

```
def min_length_subarr(s s, A A) :
```

```
    w_sum, w_st = 0, 0
```

```
    min_l = math.inf
```

```
    for w_end in range(len(A)):
```

```
        w_sum += A[w_end]
```

```
        while w_sum >= s:
```

```
            min_l = min(min_l, w_end - w_st + 1)
```

```
            w_sum -= A[w_st]
```

```
            w_st += 1
```

```
    return min_l
```

```
def max_subarray_of_size_k(k, arr):  
    max_sum, window_sum = 0, 0  
    window_start = 0    #index val.  
    for window_end in range(len(arr)):  
        window_sum += arr[window_end]  
        if window_end >= k - 1:  
            max_sum = max(max_sum, window_sum)  
            window_sum -= arr[window_start]  
            window_start += 1  
    return max_sum
```

Implement an inorder
traversal w/ constant space

9/10

7/14

Input: Bin. Tree, Output: list

Recursive Approach

def f(root: Bin. Tree) → List[int]:
res = []

def inorder(root):

if not root:
return

inorder(root.left)

res.append(root.data)

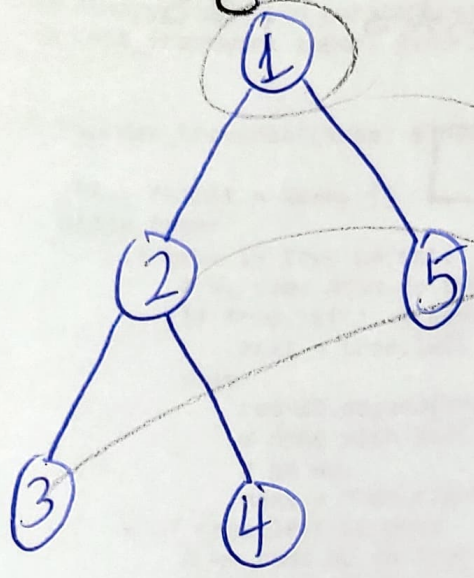
inorder(root.right)

inorder(root)

return res

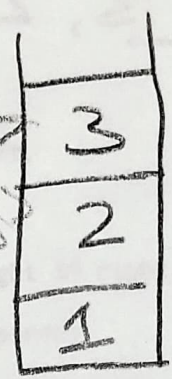
~~return res~~

Using stack



res = []

Stack



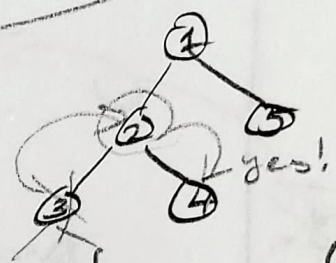
At this point, there is no left subtree. Now we pop() and add to the res. list.

function call stack

When calling func. inside a func. it has to remember to get back to the orig. func.



res = [3]

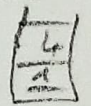


Now we check if the node has a right subtree. Since there is no right subtree, we pop again to the res.

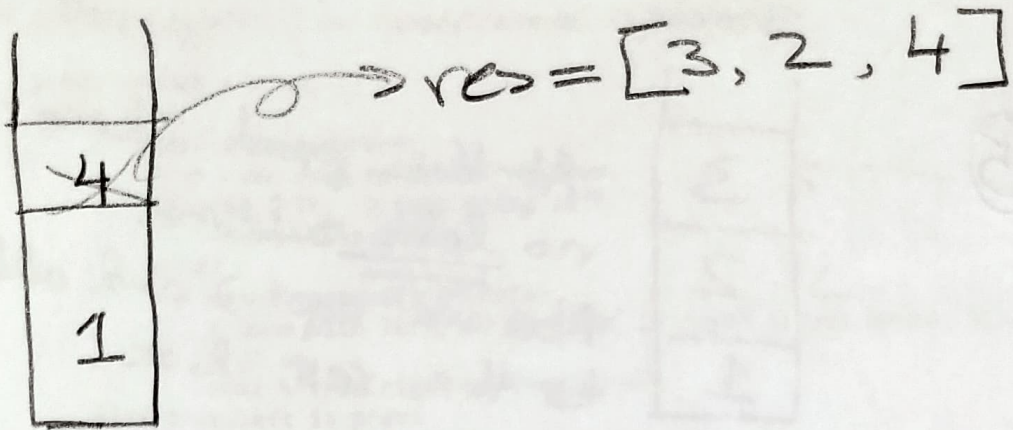


res = [3, 2]

