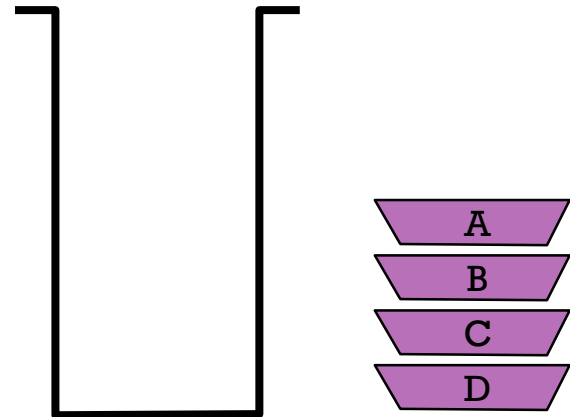


Stack

What is Stack?

■ Definition

- A collection of elements that are inserted and removed according to the **last-in first-out (LIFO) principle**.
 - The last element will be the first element to be removed.
 - Input and output are only possible at the top on the stack.

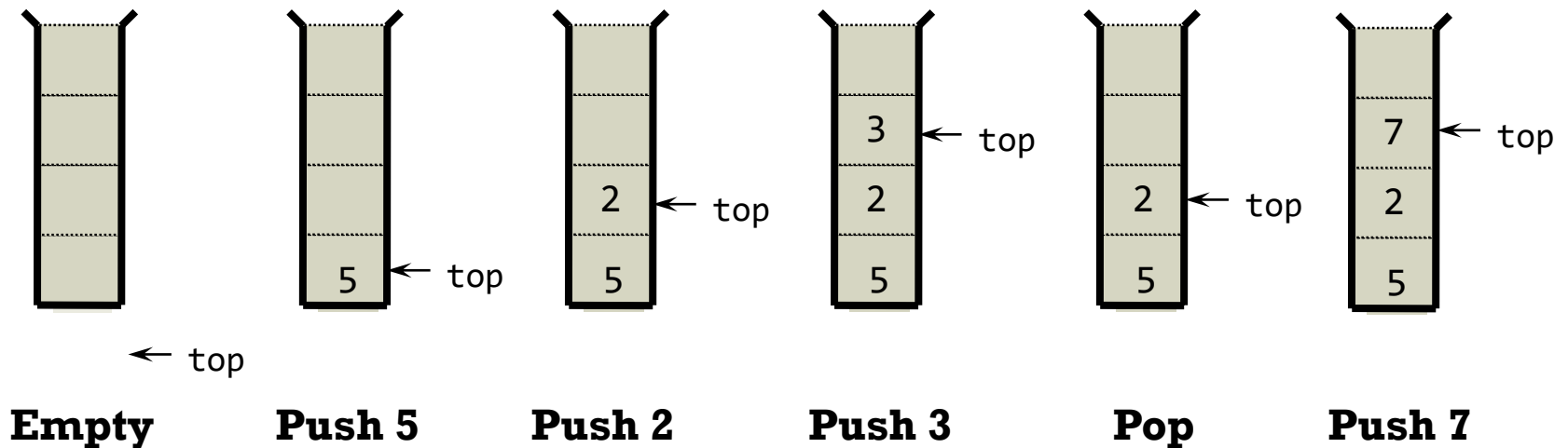


What is Stack?

■ Terminology

- **Top**: The top of stack (default = -1)
- **Push**: Insert an item on the top.
- **Pop**: Remove the item on the top.

■ How does the stack work?



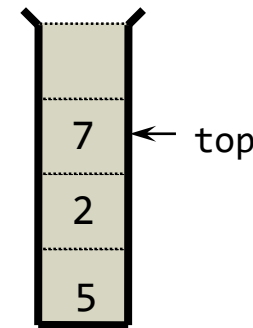
What is Stack?

■ Operations

- **InitStack**: Make stack empty.
- **IsFull**: Check whether stack is full.
- **IsEmpty**: Check whether stack is empty.
- **Peek**: Read the item at the top.
- **Push**: Insert an item at the top.
- **Pop**: Remove the item at the top.

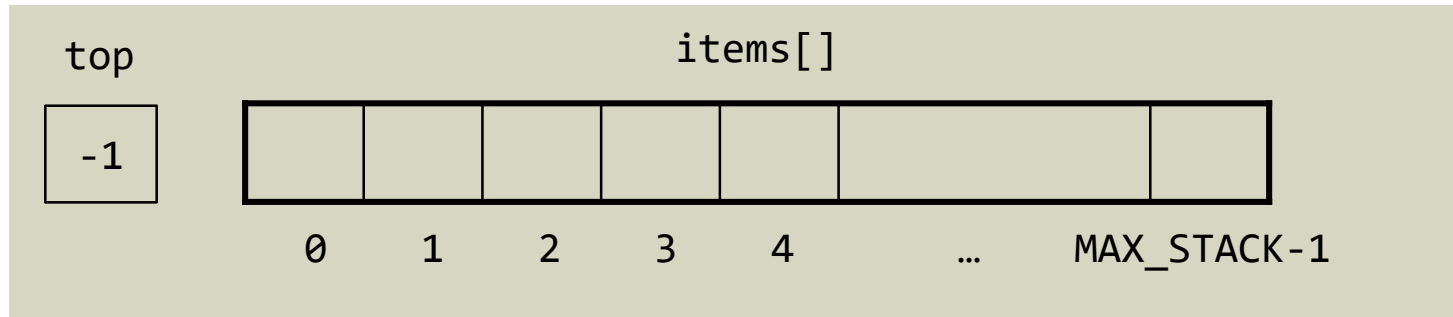
Q: Can we access items other than at the top?

A: By definition, No.



