

# LLM Watermarking: Sequential Detection



# Motivating the Problem!

### Watermark:

A hidden signal embedded in text generated by a language model (LM) to trace its origin.

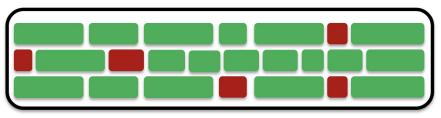
### Problem Setup:

Given an LM 'm' and a user prompt 'q', embed a watermark in the output 'y' such that:

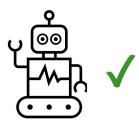
- 1. **Distortion-Free:** Output quality is unchanged (i.e.,  $P(x) \approx$  original distribution).
- 2. Model Agnostic: Detection works without access to 'm' or 'q'.
- 3. Robustness: Watermark remains detectable even after adversarial modifications.

# Kirchenbauer - Red/Green List Watermarking





Z-score = 
$$\frac{(|s|_G - \gamma T)}{\sqrt{T\gamma(1-\gamma)}}$$
 = 4



Z-score >  $\tau$  (say 3)

### **Detection**

### **Overview of KGW**

### **Solutions:**

1. The technique is model/prompt agnostic and does not need the knowledge of model. (Although this a white-box approach)

### **Problems:**

- 1. Induces distortion into the output by changing probability distribution.
- 2. The technique is susceptible to substitution attacks.



# Kuditipudi - Statistical Watermarking

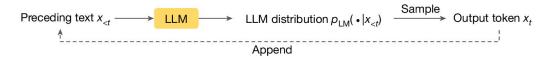
### **Solutions:**

- 1. Model/Prompt Agnostic
- 2. Since we do not manipulate the output logits of the LM, we mitigate the distortion.
- 3. The technique is robust to majority of attacks.

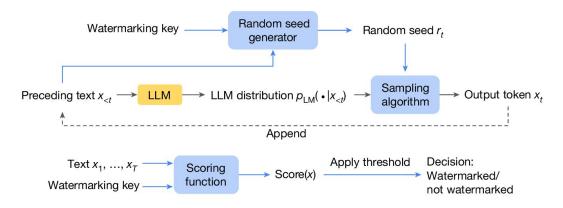
### **Problems:**

1. The algorithm has limitations on detection speeds, and is incapable of 'online' detection due to the design of the algorithm.

### **LLM** text generation



### Generative watermarking: text generation and watermark detection







**Objective:** Detect whether a given sequence of tokens  $Y_n = (Y_1, Y_2, ..., Y_n)$  was generated by a watermarked language model.

**Setup:** Given a stream of observations:

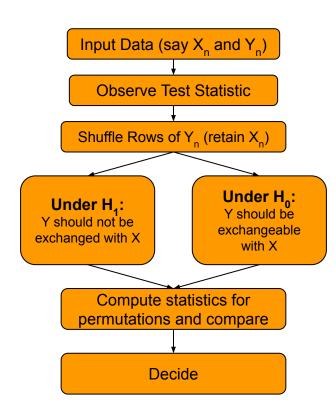
Sequence of keys  $\xi = (\xi_1, \xi_2, ..., \xi_n)$ , Output tokens  $Y_n = (Y_1, Y_2, ..., Y_n)$ 

Decide between hypotheses:

$$H_0$$
: Text  $\perp \!\!\! \perp$  Key vs.  $H_a$ : Text  $\perp \!\!\! \perp$  Key

**Goal:** For  $\alpha \in (0, 1)$ , construct a level- $\alpha$  sequential test of power one,

- → Under  $H_0$ : continue forever w.p.  $\geq 1-\alpha$
- $\rightarrow$  Under  $H_1$ : stop the test, and reject the null as soon as possible



For simplification, we modularise our watermarking scheme into 4 steps/algorithms:

- **→** Step-1: Generating the watermark
- → Step-2: Detecting the presence of watermark in a text
- → Step-3: Test statistic evaluating the misalignment between the keys and the text



### **Algorithm 1:** Watermarked text generation (generate)

Input : watermark key sequence  $\xi \in \Xi^n$ 

**Params:** generation length m, language model p, decoder  $\Gamma$ 

Output: string  $y \in \mathcal{V}^m$ 

- 1 for  $i \in {1, ..., m}$  do
- $\mathbf{2} \mid y_i \leftarrow \Gamma(\xi_{i\%n}, p(\cdot \mid y_{:i-1}))$
- $\mathbf{3}$  return y



