

All of us after next week

PSO 14

K-D Trees, Point Trees



Announcements

1. Fill out the instructor feedback surveys (ty 40% of you)
2. Last review session on Friday (time TBD, location TBD (existence TBD))
3. No OH next week
 - a. I will NOT be on duty

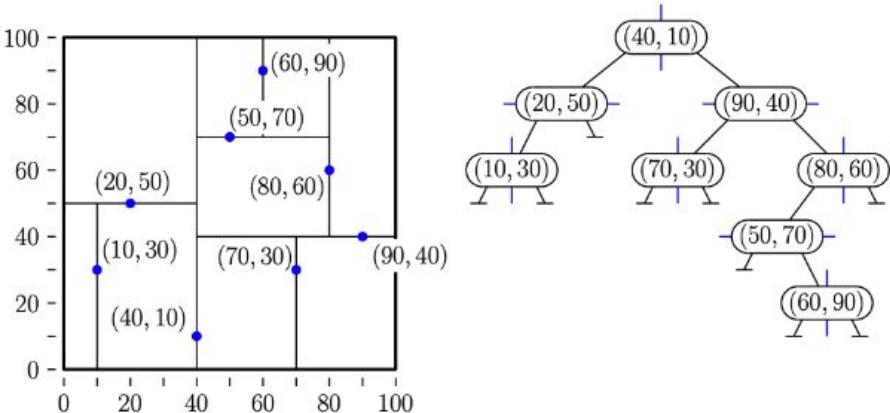
Announcements

1. Fill out the instructor feedback surveys (ty 40% of you)
2. Last review session on Friday (time TBD, location TBD (existence TBD))
3. No OH next week
 - a. I will NOT be on duty but..
 - b. I *might* happen to be sitting around the commons from 12-2PM Sat,Sun,Mon,Tues
 - c. I *might* be open to answering any questions if they *happen* to be asked
 - d. I *might* be hungover

Question 1

(kd trees)

- (1) Consider the kd-tree shown in the figure below. Assume a standard kd-tree where the cutting dimensions alternates between x and y with each level.



- (1) Show the final tree after the operation **insert((70,50))**.
- (2) Starting with the original tree, show the final tree after **delete((40,10))**.
- (3) Starting with the original tree, show the final tree after **delete((80,60))**.

Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

- (a) What is the expected depth $D(n)$ of the QuadTree?
- (b) What is the total number of **leaf nodes** in the tree in terms of n ?
- (c) If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

But first..

Question 3

(Backward pattern matching)

The Boyer-Moore algorithm is based upon backward pattern matching. Let us do a simple review via the following questions:

1. Run Boyer-Moore algorithm in the following worst-case scenario:

$$T := \underbrace{aaa \cdots a}_{9} \quad \text{and} \quad P := baaaaa.$$

Boyer-Moore: Iteratively compare pattern P with target, going backward

T	a	a	a	a	a	a	a	a	a
P	b	a	a	a	a	a			

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T[0] does not equal P[0]! Next steps..

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T[0] does not equal P[0]! Next steps.. We mismatched on target **a**

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T	a	a	a	a	a	a	a	a	a
P	b	a	a	a	a	a			

T[0] does not equal P[0]! Next steps.. We mismatched on target **a**
The last occurrence of pattern **a**

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T	a	a	a	a	a	a	a	a	a
P	b	a	a	a	a	a			

Move P (to align target **a** with pattern **a**) OR (one after target mismatch)

Whichever moves P the *least* amount – in this ex. We move one after mismatch

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T	a	a	a	a	a	a	a	a	a
P		b	a	a	a	a	a		

Fast forward..

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Boyer-Moore: Iteratively compare pattern P with target, going backward

T	a	a	a	a	a	a	a	a	a
P		b	a	a	a	a	a		

Fast forward.. Same mismatch, jump 1

Question 3

(Backward pattern matching)

The Boyer-Moore algorithm is based upon backward pattern matching. Let us do a simple review via the following questions:

1. Run Boyer-Moore algorithm in the following worst-case scenario:

$$T := \underbrace{aaa \cdots a}_{9} \quad \text{and} \quad P := baaaa.$$

Boyer-Moore: Iteratively compare pattern P with target, going backward

T	a	a	a	a	a	a	a	a	a
P			b	a	a	a	a	a	

Same thing will happen 1 more time

The example from last time

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

The example from last time

Mismatch here

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

The example from last time

We mismatched on target **X**

The last occurrence of pattern **X**

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

Move P (to align target **X** with pattern **X**) OR (one after target mismatch)
Whichever moves P the *least* amount

The example from last time

We mismatched on target **X**

The last occurrence of pattern **X**

T	O	O	X	X	X	X	O	O	O
P		O	X	X	X	X	O	O	O

Move P (to align target **X** with pattern **X**) OR (one after target mismatch)
Whichever moves P the *least* amount

Last note: If there is no last occurrence of the target mismatch, default to one after mismatch

View this at your leisure – a longer example

T	a	b	c	a	b	x	a	x	x	x	x	x	b
P	x	x	b										
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P				x	x	b							
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P				x	x	b							
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					

