

## 1) Sel Sort

From 1 to length of Array  
From  $j=i+1$  to length of Array  
if  $A[i] > A[j]$   
swap  $A[i]$  and  $A[j]$

### Invariants

- 1) The outer loop will contain sorted elements from least to greatest
- 2) The inner loop value  $A[j]$  is larger than any value in the outer loop

Why not  $n-1$

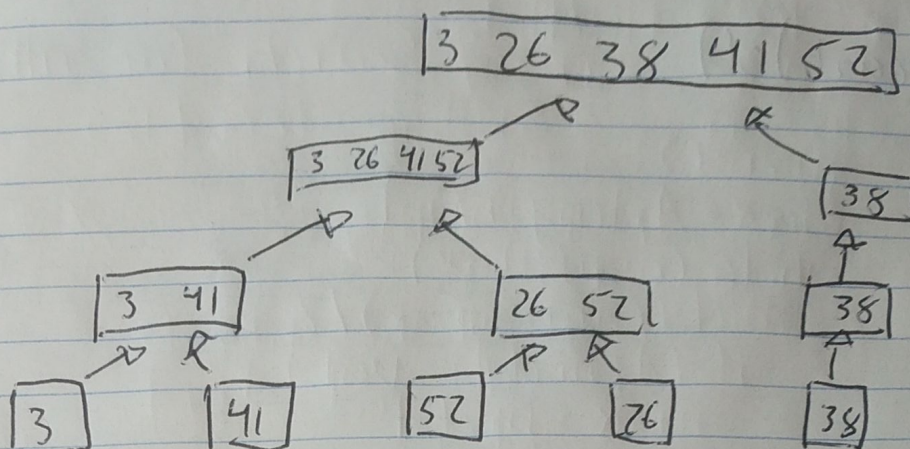
The final element will be automatically sorted because each element will be in the correct spot already

### Run times

Both best and worst case is  $\Theta(n^2)$  because both sorted and unsorted portion need to be looped through



2)



3) If there is only one element or even zero then there is nothing to ~~sort~~ sort. thus  $\Theta(1)$

The time for all elements besides  $n$  can be represented by  $T(n-1)$

The time for the  $n^{\text{th}}$  element can be represented as  $\Theta(n)$

$$T(n) = \begin{cases} \Theta(1) & n \leq 1 \\ T(n-1) + \Theta(n) & n > 1 \end{cases}$$



5 a)  $2\ 3\ 8\ 6\ 1$

j  $2\ 3\ 8\ 6\ 1$

$(2,1), (3,1), (8,6), (8,1), (6,1)$

b) The array that will have the most inversions is

$n, \dots, 3, 2, 1$

It will have  $(n-1)!$  inversions