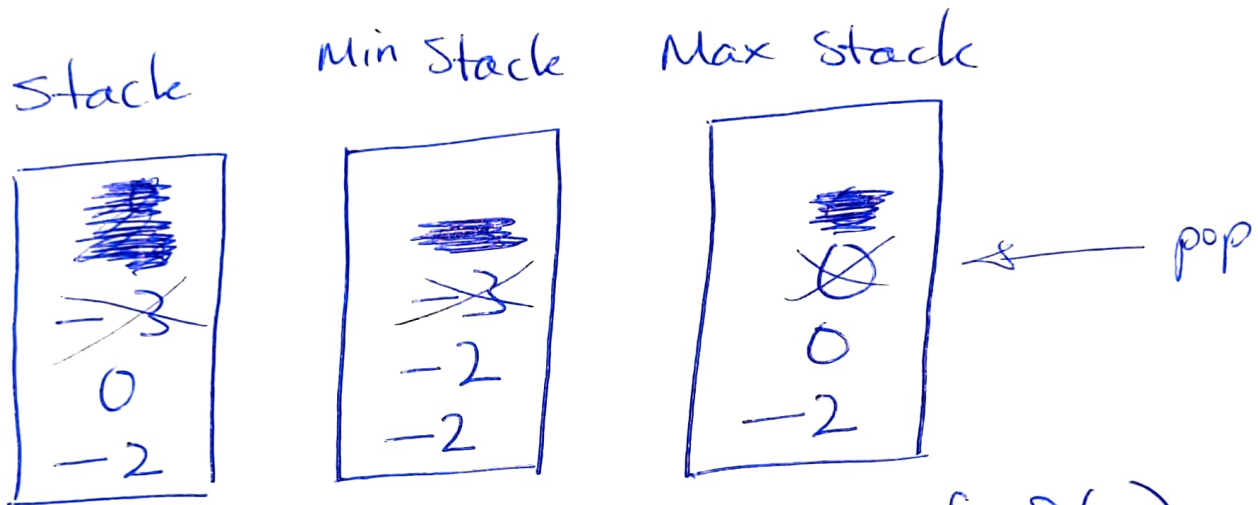


Implement a stack w/MaxAPI

8.1 7/8/22

- For this problem, we need to define another stack w/ the minimum/maximum in the order of the values pushed/popped



- This way, we can have runtime of $O(1)$ since we are comparing the pushed value w/ the ~~previous~~ min/max value, at the cost of $O(n)$ space.
- Here, we create a parallel data structure.
- This problem is less algorithmic and has more to do w/ using OOP and designing data structure.

Alternative Approaches

- Iterate through entire stack, pop every element and push them back in to check the values and return the min/max. $\rightarrow O(n), O(1)$
- $O(\log n), O(n)$ can be achieved using heap or BST, and a hash table.

```
class MaxStack (objed):
```

```
    def __init__ (self):
```

```
        self.stack = []
```

```
        self.maxes = []
```

```
    def push (self, val):
```

```
        self.stack.append (val)
```

```
        # IF the max value from the maxes stack  
        is greater, add that to stack otherwise,  
        push that value that is greater than the max
```

```
        if self.maxes and self.maxes[-1] > val:
```

```
            self.maxes.append (self.maxes[-1])
```

```
        else:
```

```
            self.maxes.append (val)
```

```
    def pop (self):
```

```
        if self.maxes:
```

```
            self.maxes.pop()
```

```
        return self.stack.pop()
```

} # Need to make
sure we also pop val
from maxes stack

```
    def max (self):
```

```
        return self.maxes[-1]
```

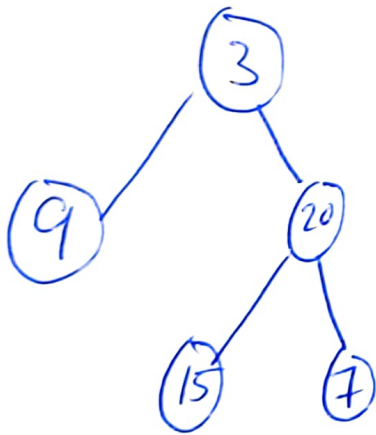
8.6 7/8/22

First set an empty list to return. Store the first node.
Then for each depth, store the children node
to a queue, ~~and~~; store the elements
in the queue to a list and append
that to the result. Do this until
all nodes are processed.

