## My Tiny BASIC Extended

# Concurrent Tasks, Irq Handling, Tokenized IL (mytb)

Version 1.1.20 IRQ/TASKING/IPC/Compiled Line Numbers Support Extended by <u>JustLostIntime@yahoo.com</u> <u>Originally by bob@corshamtech.com</u>

Ever since reading the first year of Dr. Dobb's Journal of Computer Calisthenics and Orthodontia (yes, that was the original name of the magazine), I wanted to write a tiny BASIC interpreter using the intermediate language (IL) method. The first couple years of DDJ printed source code to several BASICs but none of them used IL.

Well, the idea was always in the back of my mind, so one day I re-read the articles, found some good web pages about the topic, and started writing my own in 6502 assembly language. While it can easily be argued that this was not a good use of my time, it was fun and very satisfying, reminding me of the days when I dreamed of having a high level language on my KIM-1 computer.

Now supports both upper and lower case characters for commands and variables.

So here it is, Concurrent Tiny BASIC. It's not as tiny as it could be, but it does have some support for program storage/retrieval. It has support for the base KIM-1 computer, the xKIM monitor by Corsham Technologies, and the CTMON65 monitor by Corsham Technologies. The source is on github:

Original tiny basic @ https://github.com/CorshamTech/6502-Tiny-BASIC

For the Concurrent version supporting IRQ and Task extensions, Tokens: <a href="https://github/JustLostInTime/em6502">https://github/JustLostInTime/em6502</a>

This version of not so tiny basic is useful to learn about multi threaded and multi tasking system and the basic functionality they provide. Besides it lets you run multi task programs on your Corsham CT6502 SS-50 system and Kim system with at least 32K.

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## **Numbers and Variables**

There are 26 integer variables named A to Z., An exit code Variable for tasks ^ and finally an @[location-pointer] may be used to access any of the unused memory locations as integers while the program is running.

Variables may be subscripted var[<expr>]. Any Variable may be subscripted up to end of the variable SET. For example A may have 1 to 26 subscripts representing A-Z in the variable set. B may have subscripts 1 to 25 representing b to z and so forth up to Z which may only have a single subscript 1. Subscripts start at 1. Tasks may access other tasks Variables with the following notation <PID-Variable>!<Variable-Name>[Subscript] . Subscript is optional. A tasks exit code is available using PID!^
Examples:

```
    Access Tasks Variable A: a = task(1000):? a!a
    Access Tasks EXIT code: a = task(1000): taskw(a):? a!^
```

Each Task has its own Variable set A-Z. This version is not so small.

Numbers are signed 16 bit integers, with a range of -32768 to 32767. But may be printed as unsigned values using a % before the value to be printed. Special Variables:

```
PID - Represents the Process ID of the current task.
Currently some multiple of 25.

TRUE Represents the value -1 or hex $FFFF

FALSE Represents the value 0
```

Parameters Passed to a GOSUB or TASK may be accessed using the # variable Example:

```
1. Access parameter 0: ? #[0]
2. " 1: ? #[1]
```

See Task Section for details

The Firsts task is always PID = 0 so any task may refer to the main tasks variables using : x=0 : x!A ..... notation.

Using @ to access memory locations as integer or byte values.

(a)[0] is the first integer location after the space used by the program,(a)[0] is the first byte value after the space used by the program.

The maximum dimension is dependent upon the size of the computers memory.

The dollar sign may also be used when referencing variable such as a[10]. This return the  $10^{th}$  byte starting at the location of the a variable in memory.

## **Multi Statement lines**

A colon may be used to place more than one statement on a line.

Any line starting with an if statement execute all then and following statements on that line until CR is reached.

Examples:

<statement>[:<statement] -

Any line may contain any number of statements

if <expression> [then] <statement>::<statement>...

When true everything after the THEN is executed until the end of line is reached.

When false then move to next line of code.

As many statement as fit in 132 characters are permitted.

Putting as many statements as make sense on each line will significantly improve performance of your application.

## **Pseudo Compilation**

This version of tiny basic, turns keywords into tokens, and interprets data types. As well as translating line numbers for GoTo, GoSub and Task() into direct memory pointers. This seems to improve the performance of the Basic application as much as 60% over the original implementation.

The cost is of course some memory. The actual Basic program takes less space but the functions to compile and de-compile for listing take extra program space.

The line number conversion is done just before the program executes. And only translates line numbers stated as static values, ie not requiring computation. To this end if a line number must be calculated then it should be surrounded by () as otherwise if the first value in an expression is a number, it will be compiled resulting in an expression error for the line.

## **Expressions/Functions**

## ABS(<number>)

Returns the absolute value of the number.

