

CS6999 Web Data and Text Mining:
Parallel Text Indexer Project Report

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1 Introduction

Fundamental to the real-world application of *Information Retrieval* techniques is the ability to objectively measure the quality of a particular algorithm or system and compare it to others. Often, information retrieval algorithms lack an analytical model that would allow one to make formal statements as to the performance characteristics of a particular system. In addition, as with machine learning and classification problems, using a scalar value to represent the performance of a system as a whole is misleading and does not account for the variety of operating conditions under which the system may be expected to perform.

In measuring the performance of an information retrieval system, the following factors constitute a set of operating conditions which qualify the result:

- Accuracy of the retrieval result with respect to the query. Common measures in this respect are *Precision* and *Recall*. Search techniques employing latent semantics of document collections are subtly more difficult to quantify, though measurable nonetheless.
- Speed with which queries are answered. This metric is a function of the number of search terms and the size of the corpus.
- Size of the document collection and, if not static, the rate at which new documents are being added and/or removed.
- The computational cost of the technique. Multiple activities may contribute to this aspect: searching, indexing, preprocessing, and post processing. It is important to consider whether the technique allows indices or intermediate approximations of the document collection to be altered incrementally or whether they must be regenerated to incorporate new documents. Also significant is whether the technique can be parallelized and thereby benefit from concurrency.
- The storage requirements of the technique. How do the storage requirements evolve with the size of the collection and/or degree of concurrency?

To drive information retrieval research in one direction or another, it would be useful to have a physical implementation framework within which one could empirically explore the operating characteristics of a particular algorithm and/or

technique along the lines described above. Such a system would allow researchers to investigate how the system responds to variation in any one of these operating parameters. The aim of this project is to create such an information retrieval *test-bed*.

The initial reference implementation of this information retrieval test-bed is based on a simple *inverted index* but has been architected such that the specific algorithm could easily be changed without significant rework. This paper presents a survey of existing literature in support of our test-bed concept, as well as the specific algorithms employed in the reference implementation. In addition, a detailed architectural overview of the system is presented, giving insight into how the system would be modified for alternative algorithms. Lastly, an evaluation of our inverted index reference implementation is conducted and the results analyzed.

2 Literature Review

2.1 Indexing

Information Retrieval is a fast becoming the dominant form of information access, overtaking traditional relational database style searching [8]. The basis of most document retrieval systems is the *term-document* incidence matrix. This matrix is typically sparse and more efficiently represented as an *inverted-index* which maps terms to the parts of a document where they occur.

In order to construct an *inverted index*, documents must be scanned to determine term frequencies. The Aho-Corasick string matching algorithm[1] is a simple and efficient text scanning algorithm. The algorithm constructs a finite state machine to scan for a given set of keywords. It is, in effect, a reduced grammar regular expression parser of the type described in [2]. The algorithm is simple and efficient, construction time being proportional to the sum of the length of the keyword set, and the number of state transitions required to scan a document is independent of the number of the size of the keyword set.

2.2 Evaluation

In the case of machine learning and classification systems, the specification of a particular set of misclassification costs or false and positive misclassification rates define a particular *operating point* of the system.

This concept applies equally to *information retrieval* systems whereby it is misleading to represent the operating characteristics of a particular system with a scalar metric. In addition, the evaluation of information retrieval systems suffers from the the inherent subjectivity with respect to *relevance* [10]. Non-scalar evaluation methods try to characterize the performance of a system over a range of operating conditions. Although this makes direct performance comparisons ambiguous, we are more likely to be able to select the algorithm which will perform best in the real world by applying domain expertise to the results of a non-scalar evaluation [3].

2.3 Parallelization

Once an objective measure has been established, the performance of an information retrieval system can be improved iteratively by making various optimizations. These optimizations may take the form of algorithm enhancements, implementation efficiencies, or parallelization.

Algorithm enhancements usually represent the forefront of research in the field. Work to validate a new concept or algorithm is usually performed in a very controlled environment and on relatively small document collections. As the field matures, improvements are harder fought and often involve increasingly complex algorithms and corpus approximation methods into which it becomes more difficult to *fold-in* new documents. In addressing the considerable computation cost of re-calculating *Latent Semantic Indexing* (LSI) [4] approximations, [12] introduced new update methods which saw improved average precisions but were forced to conclude that the improvements were at a computational premium over previous methods. A semi-discrete matrix decomposition for LSI, *SDD*, performs as well as the original SVD-based LSI method for equal query times and uses less storage [6]. The decomposition time of the new method, however, is significantly larger than the original.

Enhancement of implementation efficiency in [7] saw a six-fold timing improvement in *LSI* searches on large document collections. The same study achieved nearly 180 times improvement through the use of parallelization. The study made no algorithmic improvements and the result is consistent with increased use of parallelization by large commercial systems to cope with the ever increasing bodies of data that must be processed.

Map-Reduce, popularized by Google, is emerging as an important programming model for large-scale data-parallel applications such as web indexing [5].

Fundamentally, Map-Reduce breaks up a computation into smaller tasks which are made to run in parallel on different physical machines. The *map* and *reduce* phases are inspired by the functions of the same name used in *functional programming*. The model scales well to large clusters mostly because of the decoupled nature of the tasks themselves. An initial *map* phase breaks the data into separate chunks and assigns them to a task. The *reduce* phase assembles the results of the map tasks to produce the final result.

Hadoop is an open source implementation of the MapReduce model, developed mostly by Yahoo, and now maintained by the Apache Software Foundation. Hadoop can be configured to automatically create speculative copies of slow running map tasks (stragglers). This facility is intended as a pre-emptive failure recovery mechanism as well as a means of optimizing overall throughput. The metrics used to identify stragglers as well as the algorithm for determining the target node on which speculative copies of a task should be launched, is examined in [11]. The paper introduces an alternative scheduling algorithm, *LATE*, to address excessive task speculation issues with the native Hadoop scheduler in a heterogeneous hardware environment. Experiments comparing the native Hadoop scheduler, a non-speculative scheduling strategy, and *LATE* were conducted using hosted virtual machines (VMs) in the Amazon EC2 environment as well as a dedicated local cluster. Variance in node performance was introduced by varying the distribution of VMs across physical machines and also introducing network and disk transfer loads. Results show speculative execution (native Hadoop and *LATE*) outperforming non-speculative execution and *LATE* outperforming native Hadoop in most cases.

3 Research Plan

The primary goals of this research project is to construct an information retrieval framework within which specific implementations can be evaluated under various operating conditions. To initiate the synthesis of this general framework, a reference implementation will be built, based on the notion of a simple inverted index.

The software architecture of the reference implementation will be carefully documented and follow established design patterns in order to facilitate the future construction of alternative document search schemes within the same framework. The framework will be designed to be executed in a distributed

deployment such that it may be used to evaluate the effect of parallelization on a particular algorithm.

The performance of the reference implementation will be tested with respect to the following operating parameters and compared to analytic expectations:

- Number of index keywords
- Number of documents in collection
- Average size of documents

In addition, the reference implementation will be compared against a naïve implementation of the inverted indexer. The naïve implementation provided a baseline for further evaluation.

4 Design

A black-box functional view of the indexing system is shown in figure 1. The system processes a continuous incoming flow of documents and distributes them to parallel indexing threads running on physically separate, heterogeneous nodes. Document queries are performed using an evolving search index. The operational balance is to timely index new additions to the corpus while servicing concurrent search requests. A view as to how the components of the indexer system are deployed is shown in figure 2.

As can be seen in figure 2, the indexer processes are symmetrical and deployed on multiple physical nodes. The individual indexers do not interact directly with one another, making for a simple deployment and operation model. The execution loop of each indexer is as follows:

1. Initialize the indexer by retrieving a collection of index keywords and synonyms from a relational database. These keywords are used to construct a lexical parser which will be used to scan and index documents.
2. Retrieve a batch of unprocessed documents from a relational database. The document batch will be sized according to the physical capabilities of each node. In this implementation, this tuning task is a manual exercise but future enhancements may include an adaptive loading component.
3. Parse each document retrieved and construct an inverted index representing the batch. This *delta* index, as we shall call it, is then used to augment the global index maintained in a relational database.

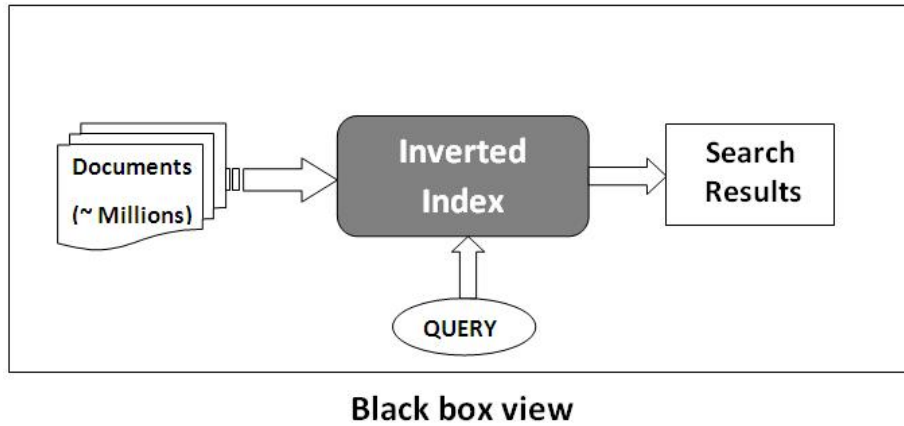


Figure 1: A black-box view of the indexer

4. Repeat from step 2.

As shown in figure 2, the searcher component can be executed from any node which has access to the relational database housing the document index. The execution path of a single search query would be as follows:

1. Canonize search terms based on keyword synonyms defined in the keyword store.
2. Execute a boolean query against the inverted index.
3. Return the list of corpus documents containing the intersection of the canons of the search terms.

