

## Unit 1: Introduction to ACT-R

ACT-R is a cognitive architecture. It is a theory of the structure of the brain at a level of abstraction that explains how it achieves human cognition. That theory is instantiated in the ACT-R software which allows one to create models which may be used to explain performance in a task and also to predict performance in other tasks. This tutorial will describe how to use the ACT-R software for modeling and provide some of the important details about the ACT-R theory. Detailed information on the ACT-R theory can be found in the paper “[An integrated theory of the mind](#)” and the book “How Can the Human Mind Occur in the Physical Universe?”. More information on the ACT-R software can be found in the reference manual which is included in the docs directory of the [ACT-R software](#).

### 1.1 Knowledge Representations

There are two types of knowledge representation in ACT-R -- **declarative** knowledge and **procedural** knowledge. Declarative knowledge corresponds to things we are aware we know and can usually describe to others. Examples of declarative knowledge include “George Washington was the first president of the United States” and “An atom is like the solar system”. Procedural knowledge is knowledge which we display in our behavior but which we are not conscious of. For instance, no one can describe the rules by which we speak a language and yet we do. In ACT-R, declarative knowledge is represented in structures called **chunks** and procedural knowledge is represented as rules called **productions**. Chunks and productions are the basic building blocks of an ACT-R model.

#### 1.1.1 Chunks in ACT-R

In ACT-R, elements of declarative knowledge are called **chunks**. Chunks represent knowledge that a person might be expected to have when they solve a problem. A chunk is a collection of attributes and values. The attributes of a chunk are called **slots**. Each slot in a chunk has a single value. A chunk also has a name which can be used to reference it, but that name is only a convenience for using the ACT-R software and is not considered to be a part of the chunk itself. Below are some representations of chunks that encode the facts that *the dog chased the cat* and that  $4+3=7$ . The chunks are displayed as a name followed by the slot and value pairs. The name of the first chunk is **action023** and its slots are **verb**, **agent**, and **object**, which have values of chase, dog, and cat respectively. The second chunk is named **fact3+4** and its slots are **addend1**, **addend2**, and **sum**, with values three, four, and seven.

```
Action023
  verb chase
  agent dog
  object cat
```

```
Fact3+4
  addend1 three
  addend2 four
  sum seven
```

### 1.1.2 Productions in ACT-R

A production is a statement of a particular contingency that controls behavior. They can be represented as if-then rules and some examples might be:

IF the goal is to classify a person

and he is unmarried

THEN classify him as a bachelor

IF the goal is to add two digits d1 and d2 in a column

and  $d1 + d2 = d3$

THEN create a goal to write d3 in the column

The condition of a production (the IF part) consists of a conjunction of features which must be true for the production to apply. The action of a production (the THEN part) consists of the actions the model should perform when the production is selected and used. The above are informal English specifications of productions. They give an overview of when the productions apply and what actions they should perform, but do not specify sufficient detail to actually implement a production in ACT-R.

## 1.2 The ACT-R Architecture

The ACT-R architecture consists of a set of **modules**. Each module performs a particular cognitive function and operates independently of other modules. We will introduce three modules in this unit and describe their basic operations. Later units will provide more details on the operations of these modules and also introduce other modules.

The modules communicate through an interface we call a **buffer**. Each module may have any number of buffers for communicating with other modules. A buffer relays requests to its module to perform actions, it responds to queries about the status of the module and the buffer itself, and it can hold one chunk at a time which is usually placed into the buffer as the result of an action which was requested. The chunk in a buffer is available for all modules to see and modify, and the set of chunks in all of the buffers is the information that is immediately available to the model.

### 1.2.1 Goal Module

The goal module is the simplest of the modules in ACT-R. It has one buffer named **goal** which is used to hold a chunk which contains the current control information the model needs for performing its current task. The only request to which the module responds is for the creation of

a new goal chunk. It responds to the request by immediately creating a chunk with the information contained in the request and placing it into the **goal** buffer.

### 1.2.2 Declarative Module

The declarative module stores all of the chunks which represent the declarative knowledge the model has which is often referred to as the model's declarative memory. It has one buffer named **retrieval**. The declarative module responds to requests by searching through declarative memory to find the chunk which best matches the information specified in the request and then placing that chunk into the **retrieval** buffer. In later units we will cover that process in more detail to describe how it determines the best match and how long the process takes. For the models in this unit, there will never be more than one chunk which matches the request and the time cost will be fixed in the models at 50 milliseconds per request.

The declarative memory in a model consists of the chunks which are placed there initially by the modeler when defining the model and the knowledge which it learns as it runs. The learned knowledge is collected from the buffers of all of the modules. The declarative module monitors all of the buffers, and whenever a chunk is cleared from one of them the declarative module stores that chunk for possible later use.

### 1.2.3 Procedural Module

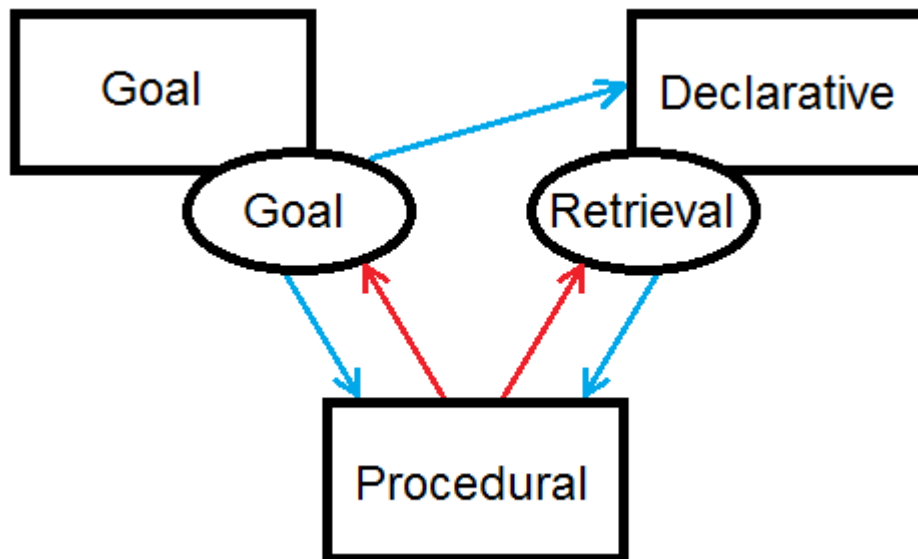
The procedural module holds all of the productions which represent the model's procedural knowledge. It does not have a buffer of its own<sup>1</sup>, and unlike other modules the procedural module does not take requests for actions. Instead, it is constantly monitoring the activity of all the buffers looking for patterns which satisfy the condition of some production. When it finds a production which has its condition met then it will execute the actions of that production, which we refer to as "firing" the production. Only one production can fire at a time and it takes 50 milliseconds from the time the production was matched to the current state until the actions happen. In later units we will look at what happens if more than one production matches at the same time, but for this unit all of the productions in the models will have their conditions specified so that at most one will match at any point in time.

### 1.2.4 Overview

These three modules are used in almost every model which is written in ACT-R, and early versions of ACT-R consisted entirely of just these three components. Here is a diagram showing how they fit together in the architecture with the rectangles representing the modules and the ovals representing the buffers:

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<sup>1</sup>Technically, there is a buffer associated with the module in the software which is useful for tracking the activities of the module, but that is only there as a convenience for the user and not a component of the ACT-R architecture itself.



The blue arrows show which modules read the information from another module's buffer, and the red arrows show which modules make requests to another module's buffer or directly modify the chunk it contains. As we introduce new modules and buffers in the tutorial, each of the buffers will have the same interface as shown for the goal buffer – both the procedural and declarative modules will read the buffer information and the procedural module will modify and make requests to the buffer.

### 1.3 ACT-R Software and Models

Now that we have described the basic components in ACT-R, we will step back and describe the ACT-R software and how one creates and runs a model using it. This tutorial is written for version 7.26 (or newer) of the ACT-R software. Current versions of the ACT-R software are distributed primarily as applications and the tutorial instructions will assume that the user is running one of those applications, but like the previous versions, the source code is available for those that prefer to run from sources. The main ACT-R software is implemented in the ANSI Common Lisp programming language, but it is not necessary to know how to program in Lisp to be able to use ACT-R because it is essentially its own language which will be described in the tutorial. In prior versions of the software it was necessary to interact with ACT-R through Lisp code, and one had to write Lisp code to build experiments or other tasks for the model to perform. However, starting with version 7.6, ACT-R provides a remote interface that can be used to interact with ACT-R from essentially any programming language, and the tutorial materials include a Python module (in the Python use of the term module not the ACT-R use as a cognitive component of the architecture) which provides functions for accessing ACT-R through that remote interface<sup>2</sup>. Using that Python module, all of the tasks for the models in the tutorial are

<sup>2</sup> The Python module included with the tutorial is sufficient for running the tasks included with the tutorial. It does not contain functions for accessing all of the commands available through the ACT-R remote interface, and it has some assumptions about how it will be used based on the needs of the tutorial. For more complex tasks or where performance is of primary importance, one might be better served by creating a custom interface instead of using the one that was built for the tutorial.

implemented in Python as well as in Lisp, and either version can be used with the models of the tutorial (the models do not depend upon which version of the task is being used). There are also examples of connecting other programming languages included with the ACT-R software, but only Lisp and Python will be used directly in the tutorial.

### 1.3.1 Starting ACT-R

To run the ACT-R software you need to run the startup script named *run-act-r* that is included with the standalone version for your operating system (versions are available for Linux, macOS, and Windows). That will open two windows for the ACT-R software. The one titled “ACT-R” is an interactive Lisp session which contains the main ACT-R system and can be used to interact with ACT-R directly. The other window, titled “Control Panel”, contains a set of GUI tools called the ACT-R Environment<sup>3</sup>. The ACT-R Environment is an optional set of tools for interacting with the ACT-R system and the use of some of those tools will be described during the tutorial. Additional information on running the software can be found in the `readme.txt` file and there are short videos on the [ACT-R software site](#) showing the typical OS protection warnings one may need to respond to the first time it is run. More information on the Environment tools can be found in the Environment’s manual included in the `docs` directory of the software. You may close the ACT-R Environment window if you do not wish to use those tools (note however that the Environment tools are necessary to perform the tutored model running exercises described in this unit and it should not be closed now). Closing the ACT-R window will exit the software.

