





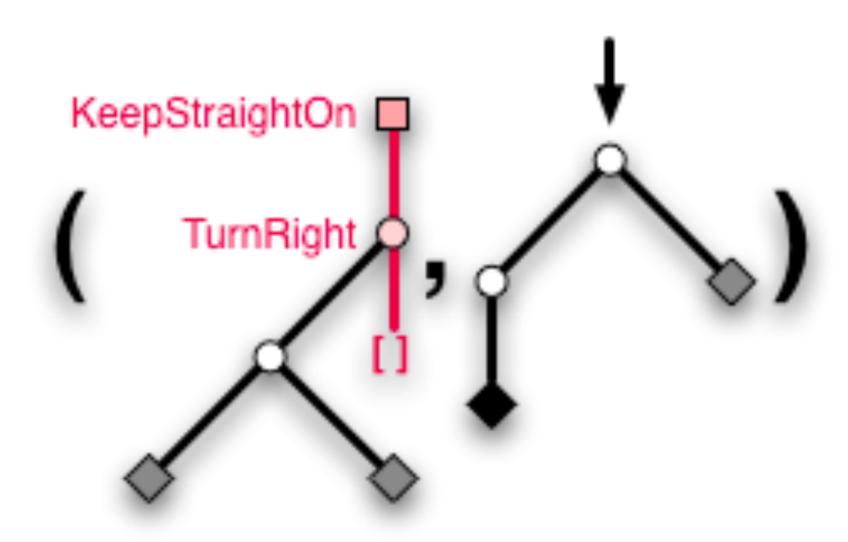
The Space Calculus in context

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Background and motivations from maths and computer science

The notion of location as a splitting has a long history in mathematics. From Dedekind cuts to Conway games this idea has born a great deal of fruit. Huet's idea of the zipper should be seen in this light.



And it generalizes to a wide range of functors (data type constructors). Formally, the zipper-based notion of location, for a given data type T, is ∂T x T where ∂T is the type of 1-holed T-contexts.

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The domain equation

The idea is to apply this notion of location to a process calculus so that names really are locations.

We begin with the domain equation which generates the rho-calculus.

$$P[X] = 1 + X \times P[X] + X \times X \times P[X] + P[X] \times P[X] + X$$

 $RP = P[RP]$

But now we want not quoted processes but locations, ∂P[X] x P[X]

$$LP = P[\partial P[LP] \times P[LP]]$$

This was the equation i wrote down in 2008, but couldn't find a solution for until 2018



Background and motivations from physics and chemistry

Oh, wait, why do we want an equation like this?

The intuition in the Einstein's field equations is that — when dynamics are taken into account — the metric is recursively intertwined with the stress-energy tensor.

$$R_{\mu
u}-rac{1}{2}Rg_{\mu
u}+\Lambda g_{\mu
u}=\kappa T_{\mu
u}.$$

We can get a discrete version of this idea using a process calculus where where the dynamics links location and synchronization.

To make this work, we need one additional ingredient, which we borrow from chemistry: the idea of a catalyst. So, a reduction can only happen in the presence of a catalyst.

In the calculus presented below, the enabling entity is more of a co-factor than a catalyst, but there are versions of the calculus where it is strictly a catalyst.



Background and motivations from computer science, physics, and biology

*X

stopped process

input-guarded continuation

output

parallel composition

deref

$$y{@Q/y} = Q$$

COMM for(x <- y)P | x!(Q) ->
$$P\{@Q/y\}$$

$$COMM_{K}$$
 for(x <- y)P | x!(Q) -> P{@K[Q]/y}

Has applications to protocol interoperability and error correction

$$COMM_{K1}$$
 o $COMM_{K2}$ for $(x < - y)P | x!(Q) -> P{@K2[K1[Q]]/y}$ composes

But is static. The context is not programmable.

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The central proposal

```
stopped process
P,Q ::= 0
                      location update
     U(x)
                      input-guarded continuation
       for (y < -x)P
                      output
       x!(Q)
       PIQ
                      parallel composition
                      deref
       *X
       COMM(K)
                      situation catalyst
x, y ::= @<K, Q>
K := [] | for(x <-y)K | x!(K) | P|K
*y{@<K,Q>/y} = K[Q]
COMM( K ) | for( x <- y )P | x!( Q ) -> P\{@<K,Q>/y\}
U(x) + *@<K,Q> -> COMM(K) + x!(Q)
```

A simple calculation

A simple calculation shows....

If K1 is of the form

If K2 is of the form

then P moves through contexts K1, K2, K3, ...



Some desiderata and future work

While this represents a discrete version of the intuitions underlying GR it is decidedly not quantum. As far as my understanding goes, there are no non-local correlations.

However, we do have about 2/3 of an encoding of quantum mechanics in rho-calculus and so it is conceivable that this encoding could be composed with the techniques developed here.