Network Optimization

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Abstract

None

1 Network Traffic Problem

We consider the problem of maximizing network traffic. A network is a graph $G = (V, E)$ with vertices $v \in V$ and edges $e = (i, j) \in E = V \times V$. Each edge e has an associated capacity limit c_e. We get a set of requests, with each request r represented as source and target pairs (s, t),[1] plus a traffic demand d_r. We would like to fulfill the demand for each request as much as possible, by splitting traffic across the K_r valid paths $p_{rk} = ((s, a), (a, b), \ldots, (z, t))$ from s to t: $\sum_{p \in r} x_p \leq d_r$, where we should not exceed the request's demand (constraint [1]). We will say that $p \in r$ if p is a valid path from $s \to t$. Additionally, we must ensure that edge traffic constraints hold. Multiple paths, as well as multiple requests, may result in overlapping traffic across particular edges, resulting in constraint [2]: $\sum_{p \in \pi(e)} x_p \leq c_e$, where $\pi(e) = \{p \mid e \in p\}$.

This yields the following optimization problem:

$$
\begin{aligned}
\text{maximize} \quad & \sum_r \sum_{p \in r} x_p \\
\text{subject to} \quad & \sum_{p \in r} x_p \leq d_r, \forall r \quad [1] \\
& \sum_{p \in \pi(e)} x_p \leq c_e, \forall e \quad [2] \\
& x_p \geq 0, \forall p.
\end{aligned}
\tag{1}
$$

While this form is compact, we would like the problem to separate over paths. This is not possible in Eqn. 1 because constraint 2 couples the traffic in paths. In order to decouple this constraint, we will transform the problem by adding dummy variables. Afterwards, we will introduce slack

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[1]We assume there is only one unique request for each source and target pair.

19 variables to turn each of the resulting inequality constraints into equality constraints, which makes
20 it more straightforward to apply the ADMM algorithm.

We add one dummy variable z_{stek} for each combination of path and edge.[2] for the traffic contributed to each edge from each path p_{stk} Adding dummy variables results in the following problem:

$$
\begin{aligned}
&\text{maximize} \quad \sum_r \sum_{p \in r} x_p \qquad &&\text{minimize} \quad -\sum_p x_p + \sum_e \sum_{p \in \pi(e)} \bar{\delta}(z_{pe}) + \sum_e \bar{\delta}(s_e) + \sum_p \bar{\delta}(s_p) \\
&\text{subject to} \quad \sum_{p \in r} x_p \le d_r, \forall r \quad &&\text{subject to} \quad c_e - \sum_{p \in \pi(e)} z_{pe} - s_e = 0, \forall e \\
&\qquad\qquad \sum_{p \in \pi(e)} x_p \le c_e, \forall e \quad &&\qquad\qquad d_p - x_p - s_p = 0, \forall p \\
&\qquad\qquad x_p \ge 0, \forall p. \quad &&\qquad\qquad x_p - z_{pe} = 0, \forall e, \forall p \in \pi(e) \\
&\quad &&\qquad\qquad x, s, z \succeq 0.
\end{aligned}
$$

(2)

We use the delta notation to indicate a function

$$
\bar{\delta}(b) = \begin{cases} \infty & \text{if condition } b \text{ does not hold,} \\ 0 & \text{o.w.} \end{cases}
$$

The augmented Lagrangian for this problem is

$$
\mathcal{L}_\rho(x, z, s, \lambda) = -\sum_p x_p + \sum_e \sum_{p \in \pi(e)} \bar{\delta}(z_{pe}) + \sum_e \bar{\delta}(s_e) + \sum_p \bar{\delta}(s_p) + \lambda^\top F(x, z, s) + (\rho/2)\|F(x, z, s)\|_2^2,
$$

(3)

where

$$
F(x, z, s) = \begin{bmatrix} F_1 \\ F_2 \\ F_3 \end{bmatrix}
$$

with $F_1 \in \mathbb{R}^{|E|}$, $F_2 \in \mathbb{R}^{K|E|}$ ($|E|$ the number of edges and K the number of paths). Each subvector is given by

$$
[F_1]_e = c_e - \sum_{p \in \pi(e)} z_{pe} - s_e, [F_2]_p = d_p - x_p - s_p, [F_3]_{pe} = x_p - z_{pe},
$$