For this unit there are no tasks for the models to interact with, and all of the interaction with the software can be done through the ACT-R Environment. For that reason, we will not describe how to use the Python module yet since there is no need for it, but it will be described in the next unit which includes interactive tasks for the models to perform.

### 1.3.2 Interacting with Lisp

Although it is not necessary to use the Lisp interface for ACT-R, it can be convenient and some familiarity with Lisp syntax can be helpful since the syntax for creating ACT-R models is based on Lisp syntax. Lisp is an interactive language and provides a prompt at which the user can issue commands and evaluate code. The prompt in the ACT-R software window is the “?” character. To evaluate (also sometimes referred to as “calling”) a command in Lisp at the prompt requires placing it between parentheses along with any parameters which it requires separated by whitespace and then hitting enter or return. As an example, to evaluate the command to add the numbers 3 and 4 requires calling the command named “+” with those parameters. That means one would type this at the prompt:

```
(+ 3 4)
```

and then hit the enter or return key. The system will then print the result of evaluating that command and display a new prompt:

---

<sup>3</sup>The ACT-R Environment is written in the Tcl/Tk language and connects through the same remote interface as the Python module.

```
? (+ 3 4)  
7  
?
```

One minor issue to note is that sometimes the output from the ACT-R system will overwrite the prompt character and it will not be visible in the ACT-R window. When that happens you can still evaluate commands by entering them at the bottom of the window and pressing enter or return.

### 1.3.3 ACT-R models

An ACT-R model is a simulated cognitive agent. A model is typically written in a text file that contains the ACT-R commands that specify how the model works, and that model file can be opened and edited in any application that can operate on text files. Because the ACT-R model syntax is based on Lisp syntax using an editor which provides additional support for Lisp formatting, like matching parentheses and automatic indenting, can be useful but is not necessary. There is a very simple text editor included with the ACT-R Environment which will do parenthesis matching, but beyond that it is a very limited text editor.

As we progress through the tutorial we will describe the ACT-R commands which one can use to create models. Later in this unit we will introduce the commands for initializing a model and creating the knowledge structures described above (chunks and productions).

### 1.3.4 Loading a model

To use an ACT-R model file it must be “loaded” into ACT-R. There are many ways to do so, but for this unit we will simply use the button in the ACT-R Environment labeled “Load ACT-R code”. In the next unit we will show how to load a model using an ACT-R command that can be called from the Lisp prompt or from an external connection, like the Python module. When the model file is loaded, the commands it contains are evaluated in order from the top down.

### 1.3.5 Running a model

Once a model has been loaded, you can run it. Models that do not interact with a task can typically be run by just calling the ACT-R **run** command. The **run** command requires one parameter, which is the maximum length of simulated time to run the model measured in seconds. Again, for this unit we will be using the tool in the ACT-R Environment instead of using the command itself. That tool is the “Run” button in the Environment and the text entry box to its right is where the time to run the model can be entered. The default time in that box is 10.0 which means pressing the “Run” button will run the model for up to 10 simulated seconds.

In later units, where the models will be interacting with various tasks, just pressing the “Run” button or calling the **run** command may not be sufficient because one might also have to run the task itself. When that is the case the tutorial will describe what is necessary to run the model and the task, and there is an additional text included with each unit which provides additional information about how the tasks are implemented and the ACT-R commands involved (those are the texts with a name that ends with “\_code”).

## 1.4 Creating an ACT-R Model

Creating an ACT-R model requires writing the text file which contains the ACT-R commands to specify the details of the model and the initial knowledge it contains. Also, in addition to the model's details, one often includes commands for controlling the general state of the ACT-R system itself.

### 1.4.1 ACT-R control commands

When creating an ACT-R model, there are two ACT-R commands for controlling the system which will almost always occur in the model file, and we will describe those commands first.

#### 1.4.1.1 *clear-all*

The **clear-all** command will usually occur at the top of every model file. This command requires no parameters and tells ACT-R that it should remove any models which currently exist and return ACT-R to its initial state. It is not necessary to call **clear-all** in a model file, but unless one is planning on running multiple models together it is strongly recommended that it occur as the first command to make sure that the model starts with the system in a properly initialized state.

#### 1.4.1.2 *define-model*

The **define-model** command is how one actually creates an ACT-R model. Within the call to **define-model** one specifies a name for the model and then includes all of the calls to ACT-R commands that will provide the initial conditions and knowledge for that model. When the model file is loaded, **define-model** will create the model with the conditions specified, and then whenever ACT-R is reset that model will be returned to that same initial state.

### 1.4.2 Chunk-Types

Before describing the commands for creating the model's initial knowledge with chunks and productions, we will first describe an additional component of the software which can be useful when creating a model. There is an optional capability available in the ACT-R software called a **chunk-type**. A chunk-type is a way for the modeler to specify categories for the knowledge in the model by indicating a set of slots which will be used together in the creation and testing of chunks. A chunk-type consists of a name for the chunk-type and a set of slot names. That chunk-type name may then be used as a declaration when creating chunks and productions in the model.

The command for creating a chunk type is called **chunk-type**. It requires a name for the new chunk-type to create and then any number of slot names. The general chunk-type specification looks like this:

```
(chunk-type type-name slot-name-1 slot-name-2 ... slot-name-n)
```

and here are some examples which could have been used in a model which created the example chunks shown earlier:

```
(chunk-type action verb agent object)
(chunk-type addition-fact addend1 addend2 sum)
```

The first creates a chunk-type named action which includes the slots verb, agent, and object. The other creates a chunk-type named addition-fact with slots addend1, addend2, and sum.

It is important to note that using a chunk-type declaration does not directly affect the operation of the model itself – the chunk-type is not a component of the ACT-R architecture. They exist in the software to help the modeler specify the model components. Creating and using meaningful chunk-types can make a model easier to read and understand. They also allow the ACT-R software to verify that the specification of chunks and productions in a model is consistent with the chunk-types that were created for that model which allows it to provide warnings when inconsistencies or problems are found relative to the chunk-types which are specified.

Although chunk-types are not required when writing an ACT-R model, most of the models in the tutorial will be written with chunk-type declarations included, and using chunk-types is strongly recommended.

### 1.4.3 Creating Chunks

The command to create a set of chunks and place those chunks into the model's declarative memory is called **add-dm**. It takes any number of chunk specifications as its arguments. As an example, we will show the chunks from the **count** model included with the tutorial that will be described in greater detail later in this unit. First, here are the chunk-type specifications used in that model:

```
(chunk-type number number next)
(chunk-type count-from start end count)
```

and here is the specification of the initial chunks which are placed into that model's declarative memory:

```
(add-dm
  (one ISA number number one next two)
  (two ISA number number two next three)
  (three ISA number number three next four)
  (four ISA number number four next five)
  (five ISA number number five)
  (first-goal ISA count-from start two end four))
```

Each chunk for **add-dm** is specified in a list – a sequence of items enclosed in parentheses. The first element of the list is the name of the chunk. The name may be anything which is not already used as the name of a chunk as long as it starts with an alphanumeric character and is a valid Lisp symbol (essentially a continuous sequence of characters which does not contain any of the symbols: period, comma, single quote, double quote, back quote, left or right parenthesis, backslash, or semicolon). In the example above the names are **one**, **two**, **three**, **four**, **five**, and **first-goal**. The purpose of the name is to provide a way for the modeler to refer to the chunk.



The name is not considered to be a part of the chunk, and it can in fact be omitted, which will result in the system automatically generating a unique name for the chunk.

The next component of the chunk specification is the optional declaration of a chunk-type to describe the chunk being created. That consists of the symbol **isa** followed by the name of a chunk-type. Note that here we have capitalized the isa symbol to help distinguish it from the actual slots of the chunk, but that is not necessary and in most cases the symbols and names used in ACT-R commands are not case sensitive.

The rest of the chunk specification is pairs of a slot name and a value for that slot. The slot-value pairs can be specified in any order and the order does not matter. When a chunk-type declaration is provided, it is not necessary to specify a value for every slot indicated in that chunk-type, but if a slot which is not specified in that chunk-type is provided ACT-R will generate a warning to indicate the potential problem to the modeler.

#### 1.4.4 Creating Productions

As indicated above, each production is a condition-action rule. Those rules are used by the procedural module to monitor the buffers of all the other modules to determine when to perform actions. The condition specifies tests for the contents of the buffers as well as the general state of the buffers and their modules. The action of a production specifies the set of operations to perform when the production is fired, and will consist of changes to be made to the chunks in buffers along with new requests to be sent to the modules.

The command for creating a production in ACT-R is called **p**, and the general format for creating a production is:

```
(p Name "optional documentation string"
  buffer tests
  ==>
  buffer changes and requests
)
```

Each production must have a unique name and may also have an optional documentation string to describe it. That is followed by the condition for the production. The *buffer tests* in the condition of the production are patterns to match against the current buffers' contents and queries of the buffers for buffer and module state information. The condition of the production is separated from the action of the production by the three character sequence ==>. The production's action consists of any buffer changes and requests which the production will make.

In separate subsections to follow we will describe the syntax involved in specifying the condition and the action of a production. In doing so we will use an example production that counts from one number to the next based on a chunk which has been retrieved from the model's declarative memory. It is similar to those used in the example models for this unit, but is slightly simpler than they are for example purposes. Here is the specification of the chunk-types used in the example production:

```
(chunk-type number number next)
(chunk-type count state current)
```

Here is the example production, which is named counting-example:

```
(P counting-example
  "example production for counting in tutorial unit 1 text"
  =goal>
    ISA      count
    state    incrementing
    current  =num1
  =retrieval>
    number   =num1
    next     =num2
==>
  =goal>
    ISA      count
    current  =num2
  +retrieval>
    ISA      number
    number   =num2
)
```

It is not necessary to space the production definition out over multiple lines or indent the components as shown above. It could be written on one line with only a single space between each symbol and still be a valid production for ACT-R. Because of that, the condition of the production is also often referred to as the left-hand side (or LHS) and the action as the right-hand side (or RHS), because of their positions relative to the ==> separator. However, since adding additional white space characters between the items does not affect the definition of a production, they are typically written spaced out over several lines to make them easier to read.

The symbols used in the production definition are also not case sensitive. The symbol ISA is only capitalized for emphasis in the example and it could have been written as isa, Isa, or any other combination of capital and lowercase letters with the same result.

