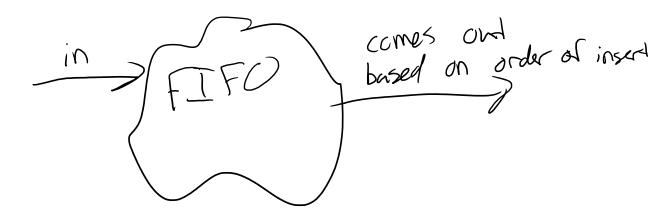
# What is a queue



# How would you write your own queue?

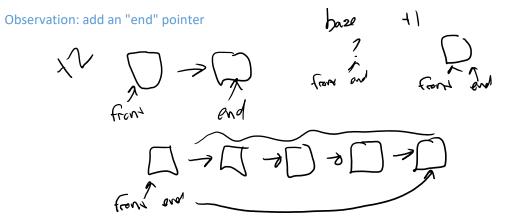
### Idea #1: Use a vector

- "push back" new items. Means that items that arrived early are in the front and that items that arrived late are in the back.
  - o Assuming that we have enough space, O(1) behavior
- Dequeue (or C++ pop) Remove the front item, shift everything over (left, or by -1) by 1.
  - o Requires FOR loop. O(N) behavior

#### Idea #2: Use a linked list

- Enqueue: add next item to the end of the list
  - o O(N) complexity
- Dequeue: remove item from front of list, alter pointer
  - O(1) complexity

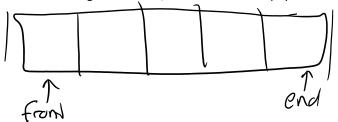




- Using an "end" pointer allows us to add items to the end of the linked list in O(1) time
- Thus, enqueue become O(1) on a linked list

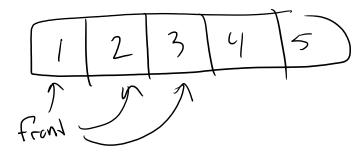
## Class Exercise: Improving the Vector-based implementation

• When working with vectors, we tend to think in physical boundaries



• Instead, consider logical or "floating" boundaries. Consider the following code:

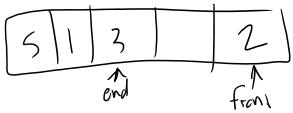
int front = 0; int value = data[front]; front++; return value;



- Downside: empty space at the front of the vector. Once we dequeue from
  a slot, we never use that space again. Thus if we do 1000 enqueues, 1000 dequeues,
  and 1000 more enqueues, our vector would be at least 2000 big.
- Solution: is to make the end of the vector a logical pointer. Consider:

Int end = vector.size() - 1;
Data[end] = value
End = (end + 1) % vector.size();

• Doing so can result in a situation in which the "front" of the queue is ahead of the "end" of the queue



• This implementation is called a Circular Queue. By using logical front and ends on a vector, a circular queue achieves O(1) enqueue and dequeue.

### "Unwinding" a circular queue when expanding

• Double the size of our data

