

PSO 9

Belt Parkway
Brooklyn, NY



WE ARE SO BACK (from Fall break)

How was the midterm?

Project due this thursday

Question 1

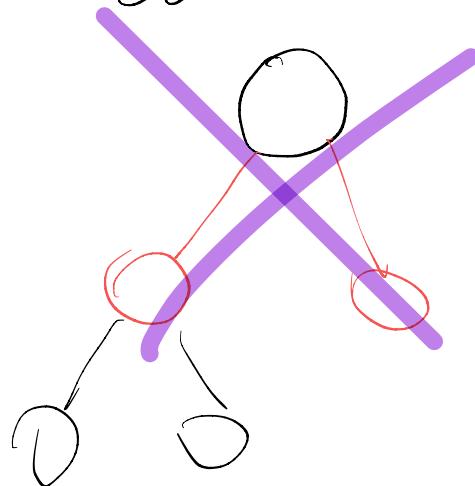
(Insertion and Deletion)

(1) Insert $\{15, 21, 7, 24, 0, 26, 3, 28, 29\}$ into an initially empty Left-Leaning Red-Black tree.

(2) Delete 7 in the final Left-Leaning Red-Black tree obtained in question (1).

LLRB Trees, what are they?

LLRB: Self balancing bst.



1) no 2 red nodes
in a row

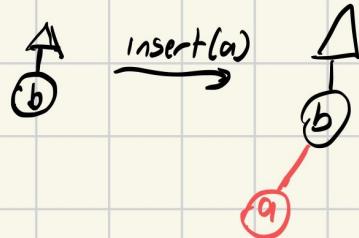
2) red node is
left-leaning

Types of LLRB inserts

(left)

Always insert in a red node

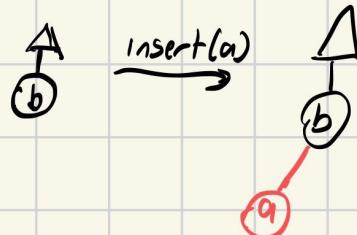
1) Left Insert, black parent



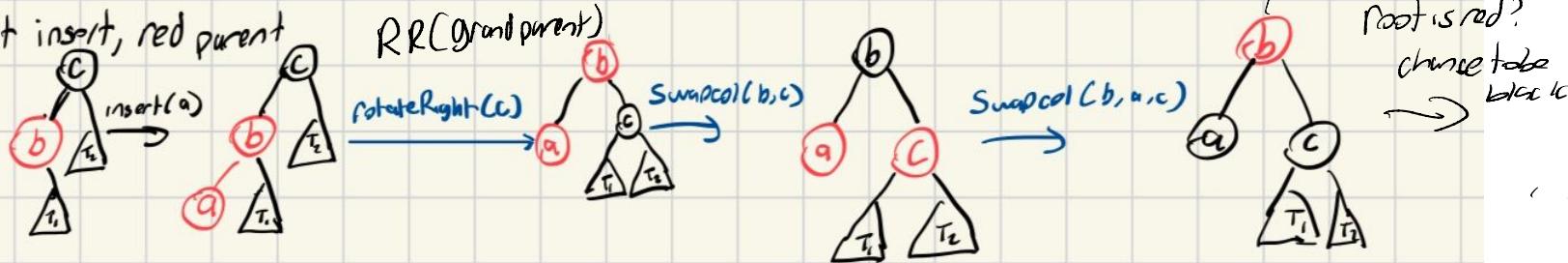
Types of LLRB inserts

Always insert in a red node

1) Left Insert, black parent



2) Left insert, red parent



Types of LLRB inserts

Always insert in a red node

3) Right Insert, any parent

(b)

Insert(c)

(b)

rc

rotate left(b)

RL(Parent)

(b)

c

Swap(c,b)

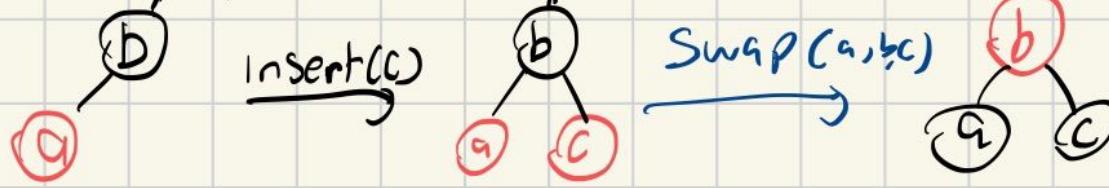
b

c

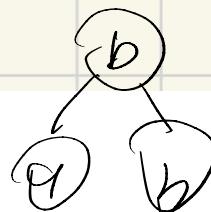
Types of LLRB inserts

Always insert in a red node

4) Right insert, sibling is also red



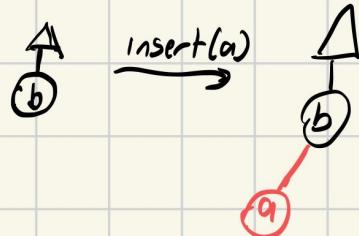
If b root, make b black



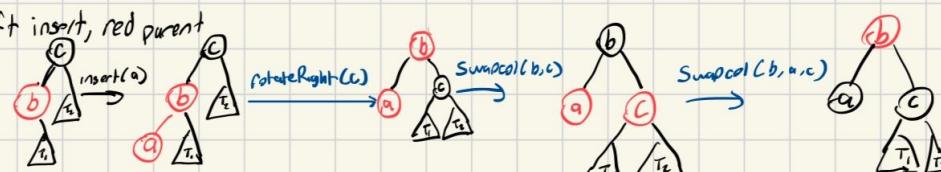
Types of LLRB inserts

Always insert in a red node

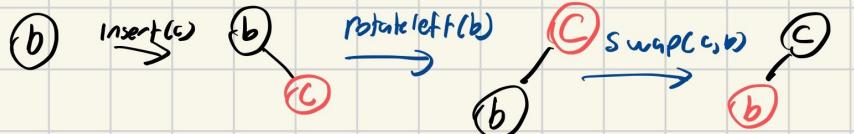
1) Left Insert, black parent



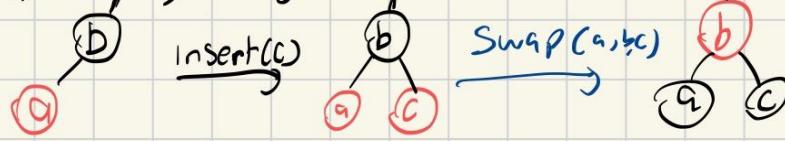
2) Left insert, red parent



3) Right Insert, any parent



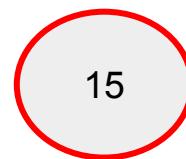
4) Right insert, sibling is also red



If b root, make b black

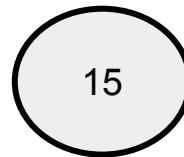
Insert: 15,21,7,24,0,26,3,28,29

Insert: 15,21,7,24,0,26,3,28,29



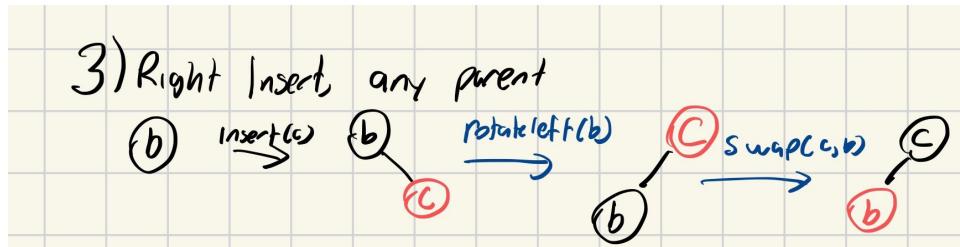
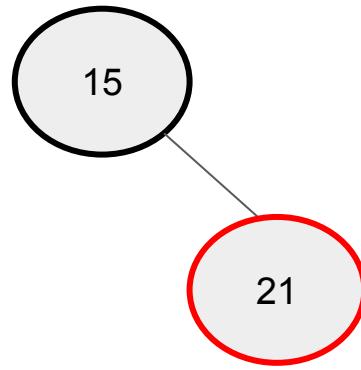
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



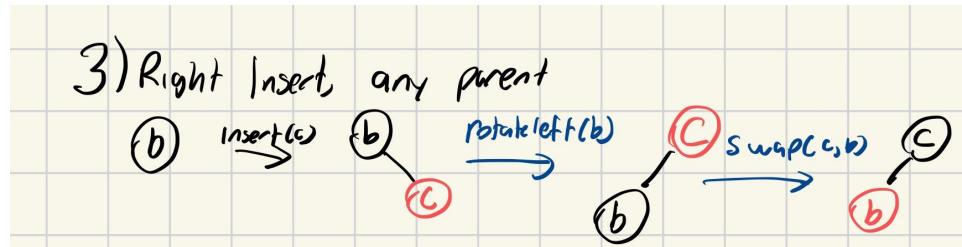
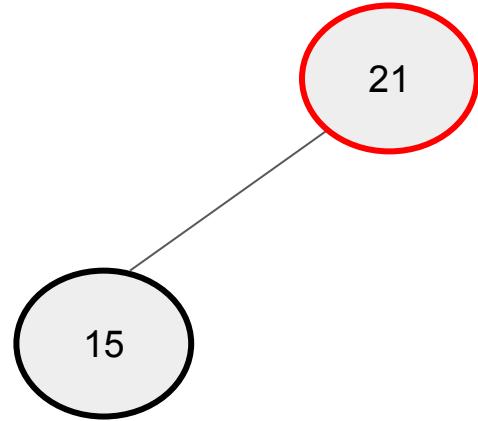
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



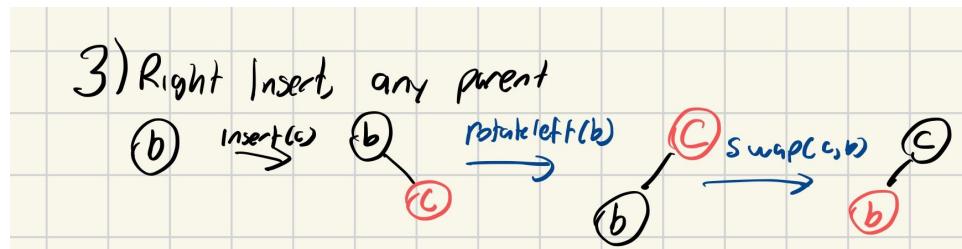
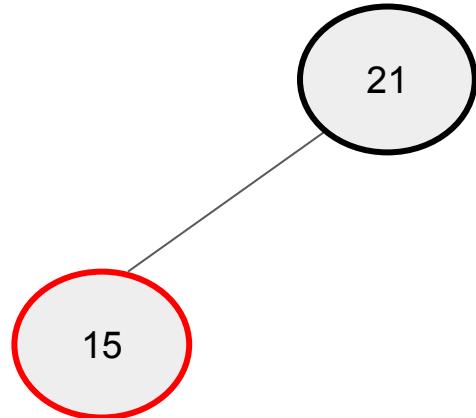
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



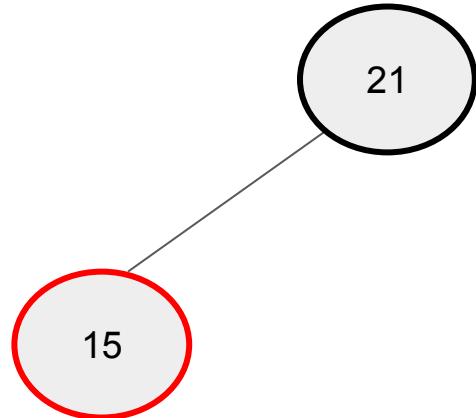
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



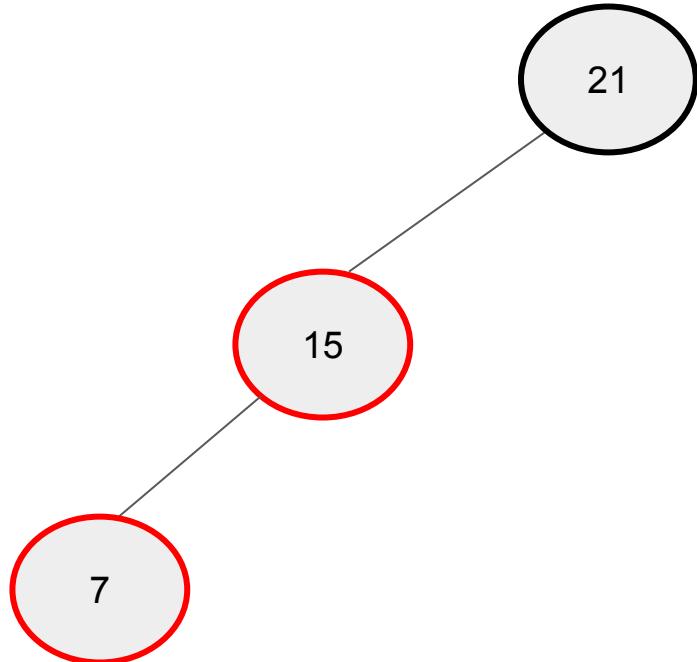
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



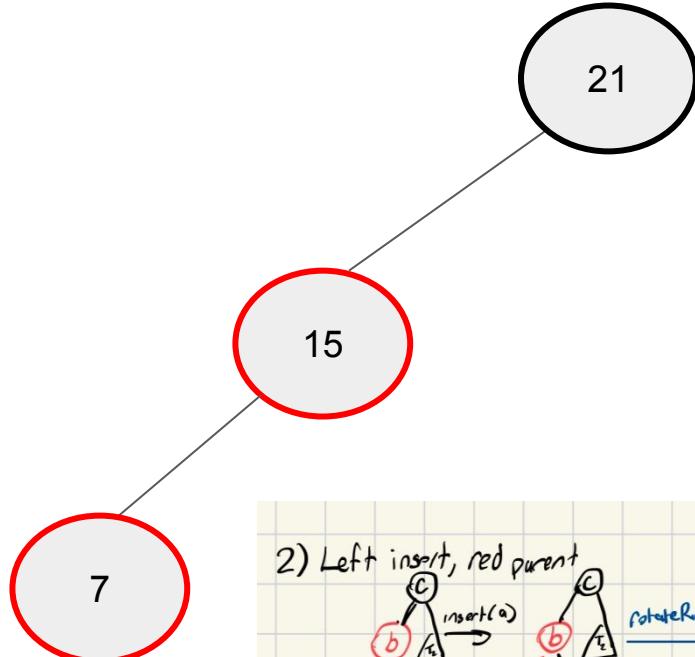
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



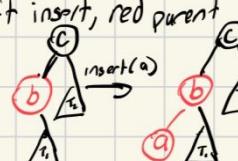
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



2)

Left insert, red parent



21

15

7

b

a

c

T₁

T₂

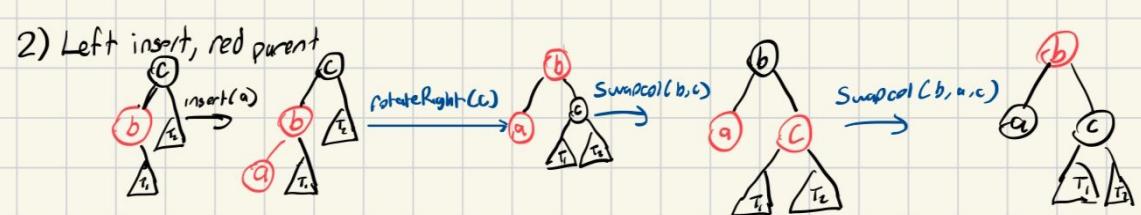
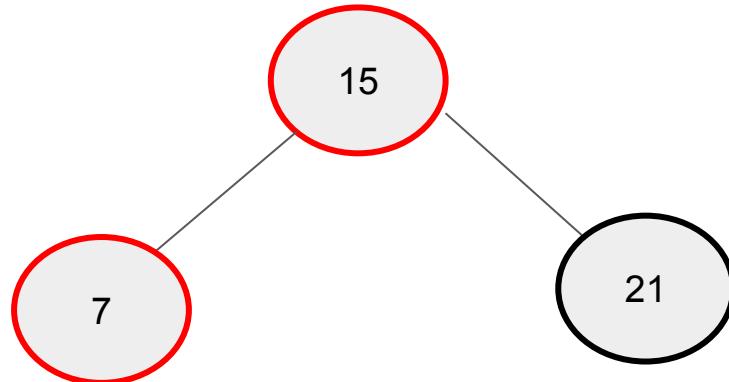
b

a

c

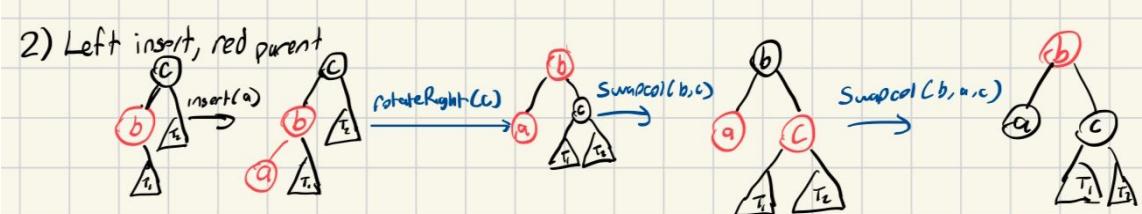
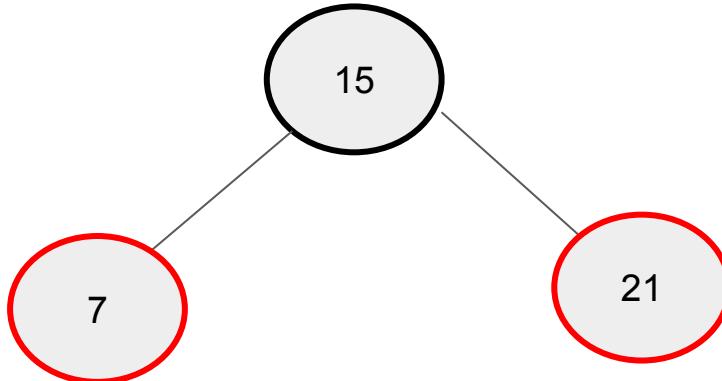
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