The green is a comparison made

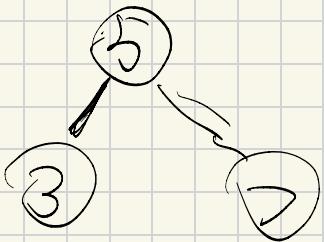
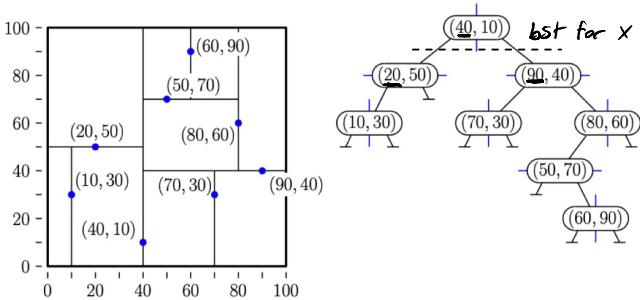
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T	a	b	c	a	b	x	a	x	x	x	x	x	b
P	x	x	b										
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P			x	x	b								
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P			x	x	b								
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
T	a	b	c	a	b	x	a	x	x	x	x	x	b
P						x	x	b					
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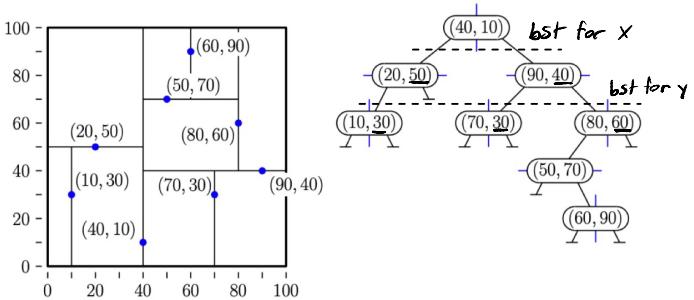
c not in the pattern
We move one after
(In this case, big jump)

The green is a comparison made

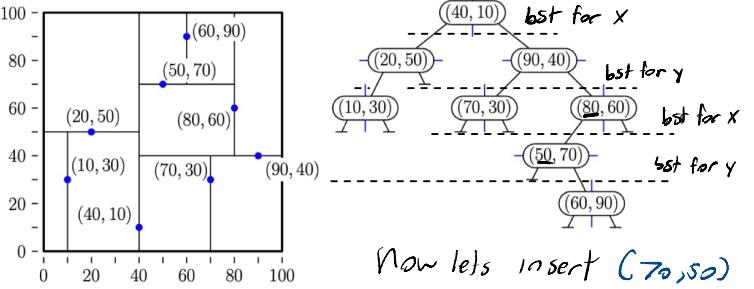
(2) (kd-trees) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation $\text{insert}((70,50))$.



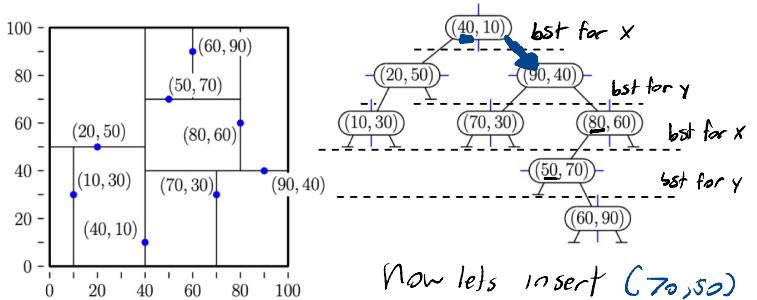
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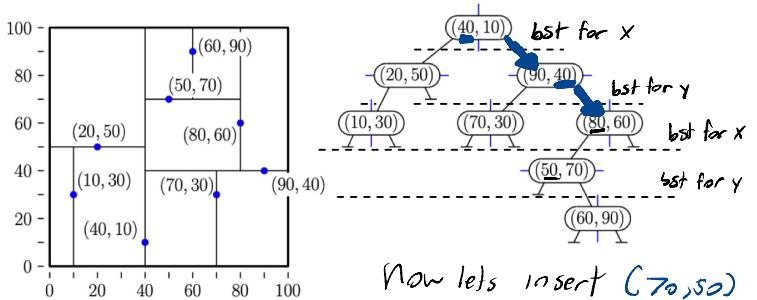


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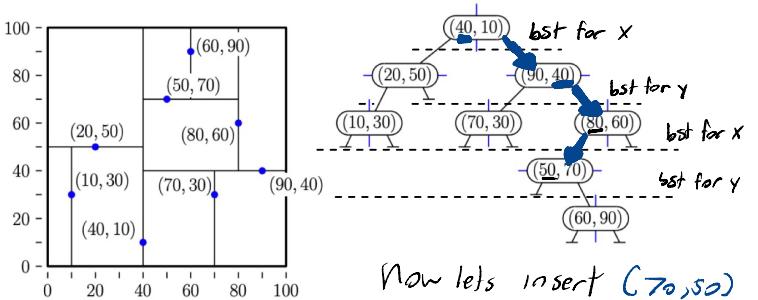


