

5.20

1/4

Input: integer ; Output: input number of lists

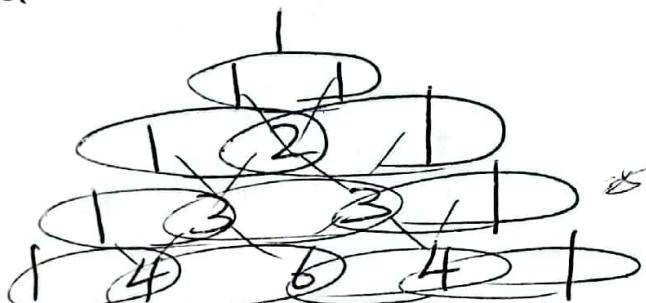
Edge Case: if  $r=1$ , return 1

Constraints:  $r > 1$  and  $r \in \mathbb{Z}$

The left and right most sides are 1 always  
There needs to be an equation to get the  
item within left and right most.

length

1  
2  
3  
4  
5



$r=1$

$r=2$

$r=3$

$r=4$

$r=5$

$$R_1[1] = R_3[0] + R_3[1]$$

$$R_4[2] = R_2[1] + R_2[2]$$

Output =  $\left[ [1], [1, 2, 1], [1, 2, 3, 1], [1, 2, 3, 4, 1] \right]$

We can use a for-loop for each row before  
the row we are calculating.

~~result = []~~, ~~new-row = [1]~~  
~~for i in range of prev. row length~~ input  
~~prev. row~~ starting with 1

$$\text{new-row}[i] = \text{result}[-1][i] + \text{result}[-1][i+1]$$

~~new-row.append(1)~~  
~~result.append(new-row)~~

this will take  $O(n)$  time  
and  $O(\frac{n(n+1)}{2}) \approx O(n^2)$  space

**Ex** input:  $r = 4$

~~result = []~~, ~~new-row = [ ]~~

#for loop

#i:

We need a while-loop first and increment upto the input r.

index = 1

# while  $i \leq 4$

new-row = [ ]

#for loop

new-row[1] = result[-1][0] + result[0][1]  $\quad i=1$

result = [1, [1, 1], [1, 2, 1]]  $\quad \text{#2}$

#for-loop

new-row[1] = ~~1~~ 1 + 2 = 3  $\quad i=1$

new-row[2] = 2 + 1 = 3  $\quad i=2$

result = [1, [1, 1], [1, 2, 1], [1, 3, 3, 1]]

2/4

```

def fun(r:int) → List[List[int]]:
    result = [[1], [1, 1]]
    if r == 1:
        return [[1]]
    if r == 2:
        return result
    if r == 0:
        return []
    index = 1
    while index <= r:
        new_row = [1]
        for i in range(len(result[-1])):
            new_row[i] = result[-1][i-1] + result[-1][i]
        new_row.append(1)
        result.append(new_row)
        index += 1
    return result

```

24 min  
 40 min to write code  
 50 min to pass all tests  
 3/4

## Evaluation:

Took too long trying figure out that I had to append to new-row

→  $\boxed{\begin{array}{l} \text{new\_row}[i] = m \times \\ \text{new\_row.append}(m) \checkmark \end{array}}$

The index had to start from the length of the already existing list (result=[[1], [1, 1]])  
→ Next time pay attention to the base case and keep track.

5.9

Input:  $n \in \mathbb{Z}$ ; Output: List [int] of primes

Constraints:

- Input is  $n \in \mathbb{Z}$  where  $n > 1$
- The items in the output list are primes so they do not have divisors other than 1 and itself.

How to check or generate primes?

~~Brute Force~~  
• Naive approach is to generate a uniform iterating list of numbers less than the input and del items that is divisible by other items in the list.

The first loop will start from the end of list and second ~~will~~ from start

Ex

~~[1, 2, 3, 4, 5]~~  
~~↳ [5, 4, 3, 2, 1]~~  
~~[1, 2, 3, 4, 5]~~

This will take  $O(n^2)$  and  $O(n)$ .

1/3

• Can I use pointers?

I] ~~Execute same approach as the~~ naive sol'n and make a new list

Note: I can use  $l = \text{list}(\text{range}(n))$

II] Or I can use an for loop, ~~and~~ check if prime, and append to empty list.

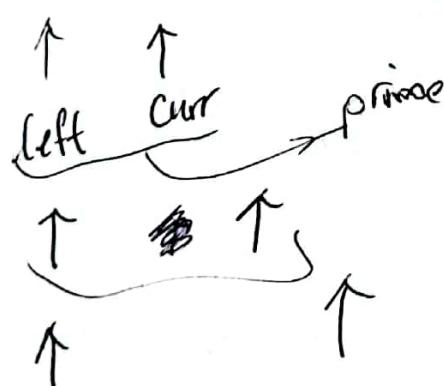
Both will be  $O(n)$   ~~$O(n)$~~  time

I] will be  $O(n)$  space

II] will be  $O(\# \text{of primes})$  space

---

Ex  $[2, 3, 4, 5, 6, 7]$



$[2]$

$[2, 3]$

$[2, 3, 5]$

$[2, 3, 5, 7]$

2/3

def fun( $n$ : int) → List[int]:

$l = \text{list}(\text{range}(n))$  or  $n$

~~for i in range(len(l)):~~

~~if l[i] <~~

while curr < len(A):

if  $l[\text{curr}] \% A[\text{left}] == 0$

~~return~~

$\text{curr} += 1$

$\text{left} = 0$

else:

}

3/3

5.5

Input: Sorted array ; Output: count of valid elem.

