

# Imp. Even-Odd Merge

7.10 7/7/22

There are two  $O(n)$ ,  $O(1)$  methods

(I) Have a left and right pointers that each points to the ends of the list.

• Then ~~change~~ make the ~~pos~~ left pointer point to the next even and make the right ptr point to the ~~next~~ left subsequent node. And make the ~~pos~~ odd pos. node point to Null.

→  $\text{left.next} = \text{left.next.next}$   
 $\text{right.next} = \text{left.next}$   
 $\text{right.next.next} = \text{None}$

(II) • Make a dummy head and even and odd pointers. Make the <sup>node</sup> pointers point to the next even/odd nodes.

• When it reached the end, make the last even node point to the first odd.

(EPI sol'n)

## 4.10 Evaluation

• I was moving back and forth w/ the two sol'n method.

~~• The idea is~~

→ I eventually stuck w/ method (I) but couldn't come up w/ the idea of having two separate even and odd dummy node and link them (last even → first odd) after the loop was finished.

• Didn't do boundary cases / edge → what if list was empty?  
one elem?  
two elem?

• If working w/ even/odd lists, consider separating them.  
• Make nodes point to Null each step, then the last elem will point to Null when the loop is finished.

7/7/22

7.10

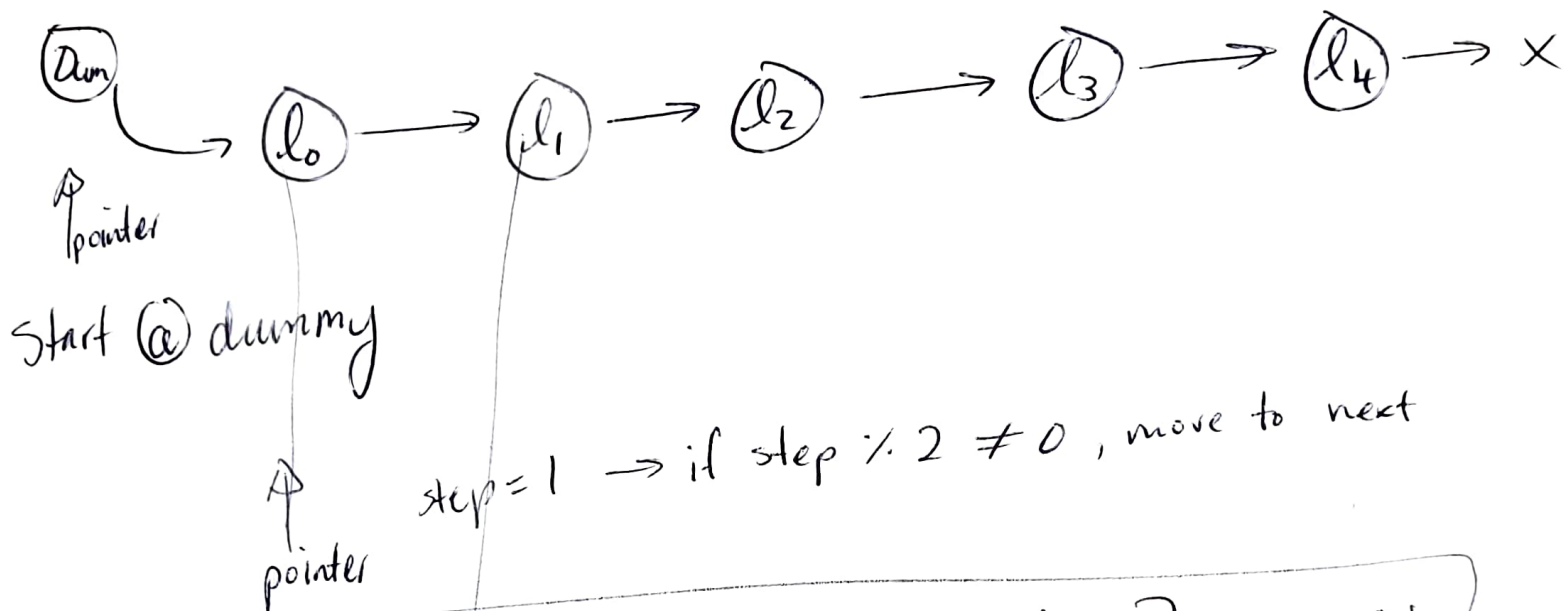
7/0 : Linked List

- Return a list w/ ~~even~~ nodes at even position first then odd.
- Starts at 0
- I probably have to use a dummy node
- The Naive approach is to keep track of the steps taken by incrementing a variable while traversing through the list.
- Then checking if a node is in even/odd position by  $(x \% 2 == 0)$ .

→ How can I rearrange the list?  
↳ sentinel / dummy

→

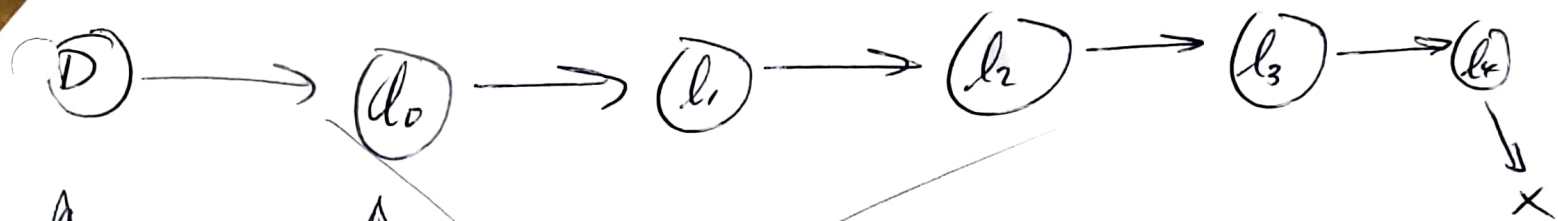
step = 0



$\rightarrow$  Should I use fast/slow pointer?  $\rightarrow$  No

$\uparrow$   
pr  
step = 2  $\rightarrow$  if  $\text{step} \% 2 = 0$ ,

$\rightarrow$  I ~~should~~ use ~~left~~ left-end pointer and another pointer that iterates through the list.



