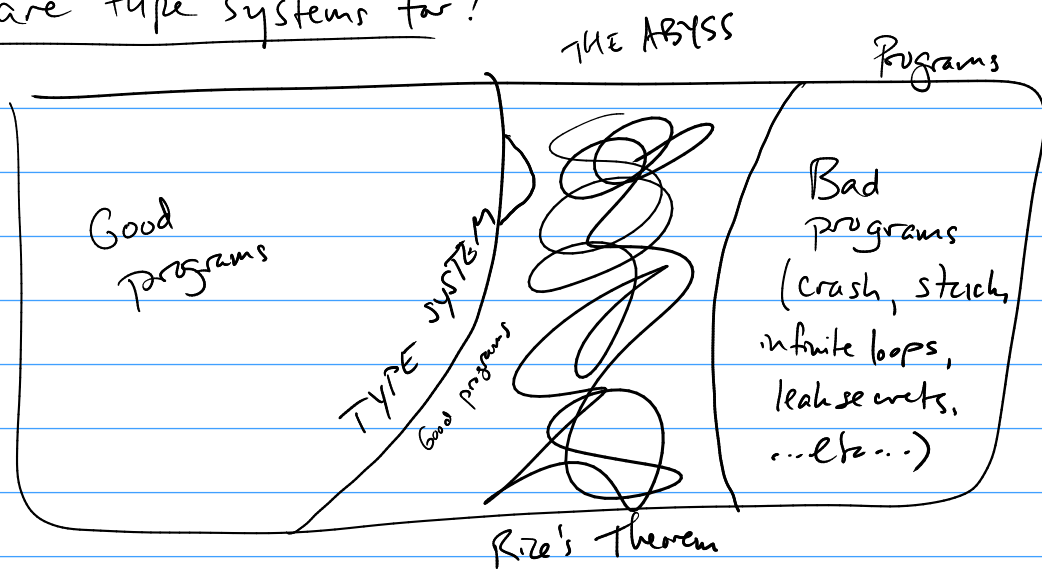


What are type systems for?



## Simply Typed $\lambda$ -calculus (STLC)

Goal: type system for  $\lambda$ -calculus that rules out nontermination (aka infinite recursion).

$t ::= x$

$| \lambda x:\alpha. t$

$| t_1 t_2$

$\alpha, \beta ::= A, B, C, \dots$

(base types)

$| \alpha \rightarrow \beta$

(function type)

$\boxed{\Gamma \vdash t : \alpha}$

$\Gamma$  = context of variable types (ie. Map Var Type)

$\frac{\Gamma \vdash t_1 : \alpha \rightarrow \beta \quad \Gamma \vdash t_2 : \alpha}{\Gamma \vdash t_1 t_2 : \beta}$

$\frac{x:\alpha \in \Gamma}{\Gamma \vdash x : \alpha}$

$\frac{\Gamma, x:\alpha \vdash t : \beta}{\Gamma \vdash \lambda x:\alpha. t : \alpha \rightarrow \beta}$

$\boxed{\Gamma \vdash t : \alpha}$

$t$  encodes a proof of the proposition  $\alpha$ .

$\frac{\Gamma \vdash t_1 : \alpha \rightarrow \beta \quad \Gamma \vdash t_2 : \alpha}{\Gamma \vdash t_1 t_2 : \beta} \rightarrow\text{-elim}$

$\frac{x:\alpha \in \Gamma}{\Gamma \vdash x : \alpha} \text{Ax}$

$\frac{\Gamma, x:\alpha \vdash t : \beta}{\Gamma \vdash \lambda x:\alpha. t : \alpha \rightarrow \beta} \rightarrow\text{-intro}$