

PROLOG DATABASE

- A Logic Database is a set of facts and rules (i.e. a logic program):

<code>father_of(john,peter) <-.</code>	<code>grandfather_of(L,M) <- father_of(L,N),</code>
<code>father_of(john,mary) <-.</code>	<code>father_of(N,M).</code>
<code>father_of(peter,michael) <-.</code>	<code>grandfather_of(X,Y) <- father_of(X,Z),</code>
<code>mother_of(mary, david) <-.</code>	<code>mother_of(Z,Y).</code>

- Given such database, a logic programming system can answer questions (queries) such as:

`<- father_of(john, peter).`

Answer: *Yes*

`<- father_of(john, david).`

Answer: *No*

`<- father_of(john, X).`

Answer: $\{X = \textit{peter}\}$

Answer: $\{X = \textit{mary}\}$

`<- grandfather_of(X, michael).`

Answer: $\{X = \textit{john}\}$

`<- grandfather_of(X, Y).`

Answer: $\{X = \textit{john}, Y = \textit{michael}\}$

Answer: $\{X = \textit{john}, Y = \textit{david}\}$

`<- grandfather_of(X, X).`

Answer: *No*

- Rules for grandmother_of(X, Y)?

Logic Programs and the Relational DB Model

Traditional → Codd's Relational Model

File	Relation	Table
Record	Tuple	Row
Field	Attribute	Column

- Example:

Name	Age	Sex
Brown	20	M
Jones	21	F
Smith	36	M

Person

Name	Town	Years
Brown	London	15
Brown	York	5
Jones	Paris	21
Smith	Brussels	15
Smith	Santander	5

Lived-in

- The order of the rows is immaterial.
- (Duplicate rows are not allowed)

Relational Database → Logic Programming

Relation Name → Predicate symbol

Relation → Procedure consisting of ground facts
(facts without variables)

Tuple → Ground fact

Attribute → Argument of predicate

• Example:

```
person(brown,20,male) <- .  
person(jones,21,female) <- .  
person(smith,36,male) <- .
```

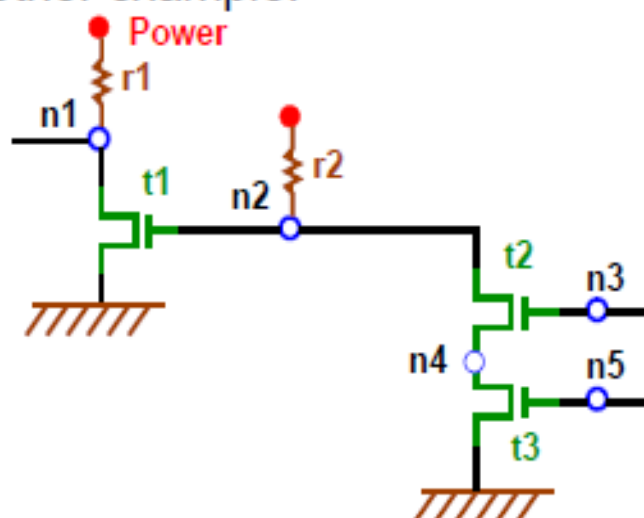
Name	Age	Sex
Brown	20	M
Jones	21	F
Smith	36	M

• Example:

```
lived_in(brown,london,15) <- .  
lived_in(brown,york,5) <- .  
lived_in(jones,paris,21) <- .  
lived_in(smith,brussels,15) <- .  
lived_in(smith,santander,5) <- .
```

Name	Town	Years
Brown	London	15
Brown	York	5
Jones	Paris	21
Smith	Brussels	15
Smith	Santander	5

- Another example:



```
resistor(power,n1) <- .
resistor(power,n2) <- .
```

```
transistor(n2,ground,n1) <- .
transistor(n3,n4,n2) <- .
transistor(n5,ground,n4) <- .
```

```
inverter(Input,Output) <-
```

```
    transistor(Input,ground,Output), resistor(power,Output).
```

```
nand_gate(Input1,Input2,Output) <-
```

```
    transistor(Input1,X,Output), transistor(Input2,ground,X),
    resistor(power,Output).
```

```
and_gate(Input1,Input2,Output) <-
```

```
    nand_gate(Input1,Input2,X), inverter(X, Output).
```

- Query `and_gate(In1,In2,Out)` has solution: $\{In1=n3, In2=n5, Out=n1\}$

TREE

C₁: pet(X) <- animal(X), barks(X).

C₂: pet(X) <- animal(X), meows(X).

C₃: animal(spot) <-.