21 The elements of λ corresponding to F_1 are $\lambda_{1,e}$, F_2 are $\lambda_{2,p}$, and F_3 are $\lambda_{3,pe}$.

The ADMM updates are as follows:

$$
\begin{aligned}
x^{k+1} &:= \operatorname*{argmin}_x \mathcal{L}_\rho(x, z^k, s^k, \lambda^k) \\
z^{k+1} &:= \operatorname*{argmin}_z \mathcal{L}_\rho(x^{k+1}, z, s^k, \lambda^k) \\
s^{k+1} &:= \operatorname*{argmin}_s \mathcal{L}_\rho(x^{k+1}, z^{k+1}, s, \lambda^k) \\
\lambda^{k+1} &:= \lambda^k + \rho(F(x^{k+1}, z^{k+1}, s^{k+1})).
\end{aligned}
$$

(4)

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[2]There are K_r paths for each request r, resulting in $|E| \sum_r \sum_{K_r}$ dummy variables.

We can compute $\text{argmin}_x \mathcal{L}_\rho(x, z^k, s^k, \lambda^k)$ by restricting our attention to terms of $\mathcal{L}_\rho$ involving x and setting the derivative equal to 0:

$$0 = \nabla_{x_p}(-x_p + \lambda_{2,p}(d_p - x_p - s_p) + \sum_{e \in p} \lambda_{3,pe}(x_p - z_{pe}) + (\rho/2)((d_p - x_p - s_p)^2 + \sum_{e \in p}(x_p - z_{pe})^2))$$

$$= -1 - \lambda_{2,p} + \sum_{e \in p} \lambda_{3,pe} + \rho(-d_p + x_p + s_p) + \rho \sum_{e \in p}(x_p - z_{pe})$$

$$x_p = \max(0, \frac{1 + \lambda_{2,p} - \sum_{e \in p} \lambda_{3,pe} + \rho(d_p - s_p + \sum_{e \in p} z_{pe})}{(1 + |p|)\rho}).$$

We perform a similar computation for $\text{argmin}_z \mathcal{L}_\rho(x^{k+1}, z, s^k, \lambda^k)$:

$$0 = \nabla_{z_{pe}} \bar{\delta}(z_{pe}) + \lambda_{1,e}(c_e - \sum_{p' \in \pi(e)} z_{p'e} - s_e) + \lambda_{3,pe}(x_p - z_{pe})$$

$$+ (\rho/2)((c_e - \sum_{p' \in \pi(e)} z_{p'e} - s_e)^2 + (x_p - z_{pe})^2)$$

$$= -\lambda_{1,e} - \lambda_{3,pe} + \rho(-c_e + \sum_{p' \in \pi(e)} z_{p'e} + s_e + z_{pe} - x_p).$$

This gives us a system of equations for each edge $0 = A_e z_e + b_e$ allowing us to solve for $z_e = -A_e^{-1} b_e$. Let $P_e = |\pi(e)|$. We then have

$$A_e = \mathbf{1}_{P_e \times P_e} + I_{P_e \times P_e}$$

$$[b_e]_p = -\lambda_{1,e} - \lambda_{3,pe} + \rho(-c_e + s_e - x_p).$$

Then, for $\text{argmin}_s \mathcal{L}_\rho(x^{k+1}, z^{k+1}, s, y^k)$:

$$0 = \nabla_{s_e} \bar{\delta}(s_e) + \lambda_{1,e}(c_e - \sum_{p \in \pi(e)} z_{pe} - s_e) + (\rho/2)((c_e - \sum_{p \in \pi(e)} z_{pe} - s_e)^2)$$

$$= -\lambda_{1,e} + \rho(-c_e + \sum_{p \in \pi(e)} z_{pe} + s_e)$$

$$s_e = \max(0, \frac{\lambda_{1,e} + \rho(c_e - \sum_{p \in \pi(e)} z_{pe})}{\rho}),$$

and

$$0 = \nabla_{s_p} \bar{\delta}(s_p) + \lambda_{2,p}(d_p - x_p - s_p) + (\rho/2)((d_p - x_p - s_p)^2)$$

$$= -\lambda_{2,p} + \rho(-d_p + x_p + s_p)$$

$$s_p = \max(0, \frac{\lambda_{2,p} + \rho(d_p - x_p)}{\rho}).$$

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