

CS 321: Assignment 2

Jared Wasinger

October 14, 2016

1. (a) $\{w \in \{a, b\}^* \mid w \text{ has an even number of b's} \}$ $(a^*ba^*ba^*)^*$
(b) $\{w \in \{a, b\}^* \mid w \text{ does not contain the substring } ab\}$
 b^*a^*
Explanation: The restriction of not allowing 'ab' means that all b's in the string must come before all a's.
(c) $\{w \in \{a, b\}^* \mid w \text{ contains the substring } ab \text{ an even number of times} \}$
 $((b^*a^*)ab(b^*a^*)ab(b^*a^*))^*$
Explanation: This is similar to (a). Each occurrence of 'ab' must match at least twice and must match an even number of times. The sub-regex (b^*a^*) represents the regex matching all strings that are not 'ab' (the only allowed filler text to occur between occurrences of 'ab'). This is in order to enforce that 'ab' only occurs an even number of times while making allowances for all other sub-strings to occur at any place in the string.
(d) $\{w \in \{a, b\}^* \mid w \text{ has an even number of a's and an even number of b's} \}$
 - i. All strings of length = 2 $\in \{a, b\}^*$
 $\{aa, ab, ba, bb\}$
 - ii. All strings with an even number of b's and a's are formed by:
 $aa, bb, 2 * (ab \text{ or } ba) = (ab + ba)(ab + ba)$
 - iii. Regex:
 $((aa + bb)^* + (ab + ba)(ab + ba))^*$
 - iv. Edge cases and assumptions:
The empty string is accepted by this regex because 0 is an even number. Thus an empty string (0 b's, 0 a's) has an even number of b's and a's.
2. (a) Let M be the language of all binary strings divisible by 3. For each string $x \in M$, $reverse(x) \in M$