depth
0

1

2



queue

[9, 20]

[15, 7]

[]

result

[[3]]

[[3], [9, 20]]

[[3], [9, 20], [15, 7]]


```
def binary_tree_depth_order(tree: BinaryTreeNode)
    → List[List[int]]:
```

```
    result: List[List[int]] = []
```

```
    if not tree:
```

```
        return result - result
```

Queue → curr_depth_nodes = [tree] ~~# Nodes @ current depth.~~

```
    while curr_depth_nodes:
```

```
        nodes_at_depth = []
```

```
        for node in curr_depth_nodes:
```

```
            nodes_at_depth.append(node)
```

```
        result.append(nodes_at_depth)
```

```
        next_depth = []
```

```
        for node in curr_depth_nodes:
```

```
            for child in node(node.left, node.right):
```

```
                if child:
```

```
                    next_depth.append(child)
```

```
        curr_depth_nodes = [] # empty the queue
```

append nodes in the depth to the result

Process the nodes in the next depth

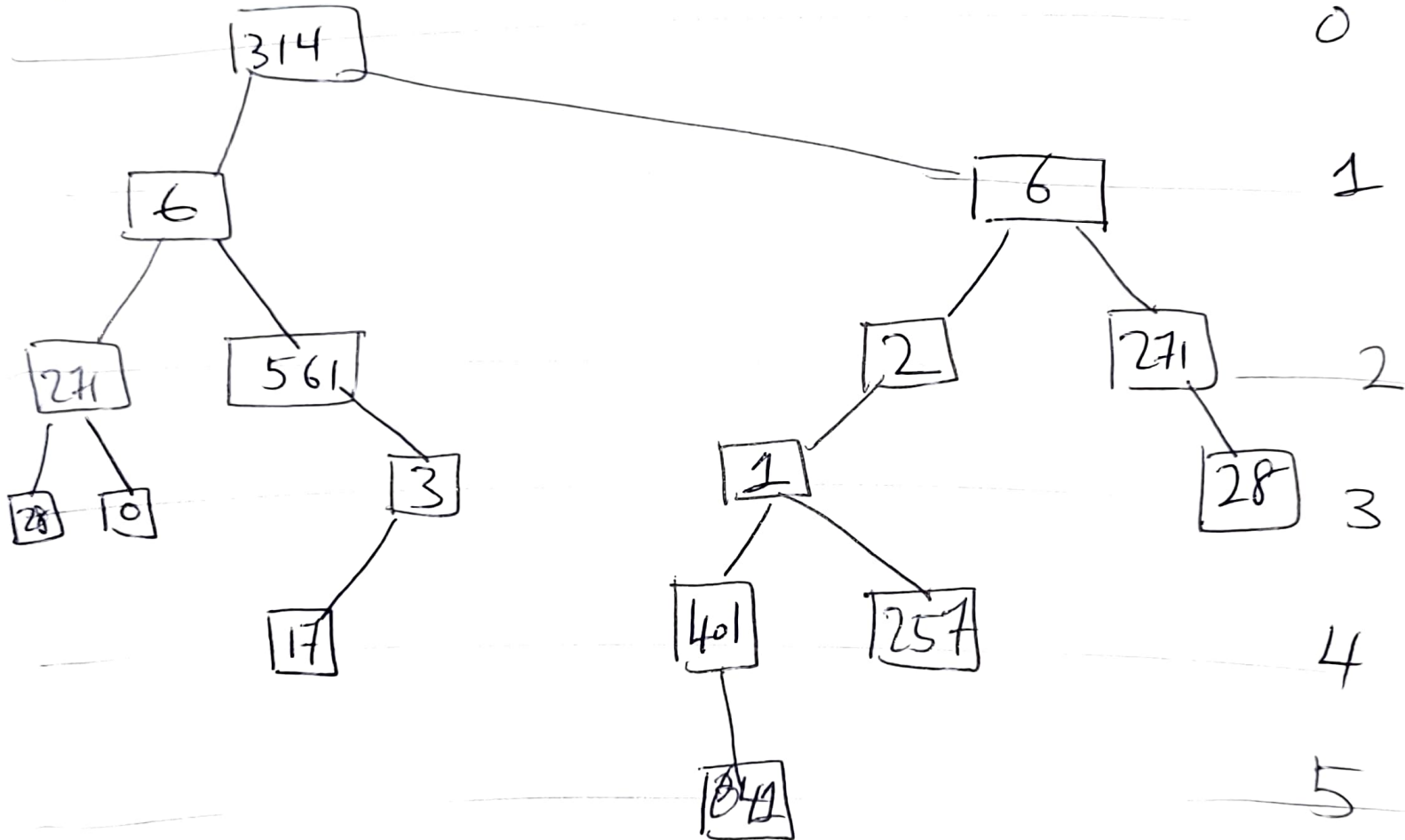
```
# prepare the next queue  
for node in next_depth:  
    curr_depth_nodes.append(node)
```

```
return result
```