6

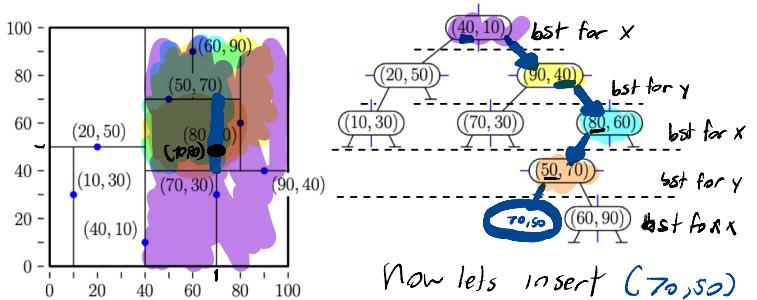
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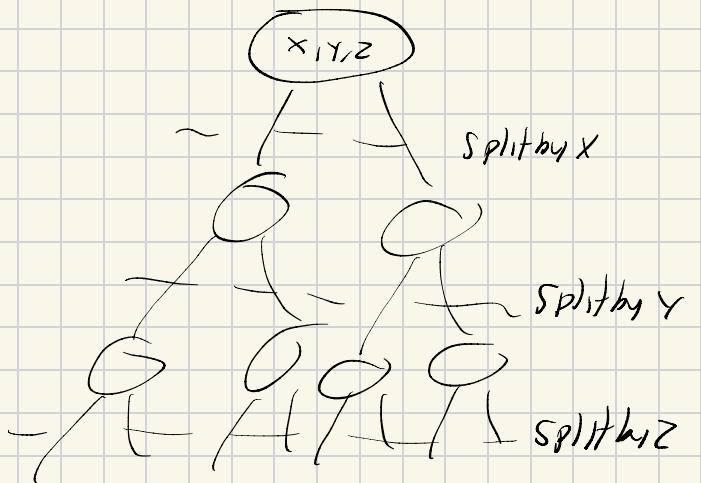


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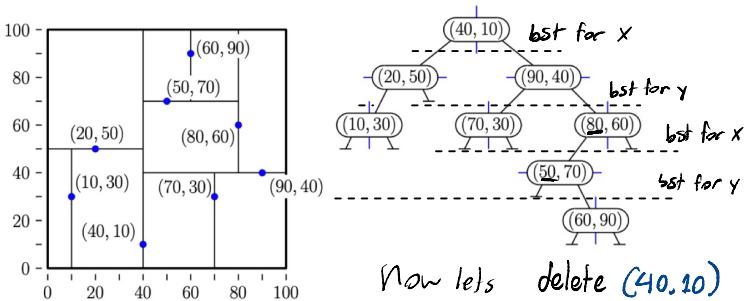


2d tree

3d tree

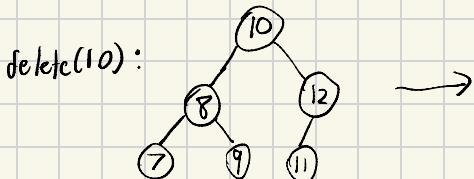


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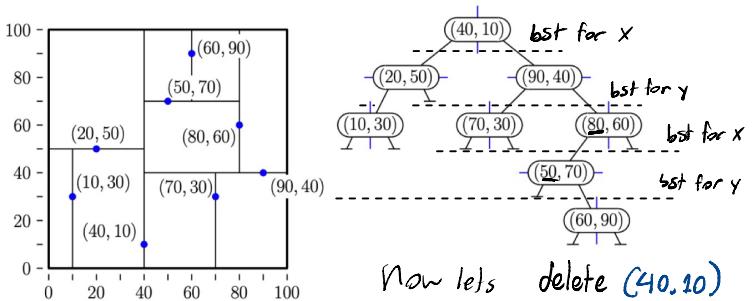


Recall deletion in normal BST

- Find Predecessor/Successor leaf to replace
ex:

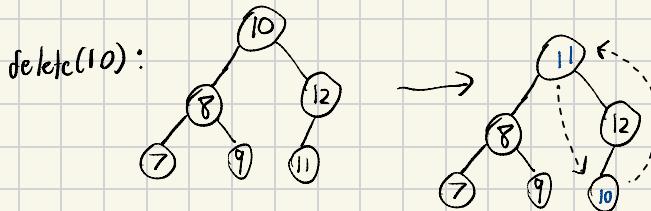


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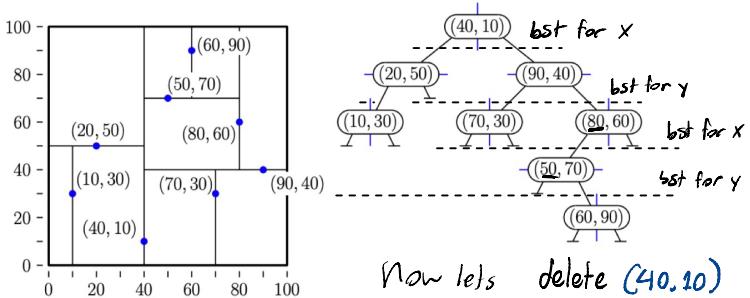