A guiding architectural principal of the indexer design is to separate the implementation into three logical layers:

- **Software Layer:** implements parsing, indexing and search algorithms.
- **Data Access Layer:** implements an *object-relational mapping*, insulating algorithm implementation from the persistence details. This layer implements classes which encapsulate persistent entities within the system.
- **Database:** the persistent store for the document collection and inverted index.

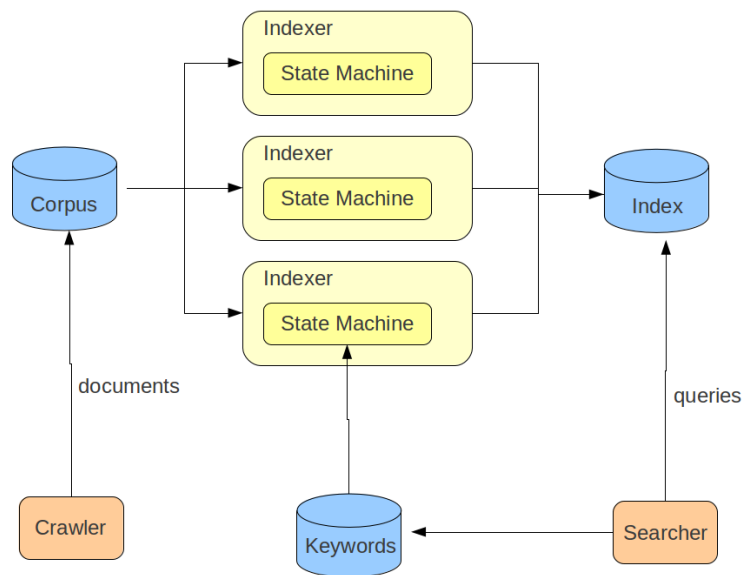


Figure 2: Indexer deployment model

The implementation classes and their relationships are represented in a UML class diagram in figure 3. Figure 4 is a UML sequence diagram showing how the layers interact at a high level. The implementation of each of these layers is described in more detail in section 5.

5 Implementation

The individual worker processes of the indexer system are implemented as Java programs. The programs make use of the following thirdparty libraries:

- **DBPool:** a database connection pooling library, used to manage connections to relational databases.
- **Apache Commons Logging:** a generic logging API used by DBPool.
- **Apache Commons CLI:** a library for parsing command line arguments passed to Java programs.
- **MySQL JDBC driver:** the MySQL Java client driver.

The indexer system maintains an *inverted index* of an evolving corpus and services concurrent search requests against the index. The current implementation is based on boolean retrieval methods described in [8]. An Aho-Corasick state machine [1] is used to scan documents for a predefined set of keywords, and construct a term index for each document (delta index). The term occurrences for each document are used to augment an existing index of documents already processed. The document collection, index, and keyword list are maintained in a single MySQL database instance. It is through this database that the parallel indexer instances are initialized and synchronized. This configuration is operationally simple and easy to implement but the database has the potential to become a throttle point as the system is scaled to larger and more rapidly evolving document collections, and is required to service a higher rate of queries. In order to scale the current implementation to large document collections, a more efficient mechanism would be required in this respect.

The following subsections detail the implementation of the architectural layers identified in section 4.

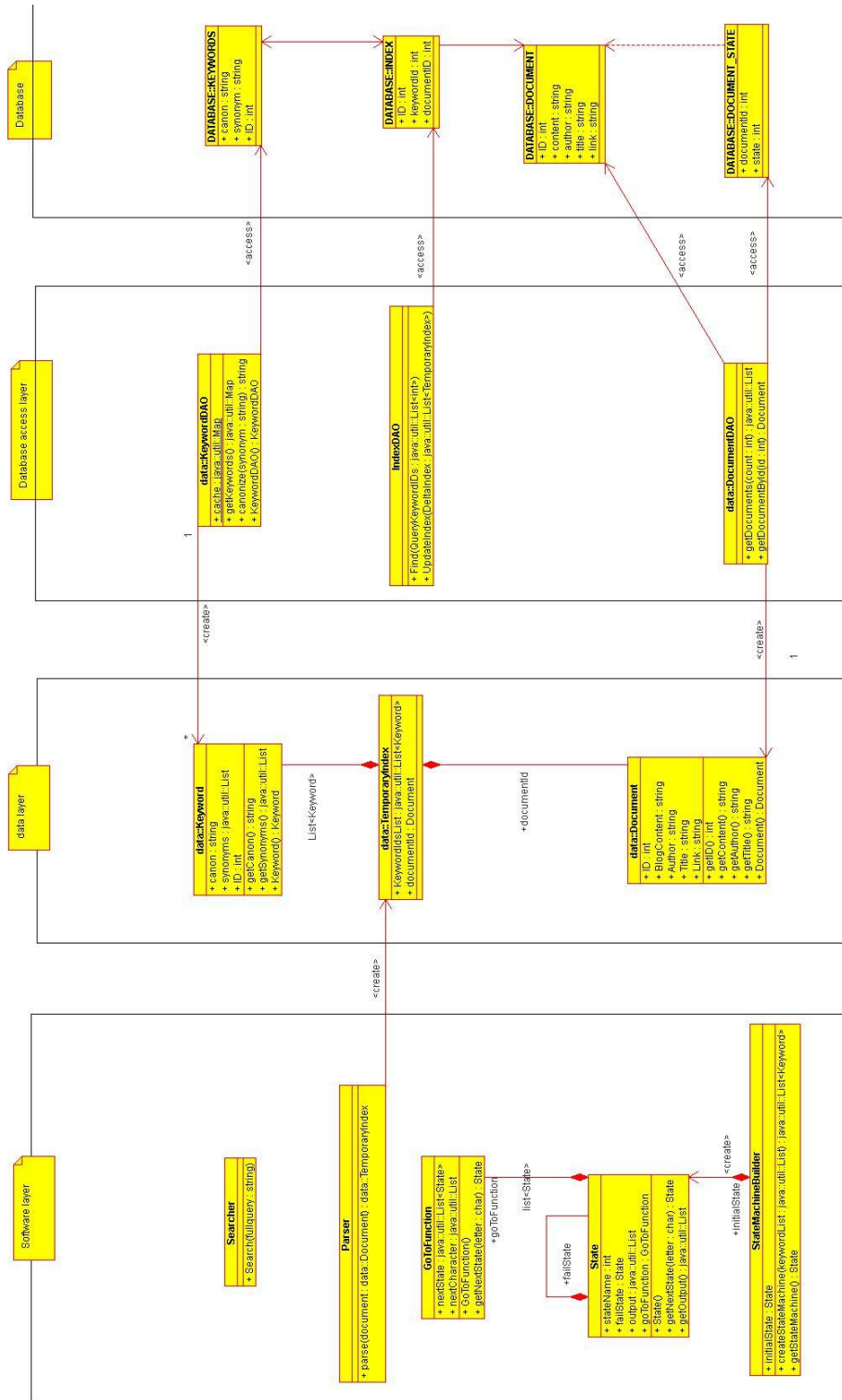


Figure 3: Overall UML class diagram design

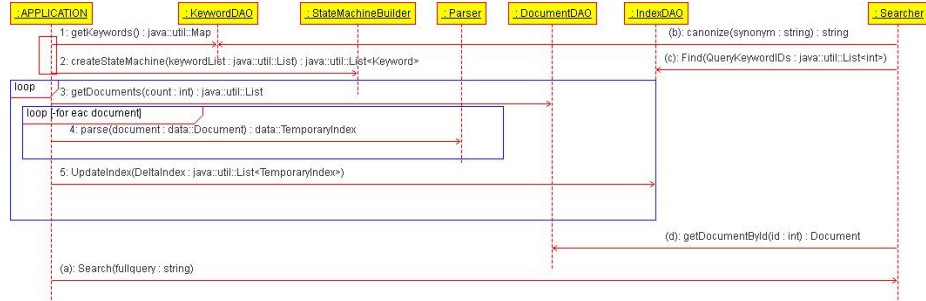


Figure 4: Overall UML sequence diagram design

5.1 Data Access Layer

Access to persistent storage is via objects of the data access layer. Confining data access to a logical grouping of classes insulates the other application components from changes in database schema. The following subsections detail the individual data entities handled by the data access layer.