## ADDR(VARIABLE)

Returns the address of the specified variable

## **CMPMEM(Length, Source 1, Source 2)**

returns:

0 s1 = s2 1 s1 > s2-1 s1 < s2

Length is unsigned 16bit value

Source1 and Source 2 are unsigned 16 bit values

## FREE()

Returns the number of free bytes for user programs.

## **CALL(<Address expression>,<Value Expression>)**

Call a system function with optionally passing a value in Accumulator.

The Call returns what ever is in the Accumulator when the system function returns.

## **GETCH()**

Returns the next character from the tty keyboard.

## FNenumber | Variable containing line number>(Parm 1, Parm 2,...)

Calls a function that returns a value.

## RND(<upper limit>)

Returns a random number from 1 to limit. If limit is not specified, it is set to 32767

## PEEK(<address expression>)

Returns the value at the specified location. Treats address value as unsigned.

ALSO SEE:

Variable = @[\$][index value]

## POKE(<address-expression>, <byte Value expression>)

Sets the memory address to the specified byte value

ALSO SEE:

@[\$][index value] = value

## Commands

#### **CLS**

Clear the screen by sending the ANSI ESC[3] sequence

## **COPYMEM(Length, Destination, Source)**

Copy a block of memory from one location to another.

Source must be before destination or results are undefined

Source: any valid expression – used as a pointer

Destination any valid expression – used as a pointer

Length any valid expression – used as unsigned length

#### **DEC <Variable-Name>**

Decrement variable by 1. Considerably faster than a=a-1

#### DIR

Lists the content of the disk.

#### **END**

Stops the currently running program, returning the user to the prompt.

#### **ERASE <File Name>**

Delete file from the disk.

#### **EXIT**

Returns back to the underlying OS/Monitor.

## **GOTO <(Line Number Expression)>**

Computes the value of the expression and then jumps to that line number, or the next line after it, if that specific line does not exist.

#### Example:

Goto(1000+10) – everything within the brackets is an expression returning a line number.

## GOTO <Valid-Line-Number | .>

GOTO < .> Goto dot "." Repeats the current line from beginning 60% Faster when repeating a line.

This type of goto is pre-computed just before the program executes. It does not search for the line number during execution. It has a direct memory transfer. Example:

Goto 1000 – this must be a valid line number. It is compiled to a direct memory transfer address just before the application executes.

## GOSUB <(Line Number-expression)>[( Parameter 1, ...)]

## GOSUB <Valid-Line-Number>[(Parameters 1, ...)]

Compute the value of the expression and then calls a subroutine at that line, or the next line after it. Return back to the calling point with the RETURN keyword.

Parameters are passed on the Stack. So are limited by the size of the stack which is also used for math. Currently 20 entries deep.

As with the goto the second form is pre-computed just before the program begins execution.

## IF <expression> [THEN] <statement>[:<statement>]

If the expression evaluates to a non-zero value (TRUE) then the statements following THEN will be executed. THEN Keyword is optional.

#### INC <Variable-name>

Increments the variable by 1. Considerably faster than a=a+1

## INPUT [prompt string ; ] <variable>,[[prompt ;]<variable>,...]

Prints a question mark, gets the user's input, converts to a number, then saves the value to the specified variable. If a string follows the keyword then it is printed as a prompt. If the variable ends with a \$ then a single character is read from the input. Example:

input "Enter a letter",a\$

This will read a letter from the keyboard and stores the value in a

## [LET] <variable> = <expression>

Assigns a value to a variable. Let is not required when assigning values to a variable.

### LOAD <"filename">

NOTE: from Version 1.0.4 Quotes are required.

Loads the specified file into memory. The file is just a text file, so you can edit programs using another editor, then load them with this command. Note that this like typing in lines at the prompt, so if there is an existing program in memory and another is loaded, they are "merged" together. Filename must match the case on the directory listing. NOTE: Quotes are not used to enclose file names.

#### **NEW**

This clears the program currently in memory. There is no mercy, no second chances, and no confirmation. The existing program is gone, instantly.

## [PRINT | PR | ?] < values > [; |, ]

Print can have quoted strings, commas, semicolons, numbers and variables.

Commas move to the next tab stop, while semicolons don't advance the cursor.

Using the ? Reduces the size of the program and speeds execution.

Print by its self prints a CR LF

A comma or semi colon at the end will not output the CR-LF.

If an expression starts with a \$ then the value is output as hex.

If an expression starts with a % then the value is displayed as an unsigned 16 bit value.

If \$ trails an expression the value is written as a character.

## Examples:

```
? free(), %free(), $free() ---- outputs: -22344 43192 A8B8
```

## **PUTCH <expression>**

Put a character to the output device. Range is 0-255

## **REM** [<comments>]

The rest of the line is ignored. It is a comment. It is not mandatory to have any text after the REM keyword. Comments made code easier to read, but they also take time to execute, so too many comments can slow down the code.