#### 1.4.4.1 Production Condition: Buffer Pattern Matching

The condition of the **counting-example** production specifies a pattern to match to the **goal** buffer and a pattern to match to the **retrieval** buffer. A buffer pattern begins with a symbol that starts with the character “=” and ends with the character “>”. Between those characters is the name of the buffer to which the pattern is applied. Thus, the symbol =goal> indicates a pattern used to test the chunk in the **goal** buffer and the symbol =retrieval> indicates a pattern to test the chunk in the **retrieval** buffer. For a production’s condition to match, the first thing that must be true is that there be a chunk in each of the buffers being tested. Thus, if there is no chunk in either the

**goal** or **retrieval** buffer, often referred to as the buffer being empty or cleared, this example production cannot match.

After indicating which buffer to test, an optional declaration may be made using the symbol **isa** and the name of a chunk-type to provide a declaration of the set of slots which are being used in the test. In the example production, the **goal** buffer pattern includes a declaration that the slots being tested are from the count chunk-type, but the **retrieval** buffer pattern does not declare a type for the set of slots specified. It is recommended that one create chunk-types and use the **isa** declarations when writing productions, but this one has been omitted for demonstration purposes. The important thing to remember is that the **isa** declaration is not a part of the pattern to be tested – it is only a declaration to allow the ACT-R software to verify that the slots used in the pattern are consistent with the chunk-type indicated.

The remainder of the pattern consists of slot tests for the chunk in the specified buffer. A slot test consists of an optional modifier (which is not used in any of the tests in this example production), the name of a slot for the chunk in the buffer, and a specification of the value that slot of the chunk in the buffer must have. The value may be specific as a constant value, a variable, or the Lisp symbol **nil**.

Here is the **goal** buffer pattern from the example production again for reference:

```
=goal>
  ISA      count
  state    incrementing
  current  =num1
```

The first slot test in the pattern is for a slot named **state** with a constant value of **incrementing**. Therefore for this production to match, the chunk in the **goal** buffer must have a slot named **state** which has the value **incrementing**. The next slot test in the pattern involves the slot named **current** and a variable value.

The “=” prefix on a symbol in a production is used to indicate a variable. The name of the variable can be any symbol, and it is recommended that the variable names be chosen to help make the production easier for a person reading it to understand. Variables are used in a production to generalize the condition and action, and they have two basic purposes. In the condition, variables can be used to compare the values in different slots, for instance that they have the same value or different values, without needing to know all of the possible values those slots could have. The other purpose for a variable is to copy a value from a slot specified in the condition to another slot specified in the action of the production.

There are two properties of variables used in productions which are important. Every slot tested with a variable in the condition must exist in the chunk in the buffer for the pattern to match. Also, a variable is only meaningful within a specific production—using the same variable name in different productions does not create any relation between those productions.

Given that, we could describe this production’s test for the **goal** buffer like this:

There must be a chunk in the goal buffer. It must have a slot named state with the value incrementing, and it must have a slot named current whose value we will refer to using the variable =num1.

Now, we will look at the **retrieval** buffer's pattern in detail:

```
=retrieval>
  number      =num1
  next        =num2
```

The first slot test it has tests the slot named **number** with the variable =num1. Since the variable =num1 was also used in the **goal** buffer test, this is testing that the **number** slot of the chunk in the **retrieval** buffer has the same value as the **current** slot of the chunk in the **goal** buffer. The other slot test for the **retrieval** buffer is for the slot named **next** using a variable named =num2.

Therefore, the test for the **retrieval** buffer can be described as:

There must be a chunk in the retrieval buffer. It must have a slot named number which has the same value as the slot named current of the chunk in the goal buffer, and it must have a slot named next whose value we will refer to using the variable =num2.

This example production did not include all of the possible items which one may need in writing the condition for a production, but it does cover the basics of the pattern matching applied to chunks in buffers. We will describe how one queries the buffer and module state later in this unit, and additional production condition details will be introduced in future units.

#### 1.4.4.2 Production Action

The action of a production consists of a set of operations which affect the buffers. Here is the RHS from the example production again:

```
=goal>
  ISA      count
  current  =num2
+retrieval>
  ISA      number
  number   =num2
```

The RHS of a production is specified much like the LHS with actions to perform on particular buffers. An action for a buffer is specified by a symbol which starts with a character that indicates the action to take, followed by the name of the buffer, and then the character ">". That is followed by an optional chunk-type declaration using isa and a chunk-type name and then the specification of slots and values to detail the action to perform. There are five different

operations that can be performed with a buffer. The three most common will be described in this unit. The other operations will be described later in the tutorial.

#### 1.4.4.2.a Buffer Modifications, the = action

If the buffer name is prefixed with the character "=" then the action is for the production to immediately modify the chunk currently in that buffer. Each slot specified in a buffer modification action indicates a change to make to the chunk in the buffer. If the chunk already has such a slot its value is changed to the one specified. If the chunk does not currently have that slot then that slot is added to the chunk with the value specified.

Here is the action for the **goal** buffer from the example production:

```
=goal>
  ISA      count
  current  =num2
```

It starts with the character "=" therefore this is a modification to the buffer. It will change the value of the **current** slot of the chunk in the **goal** buffer (since we know that it has such a slot because it was tested in the condition of the production) to the value referred to by the variable **=num2** (which is the value that the **next** slot of the chunk in the **retrieval** buffer had in the condition). This is an instance of a variable being used to copy a value from one slot to another.

An important constraint on the use of the modification action is that a production can only use it for buffers that were matched to a pattern in the condition of the production – the production must test that there is a chunk in the buffer before it can modify it.

#### 1.4.4.2.b Buffer Requests, the + action

If the buffer name is prefixed with the character "+", then the action is a request to that buffer's module, and we will often refer to such an action as a "*buffer request*" where *buffer* is the name of the buffer indicated e.g. a retrieval request or a goal request. Typically a request results in the module replacing the chunk in the buffer with a different one, but not all modules handle requests the same way. As was noted above, the goal module handles requests by creating new chunks and the declarative module uses the request to find a matching chunk in the model's declarative memory to place into the buffer. In later units of the tutorial we will also describe modules that perform perceptual and motor actions in response to requests.

Here is the retrieval request from the example production:

```
+retrieval>
  ISA      number
  number   =num2
```

It is asking the declarative module to find a chunk that has a slot named **number** that has the same value as =**num2**. If such a chunk exists in the model's declarative memory it will be placed into the **retrieval** buffer.

#### 1.4.4.2.c Buffer Clearing, the - action

A third type of action that can be performed with a buffer is to explicitly clear the chunk from the buffer. This is done by placing the "-" character before the buffer name in the action. Thus, this action on the RHS of a production would clear the chunk from the **retrieval** buffer:

```
-retrieval>
```

Clearing a buffer occurs immediately and results in the buffer being empty and the declarative module storing the chunk which was in the buffer in the model's declarative memory.

#### 1.4.4.2.d Implicit Clearing

In addition to the explicit clearing action one can make, there are situations which will implicitly clear a buffer. Any buffer request with a "+" action will also cause that buffer to be cleared. Therefore, the retrieval request from the example production will also result in the **retrieval** buffer being automatically cleared at the time the request is made. We will see another situation where implicit clearing occurs later in the unit.

## 1.5 The Count Model

The first model we will run is a simple one that counts up from one number to another, for example it will count up from 2 to 4 – 2,3,4. It is included with the tutorial files for unit 1 in the file named "count.lisp". You should now start the ACT-R software if you have not done so already, and then load the count model by pressing the "Load ACT-R code" button on the ACT-R Environment Control Panel and selecting the "count.lisp" file.

When you do that, you should see a window open up which says "Successful Load", with no other text in the window. You should press the "Ok" button on that window to continue. If there had been any problems when loading the file then details about those issues would have been shown in that window.

Now you should run the model by pressing the "Run" button in the ACT-R Environment Control Panel. That will run the model for up to 10 seconds, since the default time shown next to the button is 10.0. When you do that, you should see the following output in the ACT-R window:

0.000	GOAL	SET-BUFFER-CHUNK GOAL FIRST-GOAL NIL
0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.000	PROCEDURAL	PRODUCTION-SELECTED START
0.000	PROCEDURAL	BUFFER-READ-ACTION GOAL
0.050	PROCEDURAL	PRODUCTION-FIRED START
0.050	PROCEDURAL	MOD-BUFFER-CHUNK GOAL

	0.050	PROCEDURAL	MODULE-REQUEST RETRIEVAL
	0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
	0.050	DECLARATIVE	start-retrieval
	0.050	PROCEDURAL	CONFLICT-RESOLUTION
	0.100	DECLARATIVE	RETRIEVED-CHUNK TWO
	0.100	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL TWO
	0.100	PROCEDURAL	CONFLICT-RESOLUTION
	0.100	PROCEDURAL	PRODUCTION-SELECTED INCREMENT
	0.100	PROCEDURAL	BUFFER-READ-ACTION GOAL
	0.100	PROCEDURAL	BUFFER-READ-ACTION RETRIEVAL
	0.150	PROCEDURAL	PRODUCTION-FIRED INCREMENT
TWO			
	0.150	PROCEDURAL	MOD-BUFFER-CHUNK GOAL
	0.150	PROCEDURAL	MODULE-REQUEST RETRIEVAL
	0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
	0.150	DECLARATIVE	start-retrieval
	0.150	PROCEDURAL	CONFLICT-RESOLUTION
	0.200	DECLARATIVE	RETRIEVED-CHUNK THREE
	0.200	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL THREE
	0.200	PROCEDURAL	CONFLICT-RESOLUTION
	0.200	PROCEDURAL	PRODUCTION-SELECTED INCREMENT
	0.200	PROCEDURAL	BUFFER-READ-ACTION GOAL
	0.200	PROCEDURAL	BUFFER-READ-ACTION RETRIEVAL
	0.250	PROCEDURAL	PRODUCTION-FIRED INCREMENT
THREE			
	0.250	PROCEDURAL	MOD-BUFFER-CHUNK GOAL
	0.250	PROCEDURAL	MODULE-REQUEST RETRIEVAL
	0.250	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
	0.250	DECLARATIVE	start-retrieval
	0.250	PROCEDURAL	CONFLICT-RESOLUTION
	0.300	DECLARATIVE	RETRIEVED-CHUNK FOUR
	0.300	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL FOUR
	0.300	PROCEDURAL	CONFLICT-RESOLUTION
	0.300	PROCEDURAL	PRODUCTION-SELECTED STOP
	0.300	PROCEDURAL	BUFFER-READ-ACTION GOAL
	0.300	PROCEDURAL	BUFFER-READ-ACTION RETRIEVAL
	0.350	PROCEDURAL	PRODUCTION-FIRED STOP
FOUR			
	0.350	PROCEDURAL	CLEAR-BUFFER GOAL
	0.350	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
	0.350	PROCEDURAL	CONFLICT-RESOLUTION
	0.350	-----	Stopped because no events left to process

This output is called the trace of the model. Each line of the trace represents one event that occurred during the running of the model. Each event shows the time in seconds at which it happened, the ACT-R mechanism that generated the event (typically the name of a module), and some details describing the event. Any output generated by the model is also shown in the trace. The level of detail provided in the trace can be changed to see more or less information as needed<sup>4</sup>, and this model is set to show the most information possible. That is more information than is typically needed, but some of the steps that will help to understand how the system works are only shown at this level of detail.

