

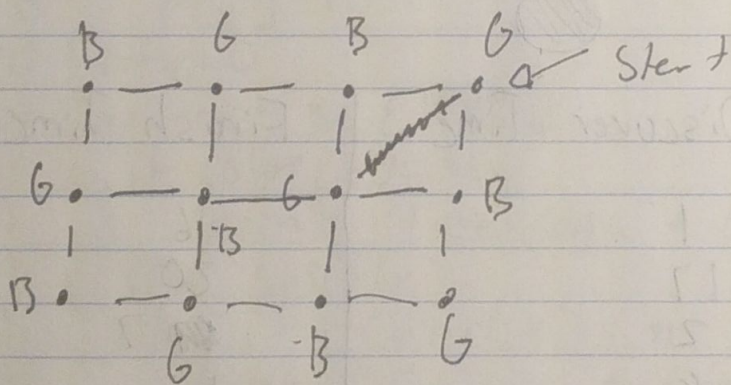
HW # 6

1) The worst case run time for BFS ~~search~~ on an adjacency matrix is:

$$\boxed{O(V^2)}$$

2) Assume graph G has n vertices and that for one vertex v its neighbors will be considered rivals.

Example



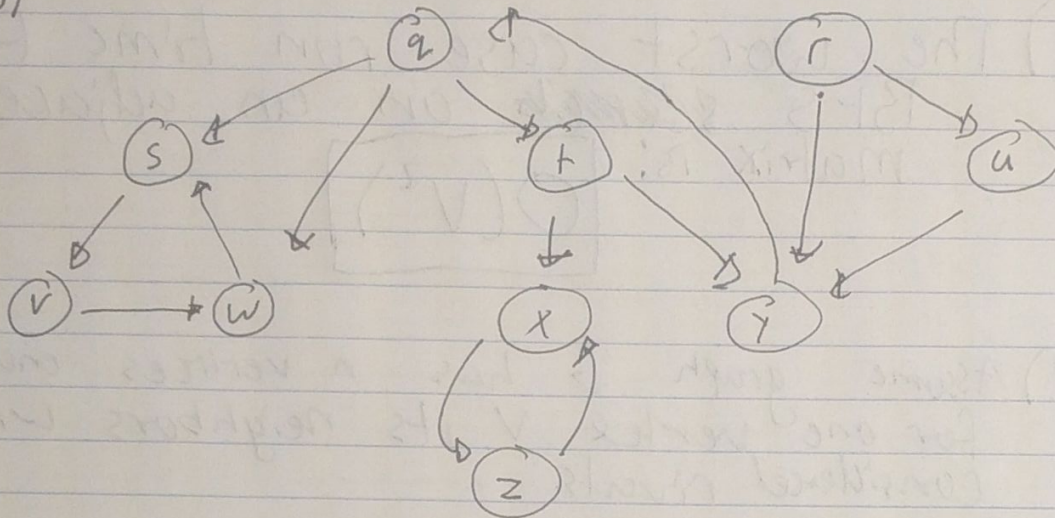
If a BFS is performed on a graph such as this it will have a time complexity: $O(n+r)$ which is desired.

$$O(n+r)$$

If a wrestler has already ~~des.~~ been designated do not designate it

If it is a rival designate neighbors as non rivals and vice versa.

3)



Vertices	Discover Time	Finish Time
q	1	16
r	17	20
s	2	17 7
t	8	15
u	18	19
v	3	6
w	4	5
x	9	12
y	13	14
z	10	11

Tree: (q, s), (s, v), (v, w), (q, t), (t, x), (x, z), (t, y), (r, u)

Forward: (q, w)

Back: (w, s), (z, x), (y, q)

Cross: (r, y), (u, y)

4)

	White		Gray		Black	
White	Tree	Back	Back	Cross	Cross	
	Forward	Cross				
Gray	Tree	Forward	Tree	Forward	Tree	Forward
			Back		Cross	
Black			Back		Tree	Forward
					Back	Cross

5) DFS

- 1) for all $x \in V$
- 2) $color[x] = w$
- 3) $P[x] = nil$
- 4) $time = 0$
- 5) for each $x \in V$
- 6) if $color[x] = w$
- 7) print "Tree edge" + coordinate
- 8) ~~else~~ visit(x)
- 9) else if $color[x] = b$
- 10) print "Forward" + coordinate
- 11) else if $d[x] > d[y]$
- 12) print "Cross" + coordinate
- 13) else
- 14) print "back" + coordinate
- 15) $color[x] = b$
- 16) $f[x] = (t + time)$

6) DFS

- 1) time = 0
- 2) $k = 1$
- 3) for all $x \in V$
- 4) $color[x] = w$
- 5) $P[x] = nil$
- 6) for all $x \in V$
- 7) if $color[x] == w$
- 8) $cc[x] = k$
- 9) $k++$
- 10) visit ~~xxxxxx~~ (x)