Array-based Implementation

■ Stack representation



```
#define MAX_STACK    100

typedef enum { false, true } bool;
typedef int Data;

typedef struct {
    Data items[MAX_STACK];
    int top;
} Stack;
```

Array-based Implementation

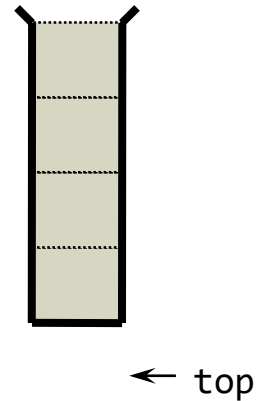
■ Operations

```
// Make stack empty.  
void InitStack(Stack *pstack);  
// Check whether stack is full.  
bool IsFull(Stack *pstack);  
// check whether stack is empty.  
bool IsEmpty(Stack *pstack);  
  
// Read the item at the top.  
Data Peek(Stack *pstack);  
// Insert an item at the top.  
void Push(Stack *pstack, Data item);  
// Remove the item at the top.  
void Pop(Stack *pstack);
```

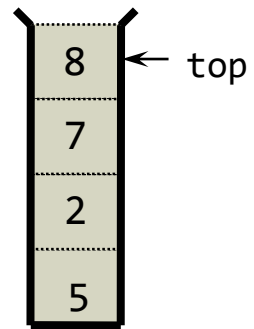
Array-based Implementation

■ Initialize and IsFull operations

```
// Make stack empty.  
void InitStack(Stack *pstack)  
{  
    pstack->top = -1;  
}
```



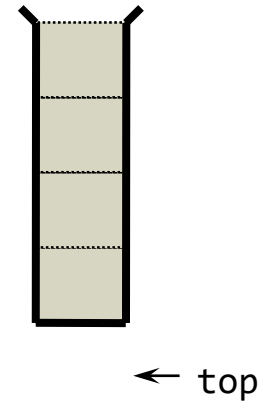
```
// Check whether stack is full.  
bool IsFull(Stack *pstack)  
{  
    return pstack->top == MAX_STACK - 1;  
}
```



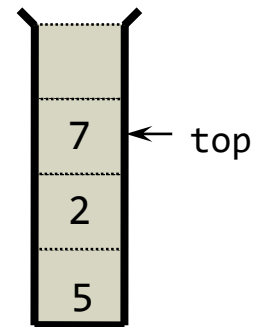
Array-based Implementation

■ IsEmpty and Peek operations

```
// check whether stack is empty
bool IsEmpty(Stack *pstack)
{
    return pstack->top == -1;
}
```



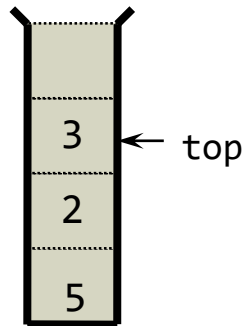
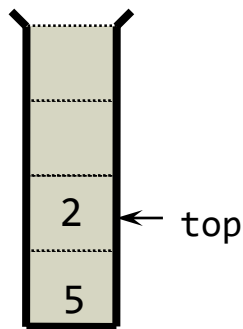
```
// Read the item at the top.
Data Peek(Stack *pstack)
{
    if (IsEmpty(pstack))
        exit(1); //error: empty stack
    return pstack->items[pstack->top];
}
```



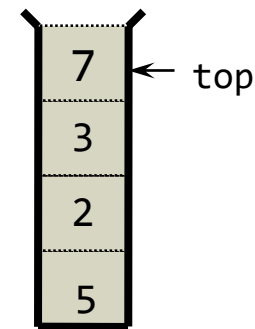
Array-based Implementation

■ Push operation

```
// Insert an item at the top.  
void Push(Stack *pstack, Data item)  
{  
    if (IsFull(pstack))  
        exit(1); //error: stack full  
    pstack->items[++(pstack->top)] = item;  
}
```



Push 3

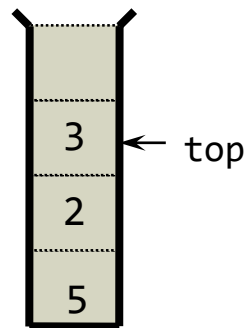
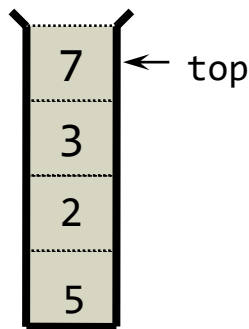


Push 7

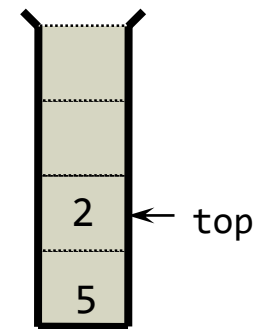
Array-based Implementation

■ Pop operation

```
// Remove the item at the top.  
void Pop(Stack *pstack)  
{  
    if (IsEmpty(pstack))  
        exit(1); //error: empty stack  
    --(pstack->top);  
}
```