C₄: animal(barry) <-.

C₅: animal(hobbes) <-.

C₆: barks(spot) <-.

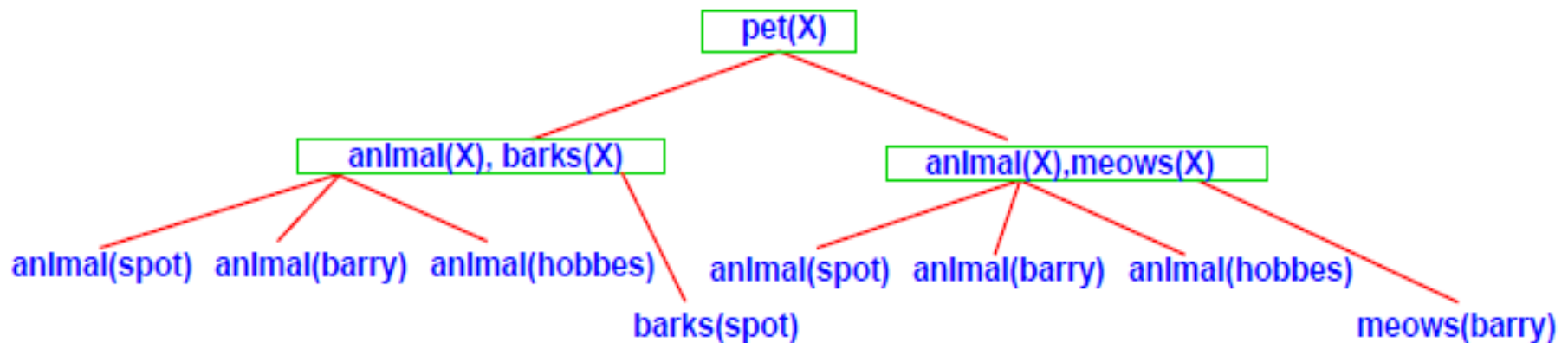
C₇: meows(barry) <-.

C₈: roars(hobbes) <-.

SEARCH TREE

- A query + a logic program together specify a search tree.

Example: query $\leftarrow \text{pet}(X)$ with the previous program generates this search tree (the boxes represent the “and” parts [except leaves]):



- The details of the operational semantics explain how the search tree will be explored during execution.
- Different query \rightarrow different tree.

Graph

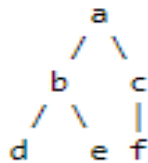
graphs are simple and understandable. Graph-theoretic structures are able to represent properties of programs. Known properties of graphs can help improve the understanding of structural properties of the programs logic.

Assume given a set of facts of the form `father(name1,name2)` (`name1` is the father of `name2`).

1. Define a predicate `brother(X,Y)` which holds iff `X` and `Y` are brothers.
2. Define a predicate `cousin(X,Y)` which holds iff `X` and `Y` are cousins.
3. Define a predicate `grandson(X,Y)` which holds iff `X` is a grandson of `Y`.
4. Define a predicate `descendent(X,Y)` which holds iff `X` is a descendent of `Y`.
5. Consider the following genealogical tree:

```
father(a,b). % 1
father(a,c). % 2
father(b,d). % 3
father(b,e). % 4
father(c,f). % 5
```

whose graphical representation is:



Say which answers, and in which order, are generated by your definitions for the queries

```
?- brother(X,Y).
?- cousin(X,Y).
?- grandson(X,Y).
?- descendent(X,Y).
```

```
brother(X,Y) :- father(Z,X), father(Z,Y), not(X=Y).      % 6

cousin(X,Y) :- father(Z,X), father(W,Y), brother(Z,W). % 7

grandson(X,Y) :- father(Z,X), father(Y,Z).               % 8

descendent(X,Y) :- father(Y,X).                           % 9
descendent(X,Y) :- father(Z,X), descendent(Z,Y).          % 10

?- brother(X,Y).
X = b    Y = c ;
X = c    Y = b ;
X = d    Y = e ;
X = e    Y = d ;
No
```


?- cousin(X,Y).

X = d Y = f ;

X = e Y = f ;

X = f Y = d ;

X = f Y = e ;

No

?- grandson(X,Y).

X = d Y = a ;

X = e Y = a ;

X = f Y = a ;

No

?- descendent(X,Y).

X = b Y = a ;

X = c Y = a ;

X = d Y = b ;

X = e Y = b ;

X = f Y = c ;

X = d Y = a ;

X = e Y = a ;

X = f Y = a ;

No

draw the SLD-tree for the first query.