5.1.1 Document

The `DocumentDAO` singleton class provides methods for retrieving `Document` objects from the document database:

- `List<Document> getDocuments (int count)`: retrieve a specified number of unprocessed documents from the database. The documents retrieved will be atomically marked as processed so that processing is not duplicated on other nodes. In this scheme, a document is considered *processed* once retrieved from the database. This may cause problems from a fault recovery perspective if the processing node fails. A proposed enhancement to this scheme is to use an intermediate state, *processing*, to indicate that the document has been retrieved and to update the state to *processed* once processing is actually complete. For the initial implementation we will use the *processed* state only.
- `Document getDocumentById (int id)`: retrieve a specific document from the database, referenced by its ID.
- `int getID ()`: accessor method.
- `String getContent ()`: accessor method.

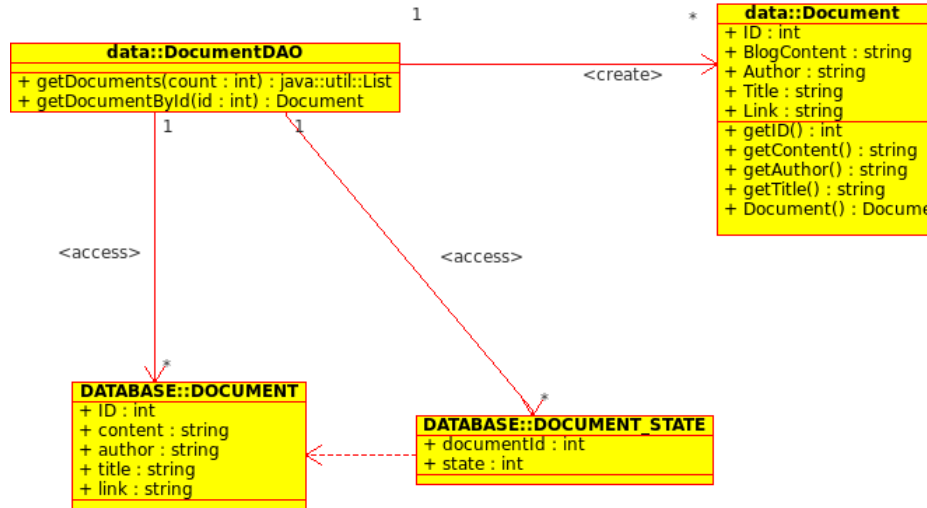


Figure 5: Document class diagram

- `String getAuthor ()`: accessor method.
- `String getTitle ()`: accessor method
- `String getLink()`: accessor method.

The `Document` class encapsulates the BLOG database table, with the addition of a state field that tracks whether a document has been processed or not. Figures 5 and 6 show class and sequence diagrams respectively for non-trivial classes and methods involved in the implementation of the Document data access layer.

5.1.2 Keyword

The `KeywordDAO` singleton class provides methods for retrieving `Keyword` objects from the document database. Because the keyword set is relatively static and accessed often, the full set of keywords will be loaded from the database and cached when the class initializes. The class also provides methods for converting search terms to their canonical form.

- `List<String> getKeywords ()`: this method will returns a list of all `Keyword` objects defined in the keyword database.

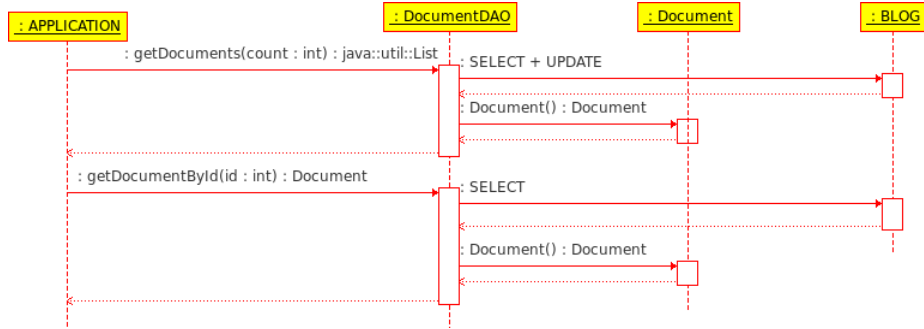


Figure 6: Document sequence diagram

- `List<Integer> canonize (String query)`: return the list of IDs corresponding to a comma-separated list of query terms.
- `Keyword getKeywordById (int id)`: return the `Keyword` object with the given identifier.

The `Keyword` field encapsulates a canonical search term and a list of its synonyms. Figures 7 and 8 show UML class and sequence diagrams respectively for non-trivial classes and methods involved in the implementation of the `Keyword` data access layer.

5.1.3 Index

The `IndexDAO` class provides methods for updating and searching the inverted index table stored in the database. There are two methods defined in this class:

- `void UpdateIndex (List<TempIndex> input)`: this method takes input of type `List<TempIndex>` (Note: Temporary index is a user defined data class, which contains Document ID and List of `keyword ids` as its attributes) and generates an SQL query which updates the `INDEX` table in the Database.
- `List<Integer> Find (List<Integer> QueryKeywordIds)`: this method takes a list of `Integers` (which corresponds to mapped `String` query into keyword ids) as input and searches the `INDEX` table in DB to find all the document ids which are common between for the given query keywords.

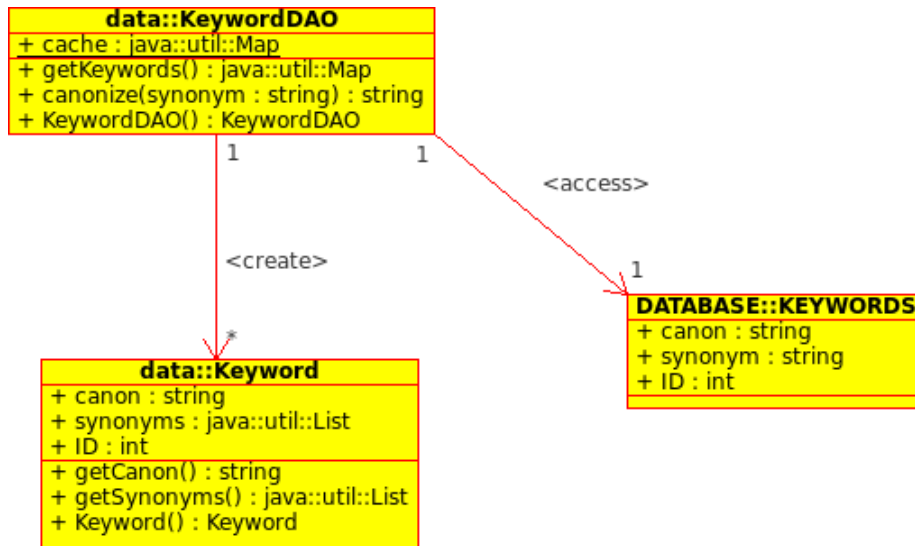


Figure 7: Keyword class diagram

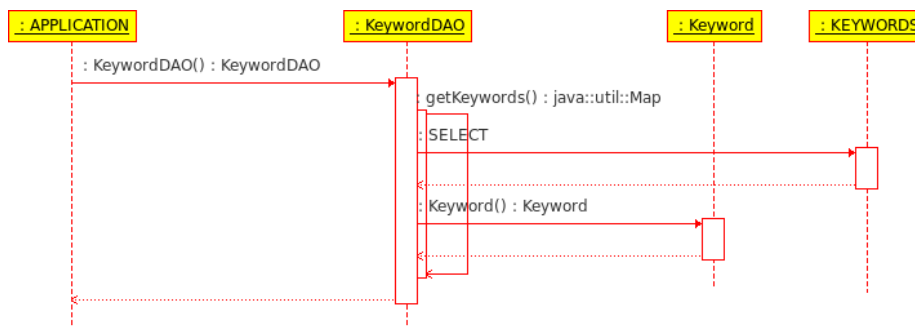


Figure 8: Keyword sequence diagram

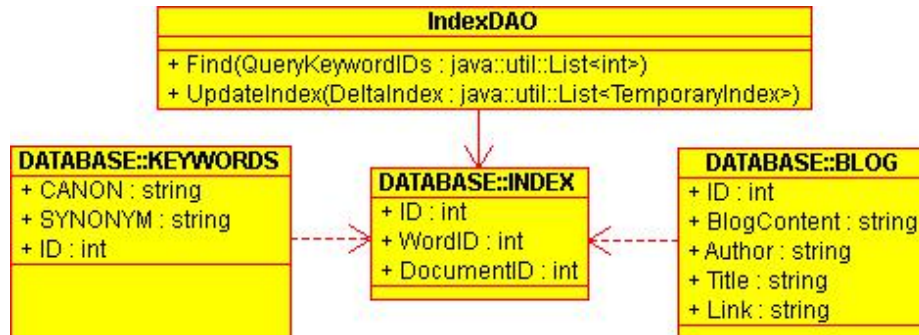


Figure 9: Index class diagram

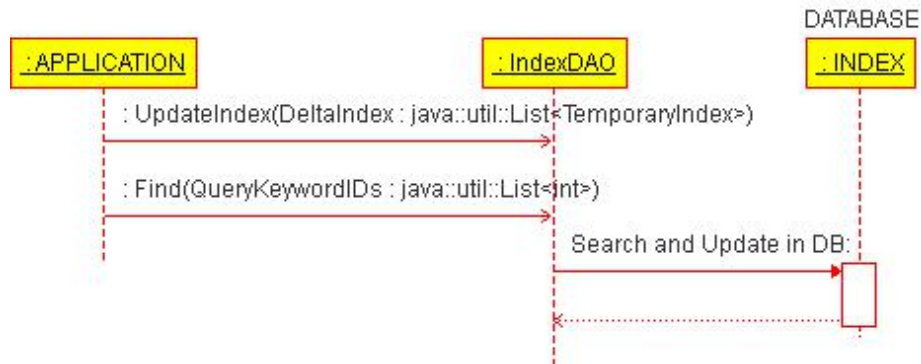


Figure 10: Index sequence diagram

Figures 9 and 10 show UML class and sequence diagrams respectively for classes and methods involved in the implementation of the Index data access layer.