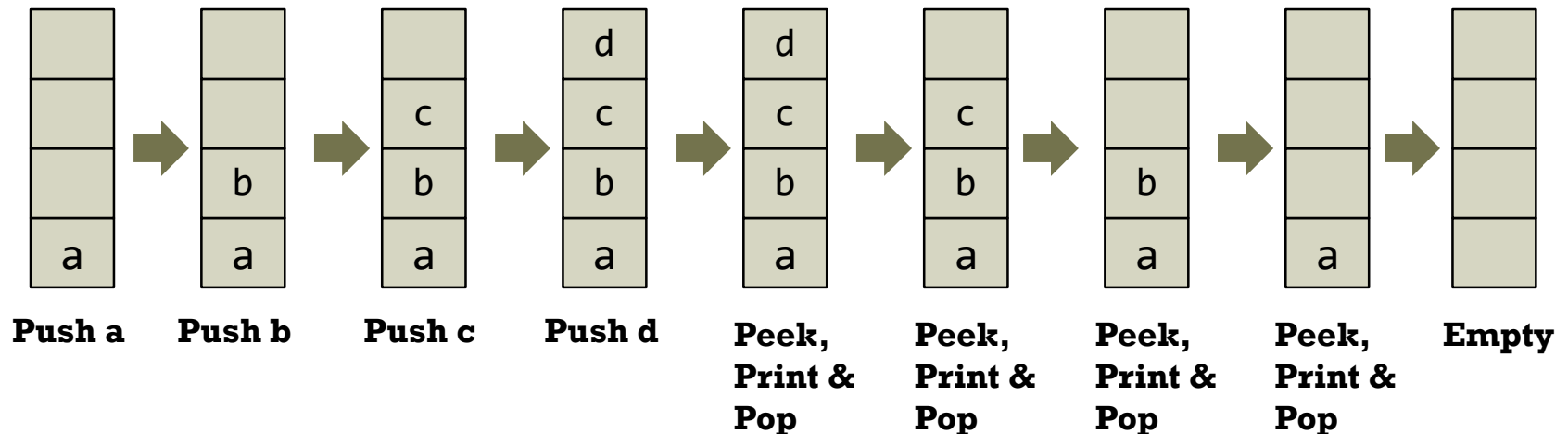
Pop
10



Pop

Print a Reverse String

- Print a string in the reverse order.
 - E.g., abcd → dcba
- How to do this with a stack?
 - Push all characters into stack.
 - Read the top, print it and pop until stack is empty.



Print a Reverse String

■ Implementation

```
void ReversePrint(char* s, int len)
{
    Stack stack;
    char ch;

    InitStack(&stack); // Make a stack empty.
    for (int i = 0; i < len; i++) // Push characters.
        Push(&stack, s[i]);

    while (!IsEmpty(&stack)) // Pop characters.
    {
        ch = Peek(&stack);
        printf("%c", ch);
        Pop(&stack);
    }
}
```

Parenthesis Matching

■ Problem

- Check if each opening symbol has a corresponding closing symbol and the pairs of parentheses are nested properly.

First open waits until last close

(() (()) ())

Most recent open matches first close

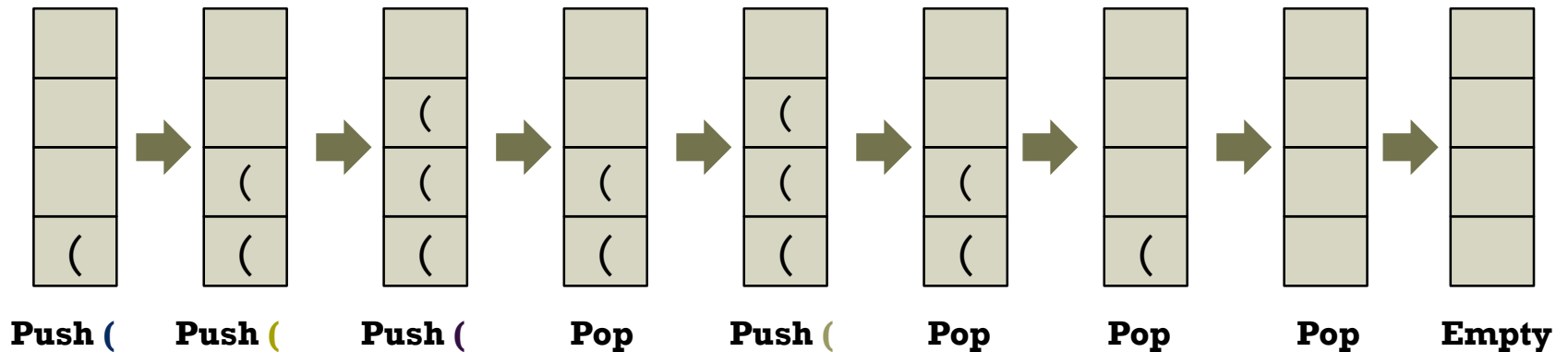
■ Example

- Balanced: $((()()))$, $((()()))$, $((()()))$
- Unbalanced: $(((((()$), $()$), $((()()$

Parenthesis Matching

- How to do this with a stack?
 - Push all open symbols into stack.
 - Whenever finding the close symbol, pop the open symbol.
 - If the stack is empty, it is not balanced.
 - After reading all parentheses, check the status of the stack.
 - If the stack is empty, it is balanced.
 - Otherwise, it is not balanced.

■ Example: ((() ()))



Parenthesis Matching

■ Implementation

```
bool IsParanBalanced(char* exp, int len)
{
    Stack stack;
    InitStack(&stack);           // Make a stack empty.
    for (int i = 0; i < len; i++) {
        if (exp[i] == '(')       // Check open symbol
            Push(&stack, exp[i]);
        else if (exp[i] == ')') { // Check close symbol
            if (IsEmpty(&stack))
                return false;    // Unbalanced case
            else
                Pop(&stack);
        }
    }
    if (IsEmpty(&stack))
        return true;            // Balanced case
    else
        return false;           // Unbalanced case
}
```