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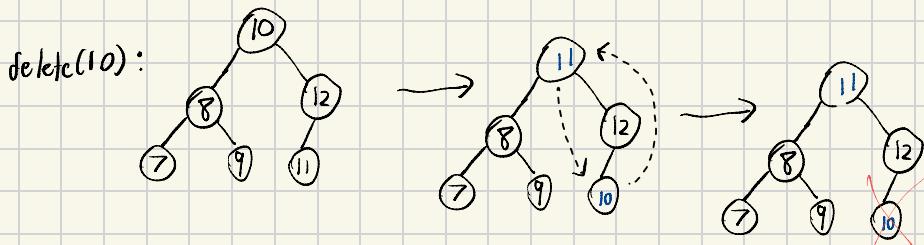


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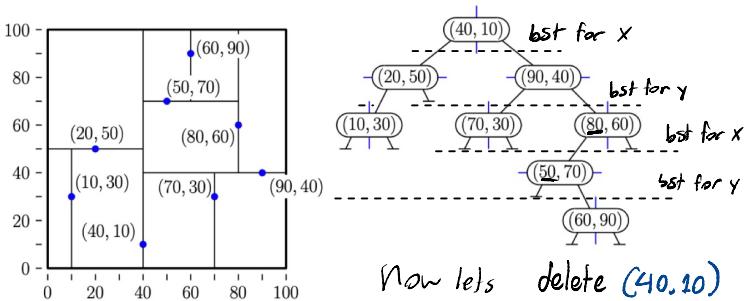
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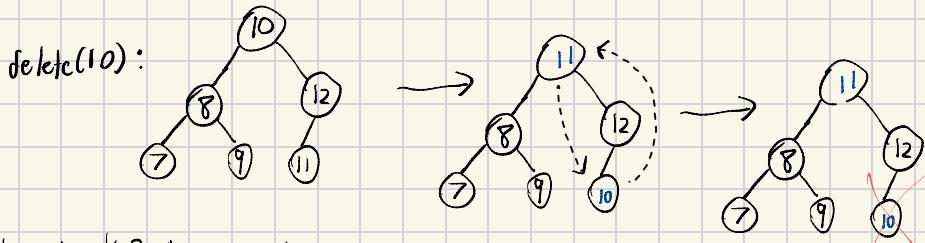
(If not a leaf continue recursively)

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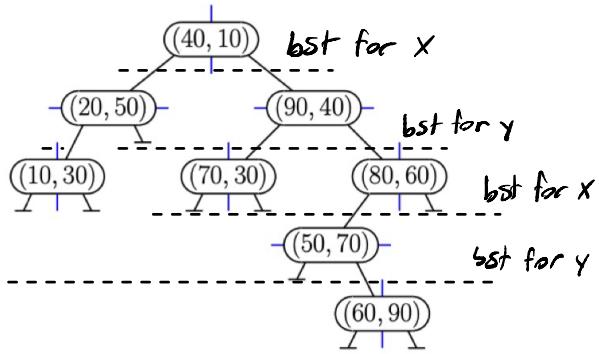
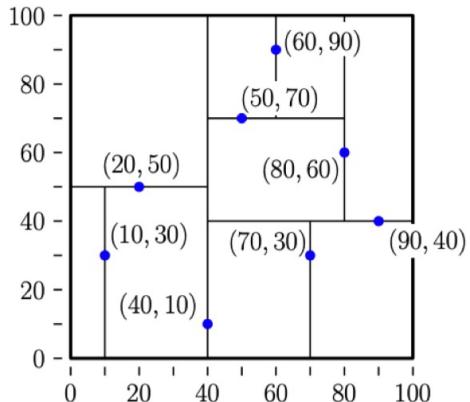
Recall deletion in normal BST

- Find Predecessor/Successor leaf to replace
ex:



Deletion in KD-trees is the same,
except you choose successor/predecessor based on dimension

(2) (kd-trees) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation `insert((70,50))`.

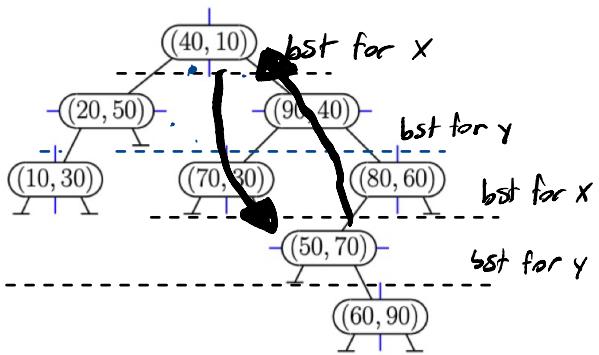
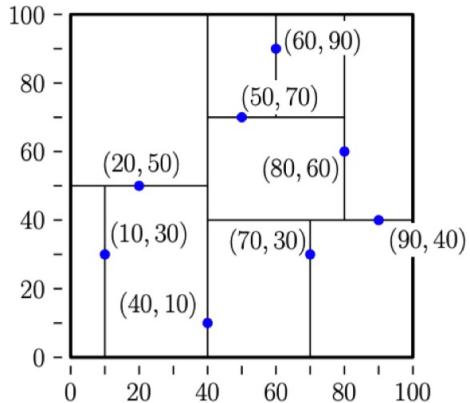


Now lets delete (40, 10)

1) Find successor for $X=40$

40, 90
90, 40

(2) (kd-trees) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation `insert((70,50))`.