You should now open the count model in a text editor (if you have not already) to begin looking at how the model is specified. Since we will not be editing the file, the simple editor provided by the ACT-R Environment is sufficient to see the model definition, and you can use that by pressing the “Open File” button on the Control Panel.

<sup>4</sup> How to change the amount of detail shown in the trace is described in the unit 1 code description document (the “unit1\_code” text).

The first two commands used in the model file are **clear-all** and **define-model** as described above, and the rest of the file contains the model definition within the **define-model** call.

The first item in the model definition is a call to the **sgp** command which is used to set parameters for the model. We will not describe that here, but it is covered in the code description document. The rest of the model definition contains the chunks and productions for performing this task, and we will look at those in detail.

### 1.5.1 Chunk-types for the Count model

The model definition has two specifications for chunk-types used by this model:

```
(chunk-type number number next)
(chunk-type count-from start end count)
```

The **number** chunk-type specifies the slots that will be used to encode the ordering of numbers. It contains a slot named **number** which indicates the current number and a slot named **next** which indicates the next number in order.<sup>5</sup> The **count-from** chunk type specifies the slots that will be used for the **goal** buffer chunk of the model, and it has slots to hold the starting number, the ending number, and the current count.

### 1.5.2 Declarative Memory for the Count model

After the chunk-types we find the initial chunks placed into the declarative memory of the model using the **add-dm** command:

```
(add-dm
  (one ISA number number one next two)
  (two ISA number number two next three)
  (three ISA number number three next four)
  (four ISA number number four next five)
  (five ISA number number five)
  (first-goal ISA count-from start two end four))
```

Each of the lists in the **add-dm** command specifies one chunk. The first five define chunks named **one**, **two**, **three**, **four**, and **five**. Each chunk represents a number, and contains slots indicating the number it represents and the next number in order. This is the knowledge that enables the model to count.

The last chunk created, **first-goal**, encodes the goal of counting from two (in slot **start**) to four (in slot **end**). Note that the chunk-type **count-from** has another slot called **count** which is not used when creating the chunk **first-goal**. Because the **count** slot is not included in the definition of the chunk **first-goal** that chunk does not have a slot named **count**.

---

<sup>5</sup> There are many ways which one could represent that information using either a single chunk or spread across multiple chunks. We have chosen a single chunk representation of numbers for the tutorial to keep things simple, and we will use that same representation throughout the tutorial (except for the final model used in this unit which does not represent numbers as chunks and instead just uses digits), extending it when necessary to include more information.



### 1.5.3 Setting the Initial Goal

The next thing we see is a call to the command **goal-focus** with the chunk name first-goal:

```
(goal-focus first-goal)
```

The **goal-focus** command is provided by the goal module to allow the modeler to specify a chunk to place into the **goal** buffer when the model starts to run. Therefore, in this model the chunk named **first-goal** will be placed into the **goal** buffer when the model starts to run, and the results of that command can be seen in the first line of the trace shown above and copied here with color coding for reference:

```
0.000    GOAL      SET-BUFFER-CHUNK GOAL FIRST-GOAL NIL
```

It shows that at time 0.000 the **goal** module performed the **set-buffer-chunk** action (the ACT-R command for placing a chunk into a buffer) for the **goal** buffer with the chunk named **first-goal**. It also indicates that this action was not requested by a production (the “nil” at the end of the event details).

### 1.5.4 The Productions

The rest of the model definition is the productions which can use the chunks from declarative memory to count, and we will look at each of the production specifications in detail along with the selection and firing of those productions as shown in the trace of the model run above.

#### 1.5.4.1 The Start Production

```
(p start
  =goal>
    ISA      count-from
    start    =num1
    count    nil
  ==>
  =goal>
    ISA      count-from
    count    =num1
  +retrieval>
    ISA      number
    number   =num1
)
```

The LHS of the start production tests the **goal** buffer. It tests that there is a value in the start slot which it now references with the variable **=num1**. This is often referred to as binding the variable, as in, **=num1** is bound to the value that is in the start slot. It also tests the count slot for the value **nil**. The value **nil** is a special value that can be used when testing and setting slots. **Nil** is the Lisp symbol which represents both the boolean false and null (the empty list). Its use in a slot test is to check that the slot does not exist in the chunk. Thus, that is testing that the chunk in the **goal** buffer does not have a slot named count.

The RHS of the start production performs two actions. The first is a modification of the chunk in the **goal** buffer. That modification will add the slot named count to the chunk in the **goal** buffer (since the condition tested that it does not have such a slot) with the value that is bound to =num1. The other action is a request of the retrieval buffer. All requests to the declarative module (the retrieval buffer's module) perform a search of the chunks in declarative memory to find one that matches the information provided and then place that chunk into the **retrieval** buffer. This request is asking the declarative module to find a chunk that has the value which is bound to =num1 in its number slot.

Looking at the model trace, we see that the first production that gets selected and fired by the model is the production **start**:

0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.000	PROCEDURAL	PRODUCTION-SELECTED START
0.000	PROCEDURAL	BUFFER-READ-ACTION GOAL
0.050	PROCEDURAL	PRODUCTION-FIRED START

The first line shows that the procedural module is performing an action called conflict-resolution. That action is the process which the procedural module uses to select which of the productions that matches the current state (if any) to fire next. The next line shows that among those that matched, the **start** production was selected. The important thing to remember is that the production was selected because its condition was satisfied. It did not get selected because it was the first production listed in the model or because it has the name start. The third line shows that the selected production's condition tested the chunk in the **goal** buffer. That last line, which happens 50 milliseconds later (time 0.050), shows that the **start** production has now fired and its actions will take effect. The 50ms time is a parameter of the procedural module, and by default every production will take 50ms between the time it is selected and when it fires.

The actions of the production are seen as the next two lines of the trace:

0.050	PROCEDURAL	MOD-BUFFER-CHUNK GOAL
0.050	PROCEDURAL	MODULE-REQUEST RETRIEVAL

Mod-buffer-chunk is the action which modifies a chunk in a buffer, in this case the **goal** buffer, and module-request indicates that a request is being sent to a module through the indicated buffer.

The next line of the trace is also a result of the RHS of the **start** production:

0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
-------	------------	------------------------

That is the implicit clearing of the buffer which happens because of the request that was made.

The next line in the trace is a notification from the declarative module that it has received a request and has started the chunk retrieval process:

0.050	DECLARATIVE	start-retrieval
-------	-------------	-----------------

That is followed by the procedural system now trying to find a new production to fire:

```
0.050    PROCEDURAL          CONFLICT-RESOLUTION
```

However, it is not followed by a notification of a production being selected, because there is no production whose condition is satisfied at this time.

The following two lines show the successful completion of the retrieval request by the declarative module after another 50ms have passed and then the setting of the **retrieval** buffer with that chunk.

```
0.100    DECLARATIVE          RETRIEVED-CHUNK TWO
0.100    DECLARATIVE          SET-BUFFER-CHUNK RETRIEVAL TWO
```

Now we see conflict resolution occurring again and this time the **increment** production is selected.

```
0.100    PROCEDURAL          CONFLICT-RESOLUTION
0.100    PROCEDURAL          PRODUCTION-SELECTED INCREMENT
```

#### 1.5.4.2 The Increment Production

```
(p increment
=goal>
  ISA          count-from
  count        =num1
  - end        =num1
=retrieval>
  ISA          number
  number       =num1
  next         =num2
==>
=goal>
  ISA          count-from
  count        =num2
+retrieval>
  ISA          number
  number       =num2
!output!      (=num1)
)
```

On the LHS of this production we see that it tests both the **goal** and **retrieval** buffers. In the test of the **goal** buffer it uses a modifier in the testing of the **end** slot:

```
=goal>
  ISA          count-from
  count        =num1
  - end        =num1
```

The “-” in front of the slot is the negative test modifier. It means that the following slot test must **not** be true for the test to be successful. The test is that the **end** slot of the chunk in the **goal**

buffer have the value bound to the **=num1** variable (which is the value from the **count** slot of the chunk in the buffer). Thus, this test is true if the **end** slot of the chunk in the **goal** buffer does **not** have the same value as the **count** slot since they are tested with the same variable **=num1**. Note that the negation of the **end** slot test would also be true if the chunk in the **goal** buffer did not have a slot named **end** because a test for a slot value in a slot which does not exist is false, and thus the negation of that would be true.

The **retrieval** buffer test checks that there is a chunk in the **retrieval** buffer which has a value in its **number** slot that matches the current **count** slot from the **goal** buffer chunk and binds the variable **=num2** to the value of its **next** slot:

```
=retrieval>
  ISA      number
  number   =num1
  next     =num2
```

We can see that these two buffers were tested by the production in the next two lines of the trace:

```
0.100  PROCEDURAL      BUFFER-READ-ACTION GOAL
0.100  PROCEDURAL      BUFFER-READ-ACTION RETRIEVAL
```

Now we will look at the RHS of this production:

```
=goal>
  ISA      count-from
  count    =num2
+retrieval>
  ISA      number
  number   =num2
!output!   (=num1)
```

The first two actions are very similar to those in the **start** production. It modifies the chunk in the **goal** buffer to change the value of the **count** slot to the next number, which is the value of the **=num2** variable, and it makes a retrieval request to get the chunk representing that next number. The third action is a special command that can be used in the actions of a production:

```
!output!   (=num1)
```

**!output!** (pronounced bang-output-bang) can be used on the RHS of a production to display information in the trace. It may be followed by a single item or a list of items, and the given item(s) will be printed in the trace when the production fires. In this production it is used to display the numbers as the model counts. The results of the **!output!** in the first firing of **increment** can be seen in the trace after the production fires:

```
0.150  PROCEDURAL      PRODUCTION-FIRED INCREMENT
TWO
```

The output is displayed on one line in the trace after the notice that the production has fired. The items in the list to output can be variables as is the case here (=num1), constant items like (stopping), or a combination of the two e.g. (the number is =num). When a variable is included in the items to output the value to which it was bound in the production will be used in the output which is displayed. That is why the trace shows TWO instead of =num1.