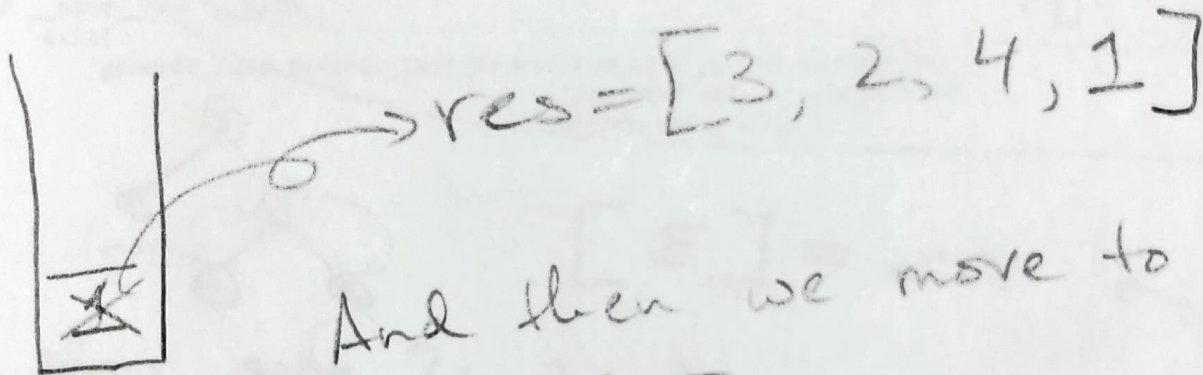
From node 2, we have right child. So push that node to the call stack and check it's children. →



Since at node w/ value 4 has no left children, we pop that into res.



Since node w/ value 4 does not have any right child we pop the ~~1~~ elem. on the stack to res.



And then we move to node w/ 5.

... $\Rightarrow res = [3, 2, 4, 1, 5]$

9/10

```
def inorder_traversal(tree: BinaryTreeNode) -> List[int]:
```

```
while tree:
```

```
# We came down to tree from prev.
```

```
next = tree.left
```

```
result.append(tree.data)
```

```
result.append(tree.data)
# Done with left, so go right if right is not empty. Otherwise,
```

```
# go up.
```

```
next = tree.right or tree.parent
```

```
# We came up to tree from its left child.
```

```
result.append(tree.data)
```

```
result.append(tree.data)
# Done with left, so go right if right is not empty. Otherwise, go
```

up.

```
next = tree.right or tree.parent
```

```
else: # Done with both children, so move up.
```

```
next = tree.parent
```

```
prev, tree = tree, next
```

```
return result
```

```
exit(
```

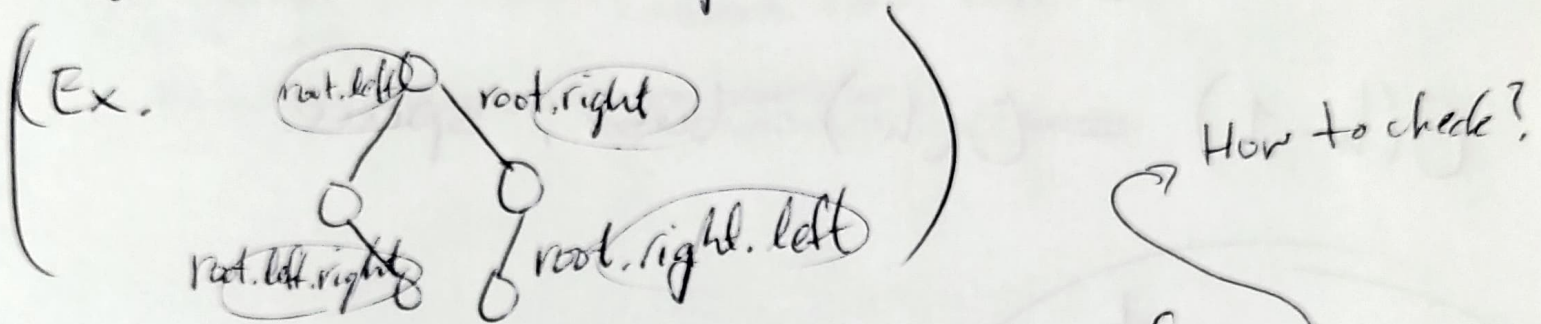
```
generic_test.generic_test_main('tree_with_parent_inorder.py',
                               'tree_with_parent_inorder.tsv',
                               inorder_traversal))
```