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- → Step-1: Generating the watermark
- → Step-2: Detecting the presence of watermark in a text
- → Step-3: Test statistic evaluating the misalignment between the keys and the text

# Algorithm 2: Watermarked text detection (detect) Input : string $y \in \mathcal{V}^*$ , watermark key sequence $\xi \in \Xi^n$ Params: test statistic $\phi$ ; watermark key sequence distribution $\nu \in \Delta(\Xi^n)$ ; resample size TOutput: p-value $\widehat{p} \in [0,1]$ 1 for $t \in 1, \ldots, T$ do 2 $\left| \begin{array}{c} \xi^{(t)} \sim \nu \\ 3 & \phi_t \leftarrow \phi(y, \xi^{(t)}) \end{array} \right|$ 4 $\widehat{p} \leftarrow \frac{1}{T+1} \left( 1 + \sum_{t=1}^T \mathbf{1} \{ \phi_t \leq \phi(y, \xi) \} \right)$ 5 return $\widehat{p}$

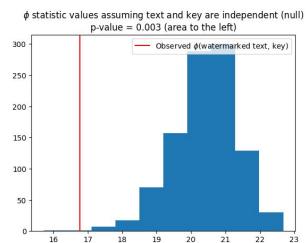
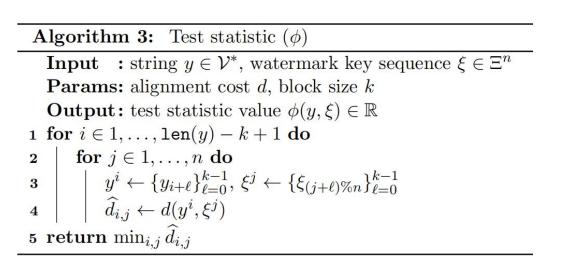


Fig: Distribution of test statistics (under the Null)



- → Step-1: Generating the watermark
- → Step-2: Detecting the presence of watermark in a text
- → Step-3: Test statistic evaluating the misalignment between the keys and the text







# Proposed Solution

Main idea: Replace Kuditipudi et. al's permutation test with a sequential test

Sequential tests allow for gathering evidence against the null hypothesis in an online fashion and stop when it becomes significant. Compared to traditional ("batch") tests:

- → They often reach decisions much earlier (saving resources), and
- → can make the same guarantees on their false positive rates

"Testing by betting" is an increasingly popular framework for designing sequential tests.



# Testing by Betting Framework

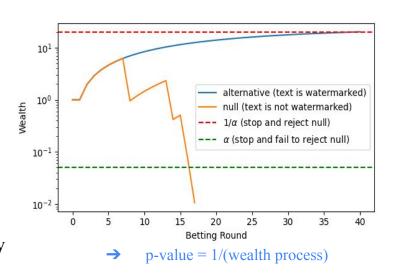
We track the wealth (a martingale by construction) of a gambler that bets against the null.

The betting function is designed such that the wealth (stochastic process):

- → Is a martingale (remains constant in expectation) under the null
- → Grows exponentially under the alternative

In our case,

- → We are testing if the text is independent of the watermarking key
- → Using a Monte Carlo permutation test, which computes T (expensive) test statistics
- → Could before reaching T if we have enough evidence against the null?





# Intuition and Guarantees of the Strategy

We focus on a log-optimal betting strategy designed specifically for hypothesis testing under a given alternative hypothesis.

### Main Properties:

- → **Log-Optimality:** Maximizes the expected log wealth under the considered alternative, ensuring statistically efficient use of evidence.
- → Finite-Time Guarantee: Achieves zero resampling risk after a finite number of permutations no need for infinite resampling to maintain validity.
- → Any-Time Valid: The method maintains type-I error control at any stage, enabling real-time, sequential analysis without needing a fixed sample size.