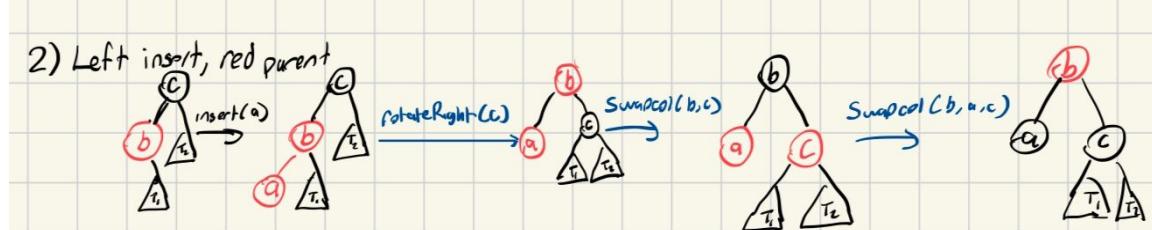
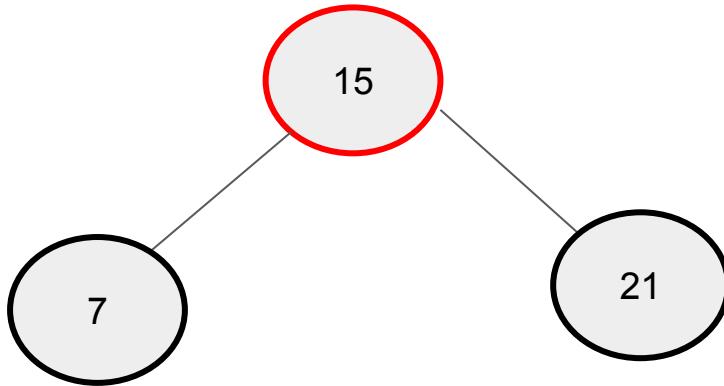
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



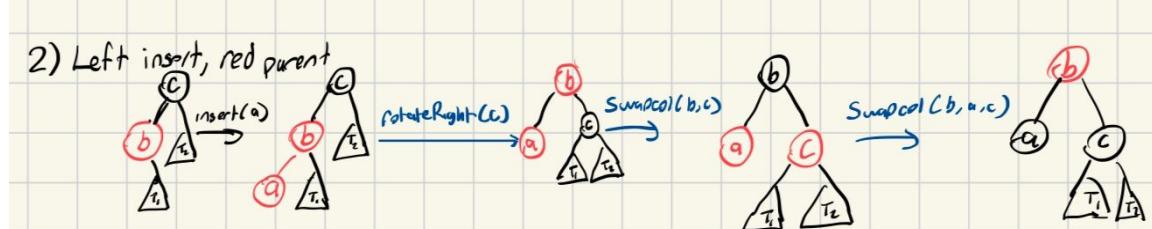
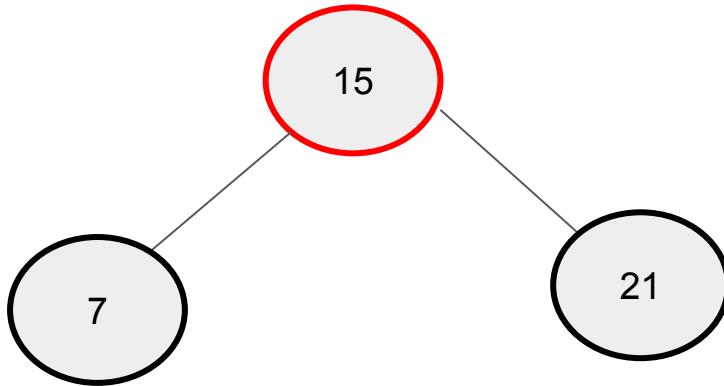
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



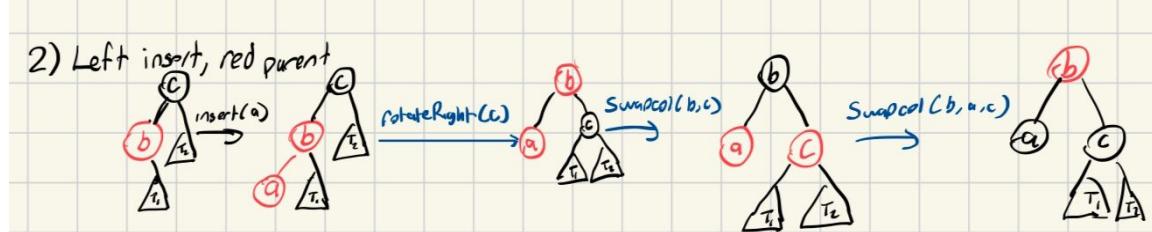
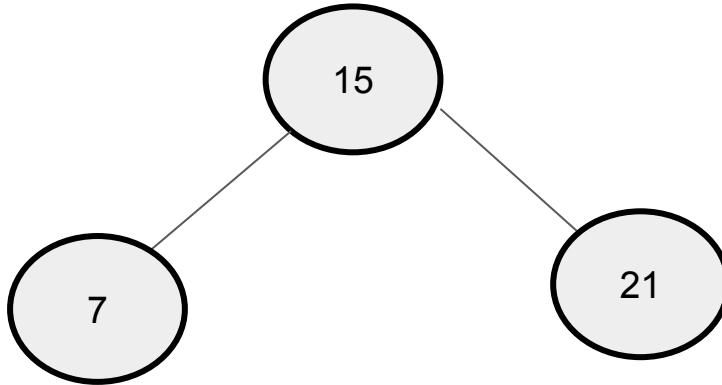
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



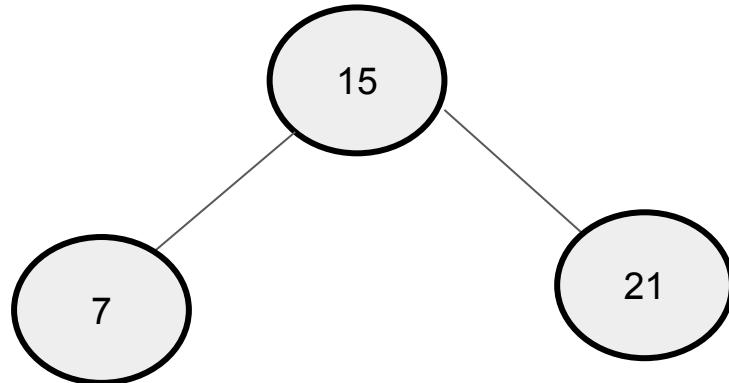
Insert: 15, 21, 7, 24, 0, 26, 3, 28, 29

If root red, make it black



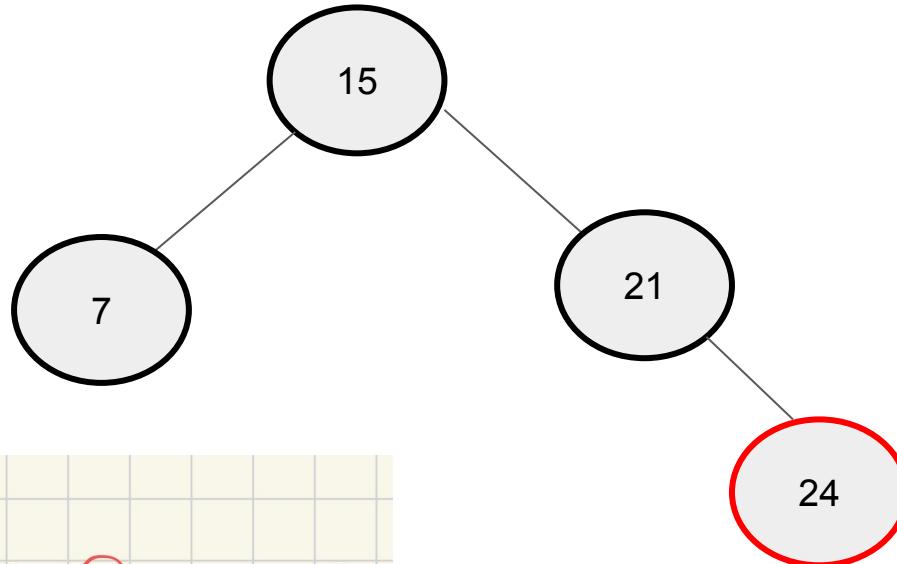
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black

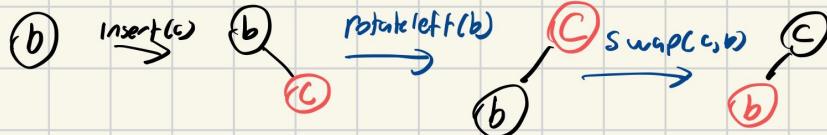


Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black

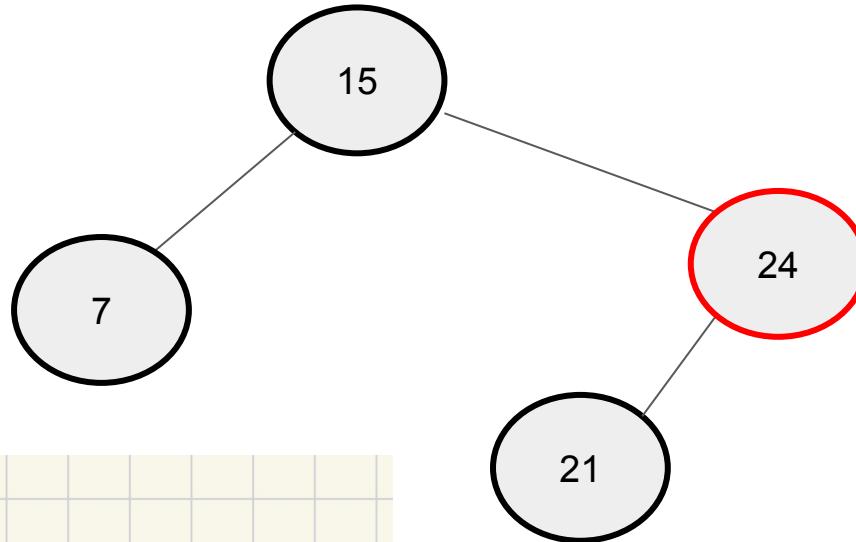


3) Right Insert, any parent

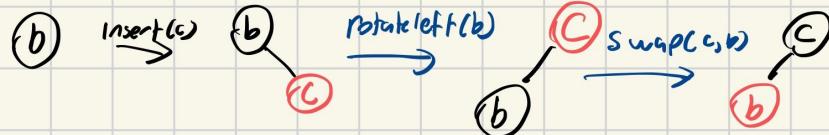


Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black

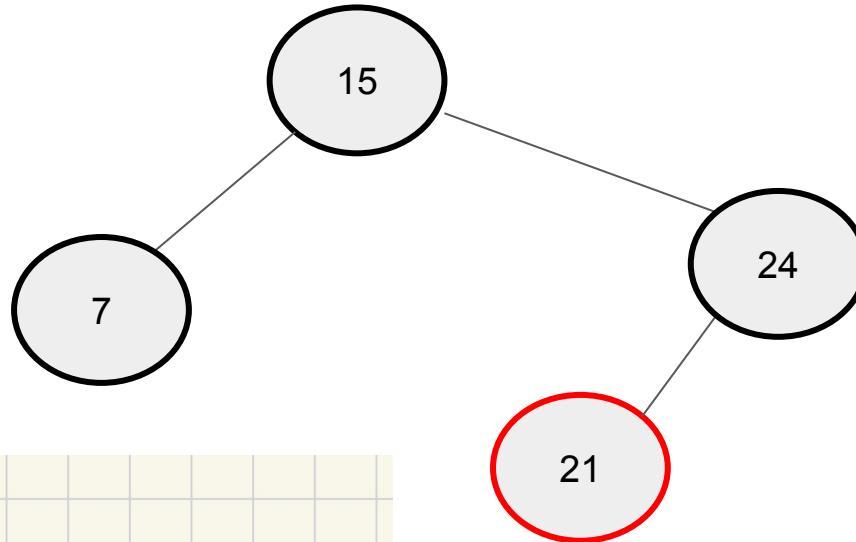


3) Right Insert, any parent

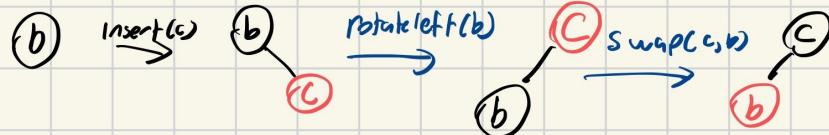


Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black

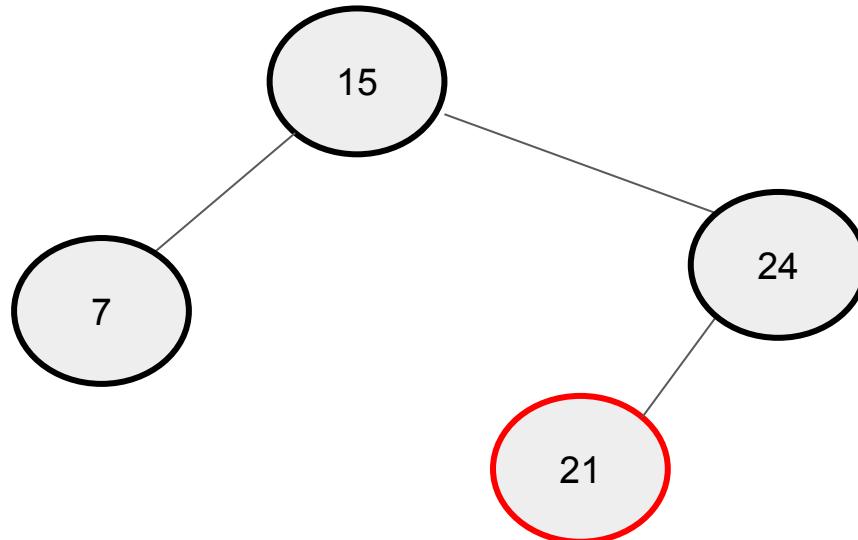


3) Right Insert, any parent



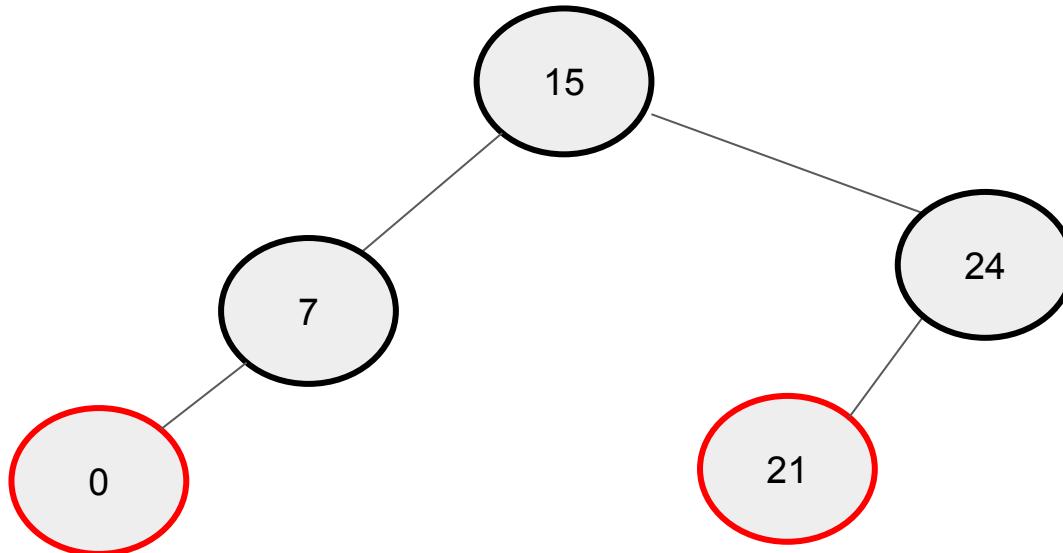
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