9.2 7/14

Input: Bin. Tree

Output: boolean

- Need to check if ~~the~~ search path is same from both sides.
- Need to have nodes points to opposite?



↳ Not sure if I need to check this.

- First need to check if depth is same
↳ return False, if no.

- Starting from the leaves, we need to check if nodes value ~~are~~ are the same for both sides.

- Empty \Rightarrow return ~~the~~ False

- Only root \Rightarrow return True

→ Do I need to check levels?

- Naive Approach

- ↳ Use hash map.

- ↳ Store all nodes into hash except root of one side

- ↳ Check other side

$O(h-1) = O(h)$ time, space

11 min quit

Use Recursion

Given binary tree root node return binary tree root
 Given root of binary tree return root

```

def isSymmetric(root: BinaryTreeNode) -> Boolean:
    def checkSymmetric(left: BinaryTreeNode, right: BinaryTreeNode):
        if left == None and right == None:
            return True
        if left == None or right == None:
            return False
        if left.val != right.val:
            return False
        return checkSymmetric(left.left, right.right) and checkSymmetric(left.right, right.left)
    return checkSymmetric(root, root)

// Test for root
// Test for left
// Test for right

```

```
def is_symmetric(tree):
```

```
    def check_symmetric(subtree0, subtree1):
```

```
        if not subtree0 and not subtree1:  
            return True
```

```
        elif subtree0 and subtree1:
```

```
            return (subtree0.data == subtree1.data
```

```
                    and check_symmetric(subtree0.left, subtree1.right)
```

```
                    and check_symmetric(subtree0.right, subtree1.left))
```

```
            return False
```

```
    return check_symmetric(tree.left, tree.right)
```

```
    or not not tree
```

Compute LCA when
nodes have parent pointers

[9.4]

7/14

A brute force approach \rightarrow is to traverse one node and store them in a hash table and then traverse the other node and ~~return the~~ check for the first common node by using the hash table. $O(h)$ time, space.

Better approach.

- If two nodes are in the same depth, we can move ~~the~~ (traverse) the two nodes at the same time and return the ~~the~~ first common node.
- The problem lies when the two input nodes have different depths, then we need to move ~~one of~~ the node w/ greater depth to the same depth as the other node.


```
def lca(n0, n1):
```

```
    def get_depth(node):
```

```
        depth = 0
```

```
        while node.parent:
```

```
            depth += 1
```

```
            node = node.parent
```

```
        return depth
```

```
    depth0, depth1 = map(get_depth, (n0, n1))
```

```
    if depth1 > depth0:
```

```
        node
```

```
        n0, n1 = n1, n0
```

```
    depth_diff = abs(n1 - n0)
```

```
    while depth_diff:
```

```
        n0 n1 parent
```

```
        n0
```

```
        depth_diff -= 1
```

```
    while n0 is not n1:
```

```
        n0, n1 = n0.parent, n1.parent
```

```
    return n0
```


9.4

```
import functools
from typing import Optional
```

```
from binary_tree_with_parent_prototype import BinaryTreeNode
from test_framework import generic_test
from test_framework.binary_tree_utils import must_find_node
from test_framework.test_failure import TestFailure
from test_framework.test_utils import enable_executor_hook
```

```
def lca(node0: BinaryTreeNode,
        node1: BinaryTreeNode) -> Optional[BinaryTreeNode]:
```

```
    def get_depth(node):
```

```
        depth = 0
```

```
        while node.parent:
```

```
            depth += 1
```

```
            node = node.parent
```

```
        return depth
```