The next few lines of the trace show the actions initiated by the **increment** production and they look very much like the actions that the **start** production generated. The **goal** buffer is modified, a retrieval request is made, a chunk is retrieved, and then that chunk is placed into the **retrieval** buffer:

0.150	PROCEDURAL	MOD-BUFFER-CHUNK GOAL
0.150	PROCEDURAL	MODULE-REQUEST RETRIEVAL
0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.150	DECLARATIVE	start-retrieval
0.150	PROCEDURAL	CONFLICT-RESOLUTION
0.200	DECLARATIVE	RETRIEVED-CHUNK THREE
0.200	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL THREE

We then see that the **increment** production is selected again:

0.200	PROCEDURAL	CONFLICT-RESOLUTION
0.200	PROCEDURAL	PRODUCTION-SELECTED INCREMENT

It will continue to be selected and fired until the value of the **count** and **end** slots of the goal chunk are the same, at which time its test of the **goal** buffer will fail. We see that it fires twice in the trace and then a different production is selected at .3 seconds into the run:

0.300	PROCEDURAL	CONFLICT-RESOLUTION
0.300	PROCEDURAL	PRODUCTION-SELECTED STOP

### 1.5.4.3 The Stop Production

```
(p stop
  =goal>
    ISA      count-from
    count    =num
    end      =num
  =retrieval>
    ISA      number
    number   =num
  ==>
  -goal>
  !output!   (=num)
)
```

The stop production matches when the values of the **count** and **end** slots of the chunk in the **goal** buffer are the same and they also match the value in the **number** slot of the chunk in the

**retrieval** buffer. The action it takes is to again print out the current number and to also clear the chunk from the **goal** buffer:

```

0.350    PROCEDURAL          PRODUCTION-FIRED STOP
FOUR
0.350    PROCEDURAL          CLEAR-BUFFER GOAL

```

The final event that happens in the run is conflict-resolution by the procedural module. No productions are found to match and no other events of any module are pending at this time (for instance a retrieval request being completed) so there is nothing more for the model to do and it stops running.

```

0.350    PROCEDURAL          CONFLICT-RESOLUTION
0.350    -----          Stopped because no events left to process

```

## 1.6 Basic Model Operations Exercise

To help you understand how an ACT-R model runs, this section contains an exercise in which you will be asked to perform the operations of the procedural and declarative modules. You will be looking at the contents of the buffers to perform the conflict resolution process (determining which productions match the current state, if any), assigning the variables in the selected production to the values they have from the buffer chunks (called instantiating the production), and determining which chunks in the model's declarative memory, if any, match the **retrieval** buffer requests that the productions makes. This will require using two of the tools in the Control Panel while running the count model described in the previous section.

### 1.6.1 Load or reset the model if necessary

If you have not yet loaded the count model then you need to press the “Load ACT-R code” button on the ACT-R Environment Control Panel and select the “count.lisp” file in the tutorial/unit1 directory. If you have already loaded the model and run it, then you need to reset it to its initial conditions so that you can run it again. To reset the model you should press the “Reset” button on the Control Panel.

### 1.6.2 The Stepper

Press the “Stepper” button on the Control Panel to open the Stepper tool. When the model is run, for each event that shows as a line in the trace, the stepper will pause the model to wait for your confirmation before handling that event. That provides you with the opportunity to inspect all of the components of the model as it progresses and can be a very useful tool when developing and debugging models. For some events, the stepper will also show additional information after it happens. The production-selected and production-fired events will show the text of the production, the bindings for the variables as they matched that production, and a set of parameters for that production (which we will describe later in the tutorial). After a retrieval request completes, the Stepper will show the details of the chunk it retrieved and the corresponding parameters that lead to that chunk being the one retrieved (more on that in a later unit).

For this exercise, you should click the “Tutor Mode” checkbox at the top of the Stepper window. That will enable the additional interactions which you will be asked to perform for the conflict-resolution, production-selected, and start-retrieval events.

### 1.6.3 The Buffers tool

Press the “Buffers” button in the Control Panel to bring up a new buffer inspection window. That will display a list of all the buffers for the existing modules. Selecting one from the list will display the chunk that is in that buffer in the text window to the right of the list by default. If you press the “Status” button the display will switch to show the information about queries for the buffer, and pressing the “Contents” button will switch it back to showing the chunk. At this point all of the buffers are empty, but that will change as the model runs.

### 1.6.4 Running with the Stepper

Now you should run the model by pressing the “Run” button in the Control Panel. When you do that you will notice that unlike before nothing shows up in the ACT-R window. Instead, the first action that will occur, the setting of the chunk in the **goal** buffer, now shows at the top of the Stepper window after “Next Step:”. That is the next action which will occur, but it is delayed until you press the “Step” button in the Stepper window to allow that action to occur. You should press it now, and then you will see that line printed out in the model trace. The Stepper now shows that event as “Last Stepped:” and a new event, conflict-resolution, is shown as the next step. In the Buffers tool window you can now see that there is a chunk in the **goal** buffer.

### 1.6.5 The goal-chunk0 Chunk and Buffer Chunk Copying

The chunk in the **goal** buffer looks like this:

```
GOAL : GOAL-CHUNK0 [FIRST-GOAL]
GOAL-CHUNK0
  START TWO
  END FOUR
```

The goal-focus command in the model made the goal module set the **goal** buffer to the chunk named **first-goal**. The trace indicated that the chunk **first-goal** was used to set the buffer, but this output is displaying that a chunk named **goal-chunk0** is currently in the **goal** buffer. Why is that? When a chunk is placed into a buffer the buffer always makes its own copy of that chunk which is then placed into the buffer. The name of the chunk that was copied is shown in the buffer inspection window in square brackets after the name of the chunk actually in the buffer. The reason for copying the chunk is to prevent the model from being able to alter the original chunk – it cannot directly manipulate the information that it has learned previously, but it may use the contents of that knowledge to create new chunks.

### 1.6.6 Conflict-resolution

You should press the Step button again to allow the conflict-resolution action to occur. When you do so, the tutoring mode will open another window. In that window the conditions of the productions in the model are shown, and it is your task to pick which, if any, match the current state. To determine that you will need to use the Buffers tool to look at the chunks in the buffers

and compare them to the patterns that are specified in the productions. Once you have chosen the item or items that you think match the current state you should press the “Check” button to have your choices compared to what the procedural model did during conflict resolution. If all of your choices were correct the window will close and you can continue stepping through the run. If any of your choices were incorrect then the conflict resolution window will show those incorrect choices, and if one of the choices does not match there will be a short description of the reason that it does not match. You will then need to press the “Ok” button to close the window and continue.

### 1.6.7 Production Selection

The next step shown is now production-selected, and you should press the Step button to allow that to happen. When a production-selected event happens the Stepper shows additional information about the selected production, which in this case, is the **start** production. In tutor mode, the production is shown in the bottom right pane of the Stepper window with all of the variables highlighted. Your task is to replace all of the variables with the values to which they are bound in the matching of this production. When you click on a highlighted variable in the production display of the Stepper a new window will open in which you can enter the value for that variable. You must enter the value for every variable in the production (including multiple occurrences of the same variable) before it will allow you to progress to the next event.

You will again need the information from the Buffers tool to be able to instantiate the production. One thing to remember is that a variable which names the buffer, for example the variable **=goal** in the **=goal>** condition, will always bind to the name of the chunk in the corresponding buffer. The bindings for the other variables will correspond to the value of the indicated slot of the tested buffer. A variable will have the same value everywhere in a production that has been selected<sup>6</sup>, and the bound values are displayed in the Stepper window as you enter them. Thus once you find the binding for a variable you can just refer to the Stepper to get the value for other occurrences of the variable in that production.

At any point in time, you can ask the tutor for help in binding a variable by hitting either the Hint or Help button of the entry dialog. A hint will instruct you on where to find the correct answer and help will just give you the correct answer.

Once the production is completely instantiated, you can continue stepping through the model run with the Step button. You should step the model to start-retrieval event which is the next part of the tutored exercise.

### 1.6.8 Declarative retrieval

The start-retrieval event is an indication from the declarative memory module that it has retrieved a request to find a chunk in memory. If you press the Step button now it will open another window which is similar to the conflict resolution selection window. In this window you are shown the retrieval request which the production made and all of the chunks in the model’s declarative memory. Your task is to select the chunks which match the request that was made.

---

<sup>6</sup> At least for the models used through most of the tutorial, but in unit 8 we will see situations where that may not always be true.



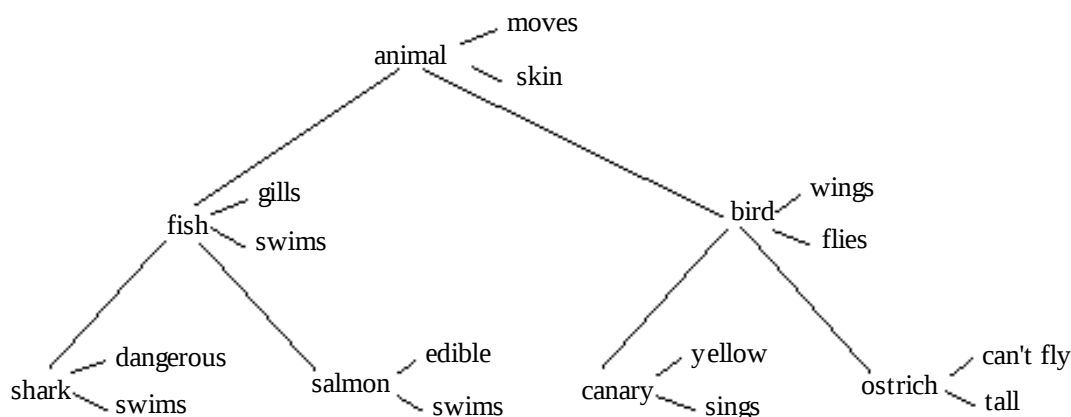
Once you have made your choices you should press the “Check” button. If you correctly chose the chunk or chunks which matched the request the window will close and you can continue stepping through the model. If you made any incorrect selections then those items will be shown and you will need to press the “Ok” button to continue.