## [RETURN | RET][(Return Value expression)

Will return to the next statement following the GOSUB which brought the program to this subroutine.

Returns the value to calling gosub in form of function using

FN<line number | variable>

Also used by tasks see Task Section.

#### **RUN**

Begins execution of the program currently in memory starting at the lowest line number.

#### SAVE <"filename">

NOTE: from Version 1.0.4 Quotes are required.

Save the current program to the specified filename. Note that the filename is used exactly as specified; nothing (like ".BAS") is automatically added.

## **SETMEMB(Value, Length, Destination)**

Set a block of memory with byte values Value any 8bit expression Length any unsigned 16bit expression Destination and unsigned 16 bit pointer expression

## SETMEMW(VALUE, LENGTH, Destination)

Set a block of memory with word value Value any 8bit expression Length any unsigned 16bit expression Destination and unsigned 16 bit pointer expression

## **SETTERM** inslot, outslot

InSlot is the index number of the 16 byte slot starting at \$E000 OutSlot is the index number of the 16 byte slot starting at \$E000 Used to configure the stdin and std out for a task. Task 0 defaults to the console. Other tasks also default to the Console.

## Trace <Switch value>

- $_{1}$  %10000000 IPPC trace the core VM
- $_2$  %01000000 Basic program trace
- 3. %01000001 Interactive Basic program trace

## TASKS and TASK MANAGEMENT

Time Sliced Circular scheduling multitasking is supported by Tiny Basic. There are 10 available task entries, The Main Task always uses the first entry leaving 9 available entries for user tasks. The following TASK management commands and functions are available.

Time slices are set by default to 512 IL instructions.. See SLICE

## KILL <Task PID - expression>

Kills the task specified by the expression should be the value returned by TASK() when a task is started.

## STAT(<Task PID - expression>)

Returns the 0 if the task has stopped, 1 otherwise.

## **SLICE <Time-Slice-Count Expression>**

Defines the number of ticks for each time slice used by the task manager. This defaults to 512.

TASK(<(Line Number expression)>[,Parameter-expression]...)

## TASK(<Line-Number>[,Parameter-Expression,...])

As a command This creates a new task starting at the specified line number. As a function it returns the PID of the new task[NOTE: support replaced after 1.0.2 for this function, replaced by task variable PID].

## TASKE[(<Exit value-Expression>)]

This may be used within an executing task to end task. If used within the MAIN Line It acts the same as and END statement. The exit value is optional and is stored in the tasks context after the task exits. This is synonymous with the use of return(exit code-expression).

The exit code is accessed using the special ^ variable.

1. Pid-expression!^

#### **TASKN**

This releases the rest of the tasks time slice to the system. Execution of the task continues at the next statement when the task receives another time slice.

## TASKW(<Task PID Expression>[,<Task PID Expression>]...

Wait for a task or group of tasks to complete.

## **Task Specific variables**

## PID

Is the PID of the current task.

## **#[Parameter index-expression]**

This is the parameter from the parameter list passed when the task was started. Basically the parameters are pushed onto the math stack when the task is started. So the stack size and the need to do math limit the number of parameters that can be passed. No checking is done...So be careful. Parameter index start at zero.

Example: a = #[0] : b = #[1]

These values are read/write and may be used as local variables.

## **Inter-process communication**

Inter process communications is supported by the system.

## ipcs(<message-expression>,<task PID-expression>)

Send a msg to another task

Write messages to the ipc message queue of another task On Return - True-good or False-failed

The message may not be sent if queue is full.

Currently 10 entries are available but this is shared by the Gosub stack.

## Example:

```
On Main task: b = task(1000): a = ipcs(100,b): ? a On Task B: a = ipcr(b): ? "Msg From "a;" Mgs is ";b
```

## ipcr(<variable name>)

Read messages from the IPS message queue

Returns

message value from message queue a message -1 is reserved meaning no entry found The provided variable contains the pid of the sending task. This is optional. This always waits for a message

before returning.

## Ipcc()

Check the message Queue for messages and return the count Returns

Number of messages waiting

## **IRQ and IRQ MANAGEMENT**

## IRQ e number -expression>

Enables the interrupts and Sets the line number to go to when an IRQ is received.

IRQ's are disabled until the IRQ subroutine completes with a ireturn statement. Setting a line number of zero stops the IRQ requests and disables interrupts.

## [IRETURN | IRET]

Returning from an interrupt service routine. Enables the IRQ interrupt.

## **Task Implementation Description.**

Tasks are implemented in the IL interpreter and are really a crude form of tasking. Allowing each task a number of IL instructions before the Task is suspended and the next task is started. Task me to some degree be cooperative and issue a task next command to release the remainder of their time slice.