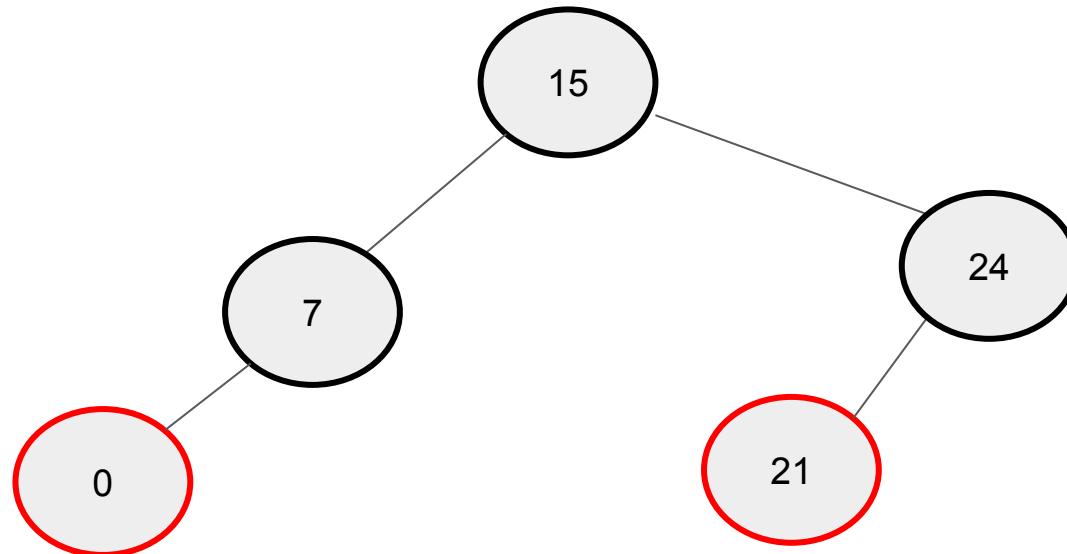
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



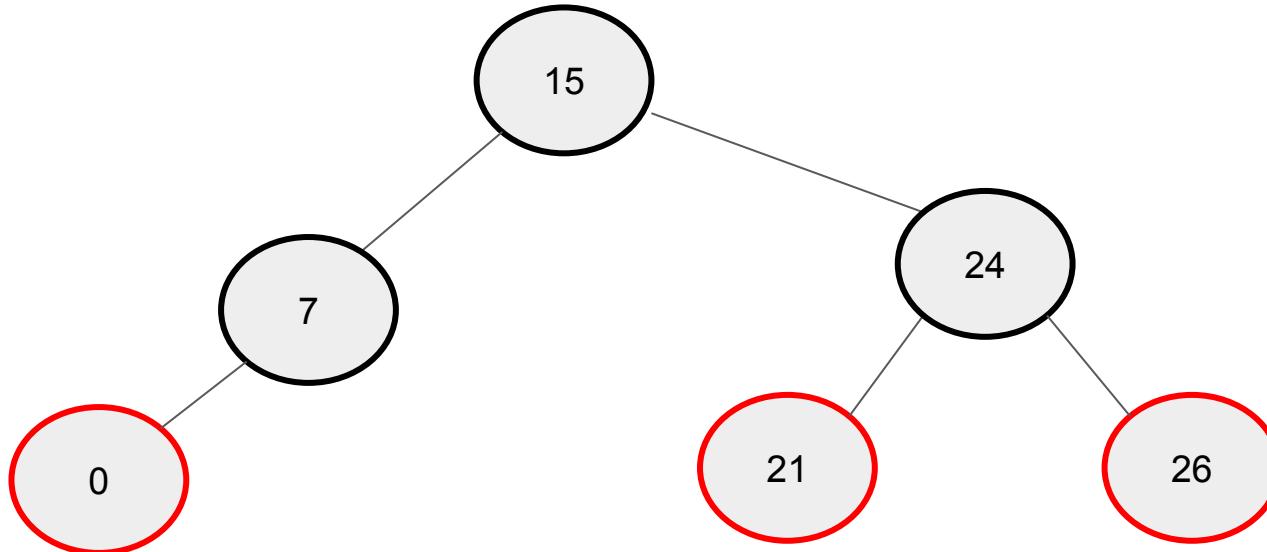
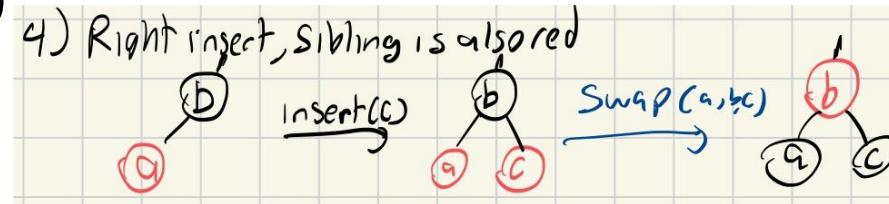
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



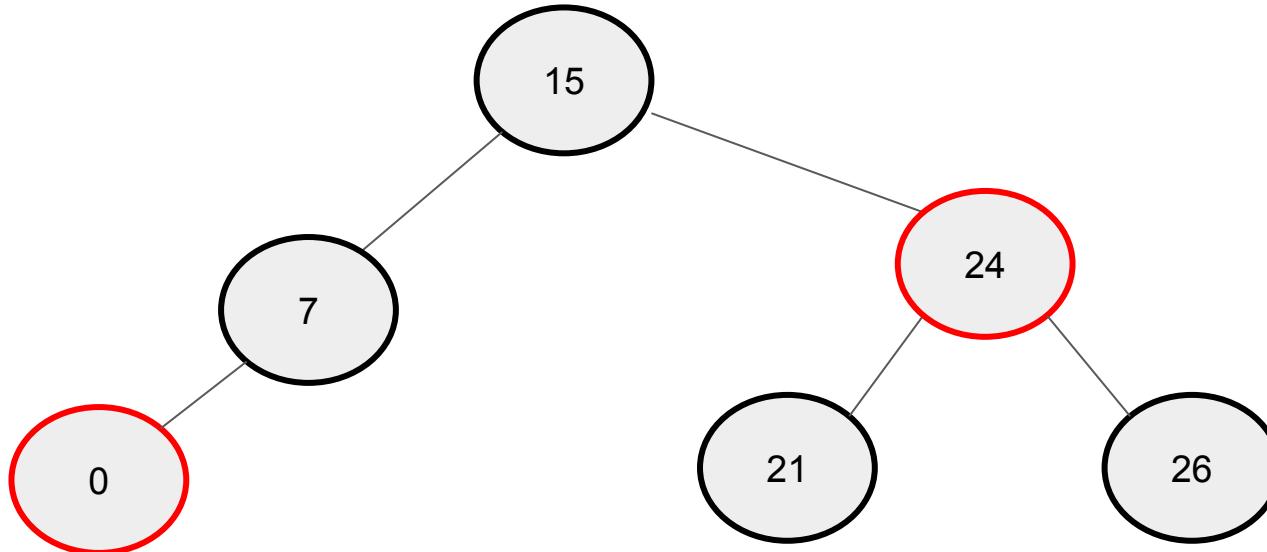
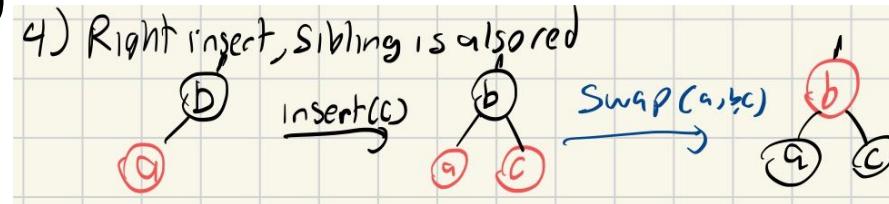
Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



Insert: 15,21,7,24,0,26,3,28,29

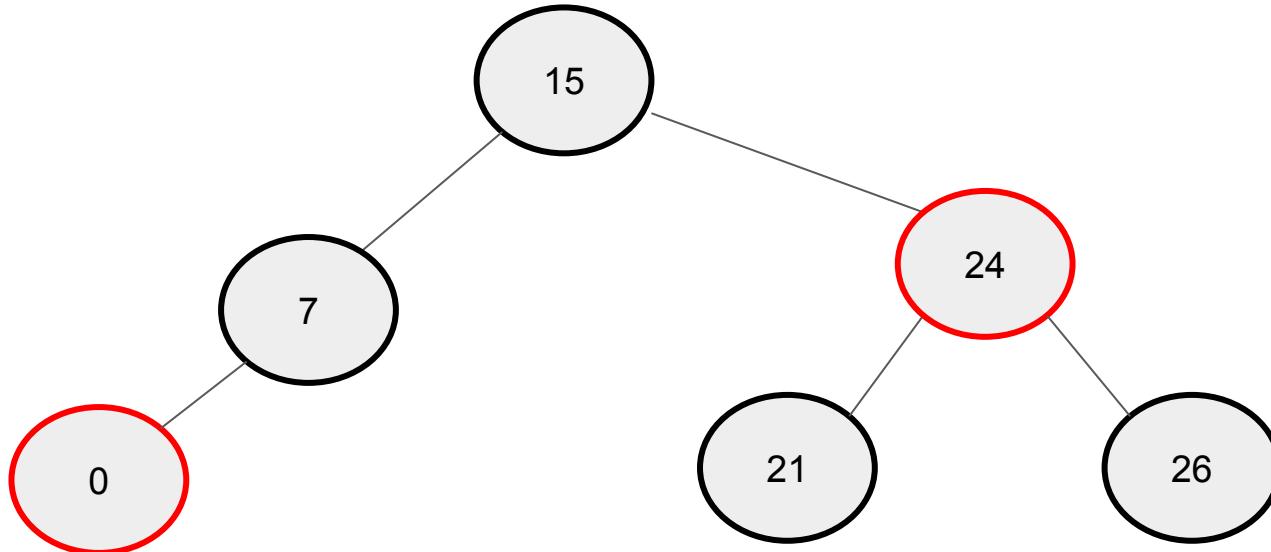
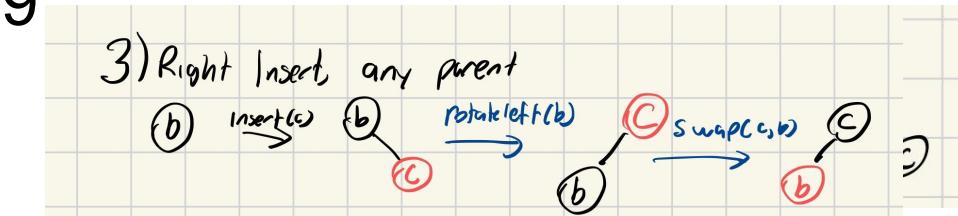
If root red, make it black



Oh no! Right red child, treat as red right insert

Insert: 15,21,7,24,0,26,3,28,29

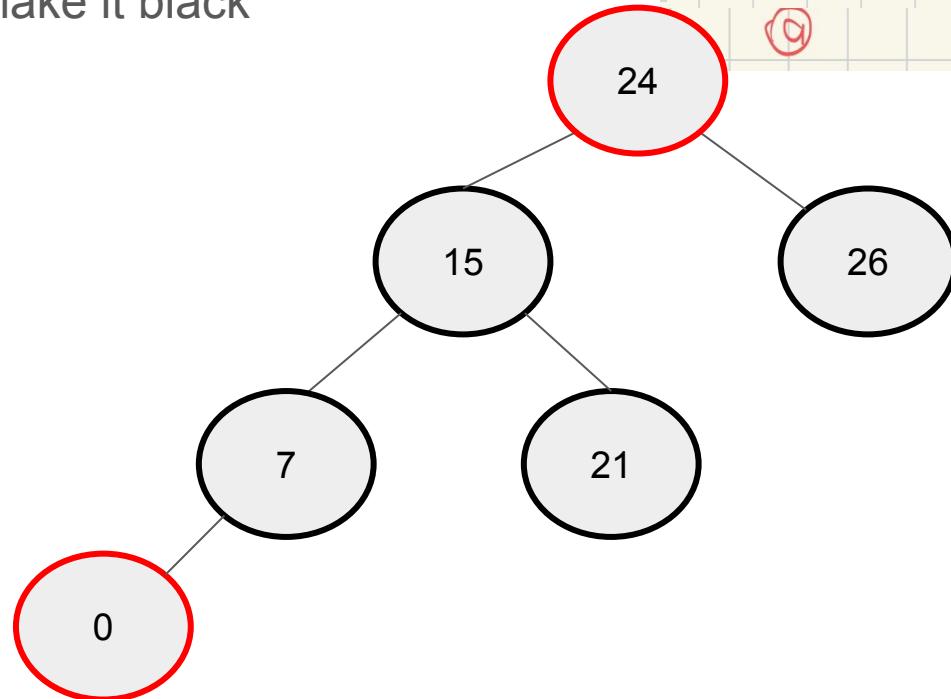
If root red, make it black



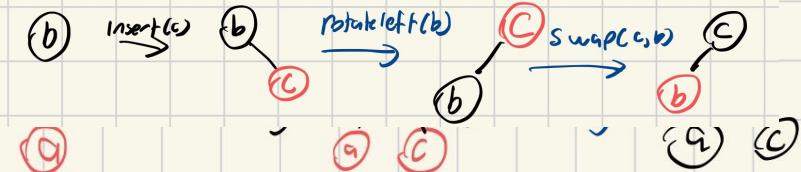
Oh no! Right red child, treat as red right insert

Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black



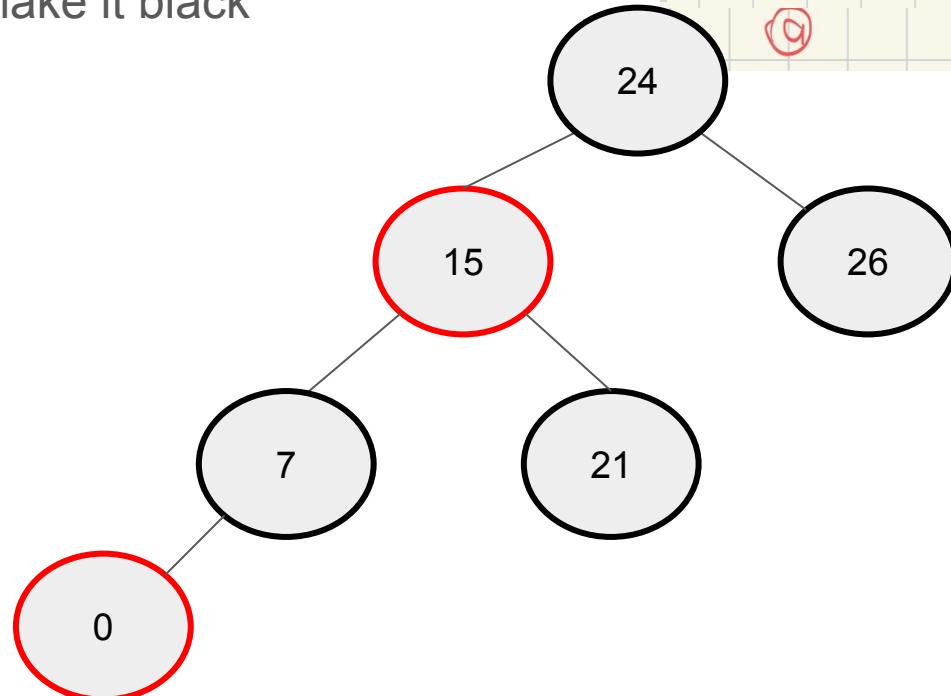
3) Right Insert, any parent



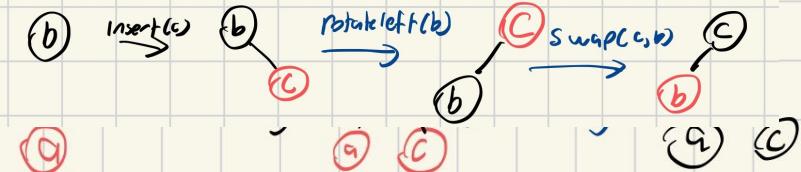
Oh no! Right red child, treat as red right insert

