



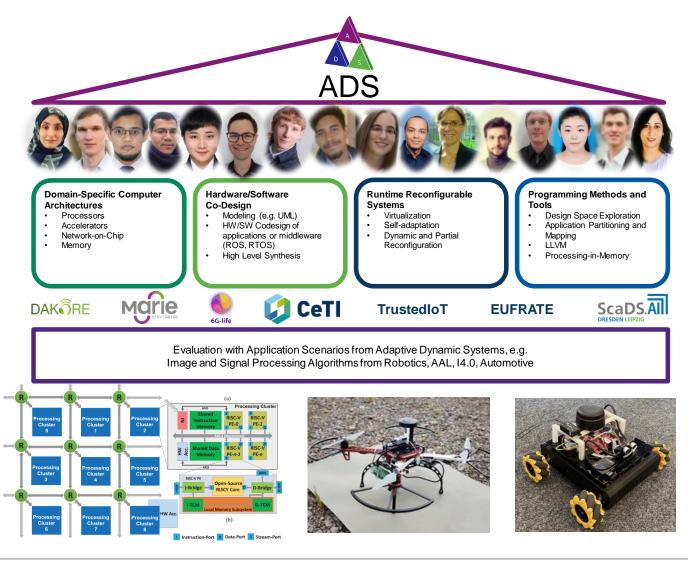
Cornelia Wulf Chair of Adaptive Dynamic Systems TU Dresden

# TrustedIoT Trusted Computing Architectures for IoT Devices

16.04.2024

**The Chair of Adaptive Dynamic Systems** 

What we do?







## The Chair of Adaptive Dynamic Systems

#### The Team



**Chair:** 

Ms Prof. Dr.-Ing. Diana Goehringer



**Postdoctoral Researcher:** 

Dr. Sergio Pertuz



**PhD Candidate:** 

Ms Dipl.-Inf. Cornelia Wulf

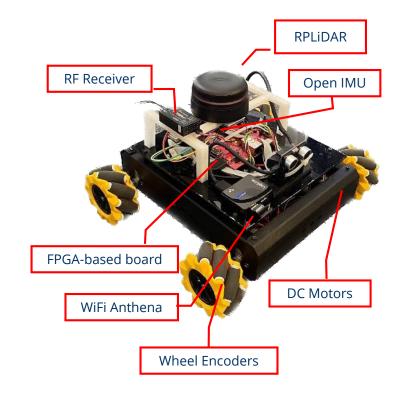




# **Low-power FPGA-based mobile robot with enhanced IoT security:**Motivation

- Software side:
   Hypervisors isolate trusted from untrusted guest operating systems.
- Hardware side:
   Fine-grained isolation mechanism for shared usage of hardware accelerators is missing.

Focus on AXI memory-mapped interfaces

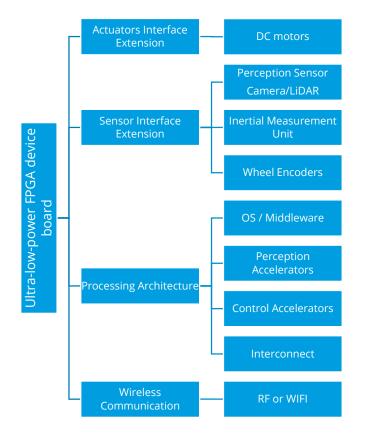


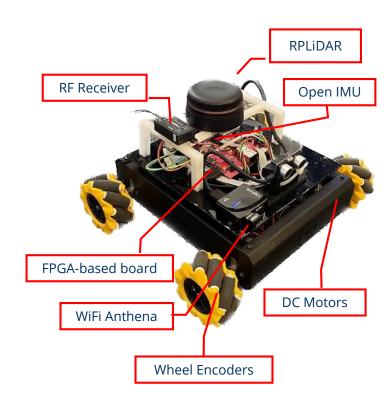




## Low-power FPGA-based mobile robot with enhanced IoT security:

## **Proposed Architecture**







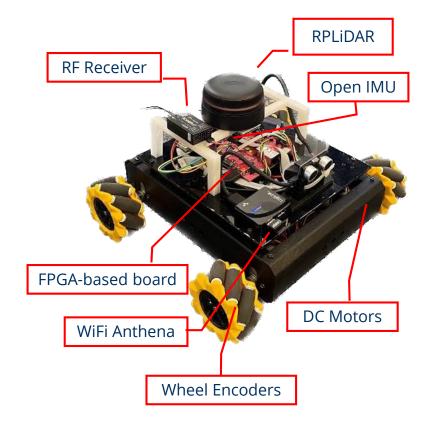


# Low-power FPGA-based mobile robot with enhanced **IoT security:** Platform



**KR260 Robotics Starter Kit:** 

Zynq UltraScale+™ MPSoC EV (XCK26)