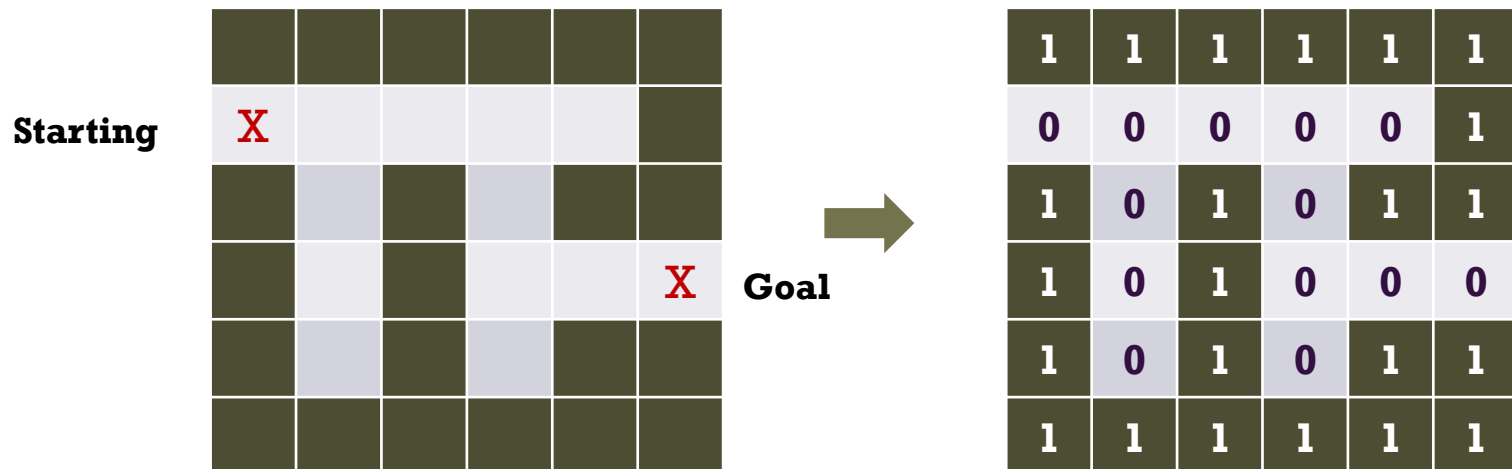
Maze Solving

■ Problem

- Find a path from the starting position to the goal point position.
- At any moment, you can only move one step in one of four directions (up, down, left, and right).

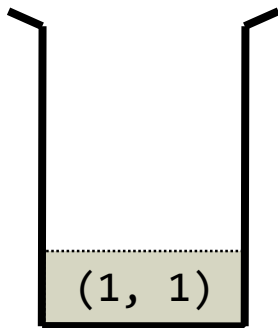
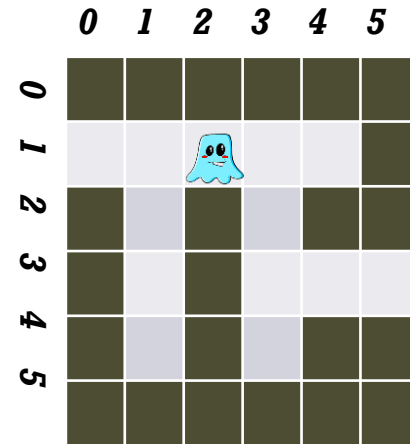
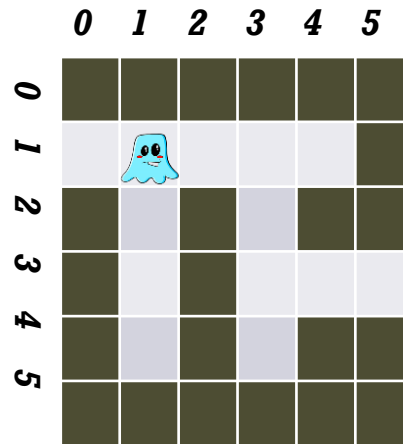
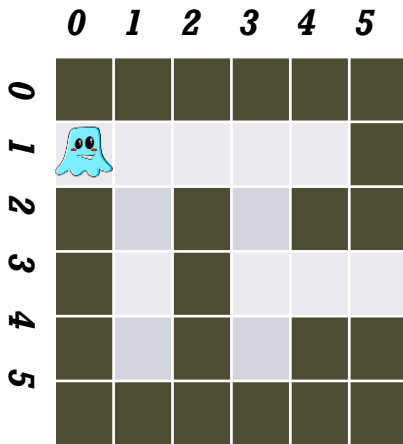
■ Representation

- The maze is represented by a 2D binary array.
 - 0: path, 1: block



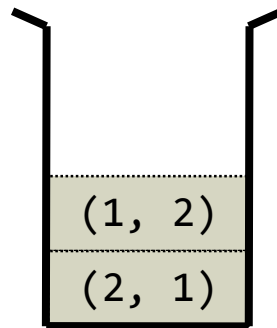
Maze Solving

- How to do this with a stack?



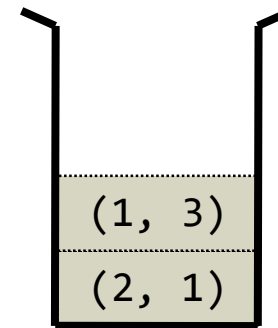
Push right direction.

Pop (1, 1)



Push down and right direction.

Pop (1, 2)

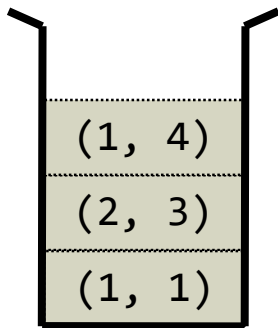
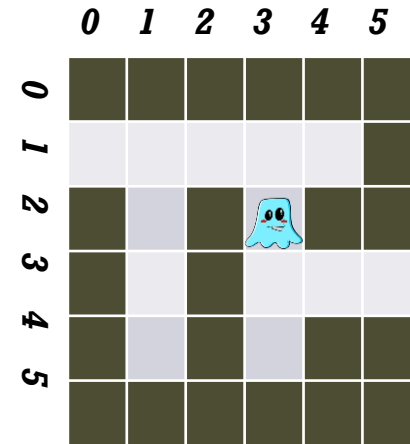
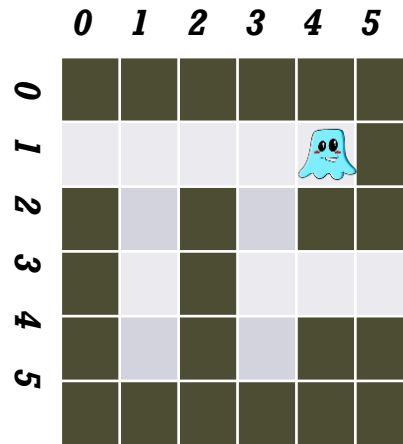
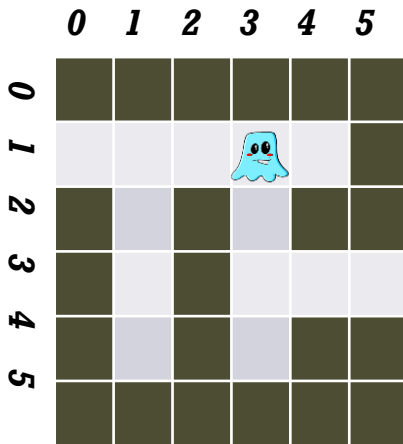