5.2 Software Layer

5.2.1 Searcher

This class is designed to perform searches for a query entered by the user. It contains a method named `search ()` which is responsible for finding all the documents which contained all query keywords.

- `List<Document> search (String)`: given a comma-separated list of search terms, returns the list of documents in which the canonical form of all

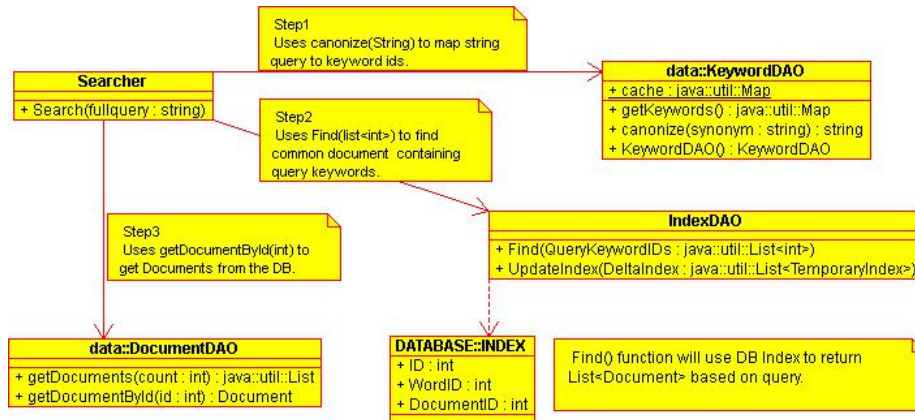


Figure 11: Searcher class diagram

search terms appear together. This method First calls `List<int> KeywordDAO.canonize (String)`, to map string query keywords into corresponding list of keyword IDs. Note: We are assuming that the query will be formed with the set of keywords, which will be a subset of the global keyword set, with which the parsing state machine is constructed.

Then it calls `List<Integer> IndexDAO.find (List<querykeywordIDs>)`, which will searches in INDEX DB for a list of Common DOC_IDs corresponding to given query keywords.

In the end, it calls `List<Document> DocumentDAO.getDocumentById (DOC_ID)`. This will map DOC_ID into a Document from the BLOG table in the database.

Figures 11 and 12 show UML class and sequence diagrams respectively for classes and methods involved in the implementation of the **Searcher** class.

5.2.2 State Machine

This class is designed to build the state machine from the list of keywords returned by the `KeywordDAO.getKeywords()` method. This class has a static attribute `initialState` that is the first state of the state machine built by the function `createStateMachine` (initializes with the value null). The first state gives us the whole state machine thanks to the architecture of a state. Once the state machine is built, the function `getStateMachine` is used to access the

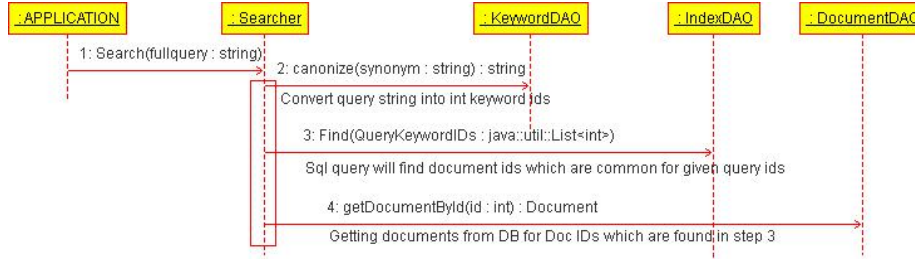


Figure 12: Searcher sequence diagram

initial state stored in the class.

Figures 13 and 14 show UML class and sequence diagrams respectively for classes and methods involved in the implementation of the **StateMachineBuilder**.

5.2.3 Parser

The **Parser** implements the document parsing algorithm within the state machine. By default, it uses the state machine contained in the static attribute of **StateMachineBuilder**. To parse a document we need to call the **parse ()** function inside this class, passing it the **Document** class encapsulating the actual text we want to parse. This class returns a **TemporaryIndex** object which encapsulates a reference to the document parsed and the keywords found to occur within the document.

Figures 15 and 16 show UML class and sequence diagrams respectively for classes and methods involved in the implementation of the **Parser**.

6 Experiment Results

The primary goal of this project is to be able to reduce the time to build a document index. As our system is building the inverted index by parsing all documents of the collection, we can expect the time complexity of the task to be proportional to the sum total of characters in the collection. As a result, changing the number of documents or the average size of each document should have a big impact in our time complexity. In addition, our implementation will incur the cost of building the Aho-Corasick state machines used to scan the documents. We can expect the time to build a state machine to be proportional to the sum of the size of the keywords.