- → Step-1: Generating the watermark
- → Step-2: Detecting the presence of watermark in a text
- → Step-3: Test statistic evaluating the misalignment between the keys and the text
- → Step-4: Sequentializing the Hypothesis Test using a Sequential-MC Test

```
Algorithm 2': Sequential Monte Carlo permutation test (seq_mc_permutation_test)
   Input: tokens y \in \mathcal{V}^*, watermark key length n, block size k, test statistic function \phi, watermark
             key sequence \xi \in \Xi^n, threshold \alpha, slack parameter c
   Output: p-value estimate \hat{p} \in [0, 1], runtime t \in \mathbb{N}
1 begin
       W \leftarrow 1:
                                                                                                     // initial wealth
       L \leftarrow 0;
                                                                                                      // success count
       \phi_0 \leftarrow \phi(y, n, k):
                                                                    // flip sign of observed test statistic
       for t = 1 to T do
           \xi^{(t)} \sim \nu;
           \phi_t \leftarrow \phi(y, \xi^{(t)});
           if \phi_t \geq \phi_0 then
            L \leftarrow L + 1;
                                                                                        // increment success count
            W \leftarrow \frac{1-\operatorname{BinomCDF}(L;t+1,c)}{c};
                                                                         // update wealth using binomial tail
10
            if W \geq \frac{1}{\alpha} or W < \alpha then
11
              break;
12
                                                                                                          // early stop
       \hat{p} \leftarrow 1/\max(W, \epsilon);
                                                                                         // final p-value estimate
13
       return \hat{p}, t
14
```



# Experiment setup



### Model and Dataset:

- → OPT-1.3B A 1.3 billion parameter open-source language model developed by Meta.
- → C4 dataset— A large-scale English-language dataset curated for language modeling tasks.

### Watermark Generation Methods:

- → ITS / ITS-edit: Inverse Transformed Sampling for watermarking and its sequential variant.
- → EXP / EXP-edit: Exponential Sampling watermarking and its sequential variant.
- → KGW-1.0: Kirchenbauer baselines for the sake of comparison.

### **Evaluation Metrics:**

- → Permutation p-value: Used to test statistical dependence between generated tokens and watermarking mechanism.
- → Number of Permutations: Reflects computational efficiency and convergence behavior of the test.

### Oracle Setup:

- → We simulate an oracle setting where the watermark detection algorithm has access to the true distribution of the watermark signal under the null hypothesis (i.e., no watermarking).
- → This setup enables us to isolate and evaluate the ideal performance of detection methods under best-case assumptions.

# Results: Watermarking the Midterm Report



### **Midterm Report**

As large language models (LLMs) continue to improve, traditional watermarking techniques—which previously depended on clear differences between machine-generated and human-written content—are becoming less reliable [9][3]. Alongside advancements in watermarking, more sophisticated detection techniques are also being developed [1][5][6].

Recently, watermarking strategies have explored statistical embedding and detection mechanisms. Notably, Kuditipudi [8] introduced a distortion-free watermarking approach coupled with a dependable detection method; however, its dependency on batch processing limits its practicality for real-time use.

To address this limitation, our project introduces a sequential watermark detection algorithm [2]. We propose an anytime-valid e-process/p-process framework that allows for real-time detection with early stopping, live access to test statistics, and reduced computational requirements.

### **Watermarked Report**

As large language models (LLMs) continue to advance, the effectiveness of traditional watermarking methods—which often rely on observable distinctions between human-authored and machine-generated content—has diminished [9][3]. In parallel with improvements in watermarking strategies, researchers have also made progress in the development of more robust detection methodologies [1][5][6].

Recent approaches in watermarking have focused on statistical techniques for both embedding and detection. One such method, proposed by Kuditipudi [8], offers a distortion-free watermarking strategy alongside a reliable detection mechanism. However, this approach is limited by its reliance on batch processing, which constrains its applicability in real-time environments

To overcome this limitation, the present work introduces a sequential watermark detection algorithm [2]. Our method leverages an anytime-valid e-process/p-process framework, which facilitates real-time detection through early stopping mechanisms, access to intermediate test statistics, and reduced computational demands.

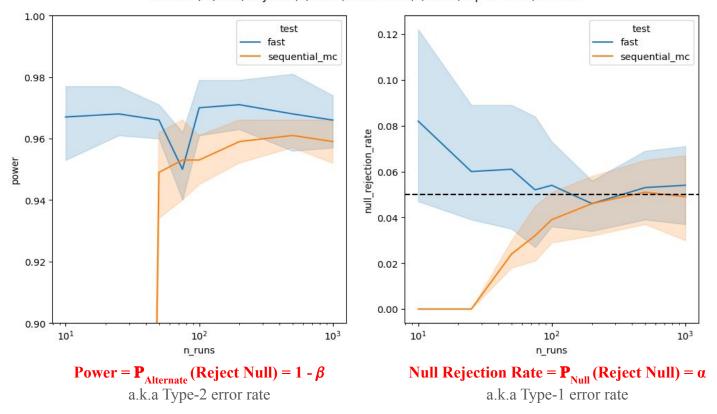
p-value:  $\sim 0.95 \Rightarrow$  Likely written by a human

p-value: 0.0494 ⇒ Likely written by LM \*(LM = OPT-1.3B here)