- Naive, brute force approach is to take an item in the list and compare it with ~~all~~ ~~of each~~ other items. Then increment, move on to the next item (and keep count of unique) This will take  $O\left(\frac{n(n+1)}{2}\right)$  time and  $O(1)$  space.   
 Grass
- We can use hashmap / hashset to iterate through the list and keep count in the dictionary. and return the number of keys in the dictionary. This will take  $O(n)$  time and space
- Another approach is to use pointers and check if the subsequent item is same or greater (since the input list is sorted) and keep count using conditionals. This will take  $O(n)$  time and  $O(1)$  space

1/4

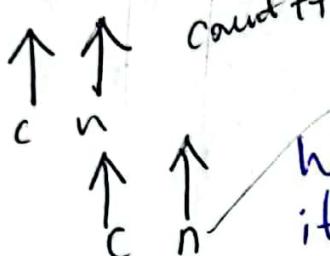
## Example

Starting @  
count = 1

[2, 3, 5, 5, 7, 11, 11, 11, 13]



count ++

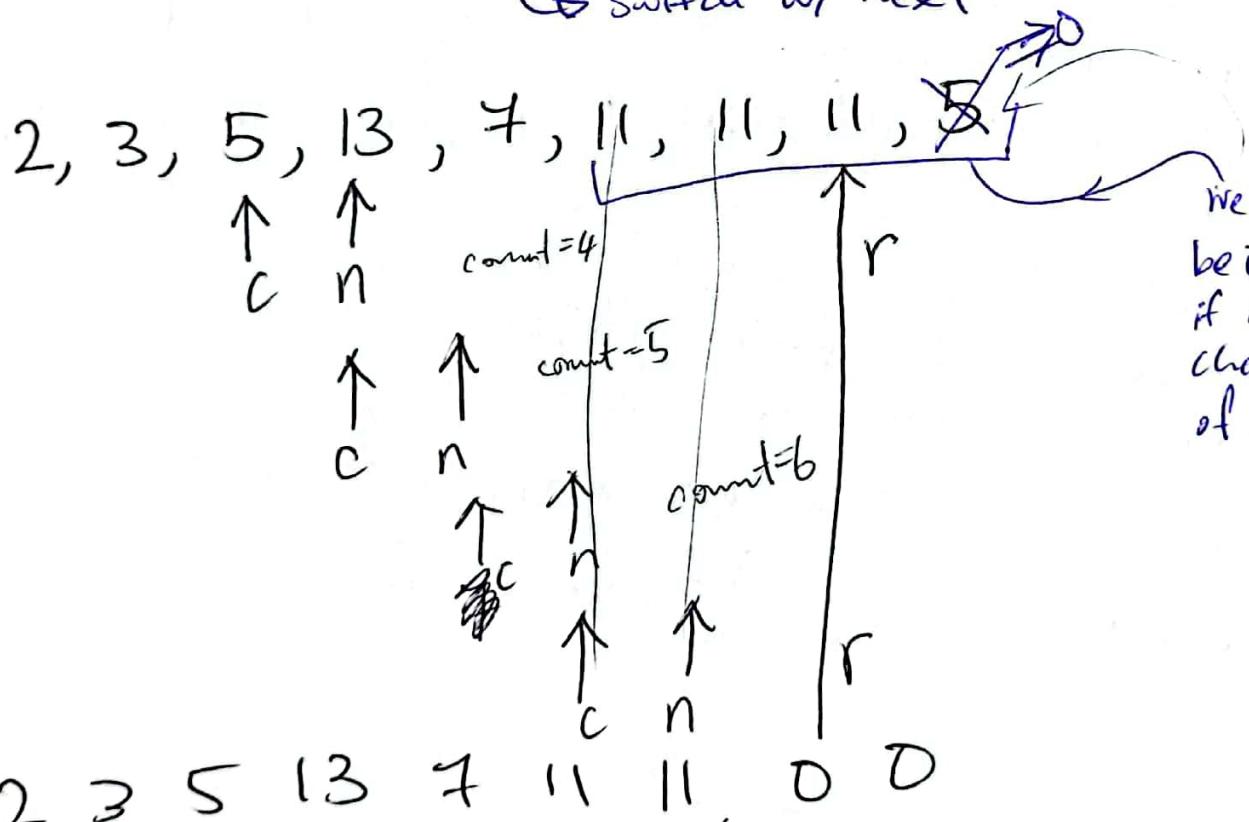


count ++

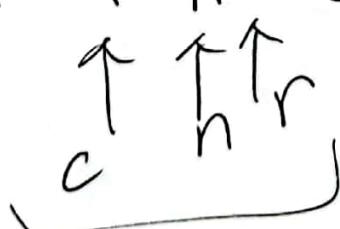
We want to switch places w/ another item if same.

↳ switch w/ current or next?

↳ switch w/ next



2 3 5 13 7 11 11 0 0



At this point, stop loop  
⇒ while curr < right

2/4

def fun (A : List[int]) → int :

count = 1

18 min I remembered the sol'n  
write-index = 1

~~curr~~ = 0

next = 1

end = len(A) - 1

while curr < end :

if A[curr] ≠ A[next] :

count += 1

curr += 1

next += 1

else :

A[next] = A[end]

A[end] = 0

end -= 1

return count

Edge case :

if len(A) == 0 :

23 min

return 0

3/4

## Evaluation / correction

- The problem was that I was unable to figure out that I shouldn't have used a while loop.
- Since I am checking two adjacent items in the list, the traversal must be uniform. So use a for loop and compare using  $i$  and  $i+1$  or  $i-1$ .
- Next time, check if I need to compare adjacent items or items that are someplace else with the current item.

4/4