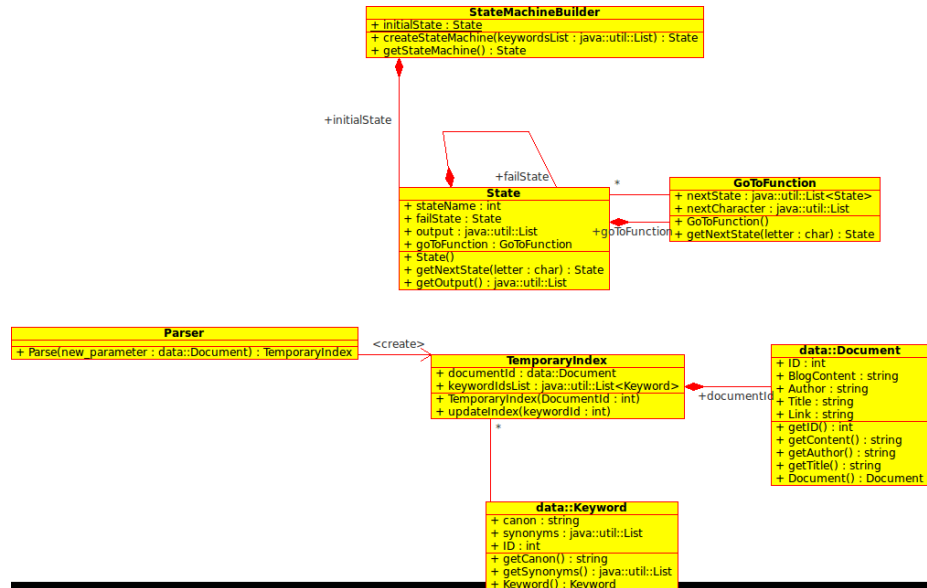


Figure 13: State machine class diagram

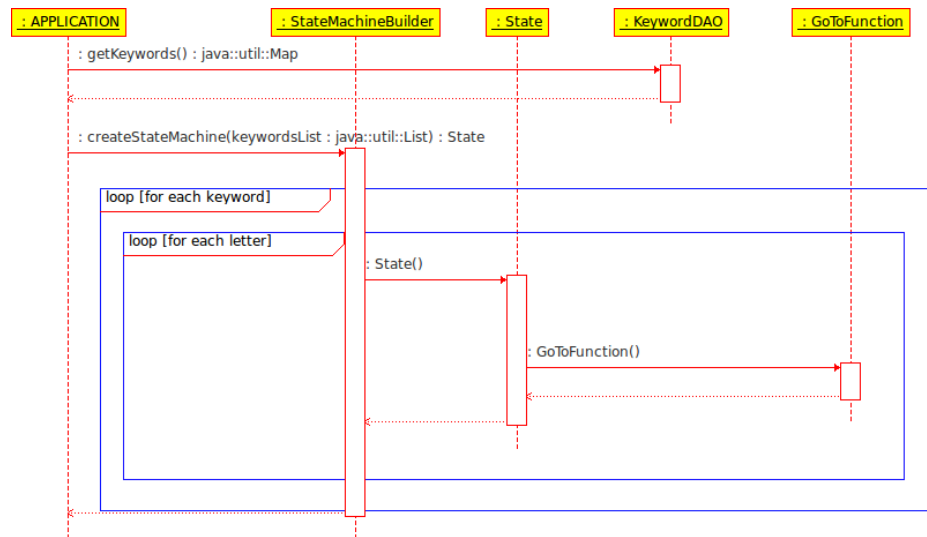


Figure 14: State machine sequence diagram

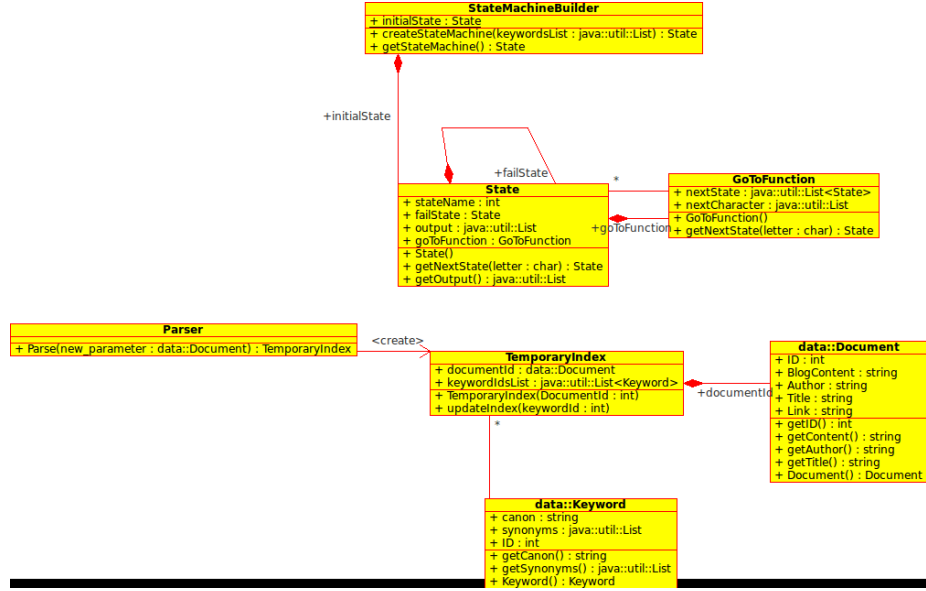


Figure 15: Parser class diagram

Finally we will do a time comparison between our indexer implementation and a naïve implementation. All the time complexities tests are done on a single computer and it is important to understand that all these processes can be parallelized. In section 6.5 we install our indexer on multiple computers and explore the performance gain of running parallel instances simultaneously.

All tests were run on a single Intel Core 2 Duo 2GHz processor, 1GB RAM, running a 64-bit Linux 2.6.35 SMP kernel. The Java Virtual Machine used was version 1.6.0-24. The database used was MySQL, running on the same node as the indexer.

6.1 Number of Keywords

To be able to study the impact of different number of keywords, we built a keywords database with 10,000 most popular English words [9]. All the experiments in this part are done with 50,000 posts randomly downloaded from the Internet. The idea is to get a good distribution of documents with respect to length and content. The experiment shown here consist of building the index for different numbers of keywords from the 10,000 keyword database and explore the time and space complexity for the building of the state machine and the

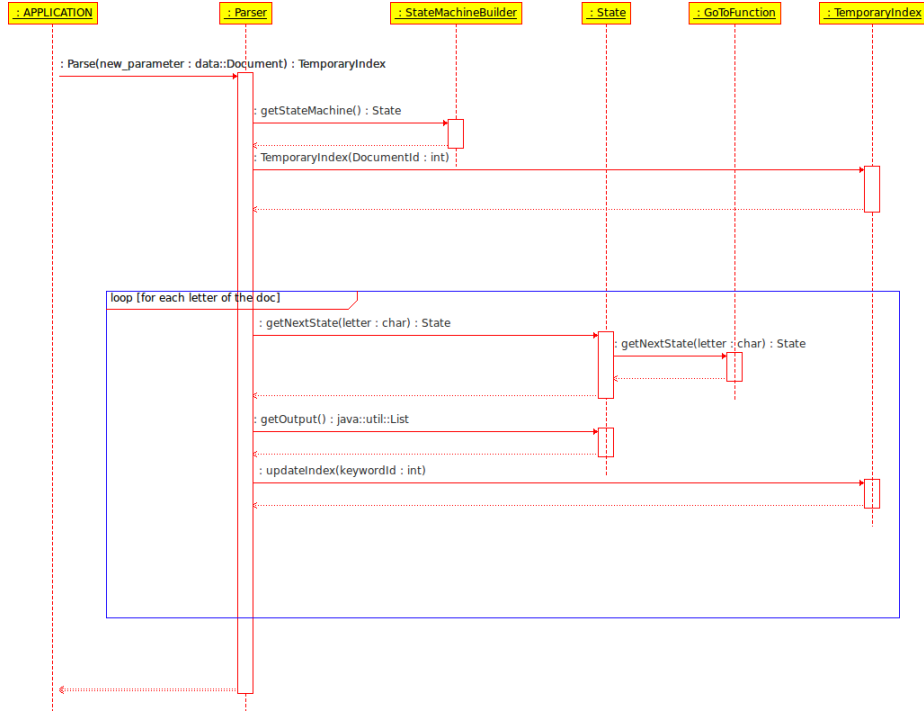


Figure 16: Parser sequence diagram

building of the index using the state machine.

6.1.1 Time Complexity

- State Machine Construction:** The time taken to construct the Aho-Corasick state machine was measured for different numbers of keywords. The results of this test is shown in figure 17 and construction time can easily be considered linear with respect to the number of keywords. This is consistent with [1] which proves that the state machine construction algorithm is linearly proportional to the sum of the lengths of the keywords used to construct the state machine.
- Index Construction:** The time taken for a state machines to build an inverted index of our corpus was measured. The test was repeated using state machines constructed with a varying number of keywords and the results plotted in figure 18. According to [1] which proves that the number

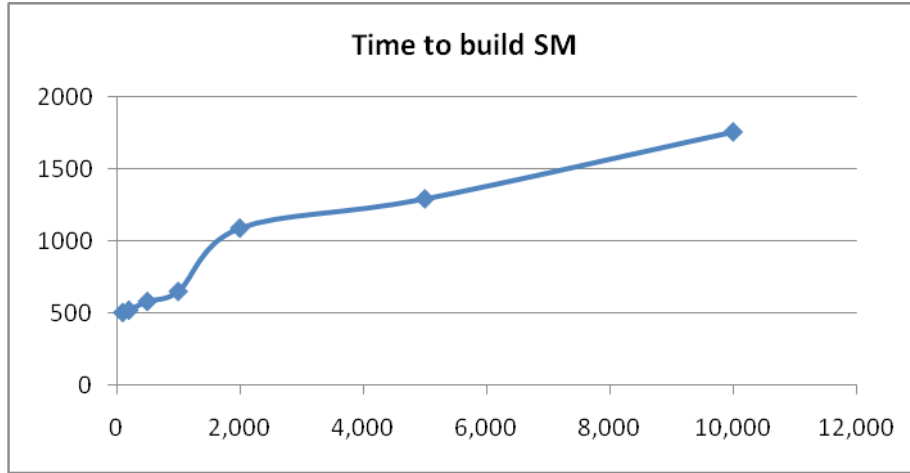


Figure 17: Time to build Aho-Corasick state machine as a function of the number of keywords. Time taken is shown in milliseconds on the y-axis and the number of keywords on the x-axis

of state transitions involved in processing an input string is independent of the number of keywords used to construct the state machine, we would not expect the time complexity to increase with the size of the keyword set. However, as can be seen in 18, the time to build the index increases logarithmically with the number of keywords. This is likely due to the fact that as the number of keywords increases, so does the size of the index and the number of database interacts required to store and update the index.

6.1.2 Space Complexity

As we would have expected it, the space used by the index in the database is proportional to the number of words we are using. It also means that our words have the same probability at the beginning or at the end of our 12,000 word dictionary. This result is reflected in figure 19.

6.2 Size of Corpus

To be able to study the impact of different number of Posts, we built a posts database with 400,000 posts randomly downloaded from the Internet. Then the keywords used to build the index are simply the first 5,000 from the previous

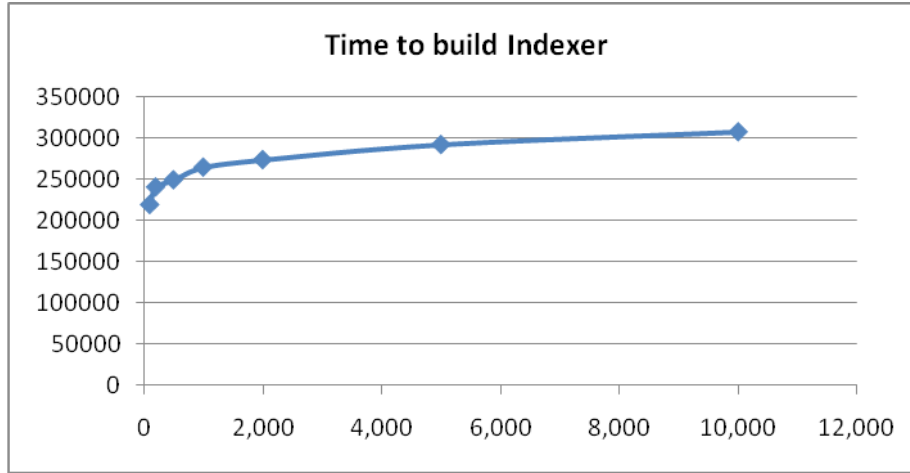


Figure 18: Time to build the inverted index as a function of the number of keywords. Time taken is shown in milliseconds on the y-axis and the number of keywords on the x-axis

keywords database.