#### **Task Control Block Definition**

```
27 private variables A-Z,^
                                                            54 Bytes
      Math stack of up to 20 entries
                                                            40 Bytes
    Gosub/For-next Stack 16 entries
                                                            64 Bytes
4. IL Interpreter stack 20 entries
                                                            40 Bytes
5. Pointers for each stack 3 Bytes6. Basic Application Instruction Pointer
                                                            03 Bytes
                                                            02 Bytes

    Basic Application Index Register
    Math Work Registers R0,R1,MQ,R1

                                                            01 Bytes
                                                            07 Bytes
9. Indirect Pointers 3
                                                            06 Bytes
                                                           216 Bytes
```

#### **Context Control Block Definition**

1.	VARIABLES		2 bytes	pointer to, 26 A-Z
2.	ILPC		2 byte	IL program counter
3.	ILSTACK		2 byte	IL call stack
4.	ILSTACKPTR		1 byte	Pointer ti current entry
5.	MATHSTACK		2 bytes	MATH Stack pointer
6.	MATHSTACKPTR		1 byte	Pointer to current stack position
7.	GOSUBSTACK		2 bytes	pointer to gosub stack
8.	GOSUBSTACKPTF	t .	1 byte	current offset in the stack, moved to task table
9.	MESSAGEPTR		2 bytes	Pointer to active message, from bottom of il stack
10.	CURPTR		2 bytes	Pointer to current Basic line
11.	CUROFF		1 byte Cu	rrent offset in Basic Line
12.	R0		2 bytes	arithmetic register 0
13.	R1		2 bytes	;arithmetic register 1
14.	MQ		2 bytes	used for some math
15.	R2		1 byte	General purpose work register(tasking)
16.		Total	25 bytes	

There are ten Task slots Allocated by default in the provided Source Code. This can be altered by the user. Therefore a total of 2140 bytes are required to support multitasking out of the box. A lot of area for some machines.

The interpreter occupies 6K of memory so for a useful Multi tasking system a minimum of 16K is required, Prefer 32 to 48K. Corsham's 6502 ss-50 system comes with 64K. They also sell memory upgraded for the KIM systems and the Rockwell systems.

This system is a good educational tool. And practical for small projects requiring multiple tasks. It is useful to explore Tasking using the em6502 emulator. Or the Corsham products.

## **Error Codes**

- 1 = Expression
- 2 = Stack underflow (expression error)
- 3 = Stack overflow (expression is too complex)
- 4 = Unexpected stuff at end of line
- 5 = Syntax error (possibly unknown command)
- 6 = Divide by zero
- 7 = Read fail loading a file
- 8 = Write fail saving a file
- 9 = No filename provided
- 10= File Not Found
- 11=Gosub Stack underflow, too many returns
- 12=Gosub stack overflow, to many nested gosub statements
- 13=Bad Line Number specified, not found
- 14=Unable to create new task, no more slots
- 15=Array Subscript out of range
- 16=Invalid Task PID provided
- 17=Out of space on queue to send new message
- 18=The expected Stack frame was not found.
- 19=Function(FN) was called and no return value was provided by Subroutine
- 20=When trying to compile a static gosub.goto or task, line number was not found
- 21=IL Stack overflow, the virtual machine has had a stack overflow.
- 22=Expected a variable name or definition
- 23=Expected a closing bracket
- 24=Expected an equal sign for assignment
- 25=ERR\_FUNCTION\_EXPECTED\_PARAMETERS
- 26=ERR\_EXPECTED\_OPENING\_BRACKET

## **Improving Speed**

Tiny BASIC on a 6502 using Hybrid IL is slower than Some implementation.

- Don't use a lot of REM statements, at least not near the beginning of the code. Every REM must be skipped at run time.
- Most other common issues have been mitigated by the use of compiled line numbers. And the tokenization of the Application program.
- When calling a subroutine with parameters try to minimize the number of parameters as each adds to the time it takes to prepare for the call and return.
- If using tasks try to use synchronization and ips rather that starting and stopping them as this takes a lot of resources.

## **Example programs**

## **Example Task program**

3040 print "Failed": return(84)

```
Sample IPC program
10 REM test ipc is working
15 cls
20 a = Task(2000) : b = task(3000,a)
30 \text{ taskw}(a,b)
40?:?: print "Test completed a's exitcode = ";a!^; " B's exit code = ";b!^
50 end
2000 Rem this task will wait for a message
2005 print "PID: ";PID, "Waiting for message":?
2010 a = ipcr(b)
2020 print "I am PID: ";PID, "Recieved msg: ";a, "From PID: ",b:?
2030 return(a)
3000 Rem Send a message to another task
3010 \text{ a} = \#[0]: Rem should be the pid start for first task
3020 Print "PID: ";PID, " Sending message to PID: ";a
3030 print "Result of sending msg = ";:if ipcs(300, a)? "Sent" : return(84)
```