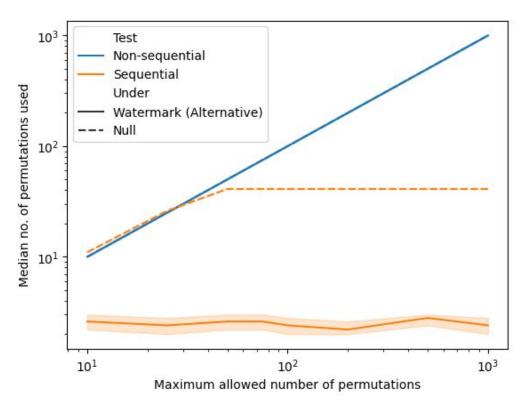


Power and Null Rejection Rate for c4 experiment without corruption text len (m)=80, key len (n)=256, # of texts (T)=200, alpha=0.05, c=0.04



# Results: Number of Permutations to Decision





Less Permutations  $\Rightarrow$  Less Time to Decision  $\Rightarrow$  Early Stopping  $\Rightarrow$  Less Computations

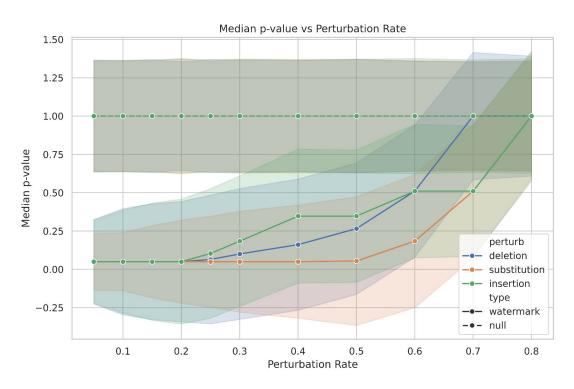
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# **Experiment Setup: Testing Robustness**

- $\rightarrow$  Do T = 200 permutations on a text of length m = 80
- → Remove/Insert/Substitute tokens from the generated output at rates ranging from 0.05 to 0.8 randomly.
- → For each corruption rate, report:
  - the Average of Median p-values
  - ◆ Null Rejection Rate
  - ◆ Empirical Power
  - ♦ Median Permutation to Decision across Runs



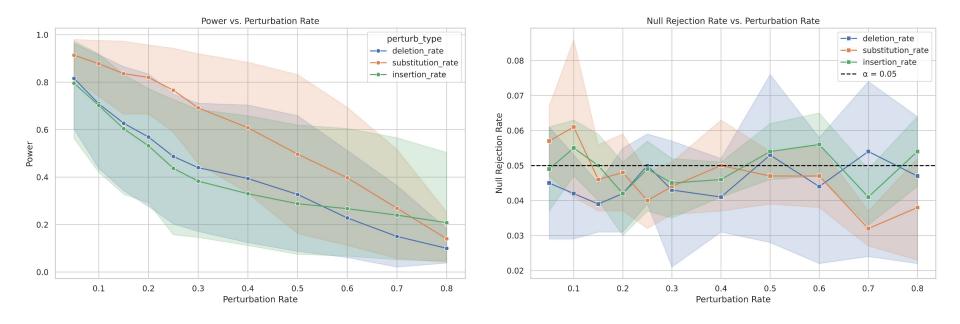
# Results: Robustness of Watermark



Median p-values remain low under perturbations ⇒ Watermark is Robust to attacks

# Results: Power & Null Rejection Rate under Attack

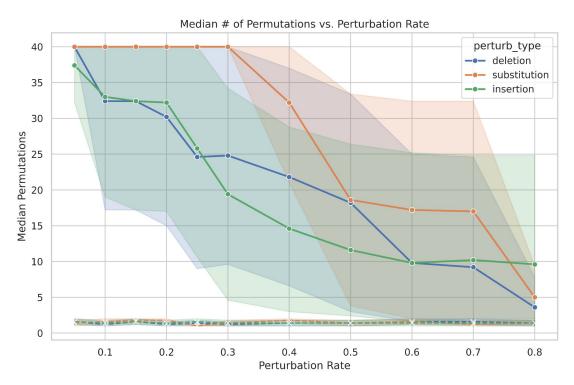




Power close to 1 + Null Rejection Rate below 0.05 ⇒ Test remains valid under attack \*Although the statistical confidence of the decision is impacted