The experiment shown here consist on building the index for different number of Posts and explore the time and space complexity for the building of the state machine and the building of the index once the state machine is built. We can expect that the time to build the State machine wont change with the number of posts and the time to build the index will increase almost linearly with the number of posts.

6.2.1 Time Complexity

As we were expecting, the time to build the index with the state machine strategy is linear with respect to the number of post, as depicted in figure 20. However, to be able to perform this experiment we had to run the indexer program with batches of 10,000 posts due to Java heap space problems. As a result, the state machine is rebuilt every 10,000 posts.

6.2.2 Space Complexity

As expected, the space used to store the index in the database is proportional to the size of the corpus. This result is shown in figure 21.

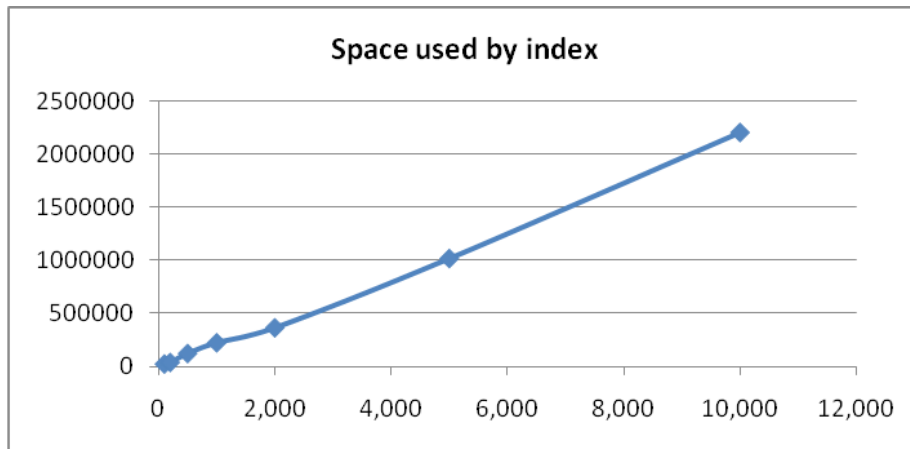


Figure 19: Size of index as a function of the number of keywords. The number of database rows used to store the index is shown on the y-axis and the number of keywords on the x-axis

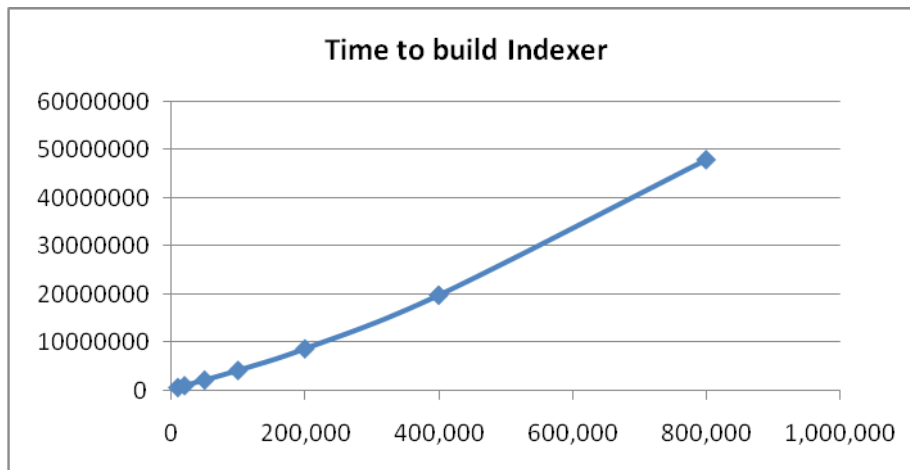


Figure 20: Time to build the inverted index as a function of the size of the corpus. Time taken is shown in milliseconds on the y-axis and the number of documents in the collection on the x-axis

6.3 Document Size

As we have seen before, the time to build the index is almost proportional with the number of post. Theoretically this time should be proportional to the total

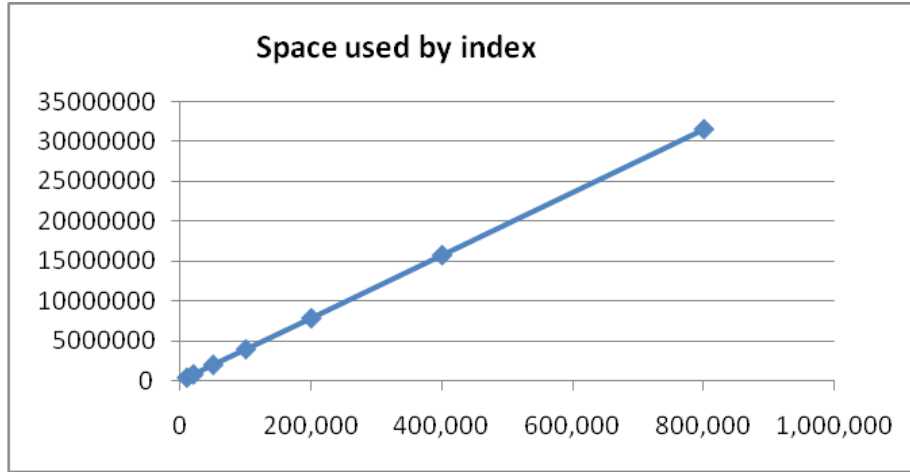


Figure 21: Size of index as a function of the size of the corpus. The number of database rows used to store the index is shown on the y-axis and the number of documents in the collection on the x-axis

length of all posts. Lets focus on this aspect and study the impact of the length of the post on the time to build the state machine. To do that we built a 5,000 posts database with posts that are all composed of 20,000 characters. Then, we ran our system on different length of substring of these posts.

6.3.1 Time Complexity

- State Machine Construction:** As expected, state machine construction time does not depend on the length of the documents. Nonetheless it is interesting to compare figures 22 and 23, showing the construction times for the state machine and inverted index relative to the average document size. The comparison shows that state machine construction time is several orders of magnitude smaller than index construction and negligible by comparison. Our state machine is built in 770 ms on average for 1000 keywords with a standard deviation of 18ms.
- Index Construction:** As expected the time to build the indexer is almost linear with respect to the average document size. This relationship is shown in figure 23.

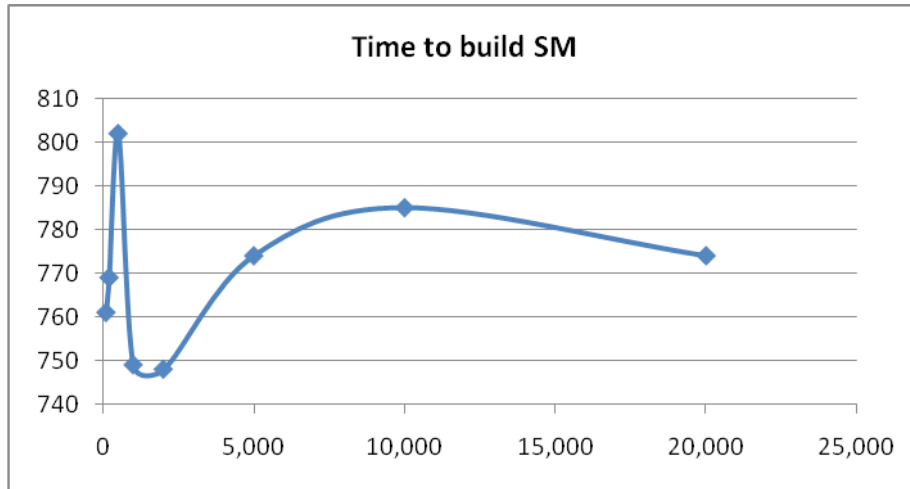


Figure 22: Time to build Aho-Corasick state machine as a function of the average document size. Time taken is shown in milliseconds on the y-axis and the average document size in characters on the x-axis

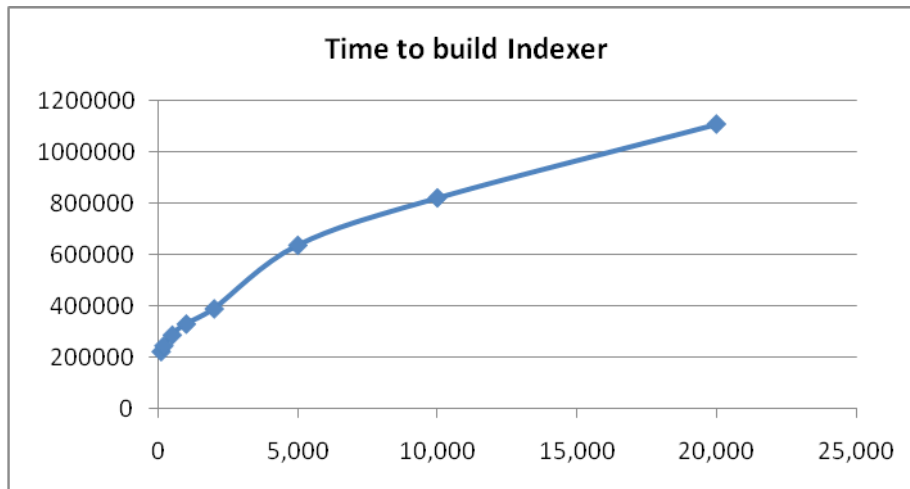


Figure 23: Time to build an inverted index as a function of the average document size. Time taken is shown in milliseconds on the y-axis and the average document size in characters on the x-axis

6.3.2 Space Complexity

Here we are studying the number of records in our index database as a function of the average document size. As expected, the larger the documents the more

records created to represent the index. However, this is not linear since if a word occurs many times in the same post we just save it once in the database. This relationship is shown in figure 24.