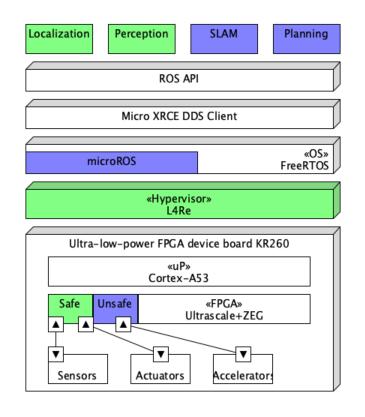


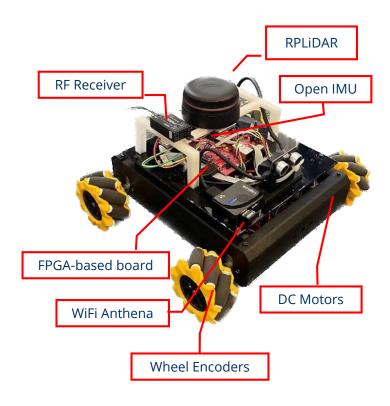




## **Low-power FPGA-based mobile robot with enhanced IoT security:**

## **Architecture Stack**



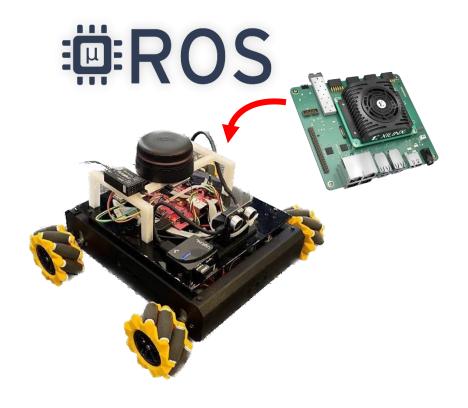






# Low-power FPGA-based mobile robot with enhanced IoT security: Software and Middleware

TUD implemented **ROS2/microROS** (novelty) on a Zynq/ZynqMP device and add a **hardware-based secure layer** to the networking and middleware to have improved security in robotics IoT.







## Threat analysis overview: Scope

The security scope here is limited to robot middleware and medium-level software vulnerabilities. In particular, **three types of attackers** are considered:

- Human attackers interact physically with the robot (Robot User),
- Another robot or system is capable of physical interaction with the robot. (Third-Party Robotic System), and
- A human teleoperation the robot or sending commands to it through a client application (e.g., smartphone app) (Teleoperator / Remote User)

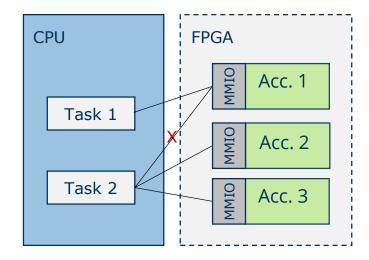






#### **Protection of hardware accelerators**

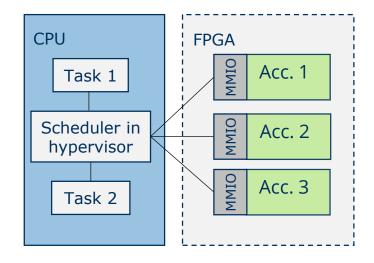
1. Fixed assignment:



#### Disadvantage:

- No flexibility
- No scalability

2. Access via software scheduler:

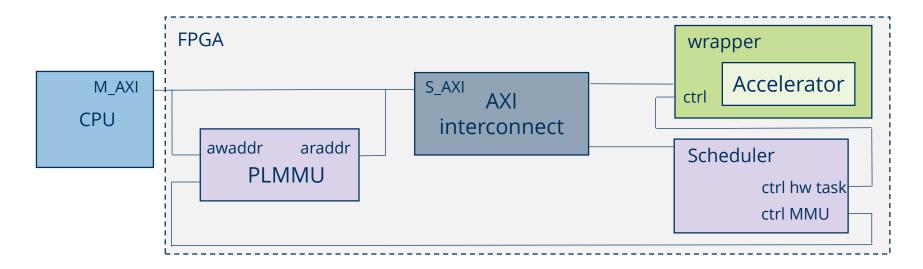


#### Disadvantage:

Latency







#### **Custom MMU**

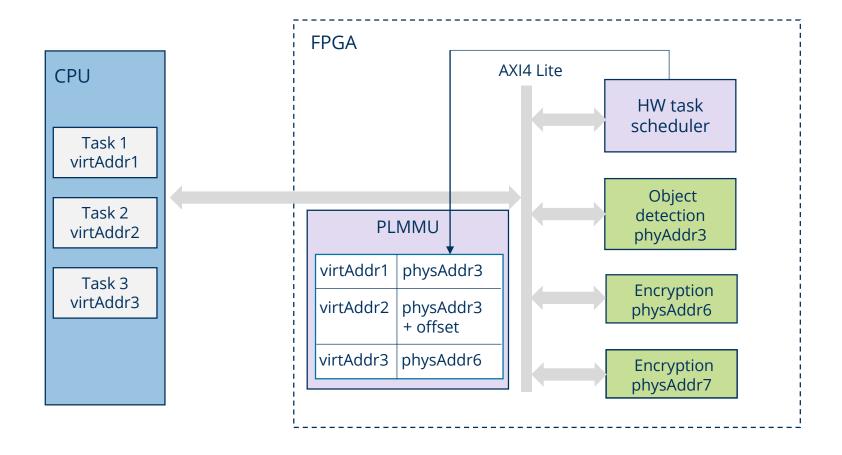
#### Scheduler

- Vitis HLS
- Input: Task ID, accelerator type, priority
- Chooses accelerator and updates the translation table