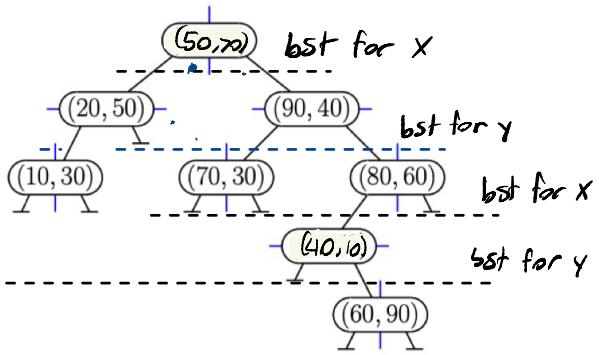
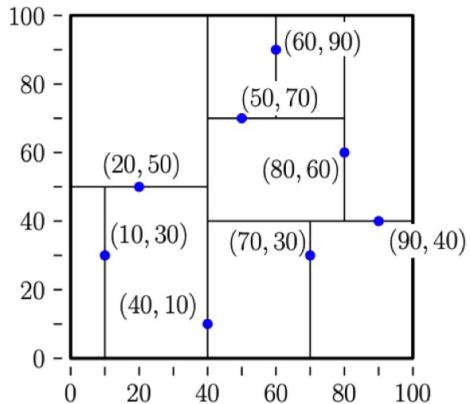


Now lets delete (40, 10)

2) Swap with (40, 10)

[Think about why this preserves k-d order]

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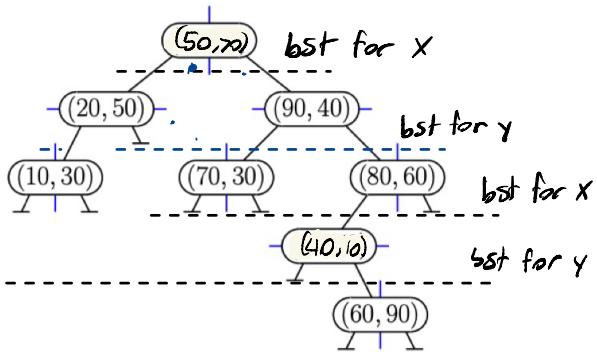
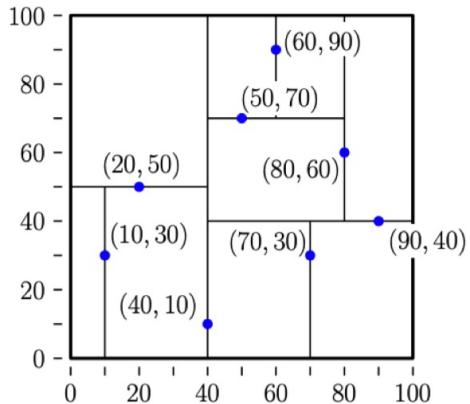


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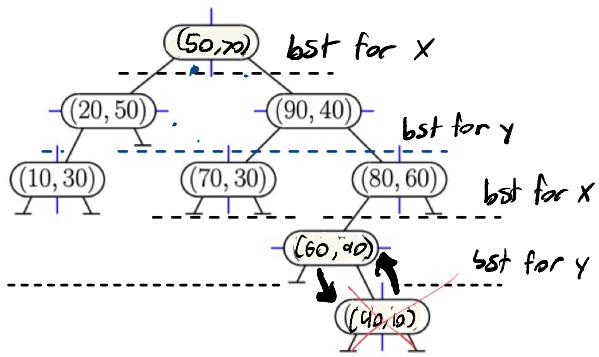
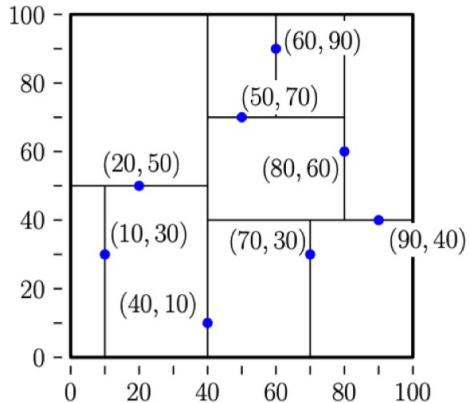
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Now lets delete (40,10)

2.2) Not at leaf yet, find successor in $y=10$ and swap.

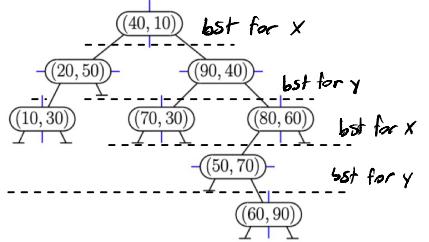
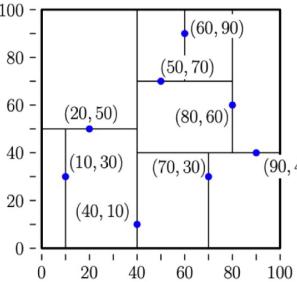
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Now let's delete (80,60).