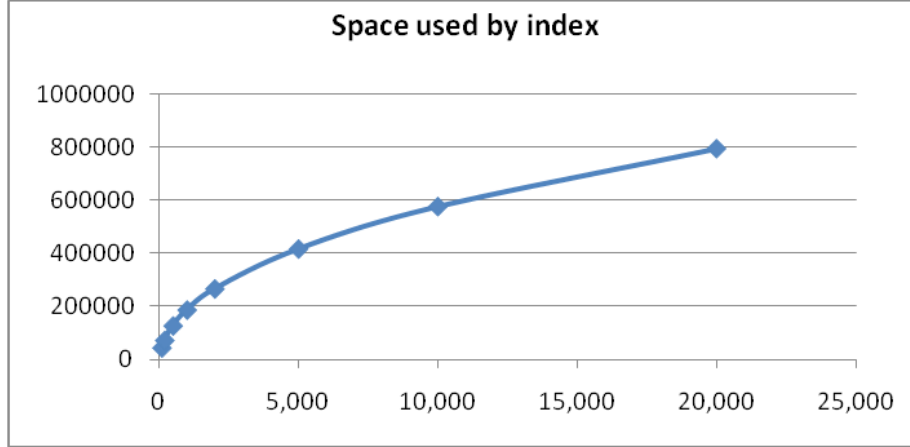


Figure 24: Size of index as a function of the average document size. The number of database rows used to store the index is shown on the y-axis and the average document size in characters on the x-axis

6.4 Comparison with Naïve Index Builder

This part focuses on the improvement that our system with respect to time complexity in comparison with a naïve index builder. The naïve indexer builder has the same input as our system i.e. a list of keywords and a list of documents, and the same output i.e. an inverted index corresponding to the keywords and documents. Algorithm 1 is used in principal for both implementations.

Algorithm 1 Build inverted index

```

for all keyword do
  Build a vector v
  for all document do
    if keyword  $\in$  document then
      Add the row (keyword.id, document.id) into v
    end for
  Save v into the database
end for

```

As we can see this algorithm has a complexity that directly depends on the

number of keywords n , the number of document m , and the length of document since we have to check in each document for each keyword. Moreover, we decided to limit, as much as possible, the accesses to the database since it is really time consuming in Java. That's why we do one access to the database for each keywords and not for each document. In the state machine we do as much accesses to the database as the number of document (in the worst case meaning we find at least one keyword per document). In the following subsections we compare the results for this algorithm with our system results for different number of keywords and posts.

6.4.1 Complexity Analysis

For the naïve algorithm described above, the theoretical worst case time complexity relationship is shown in equation 1.

$$Time \propto keyNum * (docNum * maxLength + conn) \quad (1)$$

where $keyNum$ is the number of keywords, $docNum$ is the number of documents, $maxLength$ is the maximum document length, and $conn$ is the time taken to create a connection to the database. For the state machine algorithm the theoretical complexity relationship is given by equation 2.

$$Time \propto docNum * (maxLength + conn) \quad (2)$$

6.4.2 Document Size

For this experiment we fixed the number of keywords to 1,000 and the number of documents to 5,000 and we ran our code on a database composed of document of different length. The experiment results are shown in figure 25.

As expected, the curves for both our indexer and the naïve indexer are almost linear since for the naïve implementation:

$$Time \propto (keyNum * docNum) * maxLength + conn * keyNum$$

and for our Aho-Corasick state machine method:

$$Time \propto docNum * maxLength + conn * docNum$$

We can observe that for a small length of document naïve method is faster

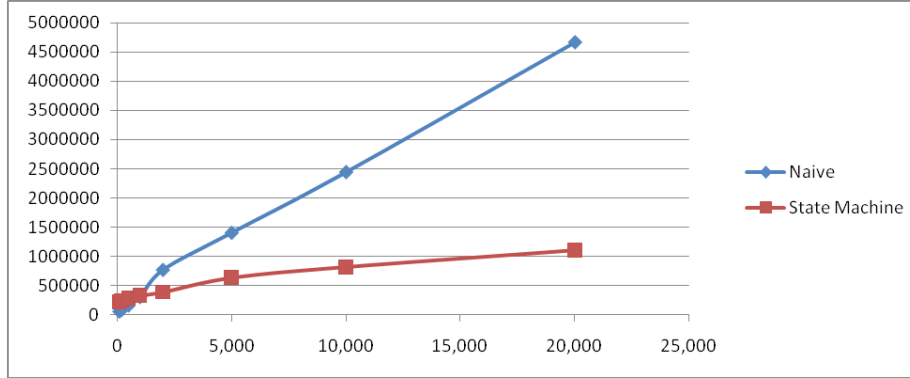


Figure 25: Time to build an inverted index as a function of the average document size. Time taken is shown in milliseconds on the y-axis and the average document size in characters on the x-axis

since $keyNum$ is smaller than $docNum$. For $docLength = 1300$ our indexer is far more efficient. To see this effect more clearly, experiment results are plotted on a logarithmic scale in figure 26.

6.4.3 Number of Keywords

As expected, the two implementations exhibit a linear response to the number of keywords used. In the naïve implementation:

$$Time \propto (keyNum * maxLength) * docNum + conn * keyNum$$

and in our Aho-Corasick state machine method:

$$Time \propto (maxLength + conn) * docNum$$

However, we would have expected the state machine method not to depend on the number of keyword. The reason why its not the case is because when we dont have any keywords in a document, we dont access the database. Knowing that the database access is very time consuming, increasing the number of keyword increases the probability of finding a keyword into each document and so increase the time complexity.

The experiment results are shown in figure 27 and on a logarithmic scale in figure 28.

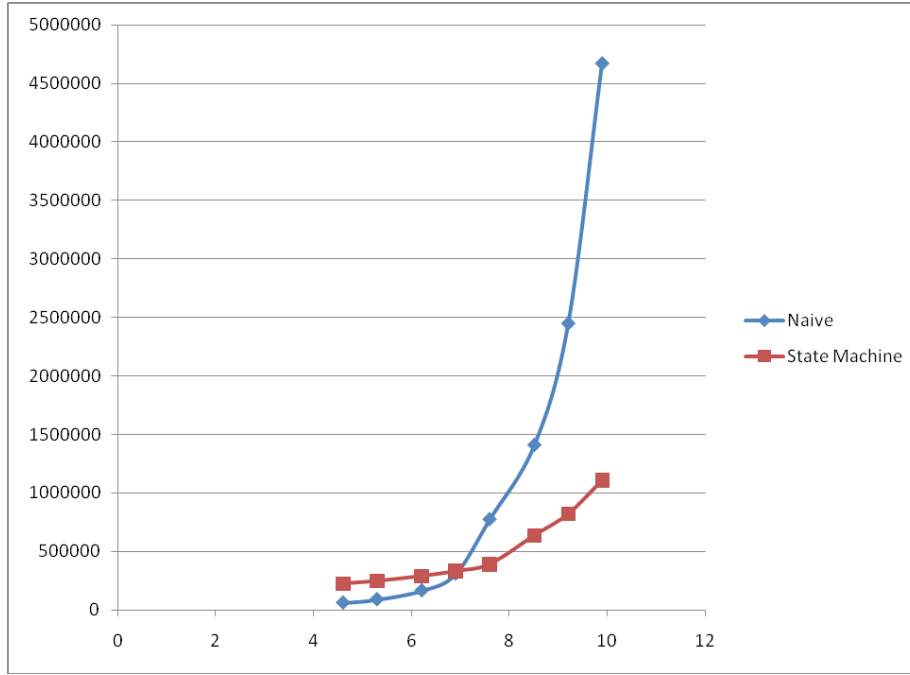


Figure 26: Time to build an inverted index as a function of the average document size. Time taken is shown in milliseconds on the y-axis and the average document size in characters on the logarithmic x-axis

6.4.4 Size of Corpus

As expected, the time complexity is linearly related to the number of documents being indexed. The experiment results are shown in figure 29 and on a logarithmic scale in figure 30. For the naïve implementation:

$$Time \propto (docNum * maxLenght + conn) * keyNum$$

and in our Aho-Corasick state machine method:

$$Time \propto docNum * maxLenght + conn * docNum$$

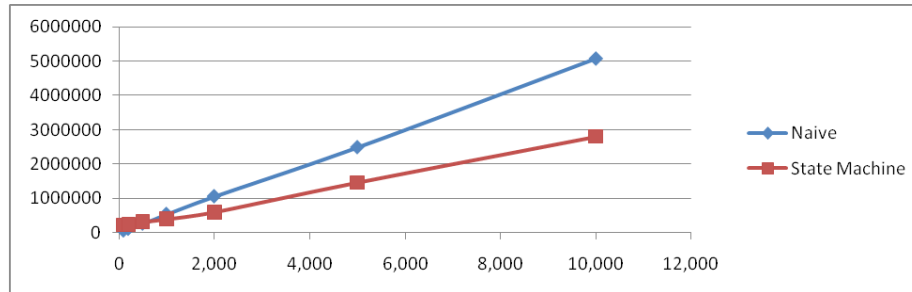


Figure 27: Time to build an inverted index as a function of the number of keywords. Time taken is shown in milliseconds on the y-axis and the number of keywords on the x-axis