Insert: 15,21,7,24,0,26,3,28,29

If root red, make it black

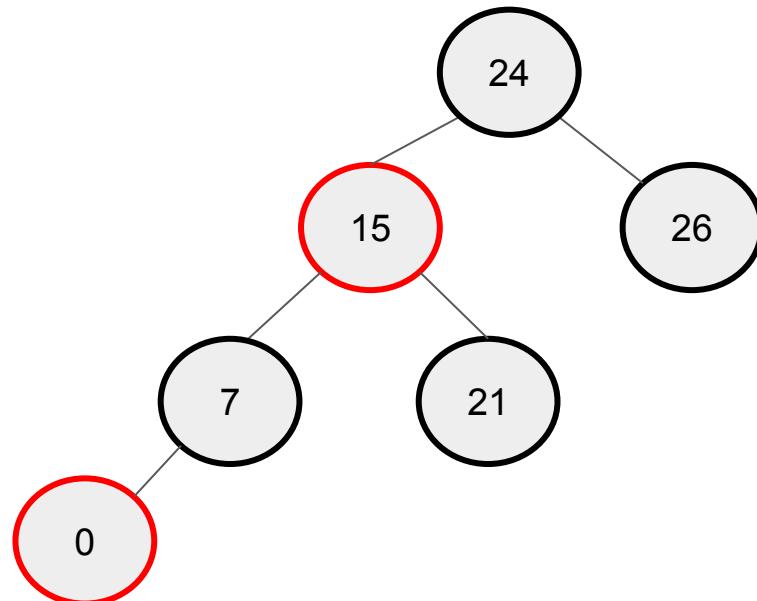


3) Right Insert, any parent

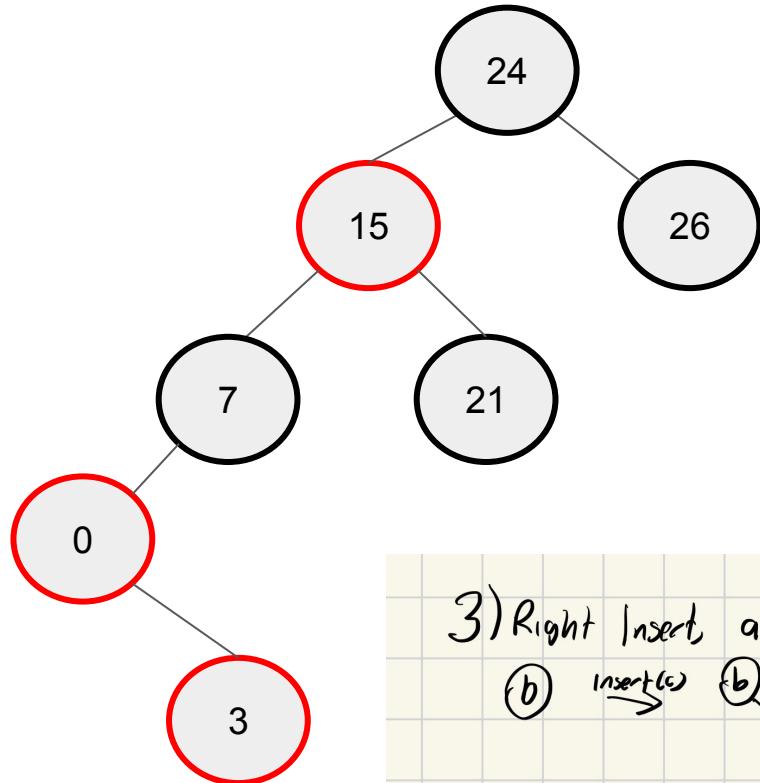


Oh no! Right red child, treat as red right insert

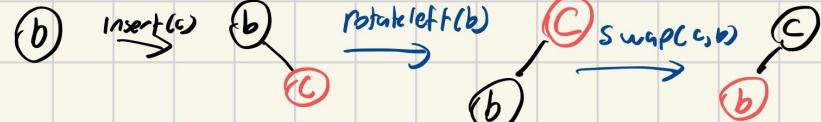
Insert: 15,21,7,24,0,26,3,28,29



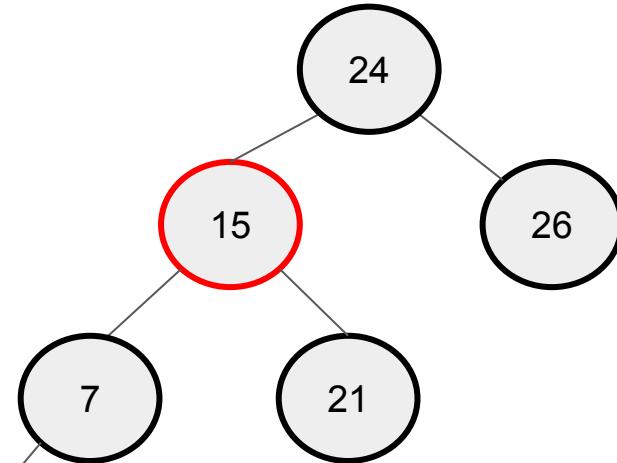
Insert: 15,21,7,24,0,26,3,28,29



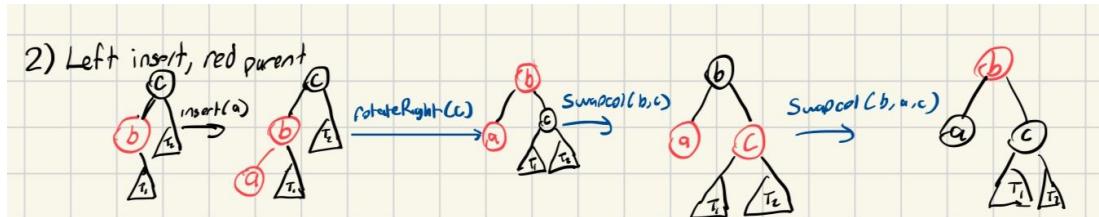
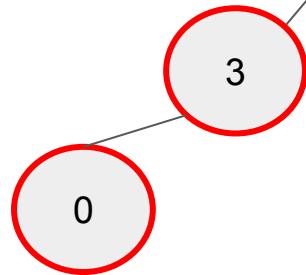
3) Right Insert, any parent



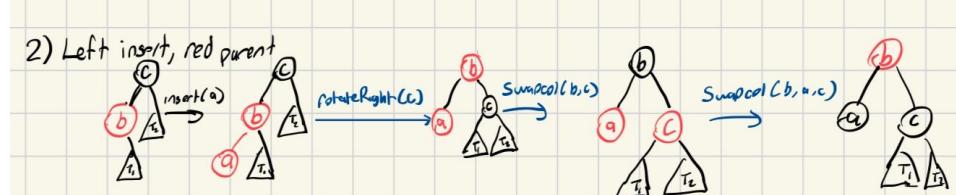
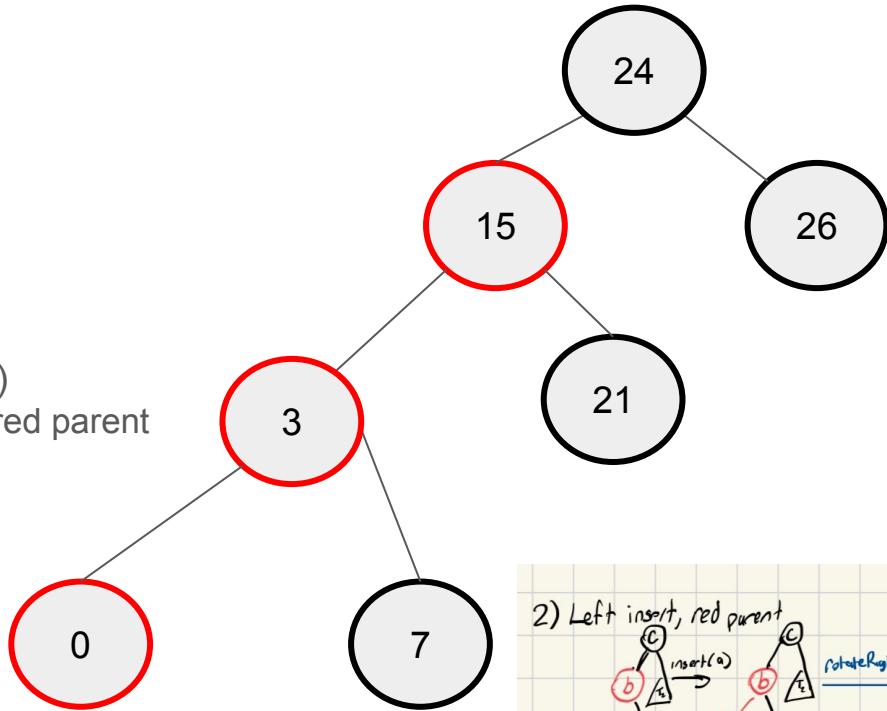
Insert: 15,21,7,24,0,26,3,28,29



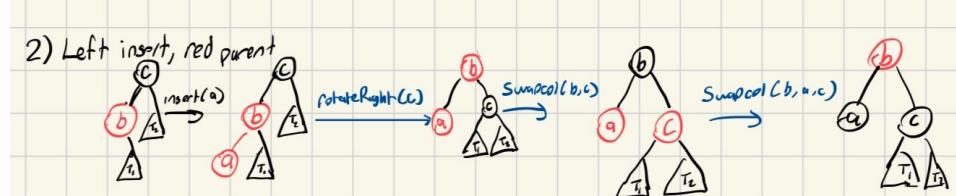
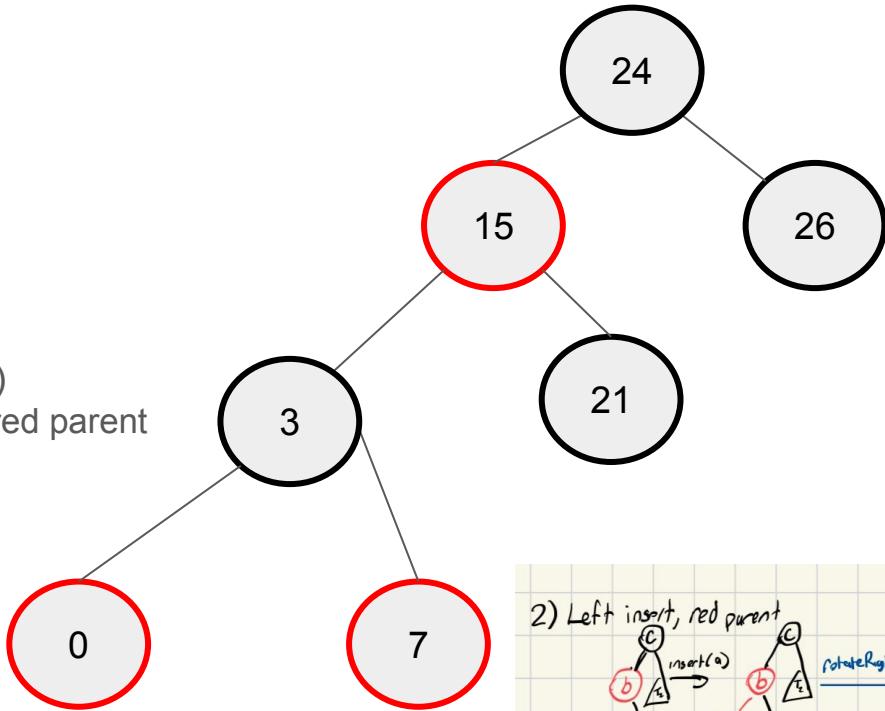
Two reds in a row (BAD)
Treat as red insert under red parent



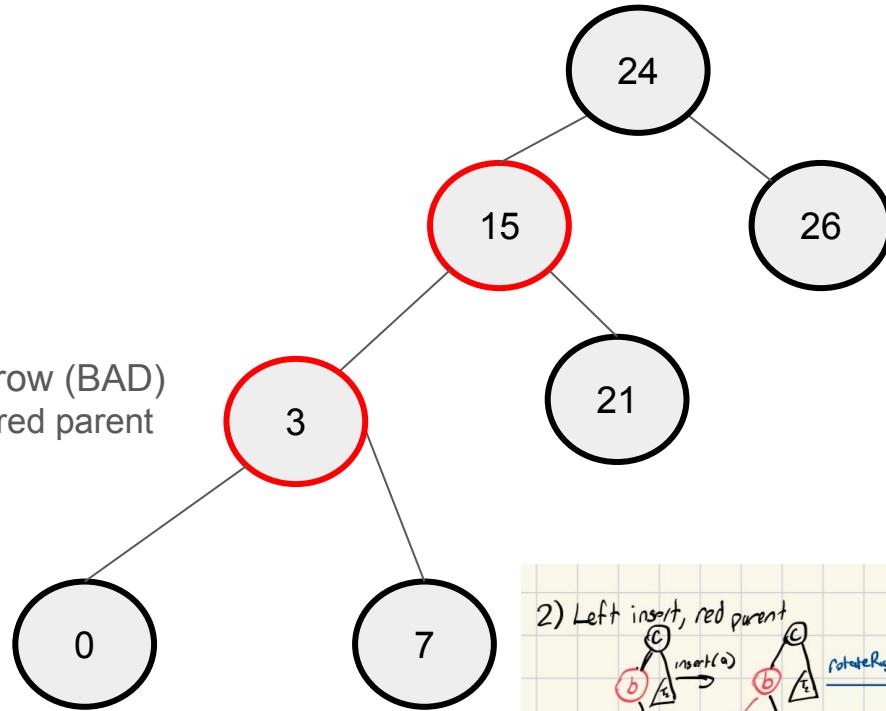
Insert: 15,21,7,24,0,26,3,28,29



Insert: 15,21,7,24,0,26,3,28,29

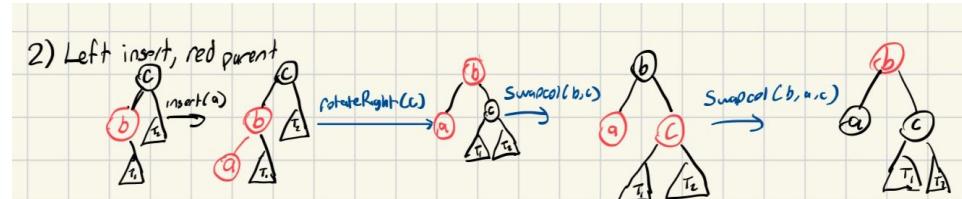


Insert: 15,21,7,24,0,26,3,28,29

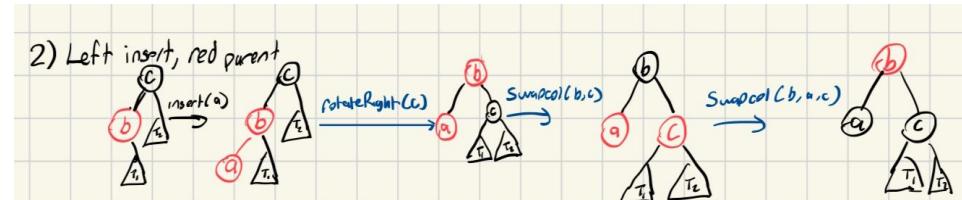
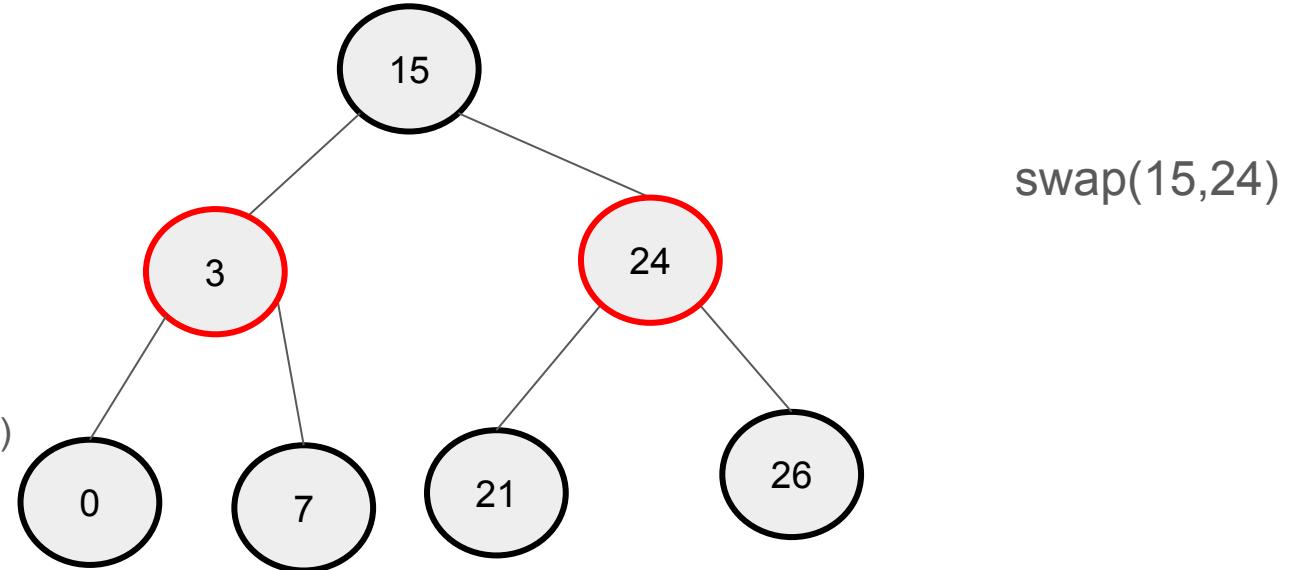


Rotate right at 24

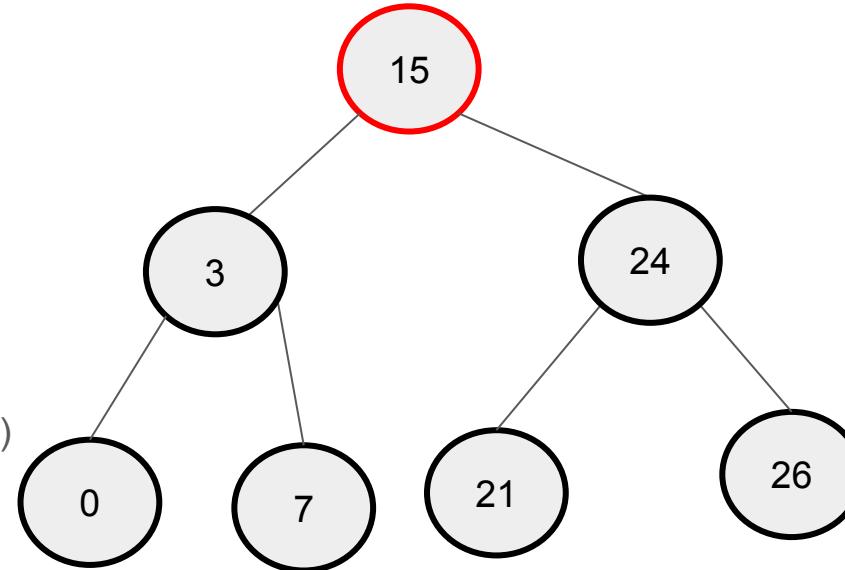
Another Two reds in a row (BAD)
Treat as red insert under red parent