Push right direction.

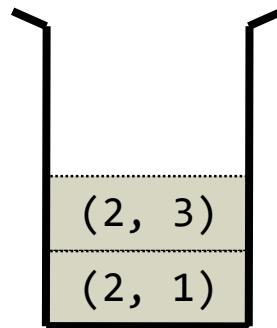
Pop (1, 3)

Maze Solving

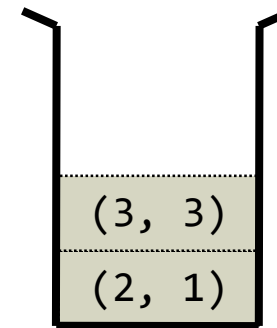
- How to do this with a stack?



Pop (1, 4)



Pop (2, 3)



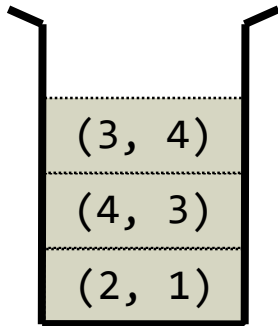
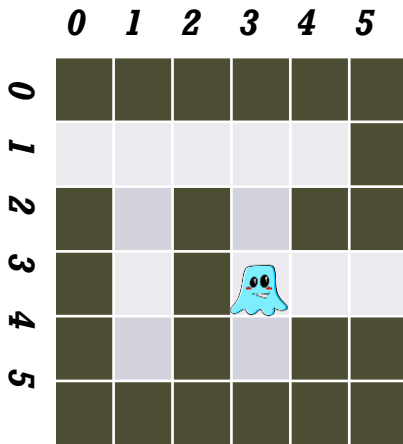
Pop (3, 3)

Push down and right direction.

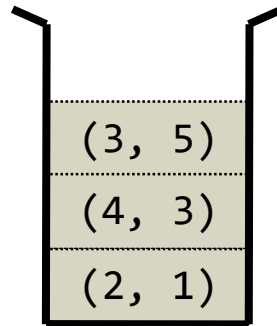
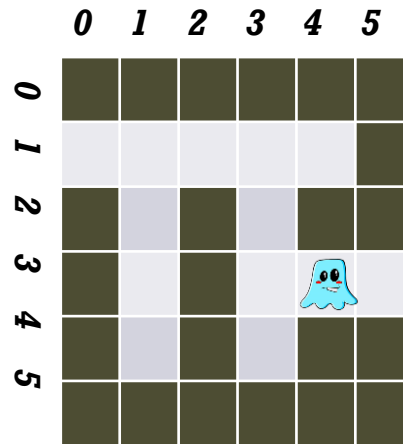
Push down direction.

Maze Solving

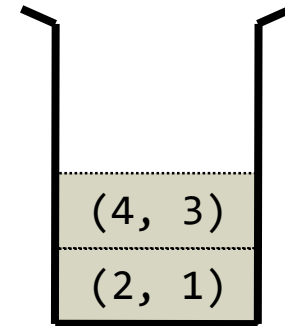
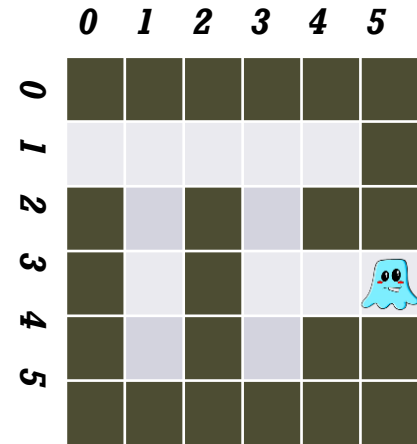
- How to do this with a stack?



Push down and
right direction.



Push right
direction.



Pop (3, 4)

Pop (3, 5)

Maze Solving

■ Overall process

1. When extending the path, **push a new position on the stack.**
 - 1.1. At the current position, extend the path one step by trying to go three directions (up, down, or right).
2. **Pop a position to move the current position.**
3. If none of the neighboring positions is available, **pop a position on the stack.**
 - 3.1. **If the stack is empty, it indicates failure.**

Repeat steps **1~3** until the goal has been reached.

Evaluation of Expression

- How to evaluate the expression?

$$3 + 4 * (5 / 2) + (7 + 9 * 3)$$

- Evaluation by human
 - Assign to each operator a priority.
 - Use parenthesis and evaluate inner-most ones.

$$((3 + (4 * (5 / 2))) + (7 + (9 * 3)))$$

- How to evaluate by computer?
 - How to evaluate the operators with parenthesis
 - How to determine the precedence of operators

Infix, Postfix, and Prefix

- Infix notation: $X + Y$
 - Operators are written in-between their operands.
 - Need extra information to make the order of evaluation of the operators clear.