You should now continue to step through the model performing the conflict resolution, production instantiation, and declarative memory matching exercises until the model finishes the task.

## 1.7 The Semantic Model

The next example model for this unit is the **semantic** model, found in the “semantic.lisp” file of tutorial unit 1. You should load that model file now. When you do so, you may notice that there is some text displayed in the window that indicates it loaded successfully and that text is also displayed in the ACT-R window. For now, you can ignore that output and just press “Ok” as you did for the previous model. We will come back to that after describing the operation of the model.

This model contains chunks which encode the following network of categories and properties.



The productions it has are capable of searching this network to make decisions about whether one category is a member of another category, for example, is a canary an animal or is a shark a bird.

### 1.7.1 Encoding of the Semantic Network

All of the links in this network are encoded by chunks with the slots **object**, **attribute**, and **value**. For instance, the following three chunks encode the links involving shark:

```
(p1 ISA property object shark attribute dangerous value true)
(p2 ISA property object shark attribute locomotion value swimming)
(p3 ISA property object shark attribute category value fish)
```

**p1** encodes that a shark is **dangerous** (set in the attribute slot) by setting the value slot to **true**. **p2** encodes that a shark can swim by setting the value slot to **swimming** with the attribute slot set as **locomotion**. **p3** encodes that a shark is a fish by setting **fish** as the value and the attribute as **category**.

There are of course many other ways one could encode this information, for example instead of using slots named attribute and value it seems like one could just use the slot as the attribute and the value as its value for example:

```
(p1 object shark dangerous true)
(p2 object shark locomotion swimming)
(p3 object shark category fish)
```

or perhaps even collapsing all of that information into a single chunk describing a shark:

```
(shark dangerous true locomotion swimming category fish)
```

How one chooses to organize the knowledge in a model can have an effect on the results, and that can be a very important aspect of the modeling task. For example, choosing to encode the properties in separate chunks vs a single chunk will affect how that information is reinforced (all the items together vs each individually) as well as how it may be affected by the spreading of activation for related information (both topics which will be discussed in later units). Typically, there is no one “right way” to write a model and one should make the choices necessary based on the objectives of the modeling effort and any related research which provides guidance.

For this model we have chosen the representation for practical reasons – it provides a useful example. It might seem that the second representation would still be better for that, and such a representation could have been used for the category searching model described below since we are only searching for the category attribute which could have been encoded directly into the productions. However, with what has been discussed so far in the tutorial, it would be difficult to use that representation to perform a more general search for information in the network. When we get to unit 8 of the tutorial we will see some additional aspects of the pattern matching in productions which could be used with such a representation to make the searching more general.

### 1.7.2 Testing for Category Membership

This model performs a test for category membership when a chunk in the **goal** buffer has object and category slot values and no slot named judgment. There are 3 different starting goals provided in the initial chunks for the model. The one initially placed in the **goal** buffer is **g1**:

```
(g1 ISA is-member object canary category bird)
```

which represents a test to determine if a canary is a bird. That chunk does not have a judgment slot, and the model will add that slot to indicate a result of yes or no. If you run the model with

**g1** in the **goal** buffer (which is already placed there by the call to goal-focus in the model definition) you will see the following trace:

0.000	GOAL	SET-BUFFER-CHUNK GOAL G1 NIL
0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.050	PROCEDURAL	PRODUCTION-FIRED INITIAL-RETRIEVE
0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.050	DECLARATIVE	start-retrieval
0.050	PROCEDURAL	CONFLICT-RESOLUTION
0.100	DECLARATIVE	RETRIEVED-CHUNK P14
0.100	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL P14
0.100	PROCEDURAL	CONFLICT-RESOLUTION
0.150	PROCEDURAL	PRODUCTION-FIRED DIRECT-VERIFY
0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.150	PROCEDURAL	CONFLICT-RESOLUTION
0.150	-----	Stopped because no events left to process

If you inspect the goal chunk after the model stops it will look like this:

```
GOAL: GOAL-CHUNK0
GOAL-CHUNK0
  OBJECT  CANARY
  CATEGORY BIRD
  JUDGMENT YES
```

This is among the simplest cases possible for this network of facts and only requires the retrieval of this property to determine the answer:

```
(p14 ISA property object canary attribute category value bird)
```

There are two productions involved. The first, **initial-retrieve**, requests the retrieval of categorical information and the second, **direct-verify**, uses that information and sets the **judgment** slot to **yes**:

```
(p initial-retrieve
  =goal>
    ISA      is-member
    object   =obj
    category =cat
    judgment nil
==>
  =goal>
    judgment pending
  +retrieval>
    ISA      property
    object   =obj
    attribute category
)
```

**Initial-retrieve** tests that there are object and category slots in the **goal** buffer's chunk and that it does not have a judgment slot. Its action is to modify the chunk in the **goal** buffer by adding a

judgment slot with the value pending and to request the retrieval of a chunk from declarative memory which indicates a category for the current object.

After that production fires and the chunk named p14 is retrieved the **direct-verify** production matches and fires:

```
(P direct-verify
  =goal>
    ISA      is-member
    object    =obj
    category  =cat
    judgment  pending
  =retrieval>
    ISA      property
    object    =obj
    attribute category
    value     =cat
==>
  =goal>
    judgment  yes
)
```

That production tests that the chunk in the **goal** buffer has values in the object and category slots and a judgment slot with the value pending, and that the chunk in the **retrieval** buffer has an object slot with the same value as the object slot of the chunk in the **goal** buffer, an attribute slot with the value category, and a value slot with the same value as the category slot of the chunk in the **goal** buffer. Its action is to modify the chunk in the **goal** buffer by setting the judgment slot to the value yes.

Something to notice about this production is that after it fires we see this event in the trace from the procedural module:

```
0.150    PROCEDURAL          CLEAR-BUFFER RETRIEVAL
```

This is another instance of implicit buffer clearing. If a production matches the chunk from a buffer on the LHS and does not modify it on the RHS, then for most buffers<sup>7</sup> that chunk will be automatically cleared from the buffer.

Some terminology which is often used when discussing productions which use a chunk that is in a buffer and then clear that chunk from the buffer is to say that it “harvests” the chunk – the **direct-verify** production harvests the chunk from the **retrieval** buffer. We would not say that it harvests the **goal** buffer’s chunk because it is modified and remains in the buffer. The implicit clearing of a chunk which is matched and not modified is referred to as “strict harvesting” and it is there as a convenience so that one does not need to add lots of explicit buffer clearing actions to the productions they write.

If you would like some more experience with how the ACT-R modules work, you can reset the model, open the Stepper, enable tutor mode, and run it like you did with the count model.

---

<sup>7</sup>Of the buffers used in the tutorial models, only the goal buffer is not subject to that automatic clearing mechanism.

### 1.7.3 Chaining Through Category Links

A slightly more complex case occurs when the category is not an immediate super ordinate of the indicated object and it is necessary to chain through an intermediate category. An example where this is necessary is in the verification of whether a canary is an animal, and such a test is created in chunk **g2**:

```
(g2 ISA is-member object canary category animal)
```

You should change the goal-focus call in the model file to set the goal to **g2** instead of **g1**. After you save that change to the file you must load the model again. That can be done using the “Reload” button on the Control Panel which will load the last model file that was loaded, or you can use the “Load ACT-R code” button to select the file and load it.<sup>8</sup>

Running the model with chunk **g2** as the goal will result in the following trace:

0.000	GOAL	SET-BUFFER-CHUNK GOAL G2 NIL
0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.050	PROCEDURAL	PRODUCTION-FIRED INITIAL-RETRIEVE
0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.050	DECLARATIVE	start-retrieval
0.050	PROCEDURAL	CONFLICT-RESOLUTION
0.100	DECLARATIVE	RETRIEVED-CHUNK P14
0.100	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL P14
0.100	PROCEDURAL	CONFLICT-RESOLUTION
0.150	PROCEDURAL	PRODUCTION-FIRED CHAIN-CATEGORY
0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.150	DECLARATIVE	start-retrieval
0.150	PROCEDURAL	CONFLICT-RESOLUTION
0.200	DECLARATIVE	RETRIEVED-CHUNK P20
0.200	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL P20
0.200	PROCEDURAL	CONFLICT-RESOLUTION
0.250	PROCEDURAL	PRODUCTION-FIRED DIRECT-VERIFY
0.250	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.250	PROCEDURAL	CONFLICT-RESOLUTION
0.250	-----	Stopped because no events left to process

This trace is similar to the previous one except that it involves an extra production, **chain-category**, which will attempt to retrieve the next category in the case that an attribute has been retrieved which does not immediately allow a decision to be made.

```
(P chain-category
=goal>
  ISA      is-member
  object   =obj1
  category =cat
  judgment pending
=retrieval>
  ISA      property
  object   =obj1
  attribute category
  value    =obj2
- value    =cat
```

<sup>8</sup>It is also possible to update the **goal** buffer chunk without editing the model file and loading it again, and the unit1\_code text shows how that could be done.

```

==>
=goal>
  object      =obj2
+retrieval>
  ISA        property
  object      =obj2
  attribute   category
)

```

This production tests that the chunk in the **goal** buffer has values in the object and category slots and a judgment slot with the value of pending, and that the chunk in the **retrieval** buffer has an object slot with the same value as the object slot of the chunk in the **goal** buffer, an attribute slot with the value category, a value in the value slot, and that the value in the value slot is not the same as the value of the category slot of the chunk in the **goal** buffer. Its action is to modify the chunk in the **goal** buffer by setting the object slot to the value from the value slot of the chunk in the **retrieval** buffer and to request the retrieval of a chunk from declarative memory which has that same value in its object slot and the value category in its attribute slot.

Again, for more experience, you can reset this model and step through it in tutor model.