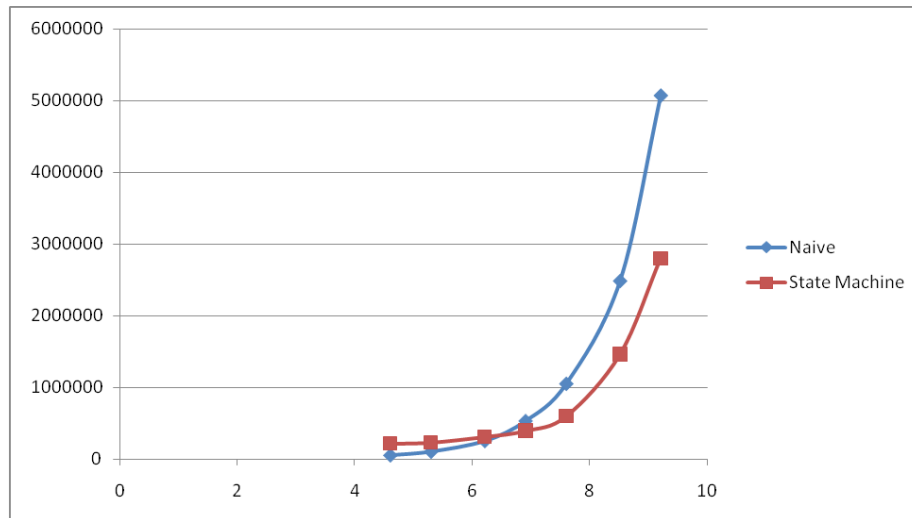


Figure 28: Time to build an inverted index as a function of the number of keywords. Time taken is shown in milliseconds on the y-axis and the number of keywords on the logarithmic x-axis

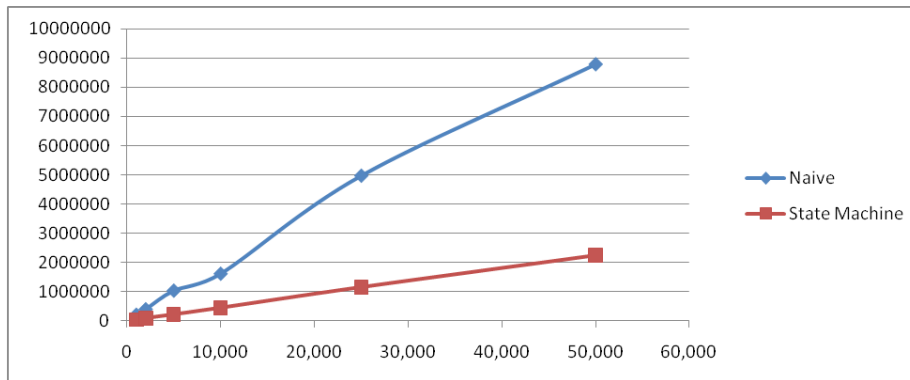


Figure 29: Time to build an inverted index as a function of the number of documents being indexed. Time taken is shown in milliseconds on the y-axis and the number of documents on the x-axis

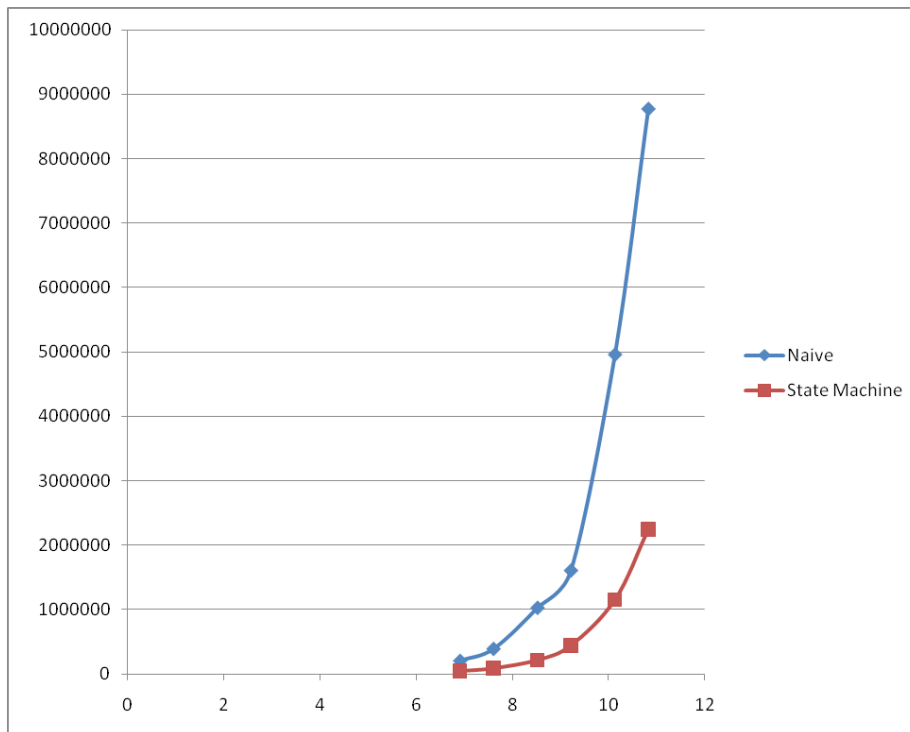


Figure 30: Time to build an inverted index as a function of the number of documents being indexed. Time taken is shown in milliseconds on the y-axis and the number of documents on the logarithmic x-axis

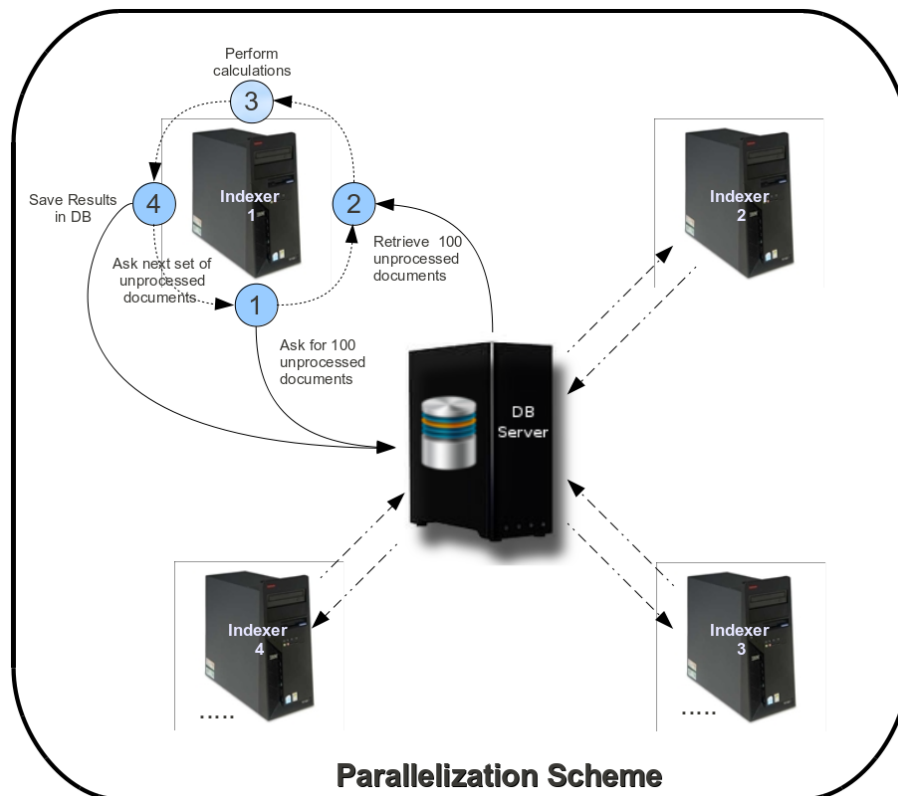


Figure 31: Deployment model of parallel indexer

6.5 Parallelization

7 Conclusions

The experimental results tend to prove what we could expect from theory. Our system is faster than the naïve indexer with respect to the number of keywords being indexed, size of the document collection, and average document size. In particular, our indexer implementation performs relatively well with respect to the number of documents being indexed. Moreover, as soon as we are dealing with a number of keywords that can be representative of real-world needs, the state machine implementation proves also to be significantly faster than the naïve indexer.

One of the primary goals of this research project was to leverage the effect of parallelization in reducing the time to index a document collection. Due to time constraints, this was not done but the resulting implementation lays a strong foundation for further research in this area. Other avenues of interest include, but are not limited to:

- Compare the Aho-Corasick state machine indexer implementation to less naïve indexer implementations.
- Use compression techniques to our algorithm in order to reduce the index size in the database.
- Study the time performance improvements possible deploying the Aho-Corasick state machine indexer within a commercial *MapReduce* framework such as *Hadoop*.

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