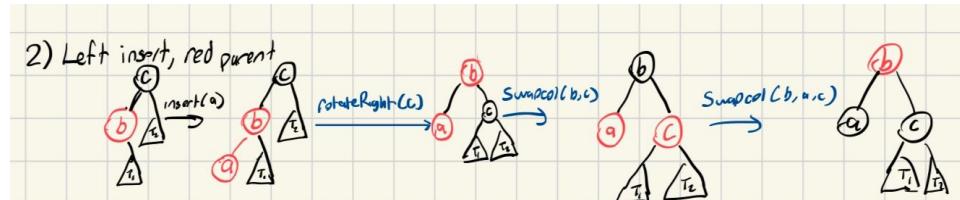
Insert: 15,21,7,24,0,26,3,28,29



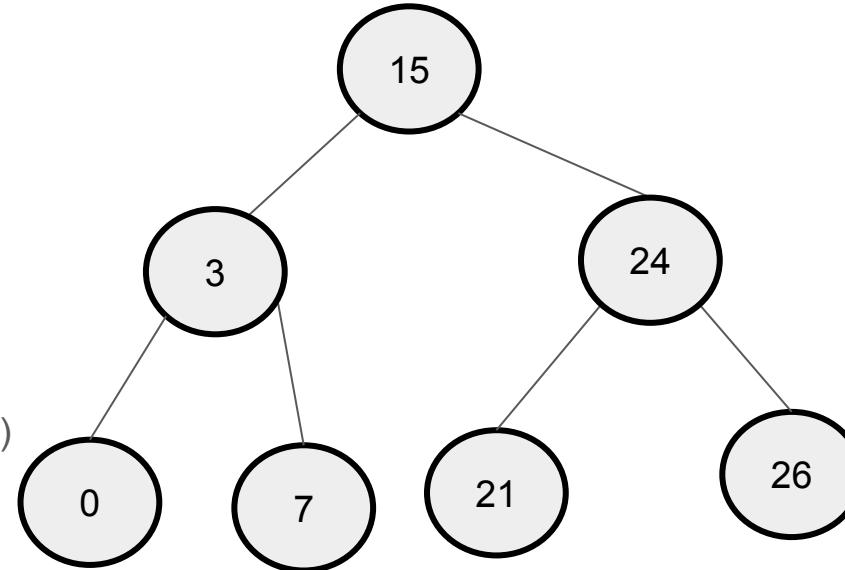
Insert: 15,21,7,24,0,26,3,28,29



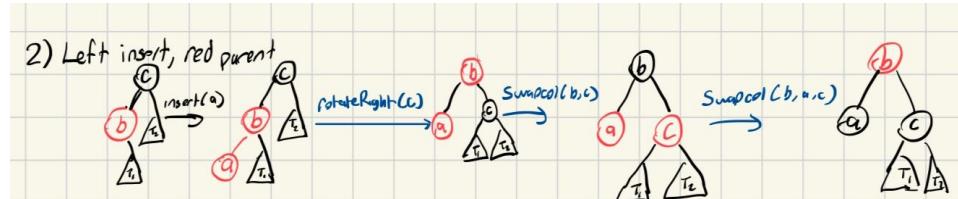
Another Two reds in a row (BAD)
Treat as red insert under red parent



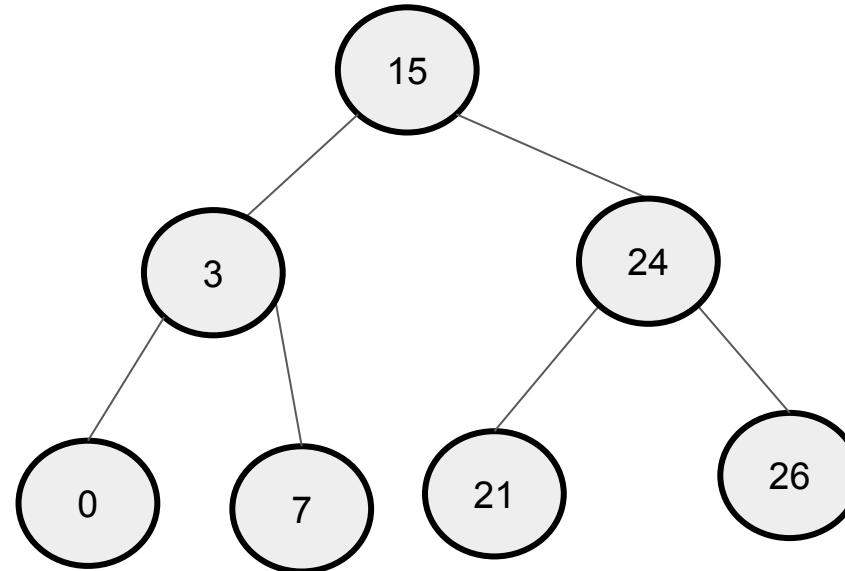
Insert: 15,21,7,24,0,26,3,28,29



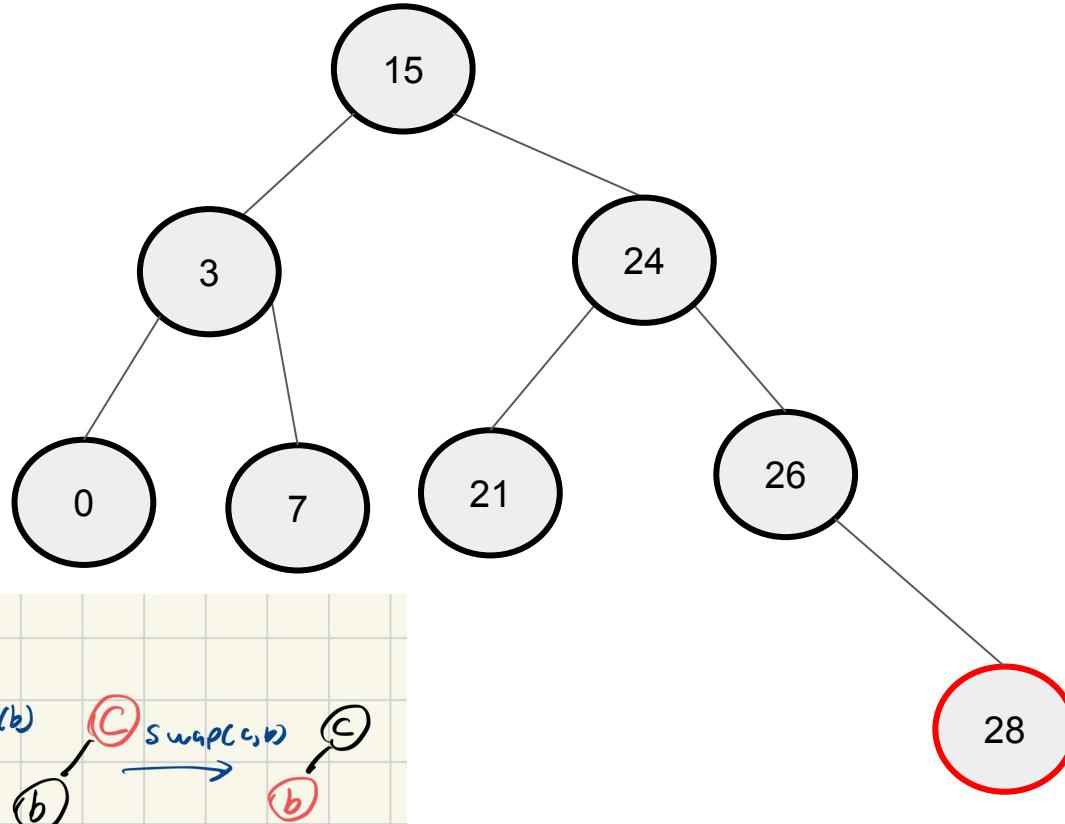
Another Two reds in a row (BAD)
Treat as red insert under red parent



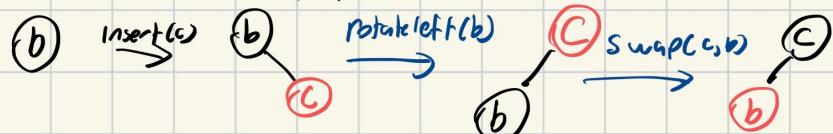
Insert: 15,21,7,24,0,26,3,28,29



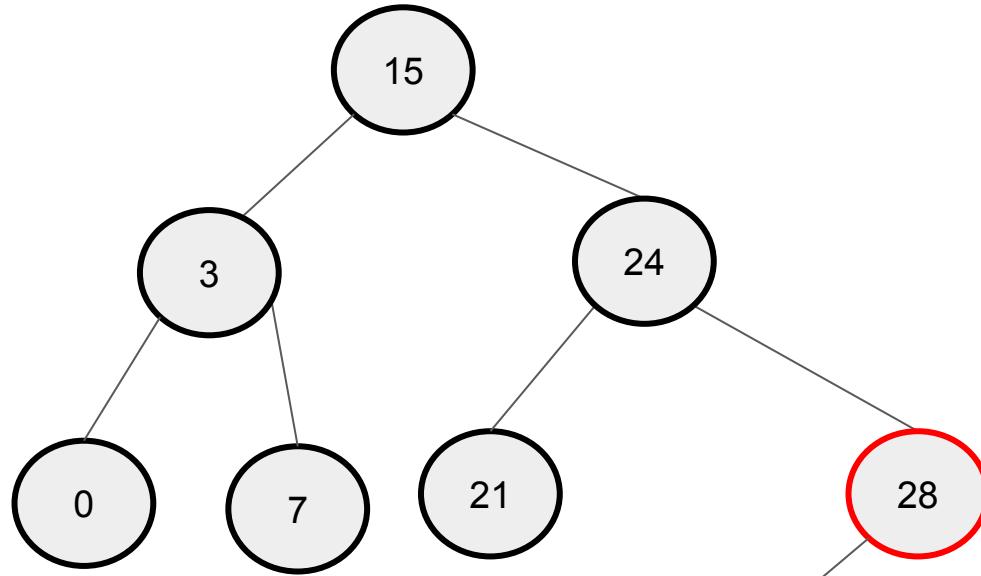
Insert: 15,21,7,24,0,26,3,28,29



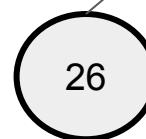
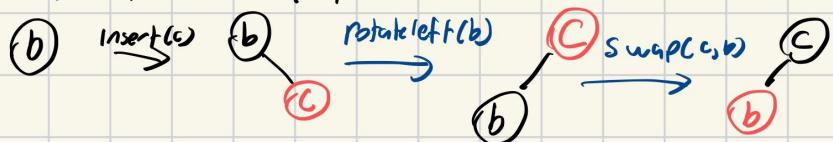
3) Right Insert any parent



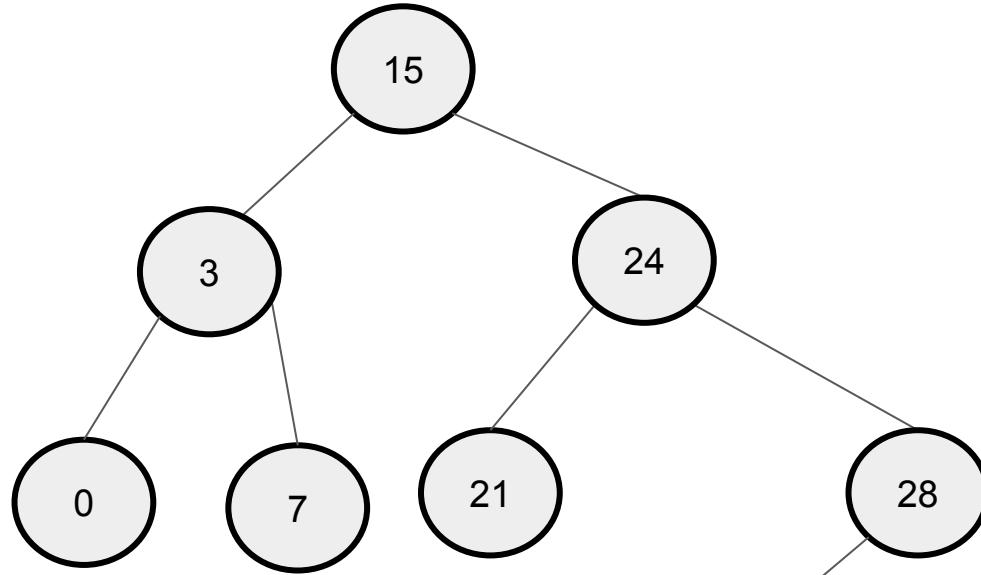
Insert: 15,21,7,24,0,26,3,28,29



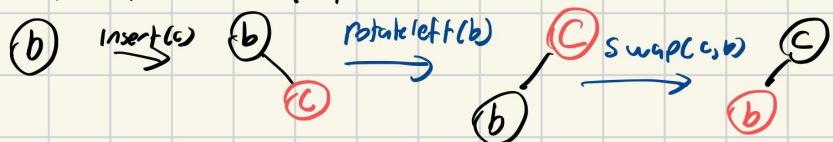
3) Right Insert any parent



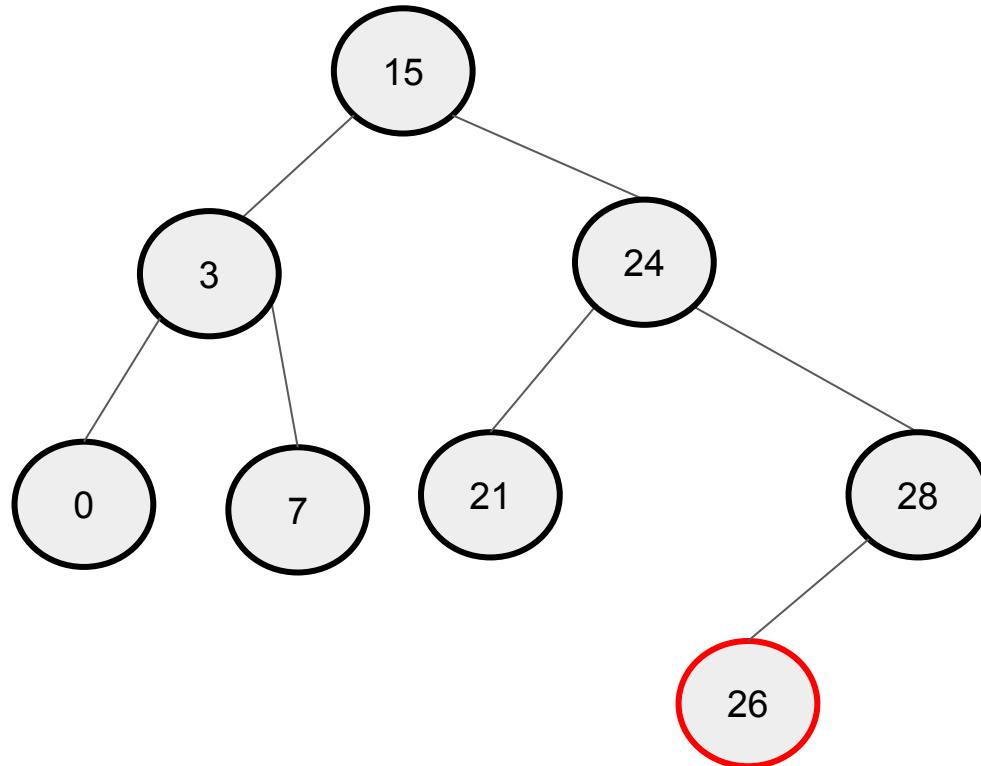
Insert: 15,21,7,24,0,26,3,28,29



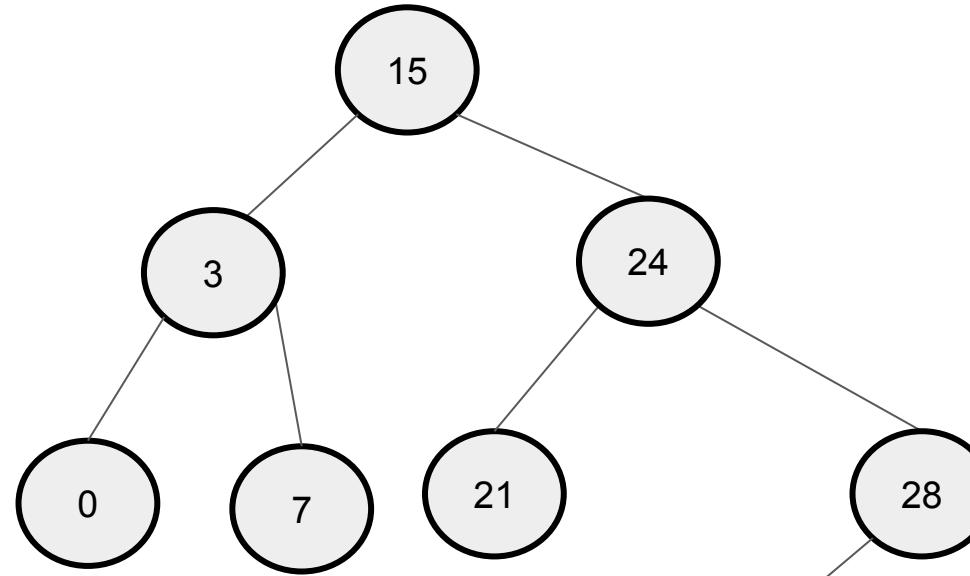
3) Right Insert any parent



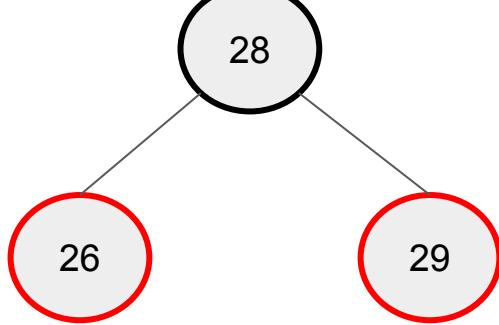
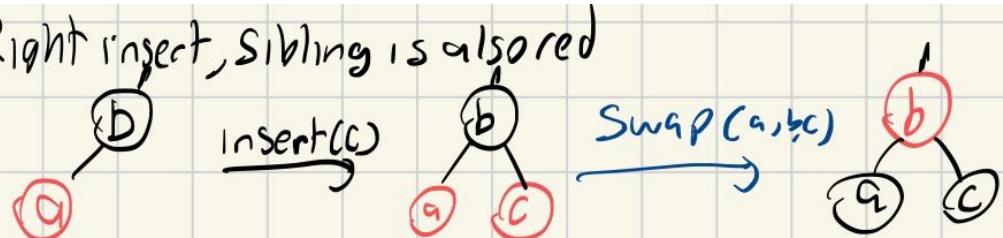
Insert: 15,21,7,24,0,26,3,28,29