#### 1.7.4 The Failure Case

Now you should change the initial goal to the chunk **g3** and reload the model file.

```
(g3 ISA is-member object canary category fish)
```

When you run the model with this goal, you will see what happens when the chain reaches a dead end:

0.000	GOAL	SET-BUFFER-CHUNK GOAL G3 NIL
0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.050	PROCEDURAL	PRODUCTION-FIRED INITIAL-RETRIEVE
0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.050	DECLARATIVE	start-retrieval
0.050	PROCEDURAL	CONFLICT-RESOLUTION
0.100	DECLARATIVE	RETRIEVED-CHUNK P14
0.100	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL P14
0.100	PROCEDURAL	CONFLICT-RESOLUTION
0.150	PROCEDURAL	PRODUCTION-FIRED CHAIN-CATEGORY
0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.150	DECLARATIVE	start-retrieval
0.150	PROCEDURAL	CONFLICT-RESOLUTION
0.200	DECLARATIVE	RETRIEVED-CHUNK P20
0.200	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL P20
0.200	PROCEDURAL	CONFLICT-RESOLUTION
0.250	PROCEDURAL	PRODUCTION-FIRED CHAIN-CATEGORY
0.250	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.250	DECLARATIVE	start-retrieval
0.250	PROCEDURAL	CONFLICT-RESOLUTION
0.300	DECLARATIVE	RETRIEVAL-FAILURE
0.300	PROCEDURAL	CONFLICT-RESOLUTION
0.350	PROCEDURAL	PRODUCTION-FIRED FAIL
0.350	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.350	PROCEDURAL	CONFLICT-RESOLUTION
0.350	-----	Stopped because no events left to process

There we see the declarative memory module reporting that a retrieval-failure occurred at time 0.3 which is then followed by the production **fail** firing. The **fail** production uses a test on the LHS that we have not yet seen.

### 1.7.5 A Query in the Condition

In addition to testing the chunks in the buffers as has been done in all of the productions we have seen to this point, it is also possible to query the status of the buffer itself and the module which controls it. This is done using a “?” instead of an “=” before the name of the buffer. There is a fixed set of queries which can be made of the buffer itself. The buffer’s module must respond to some specific queries, but may have any number of additional queries for its specific operations to which it will respond. The result of a query will be either true or false. If any query tested in a production has a result which is false, then the production does not match.

#### 1.7.5.1 Querying the buffer itself

When querying the buffer itself, it can be in one of three mutually exclusive situations: there can be a chunk in the buffer, a failure can be indicated, or the buffer can be empty with no failure noted. If there is a chunk in the buffer or a failure has been indicated, then it is also possible to test whether that was the result of a requested action or not. Here are examples of the possible queries which can be made to test the status of a buffer, using the **retrieval** buffer for the examples.

This query will be true if there is a chunk in the **retrieval** buffer and false if there is not:

```
?retrieval>
  buffer      full
```

This query will be true if a failure has been noted for the **retrieval** buffer and false if not:

```
?retrieval>
  buffer      failure
```

This query will be true if there is not a chunk in the **retrieval** buffer and there is not a failure indicated and false if there is either a chunk in the buffer or a failure has been indicated:

```
?retrieval>
  buffer      empty
```

This query will be true if there is either a chunk in the **retrieval** buffer or the failure condition has occurred for the **retrieval** buffer, and that chunk or failure was the result of a request made to the buffer’s module. Otherwise, it will be false:

```
?retrieval>
  buffer      requested
```

This query will be true if there is either a chunk in the **retrieval** buffer or the failure condition has occurred for the **retrieval** buffer, and that chunk or failure happened without a request being

made to the buffer's module, as was the case for the chunk placed into the **goal** buffer as a result of the goal-focus command in the model<sup>9</sup>. Otherwise, it will be false:

```
?retrieval>
  buffer      requested
```

### 1.7.5.2 Querying a buffer's module

Every module must respond to a query for its state which can be tested for being either free or busy. Most modules will respond to other queries that are specific to their operation, and those will be described in future units when needed. Below are some examples of querying the state using the **retrieval** buffer.

This query will be true if the **retrieval** buffer's module (the declarative module) is not currently performing an action and will be false if it is currently performing an action:

```
?retrieval>
  state      free
```

This query will be true if the **retrieval** buffer's module is currently performing an action and will be false if it is not currently performing an action:

```
?retrieval>
  state      busy
```

### 1.7.5.3 Using queries in productions

When specifying queries in a production multiple queries can be made of a single buffer. This query checks if the **retrieval** buffer is currently empty and that the declarative module is not currently handling a request:

```
?retrieval>
  buffer      empty
  state      free
```

One can also use the optional negation modifier “-” before a query to test that such a condition is not true. Thus, either of these tests would be true if the declarative module was not currently handling a request:

```
?retrieval>
  state      free

or

?retrieval>
  - state    busy
```

---

<sup>9</sup> Later units will describe modules which can perform operations without being requested to do so that are not just the result of commands specified by the modeler.



### 1.7.6 The fail production

Here is the production that fires in response to a category request not being found.

```
(P fail
  =goal>
    ISA      is-member
    object   =obj1
    category =cat
    judgment pending

    ?retrieval>
      buffer failure
==>
  =goal>
    judgment no
)
```

Note the query for a **retrieval** buffer failure in the condition of this production. When a retrieval request does not succeed, in this case because there is no chunk in declarative memory which matches the specification requested, the buffer will indicate that as a failure. In this model, this will happen when one gets to the top of a category hierarchy and there are no super ordinate categories.

Again, you can reset the model and work through its execution with the tutor mode of the Stepper tool. This time, you will also need to check the status of the **retrieval** buffer to determine when the fail production matches.

### 1.7.7 Model Warnings

Now that you have worked through the examples we will examine another detail which you might have noticed while working on this model. When you load or reload the **semantic** model there are the following warnings displayed in the Successful load window and the ACT-R window:

```
#|Warning: Creating chunk CATEGORY with no slots |#
#|Warning: Creating chunk PENDING with no slots |#
#|Warning: Creating chunk YES with no slots |#
#|Warning: Creating chunk NO with no slots |#
```

Output that begins with “#|Warning:” is a warning from ACT-R, and indicates that there is something in the model that may need to be addressed. This differs from warnings or errors which may be reported by the Lisp which implements the ACT-R system that can occur because of problems in the syntax or structure of the code in the model file. If you see ACT-R warnings when loading a model you should always read through them to make sure that there is not a serious problem in the model.

In this case, the warnings are telling you that the model uses chunks named **category**, **pending**, **yes**, and **no**, but does not explicitly define them and thus they are being created automatically. That is fine in this case. Those chunks are being used as explicit markers in the productions and there are no problems caused by allowing the system to create them automatically because they are arbitrary names that were used when writing the productions.

If there had been a typo in one of the productions however, for instance misspelling pending in one of them as “pneding”, the warnings may have shown something like this:

```
#|Warning: Creating chunk PNEDING with no slots |#
```

That would provide you with an opportunity to notice the discrepancy and fix the problem before running the model and then trying to determine why it did not perform as expected.

There are many other ACT-R warnings that may be displayed and you should always read the warnings which occur when loading a model to make sure that there are no serious problems before you start running the model. Most warnings which are not a serious problem can be eliminated with some additional declarations or changes to the model. For example, if we wanted to eliminate these warnings we could explicitly create those chunks in the model. With what we have seen so far that would be done by adding them to declarative memory with the other chunks that are created, but later in the tutorial we will show a better approach that does not require adding them to the model’s memory.

## 1.8 The Addition Model

There is another example model included with the unit materials which is similar to the count model which uses a larger set of counting facts to do a somewhat more complicated task. It will perform addition by counting up. Thus, given the goal to add 2 to 5 it will count 5, 6, 7, and report the answer 7. We will not cover that model in detail here. If you would like to examine and run that model, it is found in the “addition.lisp” file of unit 1, and you would use it the same way you loaded and ran the other models.

## 1.9 Building a Model

We would like you to now construct the pieces of an ACT-R model on your own. The “tutor-model.lisp” file included with unit 1 contains the basic code necessary for a model, but does not have any of the declarative or procedural elements defined. The instructions that follow will guide you through the creation of those components. You will be constructing a model that can perform the addition of two two-digit numbers using this process to add the numbers (which is of course not the only way one could implement the addition process):

To add two two-digit numbers start by adding the ones digits of the two numbers. After adding the ones digits of the numbers determine if there is a carry by checking if that result is equal to 10 plus some number. If it is, then that number is the answer for the ones column and there is a carry of 1 for the tens column. If it is not, then that result is the answer for the ones column and there is no carry. Then add the digits in the tens column. If there is no carry from the ones column then that sum is the answer for the tens column and the task is done. If there is a carry from the ones column then add the carry to that sum and make that the answer for the tens column to finish the task.

Once all of the pieces have been created as described below to implement that process, you should be able to load and run the model to produce a trace that looks like this:

```

0.000 GOAL SET-BUFFER-CHUNK GOAL GOAL NIL
0.000 PROCEDURAL CONFLICT-RESOLUTION
0.050 PROCEDURAL PRODUCTION-FIRED START-PAIR
0.050 PROCEDURAL CLEAR-BUFFER RETRIEVAL
0.050 DECLARATIVE start-retrieval
0.050 PROCEDURAL CONFLICT-RESOLUTION
0.100 DECLARATIVE RETRIEVED-CHUNK FACT67
0.100 DECLARATIVE SET-BUFFER-CHUNK RETRIEVAL FACT67
0.100 PROCEDURAL CONFLICT-RESOLUTION
0.150 PROCEDURAL PRODUCTION-FIRED ADD-ONES
0.150 PROCEDURAL CLEAR-BUFFER RETRIEVAL
0.150 DECLARATIVE start-retrieval
0.150 PROCEDURAL CONFLICT-RESOLUTION
0.200 DECLARATIVE RETRIEVED-CHUNK FACT103
0.200 DECLARATIVE SET-BUFFER-CHUNK RETRIEVAL FACT103
0.200 PROCEDURAL CONFLICT-RESOLUTION
0.250 PROCEDURAL PRODUCTION-FIRED PROCESS-CARRY
0.250 PROCEDURAL CLEAR-BUFFER RETRIEVAL
0.250 DECLARATIVE start-retrieval
0.250 PROCEDURAL CONFLICT-RESOLUTION
0.300 DECLARATIVE RETRIEVED-CHUNK FACT34
0.300 DECLARATIVE SET-BUFFER-CHUNK RETRIEVAL FACT34
0.300 PROCEDURAL CONFLICT-RESOLUTION
0.350 PROCEDURAL PRODUCTION-FIRED ADD-TENS-CARRY
0.350 PROCEDURAL CLEAR-BUFFER RETRIEVAL
0.350 DECLARATIVE start-retrieval
0.350 PROCEDURAL CONFLICT-RESOLUTION
0.400 DECLARATIVE RETRIEVED-CHUNK FACT17
0.400 DECLARATIVE SET-BUFFER-CHUNK RETRIEVAL FACT17
0.400 PROCEDURAL CONFLICT-RESOLUTION
0.450 PROCEDURAL PRODUCTION-FIRED ADD-TENS-DONE
0.450 PROCEDURAL CLEAR-BUFFER RETRIEVAL
0.450 PROCEDURAL CONFLICT-RESOLUTION
0.450 ----- Stopped because no events left to process

```

There is a working solution model included with the unit 1 files, and it is also described in the code description text for this unit. Thus, if you have problems you can consult that for help, but you should try to complete these tasks without looking first.