Left as
Exercise

NOTE: If no successor, find a predecessor instead and vice versa.

Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

- (a) What is the expected **depth** $D(n)$ of the QuadTree?
- (b) What is the total number of **leaf nodes** in the tree in terms of n ?
- (c) If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

for now

I'm going to ignore this question[^] and go over how we insert into a quad tree.

Question 1

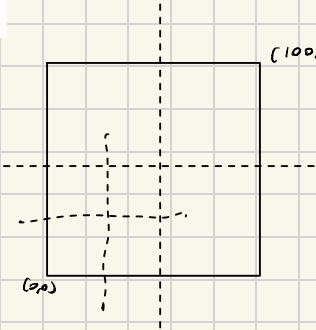
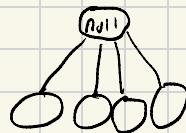
(1) (Quadtrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

Suppose $\text{root.pt} = \text{null}$

$\text{root.region} = (0, 0, 100, 100)$

$\text{root.children} = []$

1) Insert (35,40)



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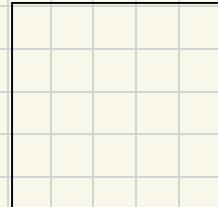
$\text{root.region} = (0, 0, 100, 100)$

$\text{root.children} = []$

1) Insert (35,40)

```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

(null)



(top)

(100, 100)

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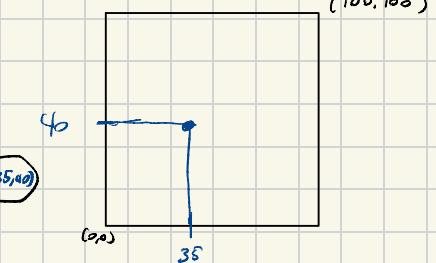
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1) Insert (35,40)

root.pt = (35,40)

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root.children = []



Question 1

(1) (Quadtrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

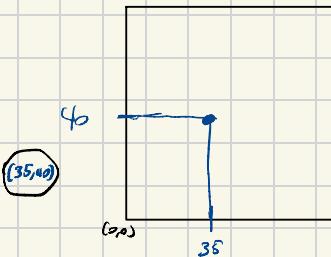
1) Insert (35,40)

2) Insert (50,10)

root.pt = (35,40)

root.region = (0,0,100,100)

root.children = []



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
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```

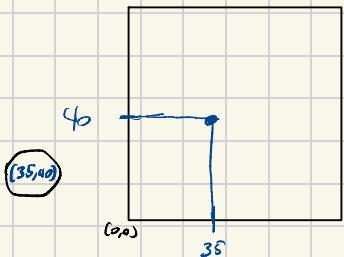
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  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm

```

$\text{root.pt} = (35, 40)$

$\text{root.region} = (0, 0, 100, 100)$

$\text{root.children} = []$



```

algorithm subdivide(root:node)
  xmin, ymin, xmax, ymax ← root.region
  xmid ← (xmin + xmax) / 2
  ymid ← (ymin + ymax) / 2
  root.children ← [
    Node(Region(xmin, ymid, xmid, ymax)), //NW
    Node(Region(xmid, ymid, xmax, ymax)), //NE
    Node(Region(xmin, ymin, xmid, ymid)), //SW
    Node(Region(xmid, ymin, xmax, ymid))] //SE
  if root.point is not Null then
    P ← root.point
    root.point ← Null
    for each quadrant in root.children do
      if insert(P, quadrant) then
        return
  end if
end algorithm

```

Question 1

(1) (Quadtree) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

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```

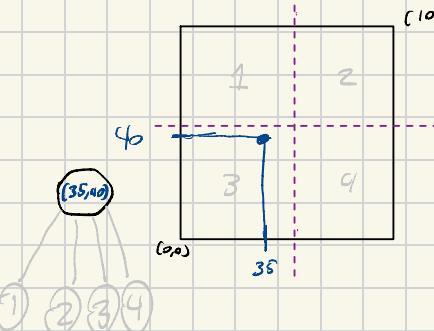
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    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm

```

root.pt = (35,40)

root.region = (0,0,100,100)

root.children = []



```

algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← 50
    ymid ← 50
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
        P ← root.point
        root.point ← Null
        for each quadrant in root.children do
            if insert(P, quadrant) then
                return

```

Question 1

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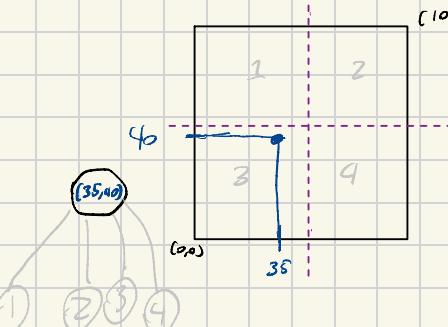
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```

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        Node(Region(xmin, ymid, xmid, ymax)), //NW
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        Node(Region(xmin, ymin, xmid, ymid)), //SW
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    if root.point is not Null then
        P ← root.point
        root.point ← Null
        for each quadrant in root.children do
            if insert(P, quadrant) then
                return

```

What does this intuitively do?