Priority queue for each accelerator

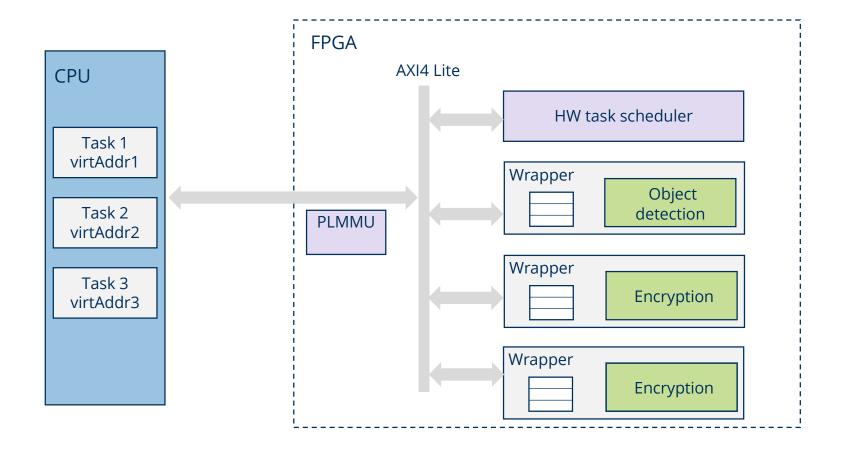
















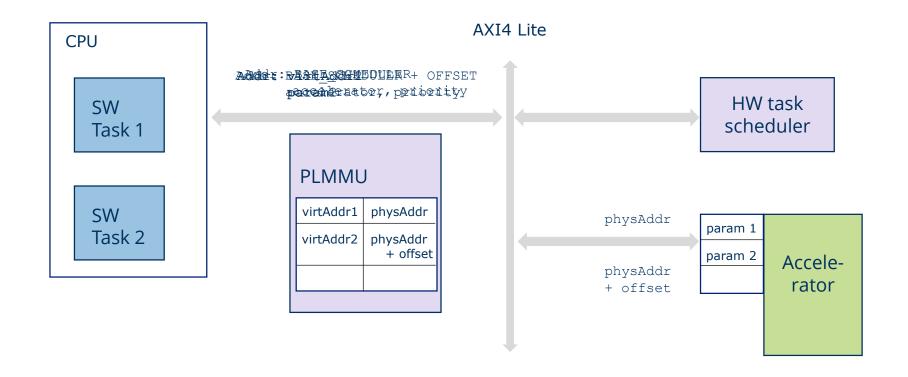
#### l4rec.io

```
Io.hw.add devices(function()
       hw scheduler = Io.Hw.Device(function()
              Resource.regs = Io.Res.mmio(0xA000000, 0xFFFFFFFF)
              Resource.irq1 € Io.Res.irq(121);
       end);
       task1 = Io.Hw.Device(function()
              Resource.regs = Io.Res.mmio(BASE SCHEDULER, BASE SCHEDULER + OFFSET)
              Resource.regs = Io.Res.mmio(BASE VIRTI, HIGH VIRTI)
       end);
       task2 = Io.Hw.Device(function()
              Resource.regs = Io.Res.mmio(BASE SCHEDULER + OFFSET,
              Resource.regs = Io.Res.mmio(BASE VIRT2, HIGH VIRT2
       end);
End)
```





## Hardware task scheduler

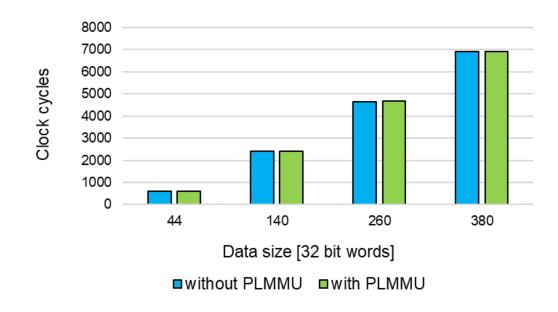






### **Evaluation**

- Hardware accelerators: Encryption
- PLMMU adds negligible overhead to the communication between software task and hardware accelerator
- For a First Come First Serve strategy, the hardware task scheduler requires on average 169 clock cycles at 100 MHz.







### Conclusion

Memory-mapped access controlled by a PLMMU and a hardware task scheduler

#### Advantages

- Prevention of unauthorized access
- Shared usage of hardware accelerators
- Preservation of priorities
- Latency reduction compared to a software approach

#### Disadvantage

Overhead





## **Future work**

- Static and / or dynamic priorities
- Virtualization of interrupts





# Questions?

## Remarks

## Discussion