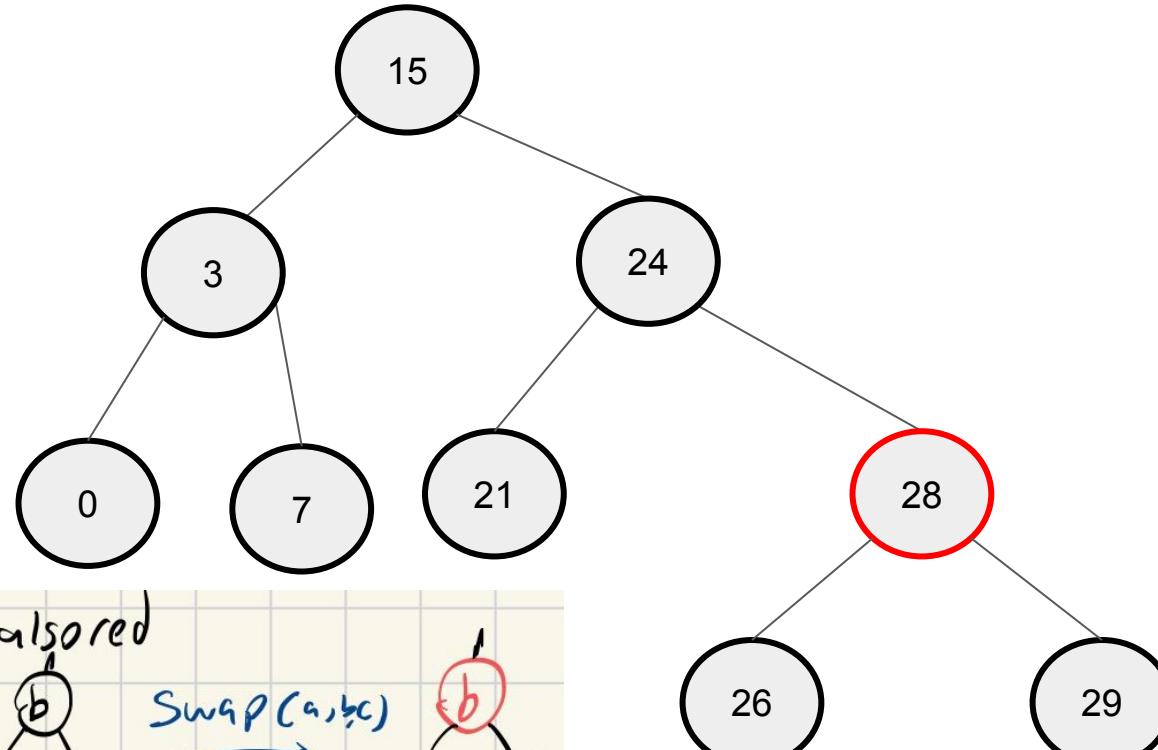
Insert: 15,21,7,24,0,26,3,28,29



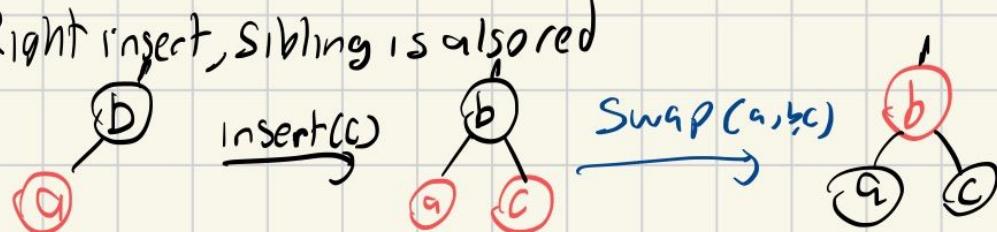
4) Right insert, sibling is also red



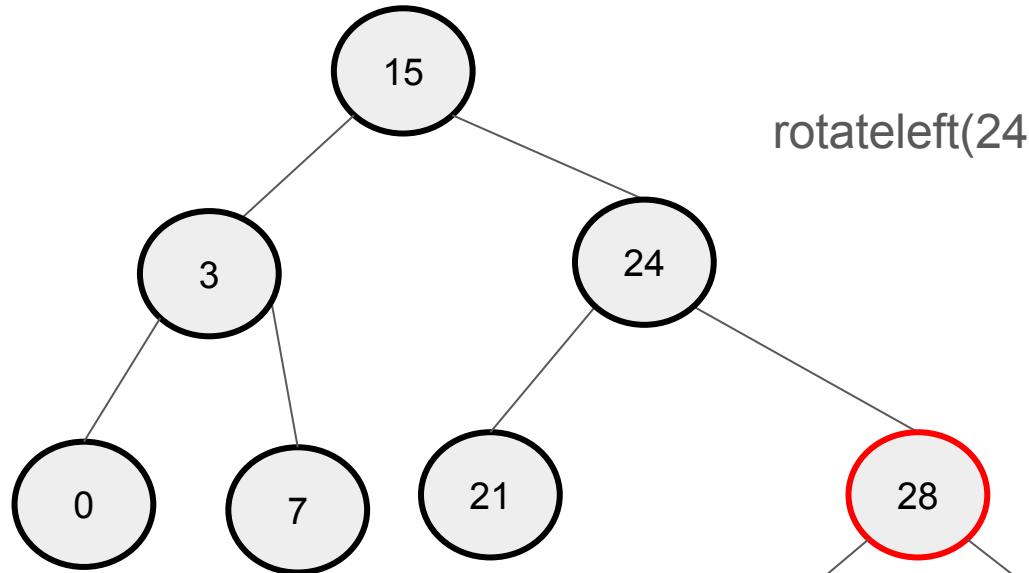
Insert: 15,21,7,24,0,26,3,28,29



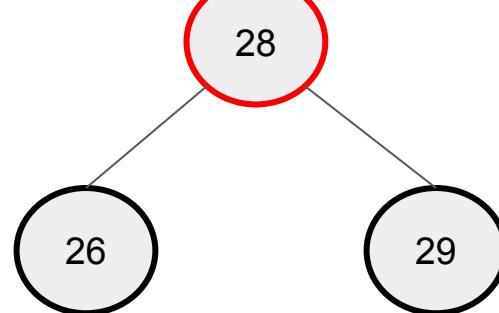
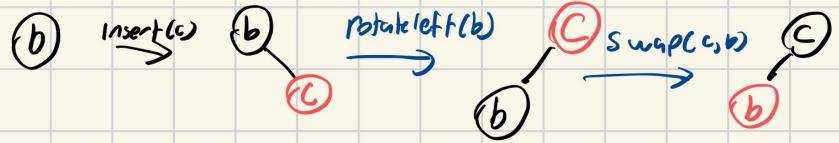
4) Right insert, sibling is also red



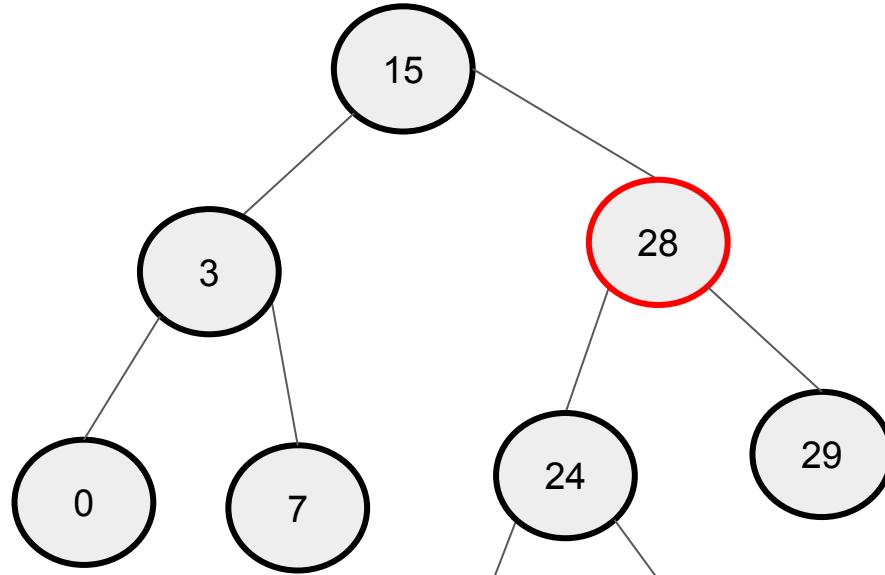
Insert: 15,21,7,24,0,26,3,28,29



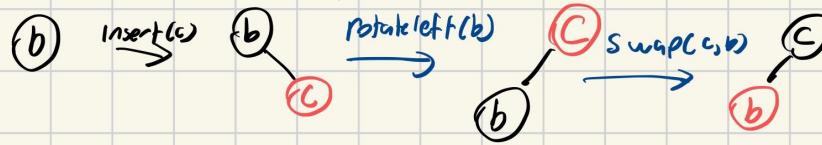
3) Right Insert any parent



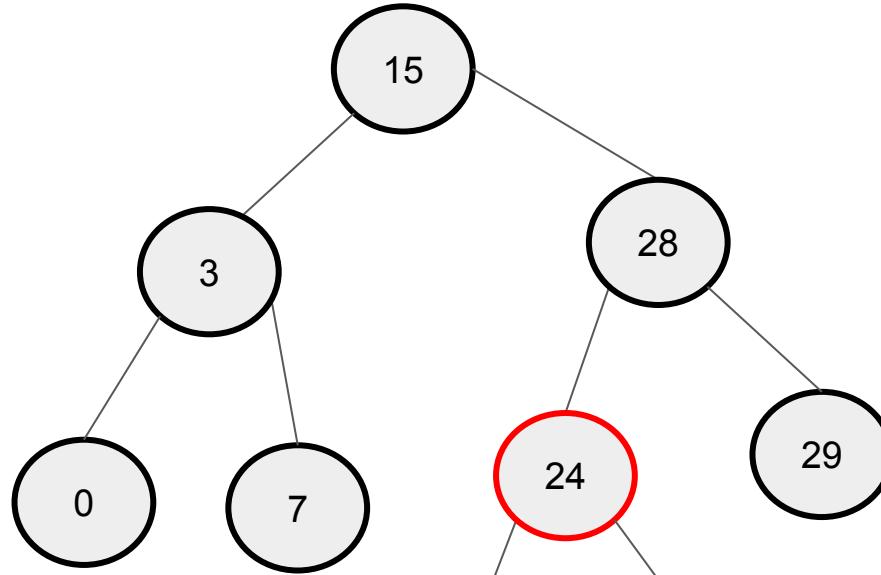
Insert: 15,21,7,24,0,26,3,28,29



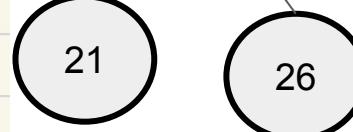
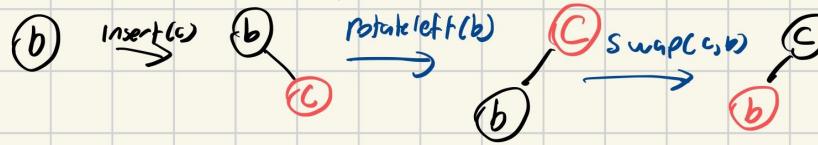
3) Right Insert, any parent