You should now open the `tutor-model.lisp` file in a text editor if you have not already. The following sections will describe the components that you should add to the file in the places indicated by comments in the model (the lines that begin with the semicolons) to construct a model for performing this task. There are of course many ways one could represent the addition facts and process them using productions to perform the addition, but the description below corresponds to the solution model which is provided for reference and thus should be followed explicitly if you want to be able to compare to that as you go along.

### 1.9.1 Chunk-types

The first thing we should do is define the chunk-types that we will use in writing the model. There are two chunk-types which we will define for doing this task. One to represent addition facts and one to represent the goal chunk which will hold the initial components of the task and the correct answer when it is done. These chunk types will be created with the **chunk-type** command as described in [section 1.4.2](#).

### 1.9.1.1 Addition Facts

The first chunk-type you should create is one to represent the addition facts. It should be named **addition-fact** and have slots named **addend1**, **addend2**, and **sum**. Chunks with these slots will represent the basic addition facts for addends from 0-10.

### 1.9.1.2 The Goal Chunk Type

The other chunk type you should create is one to represent the goal of adding two two-digit numbers. It should be named **add-pair**. It will have slots to encode all of the necessary components of the task. It should have two slots to represent the ones digit and the tens digit of the first number called **one1** and **ten1** respectively. It will have two more slots to hold the ones digit and the tens digit of the second number called **one2** and **ten2**, and two slots to hold the answer, called **one-ans** and **ten-ans**. It will also need a slot to hold any information necessary to process a carry from the addition in the ones column to the tens column which should be called **carry**.

## 1.9.2 Chunks

We are now going to define the chunks that will allow the model to solve the problem  $36 + 47$  and place them into the model's declarative memory. This is done using the **add-dm** command which was described in [section 1.4.3](#).

### 1.9.2.1 The Addition Facts

You need to create addition facts which encode the following math facts in the model's declarative memory to be able to solve this problem using the process described above and implemented with the productions which will be described later:

3+4=7  
6+7=13  
10+3=13  
1+7=8

They will use the slots of the type addition-fact and should be named based on their addends. For example, the fact that  $3+4=7$  should be named **fact34**. The addends and sums for these facts will be the corresponding numbers. Thus, **fact34** will have slots with the values 3, 4, and 7. Note that for simplicity we are just using the actual numbers as the values of the slots instead of creating separate number chunks like the other models in this unit.

### 1.9.2.2 The Initial Goal

You should now create a chunk named **goal** which encodes that the goal is to add  $36+47$  which will use slots from the add-pair chunk-type. This should be done by specifying the values for the ones and tens digits of the two numbers and leaving all of the other slots empty.

### 1.9.2.3 Checking the results so far

Once you have completed adding the chunk-types and chunks to the model you should be able to save the file and then load it to inspect the components you have created. To see the chunks you have created you can press the “Declarative” button on the Control Panel. That will open a declarative memory viewer which allows you to see all of the chunks in the model’s declarative memory (every time you press that button a new declarative viewer window will be opened so that you can view different chunks at the same time when needed). To view a particular chunk with that tool select it in the list of chunks on the left of the window. The window on the right will then show the declarative parameters for that chunk (which will be described in later units) along with the slots and values for that chunk.

Once you are satisfied with the chunks that you have created you can move on to creating the productions which will use those chunks.

### 1.9.3 Productions

So far, we have been looking mainly at the individual productions in the models. However, a model really only functions because of the interaction of the productions it contains. Essentially, the action of one production will set the state so that the condition for another production can match, which has an action that sets the state to allow another to match, and so on. It is that sequence of productions firing, each of which performs some small step, which leads to performing the entire task, and one of the challenges of cognitive modeling is determining how to perform those small steps in a way that is both capable of performing the task and consistent with human performance on that task – just “doing the task” is not usually the objective for creating a model with a cognitive architecture like ACT-R.

Your task now is to write the ACT-R productions which perform the steps described in English above to do multi-column addition. To do that you will need to use the ACT-R **p** command to specify the productions as described in [section 1.4.4](#), and to help with that we will further detail six productions which will implement that process using the chunk-types and chunks created previously.

#### **START-PAIR**

If the goal buffer contains slots holding the ones digits of two numbers and does not have a slot for holding the ones digit of the answer then modify the goal to add the slot one-ans and set it to a value of **busy** and make a retrieval request for the chunk indicating the sum of those ones digits.

#### **ADD-ONES**

If the chunk in the goal buffer indicates that it is **busy** waiting for the answer for the addition of the ones digits and a chunk has been retrieved containing the sum of the ones digits then modify the goal chunk to set the one-ans slot to that sum and add a slot named carry with a value **busy**, and make a retrieval request for an addition fact to determine if that sum is equal to 10 plus some number.

**PROCESS-CARRY**

If the chunk in the goal buffer indicates that it is **busy** processing a carry, a chunk has been retrieved which indicates that 10 plus a number equals the value in the one-ans slot, and the digits for the tens column of the two numbers are available in the goal chunk then set the carry slot of the goal to 1, set the one-ans slot of the goal to the other addend from the chunk in the retrieval buffer, set the ten-ans slot to the value **busy**, and make a retrieval request for an addition fact of the sum of the tens column digits.

**NO-CARRY**

If the chunk in the goal buffer indicates that it is **busy** processing a carry, a failure to retrieve a chunk occurred, and the digits for the tens column of the two numbers are available in the goal chunk then remove the carry slot from the goal by setting it to **nil**, set the ten-ans slot to the value **busy**, and make a retrieval request for an addition fact of the sum of the tens column digits.

**ADD-TENS-DONE**

If the chunk in the goal buffer indicates that it is **busy** computing the sum of the tens digits, there is no carry slot in the goal chunk, and there is a chunk in the retrieval buffer with a value in the sum slot then set the ten-ans slot of the chunk in the goal buffer to that sum.

**ADD-TENS-CARRY**

If the chunk in the goal buffer indicates that it is **busy** computing the sum of the tens digits, the carry slot of the goal chunk has the value 1, and there is a chunk in the retrieval buffer with a value in the sum slot then remove the carry slot from the chunk in the goal buffer, and make a retrieval request for the addition of 1 plus the current sum from the retrieval buffer.

**1.9.4 Running the model**

When you are finished entering the productions, save your model, reload it, and then run it.

If your model is correct, then it should produce a trace that looks like the one shown above, and the correct answer should be encoded in the **ten-ans** and **one-ans** slots of the chunk in the **goal** buffer:

```
GOAL : GOAL-CHUNK0
GOAL-CHUNK0
  ONE1  6
  TEN1  3
  ONE2  7
  TEN2  4
  TEN-ANS  8
  ONE-ANS  3
```

**1.9.5 Incremental Creation of Productions**

It is also possible to write just one or two productions and test them out first before you go on to try to write the rest – to make sure that you are on the right track. For instance, this is the trace you would get after successfully writing the first two productions and then running the model:

0.000	GOAL	SET-BUFFER-CHUNK GOAL GOAL NIL
0.000	PROCEDURAL	CONFLICT-RESOLUTION
0.050	PROCEDURAL	PRODUCTION-FIRED START-PAIR
0.050	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.050	DECLARATIVE	start-retrieval
0.050	PROCEDURAL	CONFLICT-RESOLUTION
0.100	DECLARATIVE	RETRIEVED-CHUNK FACT67
0.100	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL FACT67
0.100	PROCEDURAL	CONFLICT-RESOLUTION
0.150	PROCEDURAL	PRODUCTION-FIRED ADD-ONES
0.150	PROCEDURAL	CLEAR-BUFFER RETRIEVAL
0.150	DECLARATIVE	start-retrieval
0.150	PROCEDURAL	CONFLICT-RESOLUTION
0.200	DECLARATIVE	RETRIEVED-CHUNK FACT103
0.200	DECLARATIVE	SET-BUFFER-CHUNK RETRIEVAL FACT103
0.200	PROCEDURAL	CONFLICT-RESOLUTION
0.200	-----	Stopped because no events left to process

The first production, **start-pair**, has fired and successfully requested the retrieval of the addition fact **fact67**. The next production, **add-ones**, then fires and makes a retrieval request to determine if there is a carry. That chunk is retrieved and then because there are no productions which match the current state of the system and no actions remaining to perform it stops. You may find it helpful to try out the productions occasionally as you write them to make sure that the model is working as you progress instead of writing all the productions and then trying to debug them all at once.

### 1.9.6 Debugging the Productions

In the event that your model does not run correctly you will need to determine why that is so you can fix it. One tool that can help with that is the “Why not?” button in the Procedural viewer. To use that, first press the “Procedural” button on the Control Panel to open a viewer for the productions in the model. That tool is very similar to the Declarative viewer described previously – there is a list of productions on the left and selecting one will show the parameters and text of the production on the right. Pressing the “Why not?” button at the top of the Procedural viewer window when there is a production selected in the list will open another window. If the chosen production matches the current state of the buffers then the new window will indicate that it matches and display the instantiation of that production. If the chosen production does not match the current state then the text of the production will be printed and an indication of the first mismatching item from the LHS of the production will be indicated, for example, when the start-pair production described above does not match it might say “The chunk in the GOAL buffer has the slot ONE-ANS” since that production specifies that the **goal** buffer should not have a slot named one-ans.

The Stepper is also an important tool for use when debugging a model because while the model is stopped you can inspect everything in the model using all of the other tools. For more information on writing and debugging ACT-R models you should also read the “Modeling” text which accompanies this unit. That file is the “unit1\_modeling” document, and each of the odd

numbered units in the tutorial will include a modeling text with details on potential issues one may encounter and ways to detect and fix those issues.



## References

Anderson, J. R., Bothell, D., Byrne, M. D., Douglass, S., Lebiere, C., & Qin, Y . (2004). [An integrated theory of the mind.](#) *Psychological Review* 111, (4). 1036-1060.

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