Input: Binary Tree

Output: Array

In:



Should return:

$\langle \langle 314 \rangle, \langle 6, 6 \rangle, \langle 271, 561, 2, 271 \rangle \dots \rangle$
 (d0) (d1) (d2)

So we are return a `List[List[int]]`
~~where the elements are entries~~ for each depth. 1/3

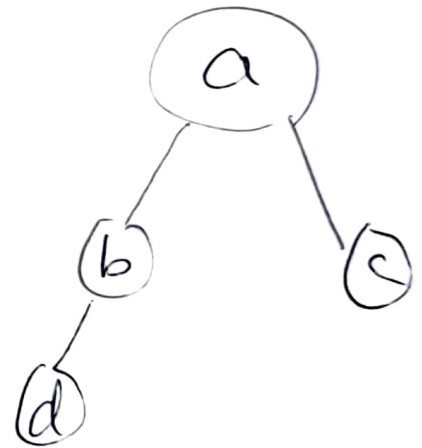
• I don't know how the bin. tree input looks like.

Bin Tree will look like...

```
class Node(object):  
    def __init__(self, val data, left=None, right=None):  
        self.data = data  
        self.left = left  
        self.right = right
```

and the ~~inputs~~ inputs:

```
root = Node('a')  
root.left = Node('b')  
root.left.left = Node('d')  
root.right = Node('c')
```



13min

- So the problem is asking for a class that has a method for getting the depth and for each depth, append all the values into a list and append to another list to return the list of values. ~~upto~~ (x depths)

~~Since~~

Boundary / edge cases:

↳ Since this is a binary tree, at each depth the max amount of elem is 2^{depth} .

⊛ How can I use queue in this?

⊛ What is the Naïve Approach?

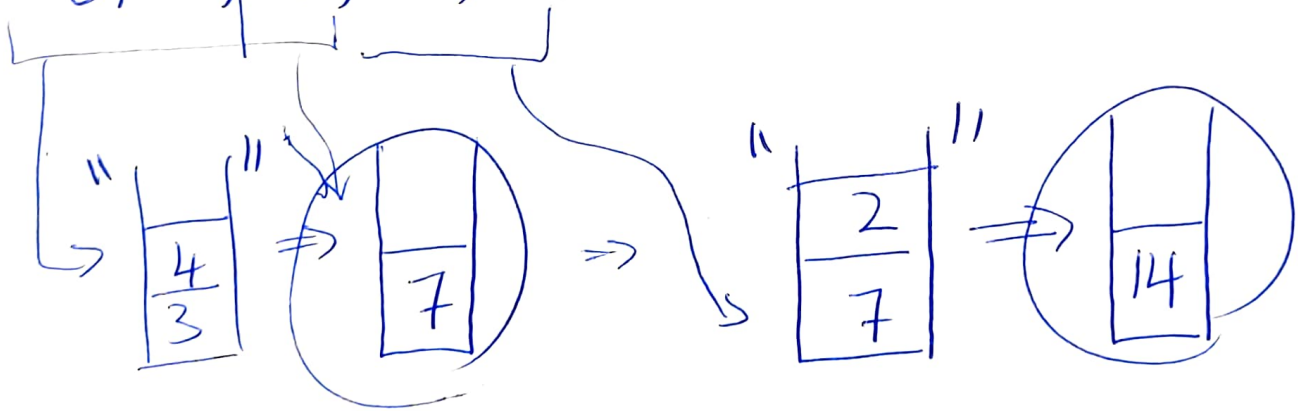
→ Maybe I can append to a list the max amount of elem, checking each branch.
Where $\text{node.left} \neq \text{null}$; ~~result.append~~ $\text{temp_list.append}(\text{node.left})$
and then $\text{result.append}(\text{temp_list})$

→ Since we are check all branches $\rightarrow O(2^{\text{max}(\text{depth})})$ time
 $O(\text{depth})$ space $\frac{3}{3}$

Evaluation

↳ The design was similar but not correct. When doing the computation, I had to observe the A and B by popping them and do the computation. Then push the value back into the list.