- Postfix notation: $X Y +$
 - Operators are written after their operands.
 - The order of evaluation of operators is always left-to-right.
 - Unnecessary to use parenthesis and precedence of operators.

- Prefix notation: $+ X Y$
 - Operators are written before their operands.
 - As for postfix, operators are evaluated left-to-right.

Infix vs. Postfix

■ Example

| Infix | Postfix | Prefix |
|-------------------|-----------------|-----------------|
| $A * B + C / D$ | $A B * C D / +$ | $+ * A B / C D$ |
| $A * (B + C) / D$ | $A B C + * D /$ | $/ * A + B C D$ |
| $A * (B + C / D)$ | $A B C D / + *$ | $* A + B / C D$ |
| $A * B / C - D$ | $A B * C / D -$ | $- / * A B C D$ |

■ Why is postfix useful?

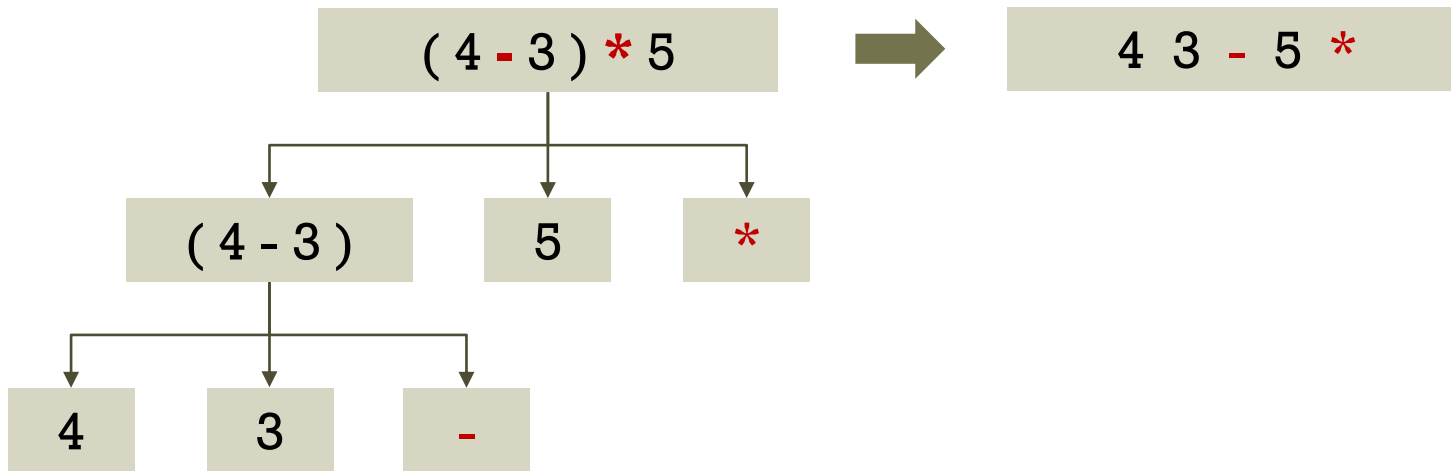
- No parenthesis is needed
- No precedence of operators is needed