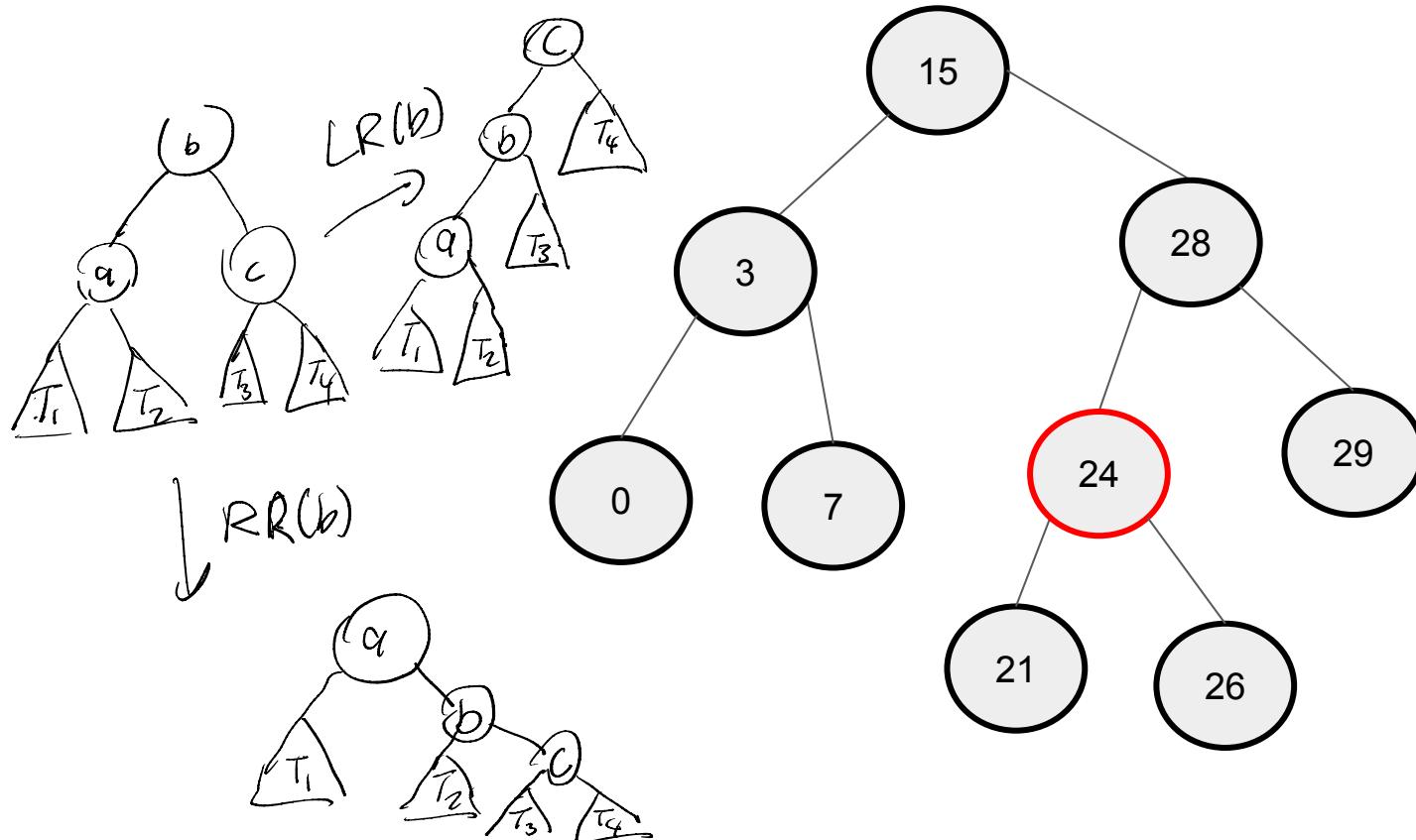
Insert: 15,21,7,24,0,26,3,28,29



3) Right Insert, any parent



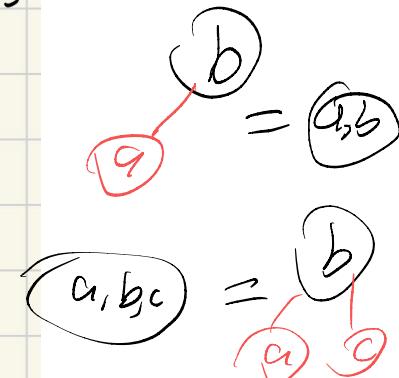
Insert: 15,21,7,24,0,26,3,28,29



Bonus: LLRB Hard to Understand

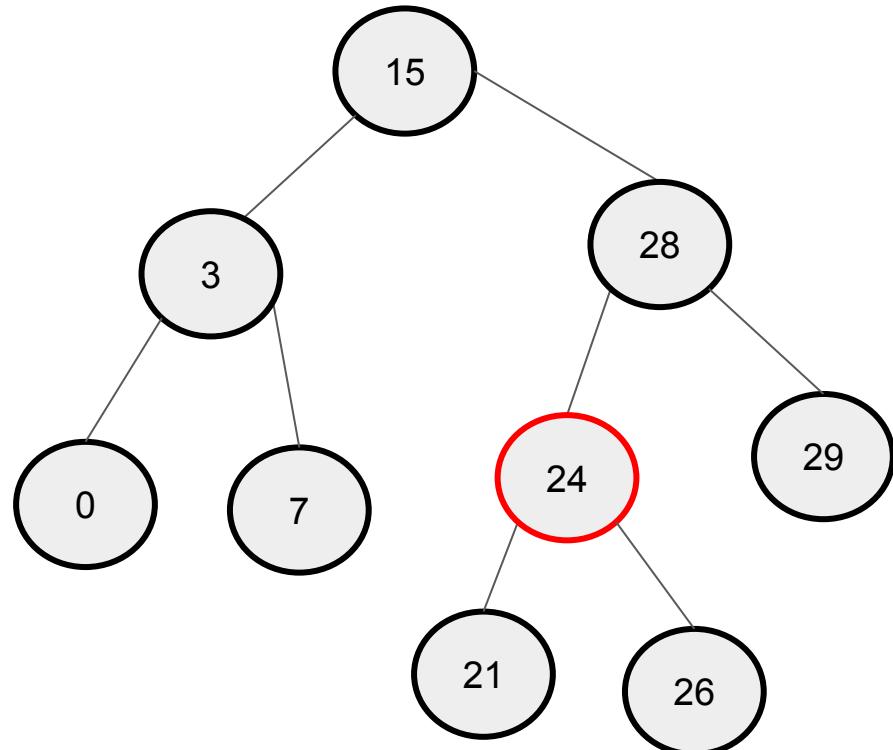
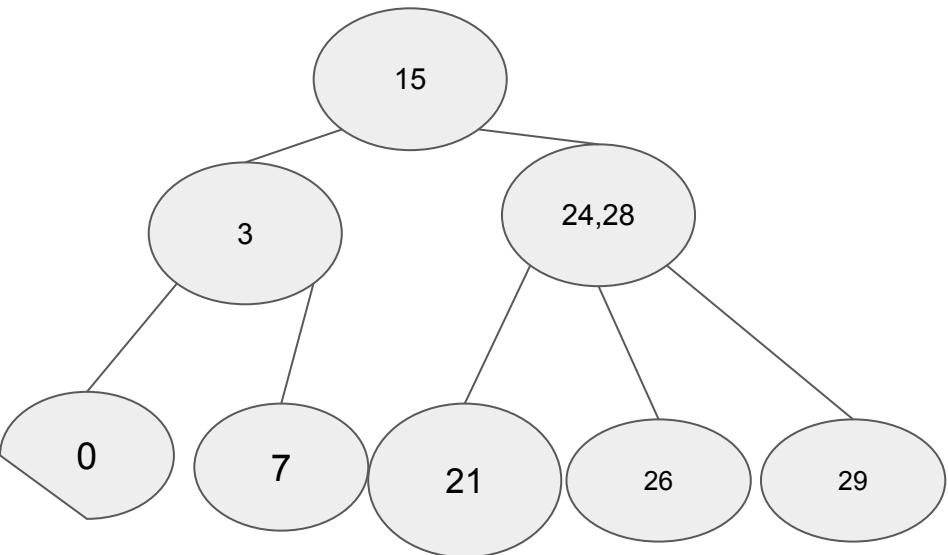
LLRB trees are “the same as” to 2-3 trees

	• binary tree equivalent of 2-3 trees
2-3 Tree	↔ LLRB Tree
• all paths have same depth	• all paths have same # black nodes
• There are 3-nodes	- red nodes, left leaning
• Inserting ‘overstuffed’ nodes	• insert red nodes



Exercise: Compare the insert trace of the 2-3 tree vs the LLRB Tree

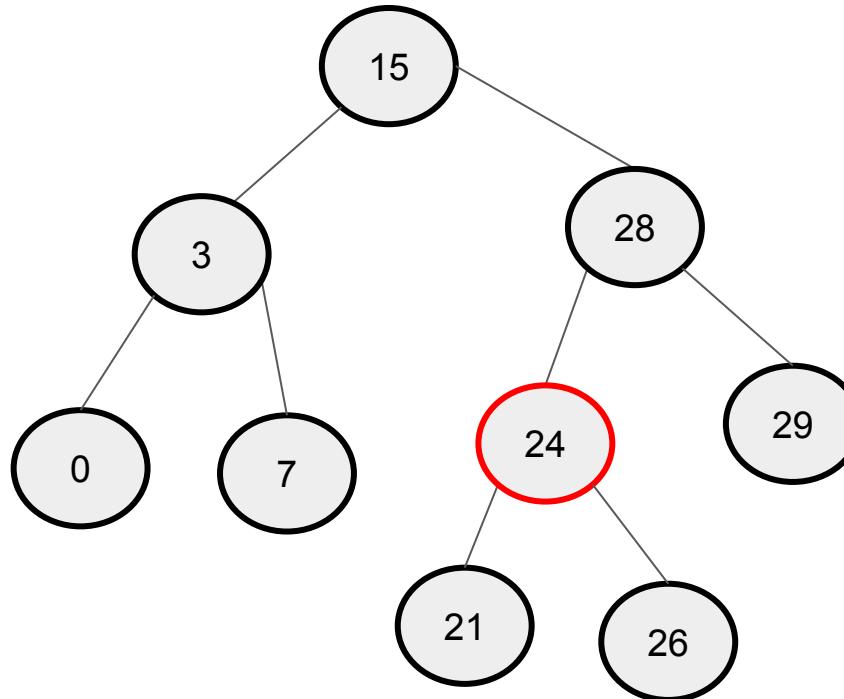
Notice how red nodes == 3-nodes!



Question 3

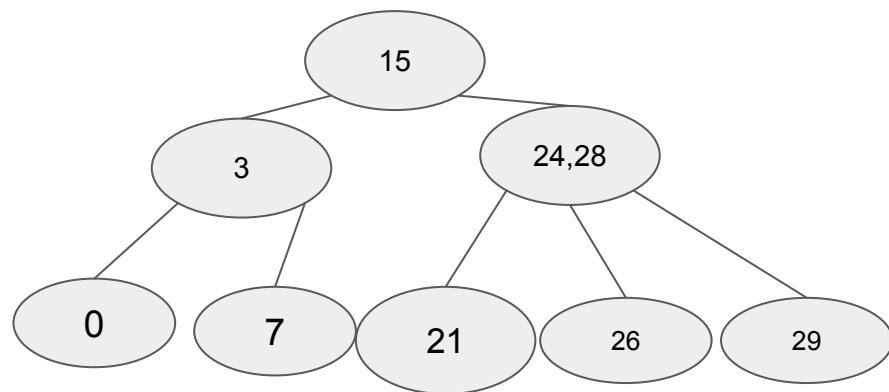
(Deletion) Show intermediate steps of the following questions:

- (1) How to delete 7 in the final 2-3 tree of Q1?
- (2) How to delete 7 in the final Left-Leaning Red-Black tree of Q1?

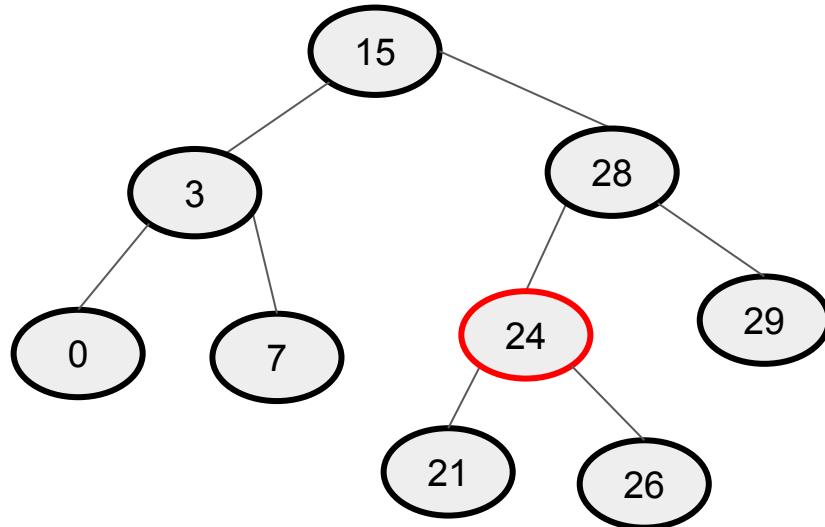


Deletion in LLRB is quite hard..

Idea: Use equivalence between LLRB and 2-3 trees



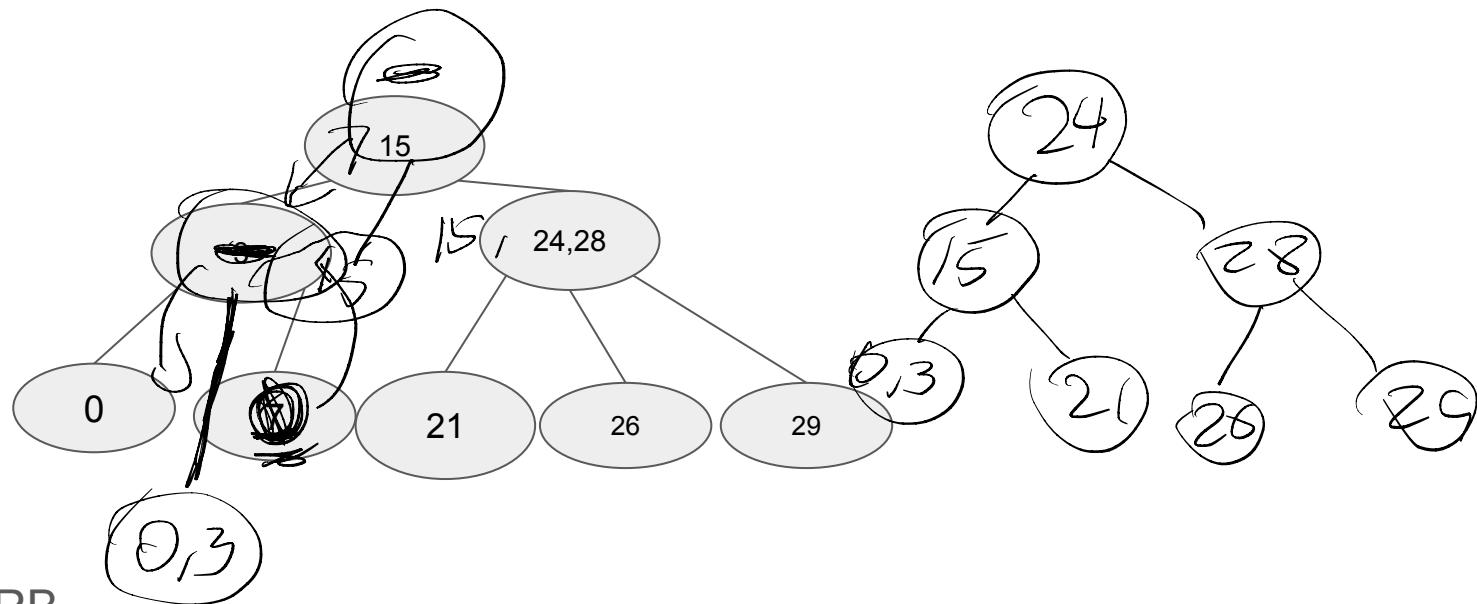
Idea: Use equivalence between LLRB and 2-3 trees



Take the LLRB

- 1) Turn it into a 2-3 tree
- 2) Run the delete algorithm
- 3) Turn it back into a LLRB tree

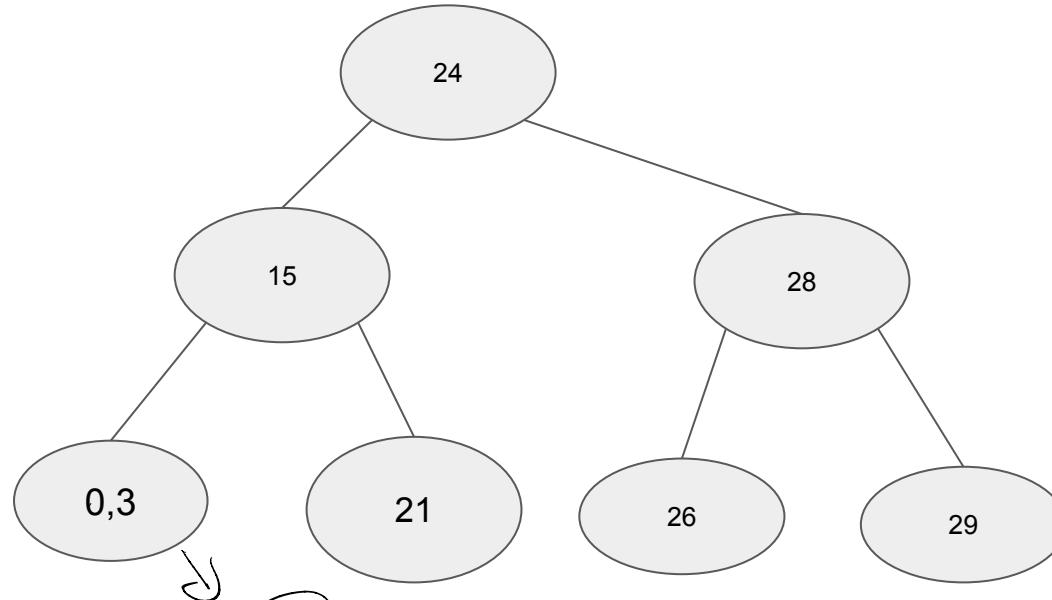
Idea: Use equivalence between LLRB and 2-3 trees



Take the LLRB

- 1) Turn it into a 2-3 tree
- 2) Run the delete algorithm
- 3) Turn it back into a LLRB tree

Idea: Use equivalence between LLRB and 2-3 trees



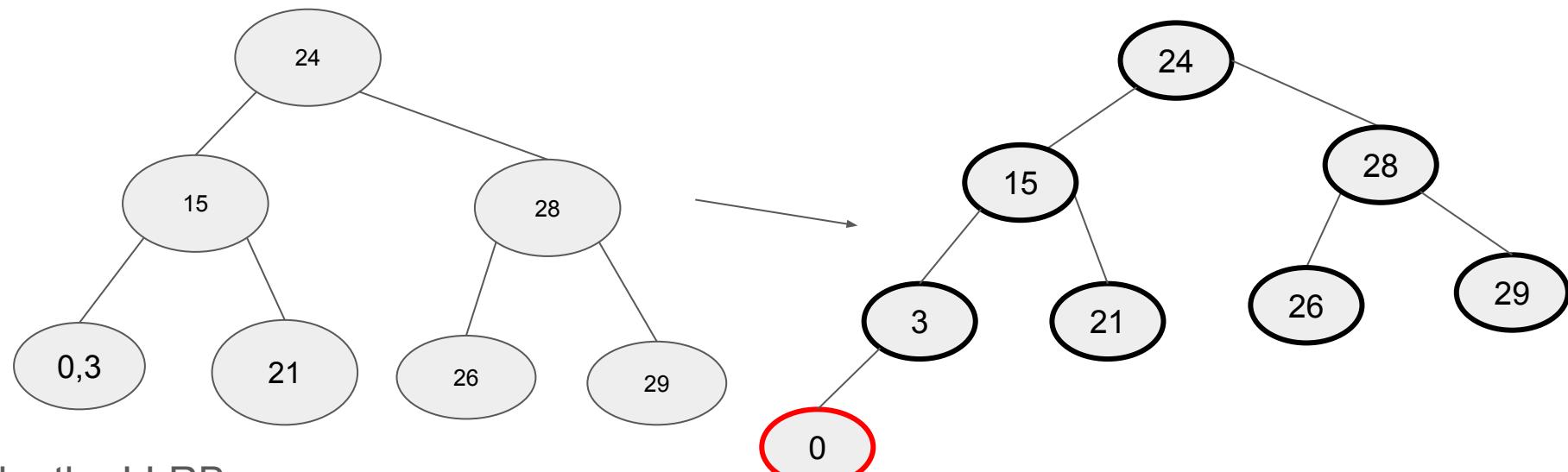
Take the LLRB

- 1) Turn it into a 2-3 tree
- 2) **Run the delete algorithm**
- 3) Turn it back into a LLRB tree

0

3

Idea: Use equivalence between LLRB and 2-3 trees



Take the LLRB

- 1) Turn it into a 2-3 tree
- 2) Run the delete algorithm
- 3) Turn it back into a LLRB tree

Keep things simple :)

Question 2

(Adjacency-matrix Representation)

1. Give an adjacency-matrix representation for a complete binary search tree on 7 vertices numbered from 1 to 7.

What is in the world in an adjacency matrix?