"3, 4, +, 2, x, 1, +"



⇒

15

⊛ For each number value (not operators) we convert and push/append as integer.

8.2

Input : RPN ; Output : ~~int~~ ~~float~~
Number

"3, 4, +, 2, x, 1, +"

$\underbrace{3, 4, +}_{=7}$
 $\underbrace{\quad\quad\quad}_{=14}$
 $\underbrace{\quad\quad\quad}_{=15}$

"-641, 6, /, 28, /"

$\underbrace{(-641, 6, /)}_{=706.833}$ ← Must be an integer, round up or down?
 $\underbrace{\quad\quad\quad}_{\approx 3.82}$

• Are we returning an \mathbb{Z} ? or float?

The Naive Approach would be to iterate through the string for a math symbol, convert the numbers before the symbol and compute the calculation.

We can take that number value and traverse the string until next math symbol or ^(the) end.

We can do this recursively.

$O(n)$, $O(1)$
 \hookrightarrow string \uparrow
 ? $\frac{1}{2}$

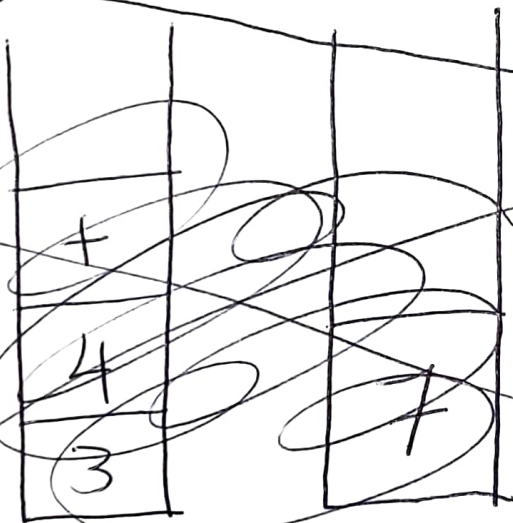
13min

We might be able to use stacks to store the computed value and compute again when there is a symbol.

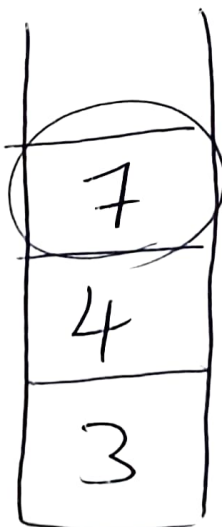
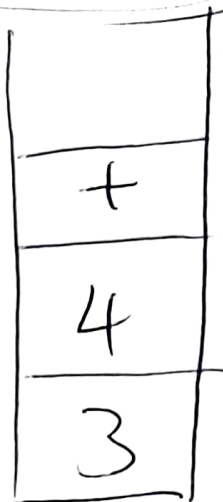
To compute ~~in~~ accordance to the symbol, we might have to use a dictionary.

This can reduce the time but use extra memory ($O(1)$, $O(n)$)

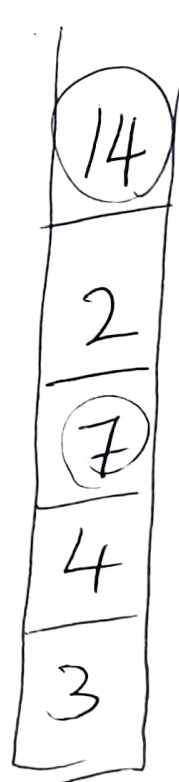
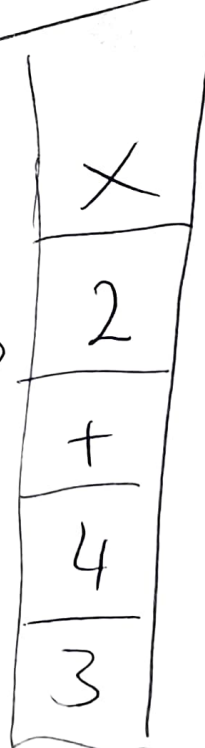
LIFO



"3, 4, +, 2, x"



=>



20min Quit

2/2