Converting Infix to Postfix

- Applying the conversion rule in a **recursive** way

$$A * B \Rightarrow AB^*$$

- Example



How to Evaluate Postfix Notation

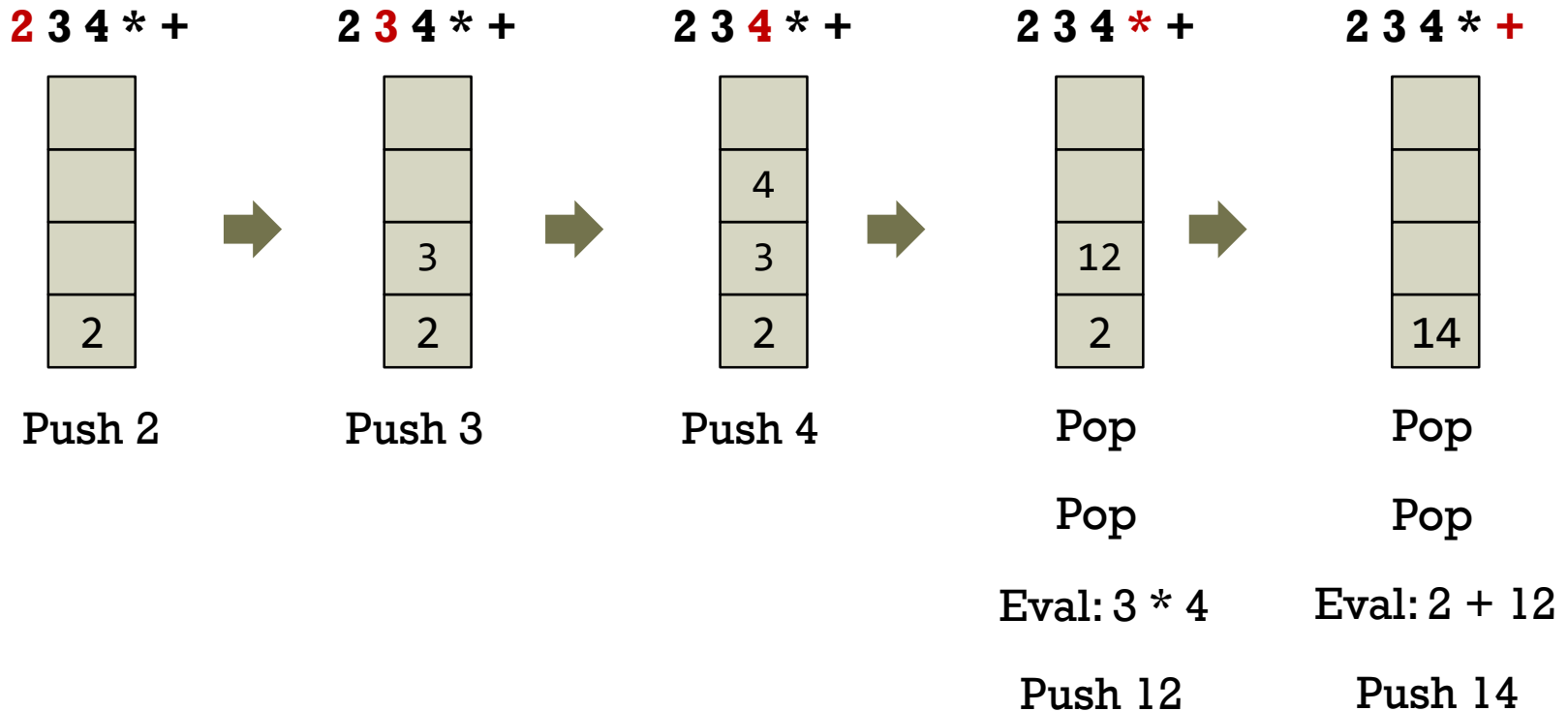
■ Overall process

1. **Push operands on the stack** until finding an operator.
2. If the operator is found, **pop two operands** and evaluate the operator. Then, **push the result** on the stack.

Repeat steps **1~2** until reading all characters in postfix notation.

How to Evaluate Postfix Notation

- Example: $2\ 3\ 4\ *\ +$
 - Push operands on the stack until finding an operator.
 - If the operator is found, pop two operands and evaluate the operator. Then, push the result on the stack.

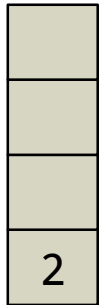


How to Evaluate Postfix Notation

■ Example: $2\ 3\ +\ 4\ *$

- Push operands on the stack until finding an operator.
- If the operator is found, pop two operands and evaluate the operator. Then, push the result on the stack.

2 3 + 4 *



Push 2



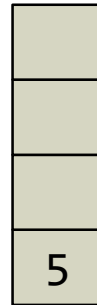
2 3 + 4 *



Push 3



2 3 + 4 *



Pop

Pop

Eval: $2 + 3$

Push 5



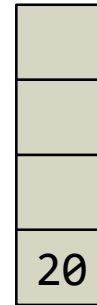
2 3 + 4 *



Push 4



2 3 + 4 *



Pop

Pop

Eval: $4 * 5$

Push 20

How to Evaluate Postfix Notation

```
int EvalPostfix(char* exp, int len)
{
    Stack stack;
    int op1, op2;
    InitStack(&stack);
    for (int i = 0; i < len; i++) {
        if (isdigit(exp[i])) // Push an operand.
            Push(&stack, exp[i] - '0');
        else {
            // Evaluate an operator.
            op2 = Peek(&stack); Pop(&stack);
            op1 = Peek(&stack); Pop(&stack);
            if (exp[i] == '+')
                Push(&stack, op1 + op2);
            else if (exp[i] == '-')
                Push(&stack, op1 - op2);
            else if (exp[i] == '*')
                Push(&stack, op1 * op2);
            else if (exp[i] == '/')
                Push(&stack, op1 / op2);
        }
    }
    return Peek(&stack);
}
```

Converting Infix to Postfix

■ Overall process

1. If the operand is found, **print it**.
2. If the operator is found, **push it into stack**, but before pushing
 - 2.1. See the operator at the top of the stack.
 - 2.2. If the priority of the incoming operator is lower than or equal to the top, **pop and print the top** and go to step 2.1.

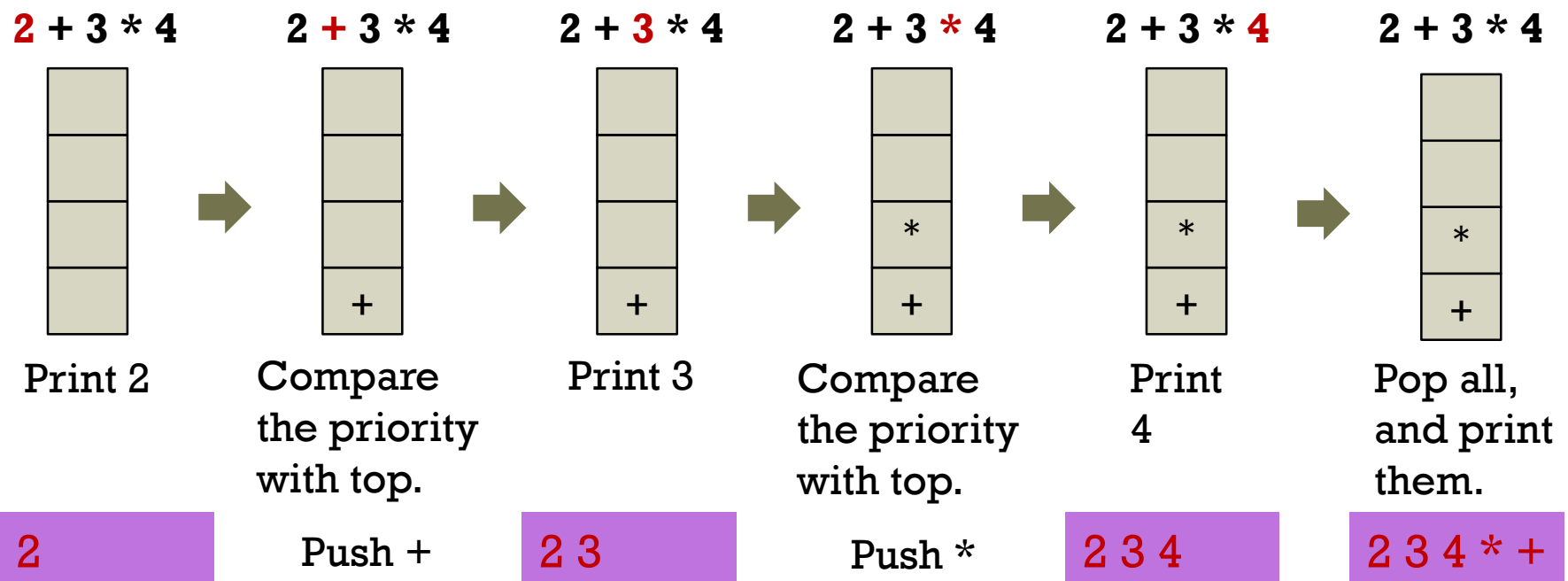
Repeat steps 1~2 until reading all characters in infix notation.