Question 2

(Adjacency-matrix Representation)

1. Give an adjacency-matrix representation for a complete binary search tree on 7 vertices numbered from 1 to 7.

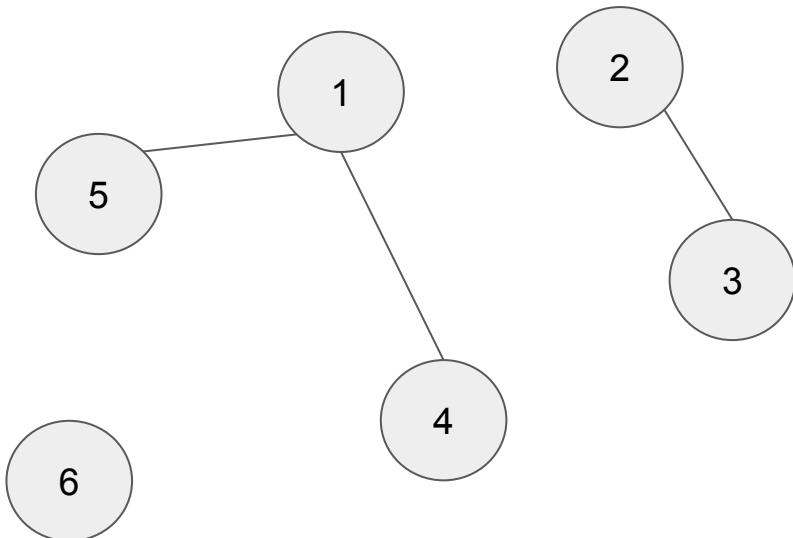
What is in the world in an adjacency matrix?

Adjacency Matrix

$n \times n$

Edges represented in a $|V| \times |V|$ matrix

E.g. if undirected..



$$A = A^T$$

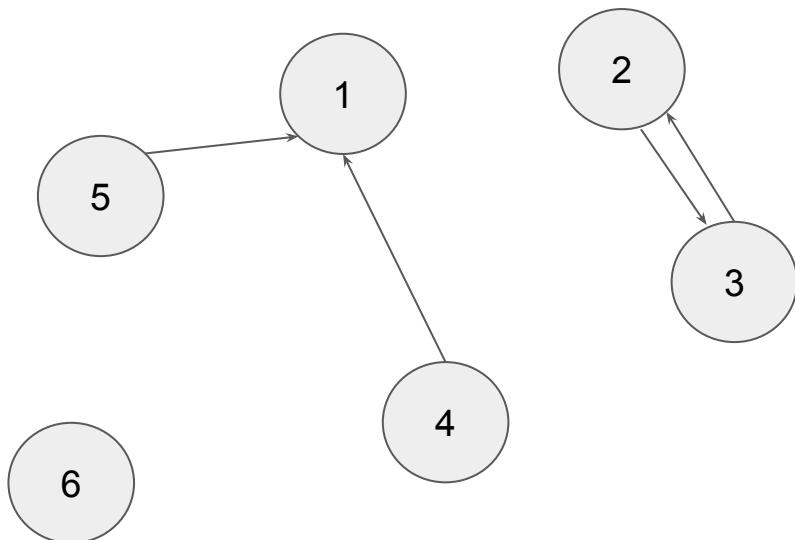
	1	2	3	4	5	6
1	0	0	0	1	1	0
2	0	0	1	0	0	0
3	0	1	0	0	0	0
4	1	0	0	0	0	0
5	1	0	0	0	0	0
6	0	0	0	0	0	0

Adjacency Matrix

$A[i][j]$: There is an edge $i \rightarrow j$
"Row goes to column"

Edges represented in a $|V| \times |V|$ matrix

E.g. if directed..



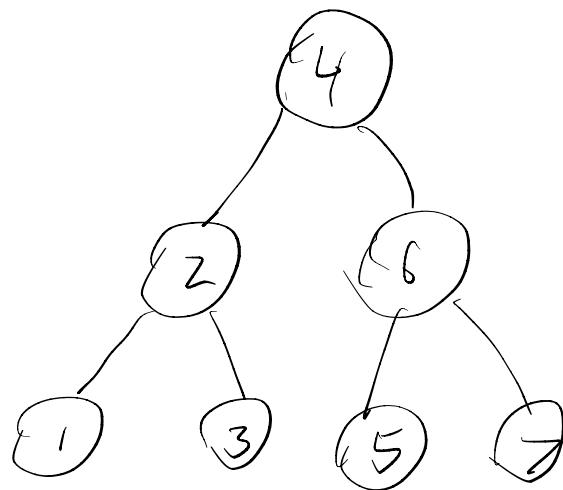
	1	2	3	4	5	6
1	0	0	0	0	0	0
2	0	0	1	0	0	0
3	0	1	0	0	0	0
4	0	0	0	0	0	0
5	1	0	0	0	0	0
6	0	0	0	0	0	0

Question 2

(Adjacency-matrix Representation)

1. Give an adjacency-matrix representation for a complete binary search tree on 7 vertices numbered from 1 to 7.

Someone give me a complete binary search tree

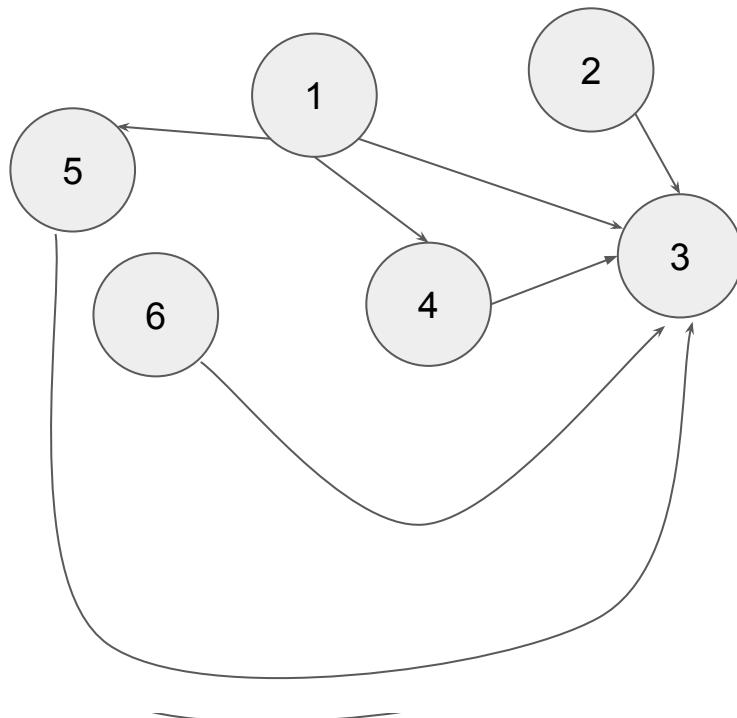


	1	2	3	4	5	6	7
1	0	1	0	0	0	1	0
2	1	0	0	0	0	0	1
3	0	0	1	0	0	0	0
4	0	0	0	1	0	0	0
5	0	0	0	0	1	0	0
6	0	0	0	0	0	1	0
7	0	0	0	0	0	0	1

Question 2

(Adjacency-matrix Representation)

1. Give an adjacency-matrix representation for a complete binary search tree on 7 vertices numbered from 1 to 7.



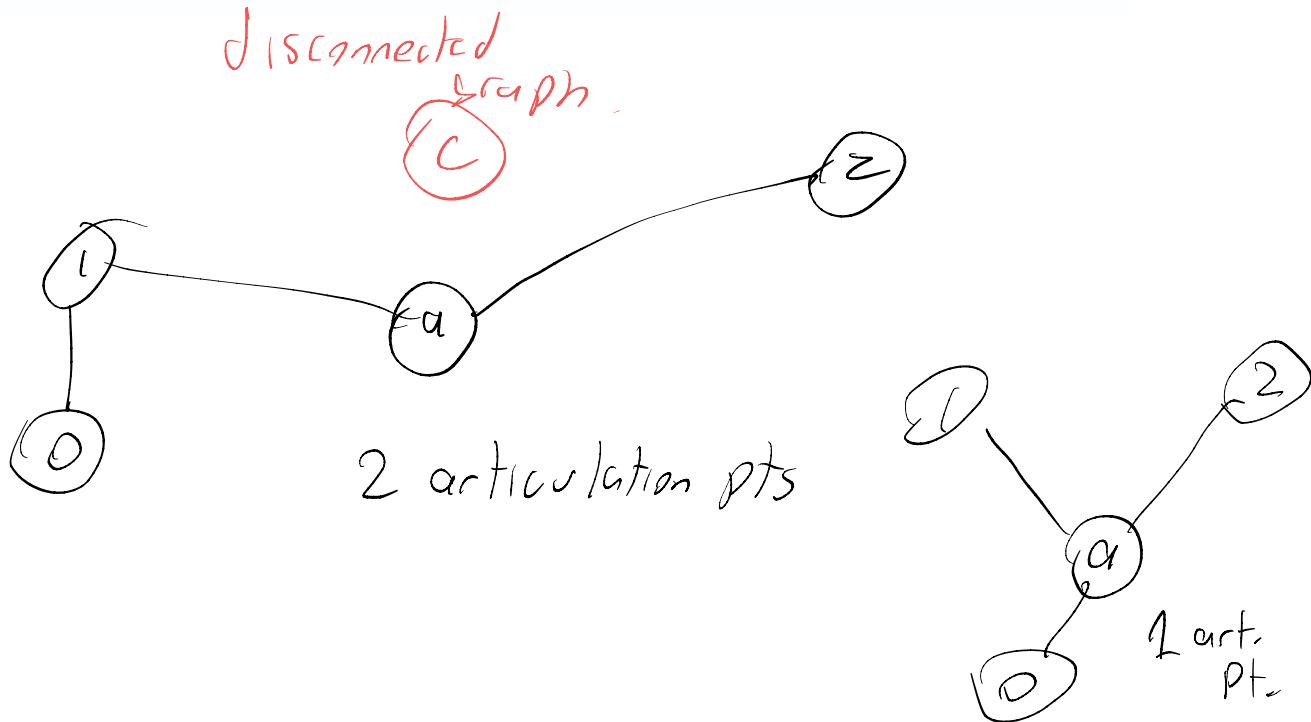
1	2	3	4	5	6
1	0	0	1	1	0
2	0	0	1	0	0
3	0	0	0	0	0
4	0	0	1	0	0
5	0	0	1	0	0
6	0	0	1	0	0

Question 1

(Articulation point)

We define an *articulation point* as a vertex that when removed causes a connected graph to become disconnected. For this problem, we will try to find the articulation points in an undirected graph G .

- (1) How can we efficiently check whether or not a graph is disconnected?
- (2) How to determine if a node u is an articulation point or not?

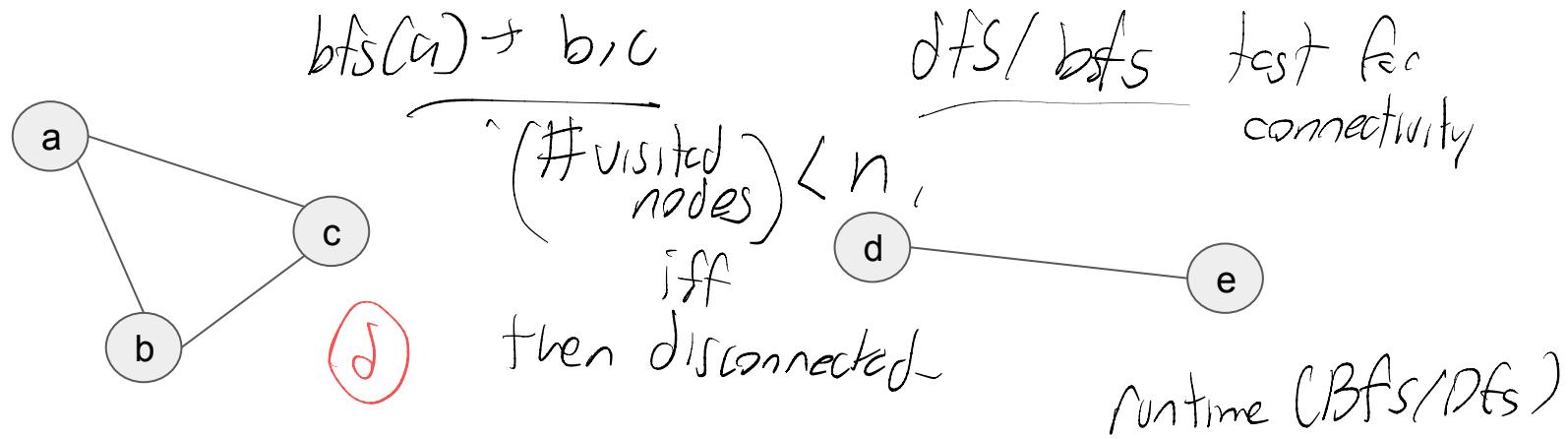


(undirected)

(1) How can we efficiently check whether or not a graph is disconnected?

Two vertices u, v are connected if there is some way to get from $u \rightarrow v$.

Graph is connected if for *all* u, v vertices, u and v are connected



Are there any algorithms that can help us here?

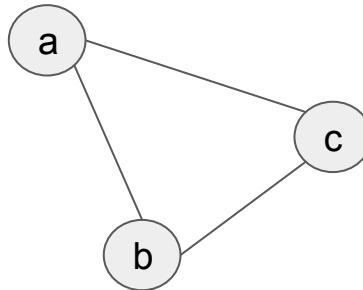
$O(n+m)$

(undirected)

(1) How can we efficiently check whether or not a graph is disconnected?

Two vertices u, v are connected if there is some way to get from $u \rightarrow v$.

Graph is connected if for *all* u, v vertices, u and v are connected



Use BFS/DFS, count the number of vertices visited

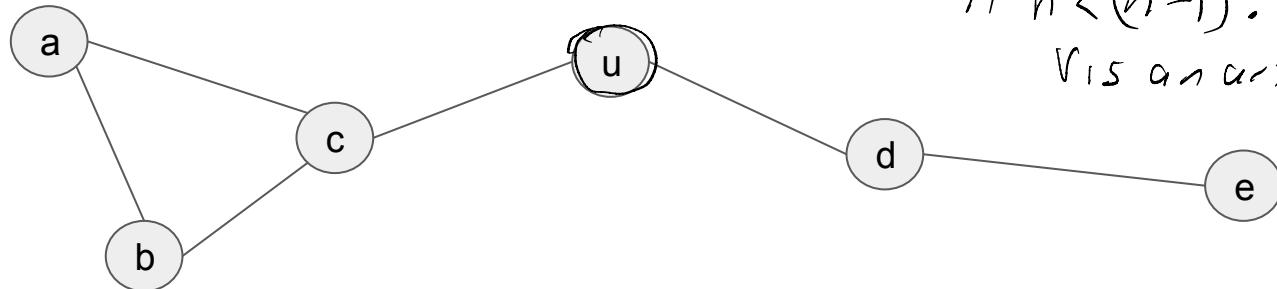
Question 1

(Articulation point)

We define an *articulation point* as a vertex that when removed causes a connected graph to become disconnected. For this problem, we will try to find the articulation points in an undirected graph G .

(2) How to determine if a node u is an articulation point or not?

Any idea from pt 1?



for v in V :

$$G' \leftarrow G - v.$$

$$n' \leftarrow \text{Dfs}(G')$$

if $n' < (n-1)$:

v is an articulation pt,

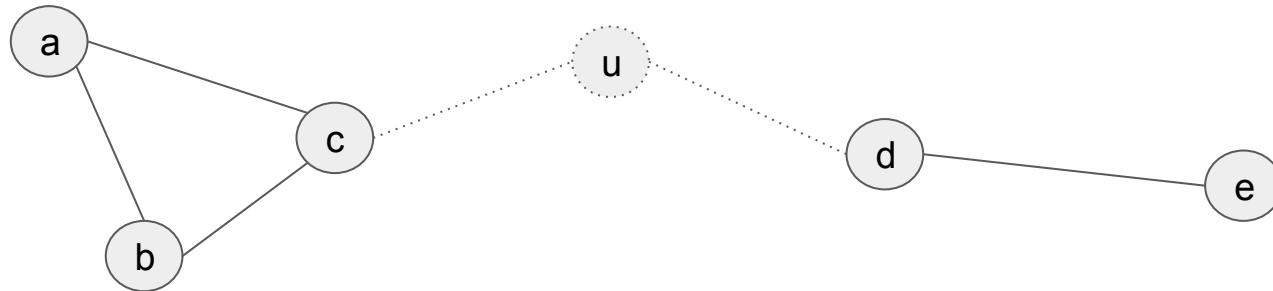
Question 1

(Articulation point)

We define an *articulation point* as a vertex that when removed causes a connected graph to become disconnected. For this problem, we will try to find the articulation points in an undirected graph G .

- (2) How to determine if a node u is an articulation point or not?

Any idea from pt 1? Check connectivity when u is deleted



Exercise: Suppose you wanted to find all articulation points. Can you do so in $O(|V| + |E|)$ time?
Hint: a point is an articulation point iff it is not in a cycle.