Question 1

(1) (Quadtree) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

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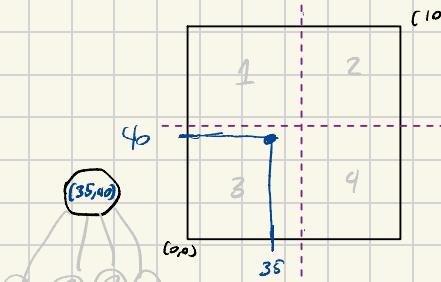
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    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm

```

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```

algorithm subdivide(root:node)
  xmin, ymin, xmax, ymax ← root.region
  xmid ← 50
  ymid ← 50
  root.children ← [
    Node(Region(xmin, ymid, xmid, ymax)), //NW
    Node(Region(xmid, ymid, xmax, ymax)), //NE
    Node(Region(xmin, ymin, xmid, ymid)), //SW
    Node(Region(xmid, ymin, xmax, ymid))] //SE
  if root.point is not Null then
    insert root into the quadrant it belongs
    (and set original root pt to null)

```

Question 1

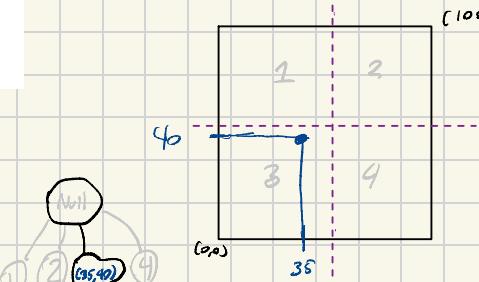
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```

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    insert root into the quadrant it belongs
```



- We do this on the second insert!
- Still need to insert (50,10)

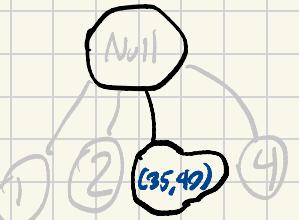
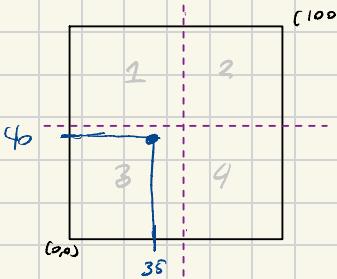
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```



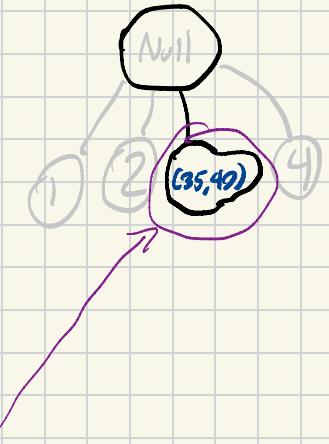
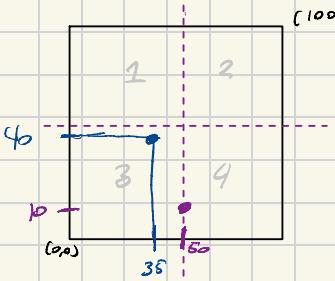
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• Insert into quadrant 3

Question 1

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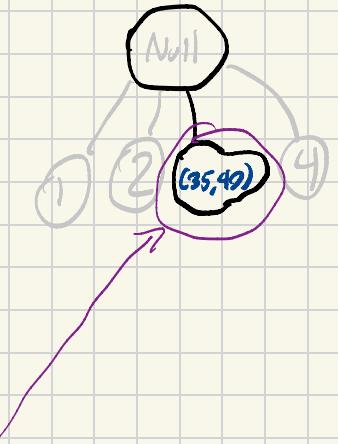
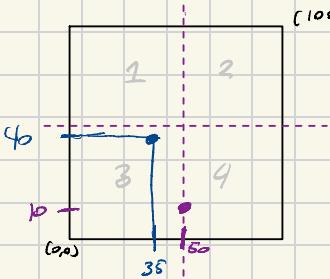
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```



- Insert into quadrant 3
- this node has no children,
need to further subdivide, then move into new subquadrant

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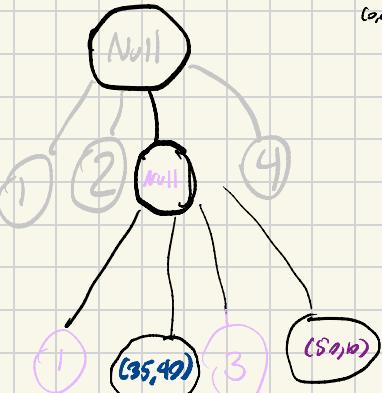
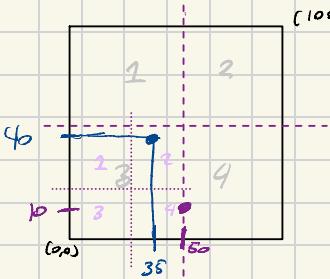
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```

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```



- Insert into quadrant 3
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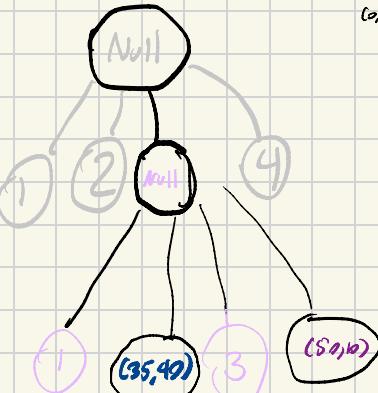
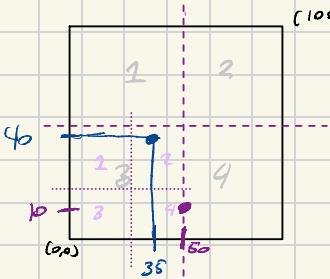
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```
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end algorithm
```



Take away

Insert always inserts at a leaf.