# Results: Number of Permutations to Decision under Attack





Perturbation Increase ⇒ Watermark becomes weak ⇒ Algorithm initiates Early Stopping \*Still taking lesser time than permutations to make a decision



# Analysis

Substitution is most robust to corruption

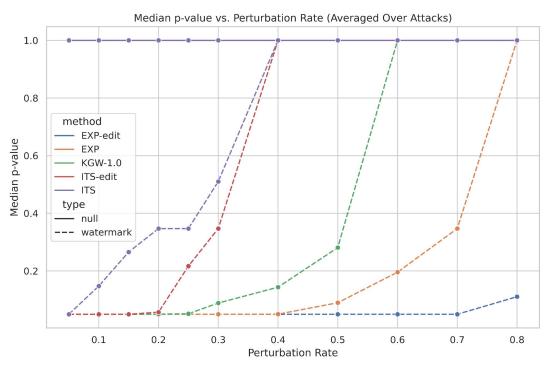
Null rejection rate stays close to  $\alpha$  = 0.05, indicating that the false positive rate is well-controlled.

Require far less permutations

Require less to test null output



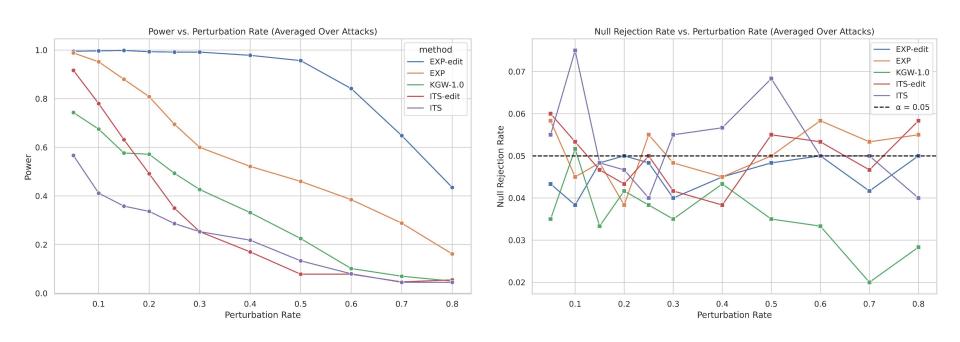
# Results: Comparing different generate algorithms



**EXP-edit method is most robust to attack outperforming EXP method!** 

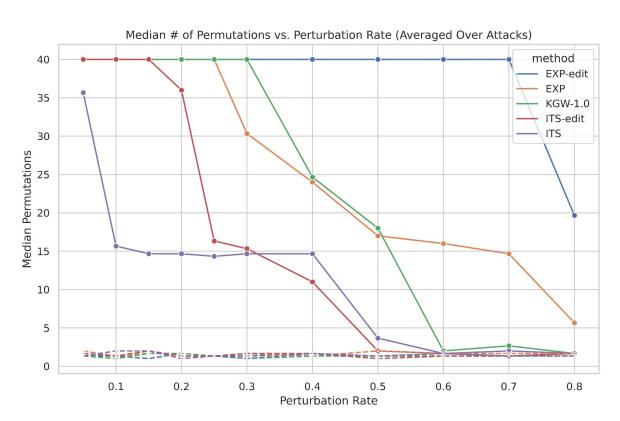


# Results: Performance under Attack





# Results: Performance under Attack





## Conclusion

- → Robust Sequential Monte Carlo Test Implemented: Developed a reliable and scalable SMC-based framework for watermark detection under real-world conditions.
- → Outperform the Permutation Tests: Demonstrates significantly higher efficiency and greater robustness, especially in limited-sample or corrupted data settings.
- → **High Statistical Power**: Power evaluations show strong and consistent detection capability, confirming the test's reliability across diverse scenarios.
- → Robust to Corruptions: Maintains performance under token-level corruptions such as substitutions and deletions, making it practical for noisy or adversarial text.