→ This is given

→ getdepth(node0), ... (node1)

```
    depth0, depth1 = map(get_depth, (node0, node1))
```

```
    # Makes node0 as the deeper node in order to simplify the code.
```

```
    if depth1 > depth0:
```

```
        node0, node1 = node1, node0
```

```
    # Ascends from the deeper node.
```

```
    depth_diff = abs(depth0 - depth1)
```

```
    while depth_diff:
```

```
        node0 = node0.parent
```

```
        depth_diff -= 1
```

```
    # Now ascends both nodes until we reach the LCA.
```

```
    while node0 is not node1:
```

```
        node0, node1 = node0.parent, node1.parent
```

```
    return node0
```

} when two nodes meet during traversal.

```
@enable_executor_hook
```

```
def lca_wrapper(executor, tree, node0, node1):
```

```
    result = executor.run(
```

```
        functools.partial(lca, must_find_node(tree, node0),
                           must_find_node(tree, node1)))
```

```
    if result is None:
```

```
        raise TestFailure('Result can\'t be None')
```

```
    return result.data
```

```
if __name__ == '__main__':
```

```
    exit(
```

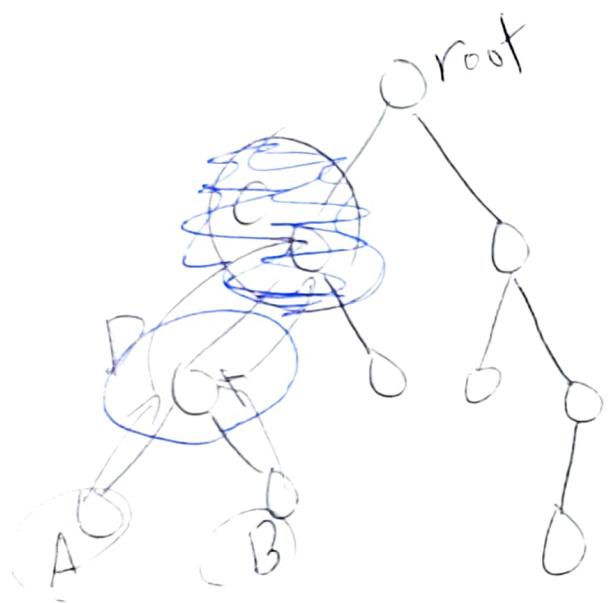
```
        generic_test.generic_test_main('lowest_common_ancestor_with_parent.py',
                                         'lowest_common_ancestor.tsv',
                                         lca_wrapper))
```

• Parent Pointer?

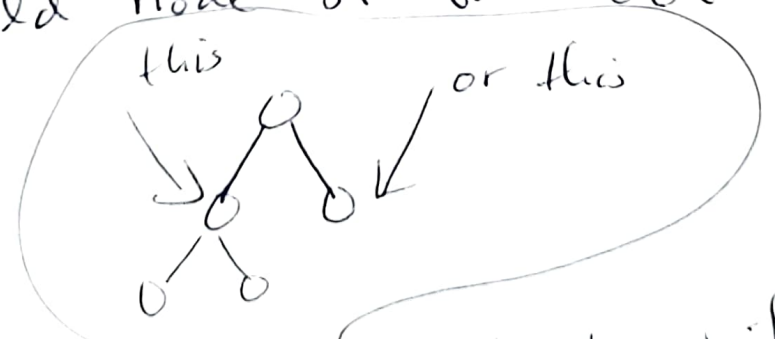
Input: Tree ~~(root)~~ nodes $(n1, n2)$

Output: Node (tree?) or some value

- Probably has to do w/ traversal.
- Maybe like 9.1 and start from the two leaves and traversal up to root.



• Wouldn't the LCA always be ~~the~~ one of the child node of the root?



10min
Quit

In this case we just need to know if the input nodes are in the left or right (or left & right)