Always for $\log_4 n$ height

Question 1

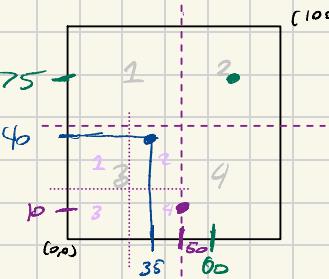
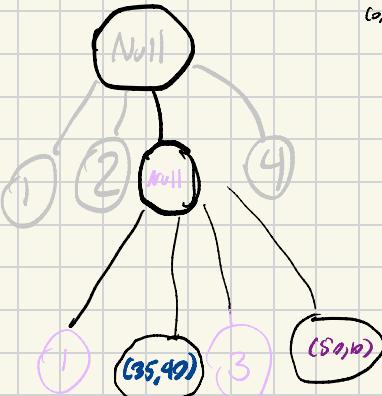
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1) Insert (35,40)

2) Insert (50,10)

3) Insert (60,75)

```
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Question 1

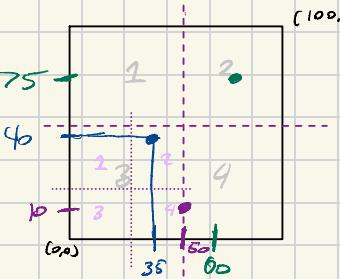
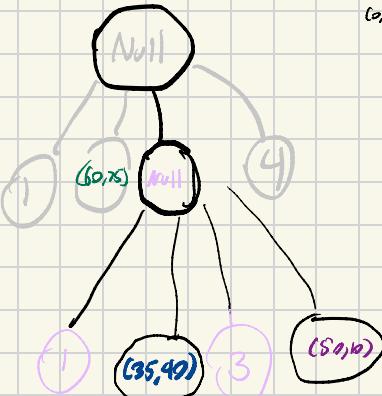
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1) Insert (35,40)

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3) Insert (60,75)

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```



Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

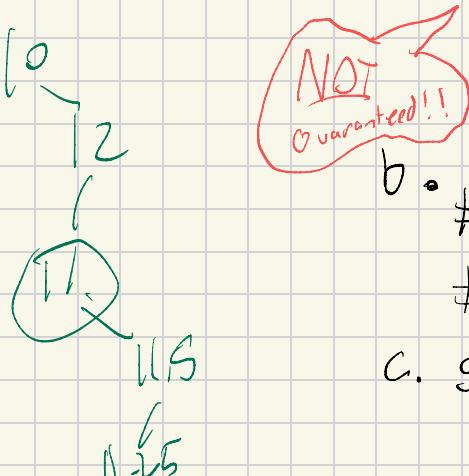
- What is the expected depth $D(n)$ of the QuadTree?
- What is the total number of **leaf nodes** in the tree in terms of n ?
- If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

bst

$2^d \approx n$

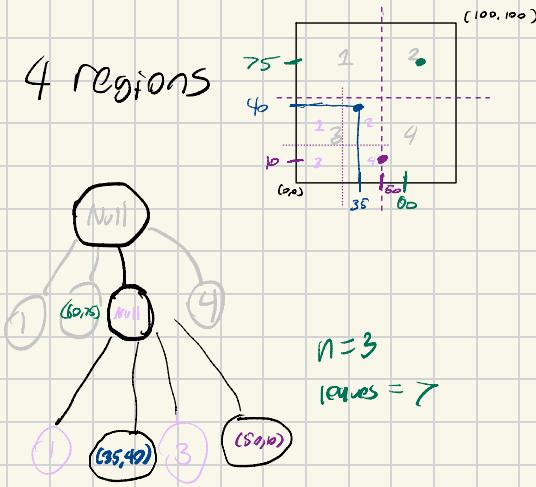
What is the takeaway of the example for this question ?

a) When we evenly subdivide into 4 regions



$4^d \approx n$

- b.
- # Non-null leaves = n
 - # leaves $\approx n$
 - c. geometry with question.



Poll

Data structure for hangout

(Guest performance by Fibonacci hang)

What do YOU want to see at the last practice session?

- Asymptotic analysis 50.1%
- Graph problems 40.1%

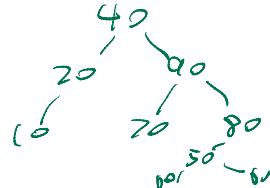
Exam 1
Sorting scenarios

That ONE topo sort question.

b trees hashtable + Z

red black trees (deletion ?) pq

Do you want us to keep doing Mult. Choice or do more Free response-y questions?



2sat reduction
dp ?