3. At the end of the infix notation, **pop all operators**.

Converting Infix to Postfix

■ Example: $2 + 3 * 4$

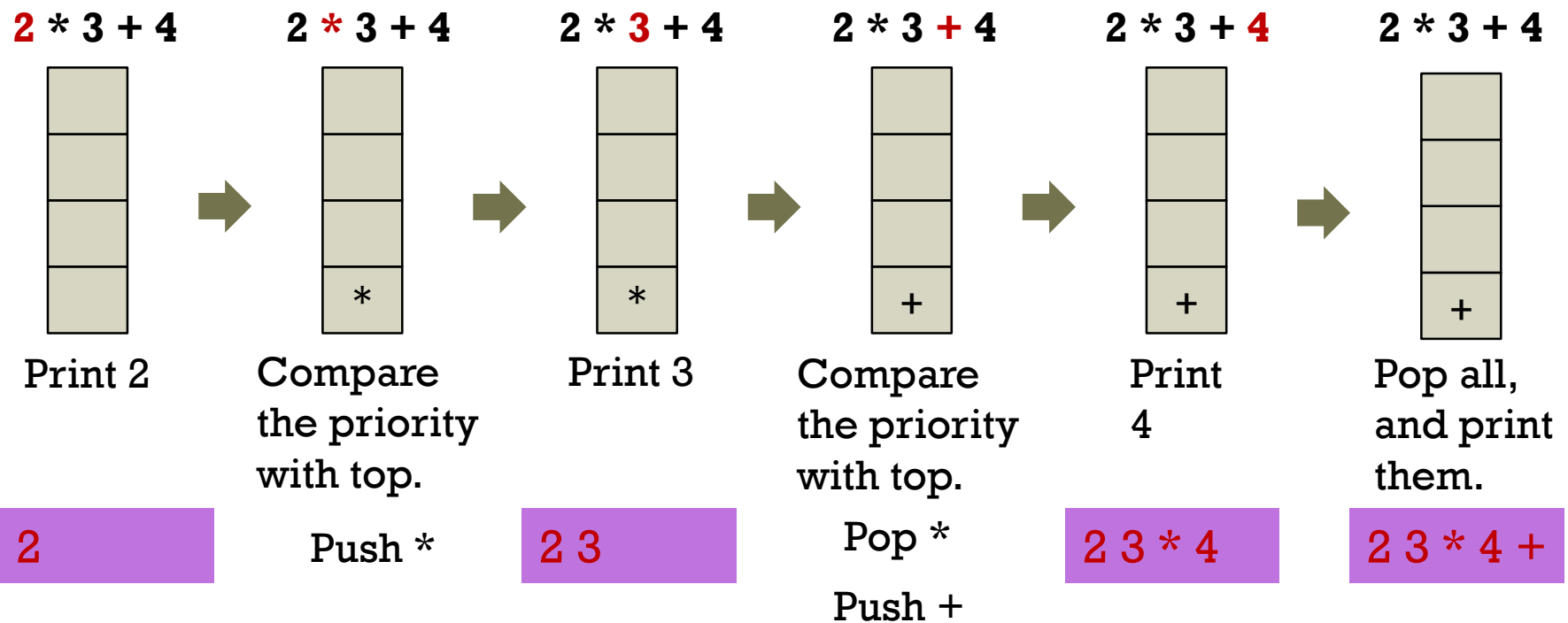
- While reading a character from infix notation
 - If we find operand, print it.
 - Otherwise, check the priority of operator and determine whether push it into stack.
- Pop all operators at the end of the infix notation.



Converting Infix to Postfix

■ Example: $2 * 3 + 4$

- While reading a character from infix notation
 - If we find operand, print it.
 - Otherwise, check the priority of operator and determine whether push it into stack.
- Pop all operators at the end of the infix notation.



Converting Infix to Postfix

■ Implementation: $2 * 3 + 4$

```
void ConvInfixToPostfix(char* exp, char* convExp, int len)
{
    Stack stack;
    int idx = 0;
    InitStack(&stack);
    for (int i = 0; i < len; i++)
    {
        if (isdigit(exp[i]))
            convExp[idx++] = exp[i]; // Print an operand.
        else {
            while (!IsEmpty(&stack) && ComPriority(Peek(&stack), exp[i])) {
                convExp[idx++] = Peek(&stack); // Print an operator.
                Pop(&stack); // Pop an operator.
            }
            Push(&stack, exp[i]); // Push an operator.
        }
    }
    while (!IsEmpty(&stack)) {
        convExp[idx++] = Peek(&stack); // Print an operator.
        Pop(&stack); // Pop an operator.
    }
}
```


Converting Infix to Postfix

■ Implementation

```
int GetPriority(char op)
{
    if (op == '*' || op == '/')
        return 2;
    else if (op == '+' || op == '-')
        return 1;
    else
        return 0;
}

bool ComparePriority(char op1, char op2)
{
    int op1_pr = GetPriority(op1);
    int op2_pr = GetPriority(op2);

    if (op1_pr >= op2_pr)
        return true;
    else
        return false;
}
```