↑  
l

↑  
r

is ~~a~~ in even position?

↳ yes, move to next

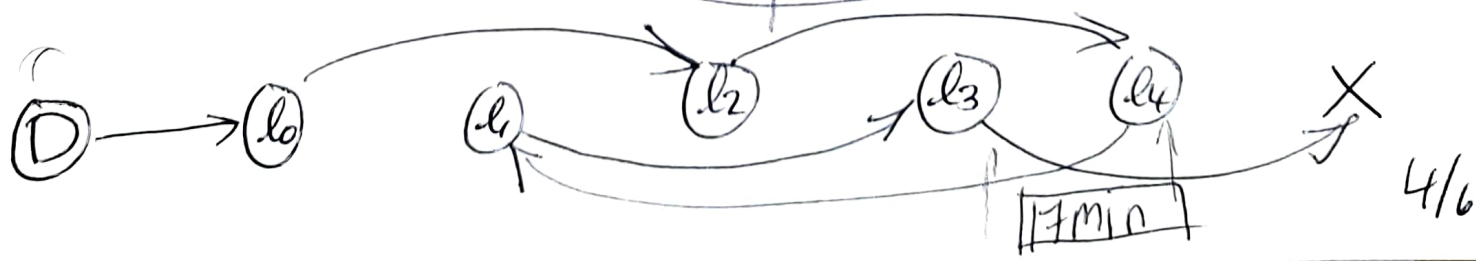
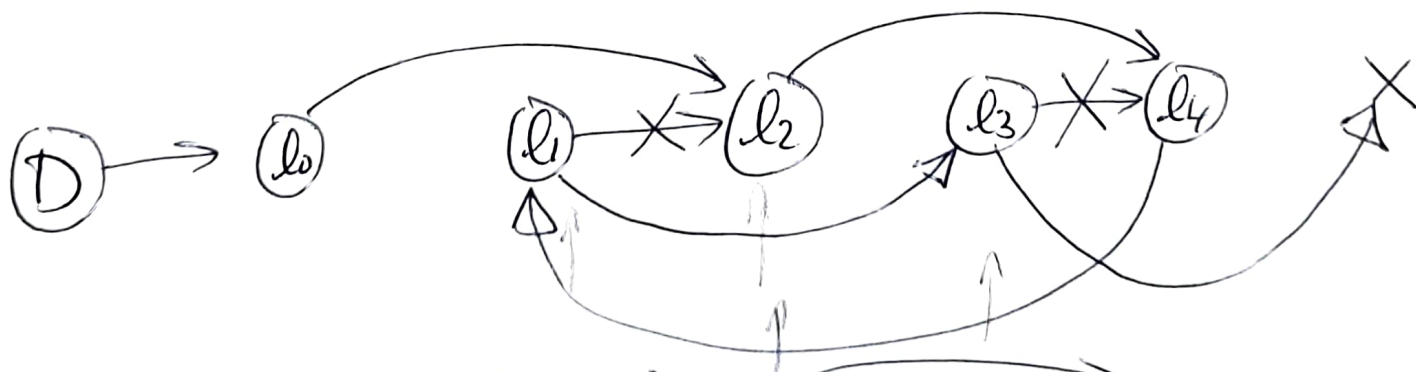
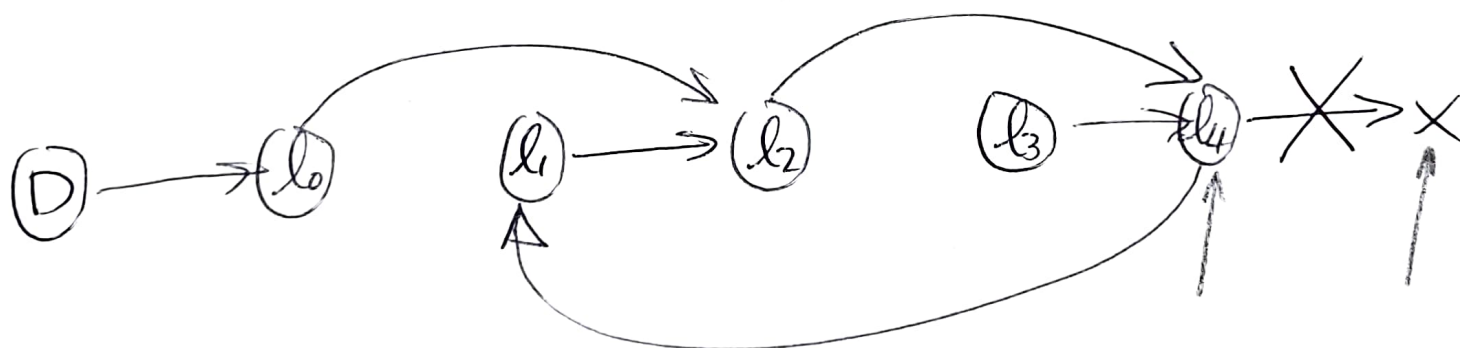
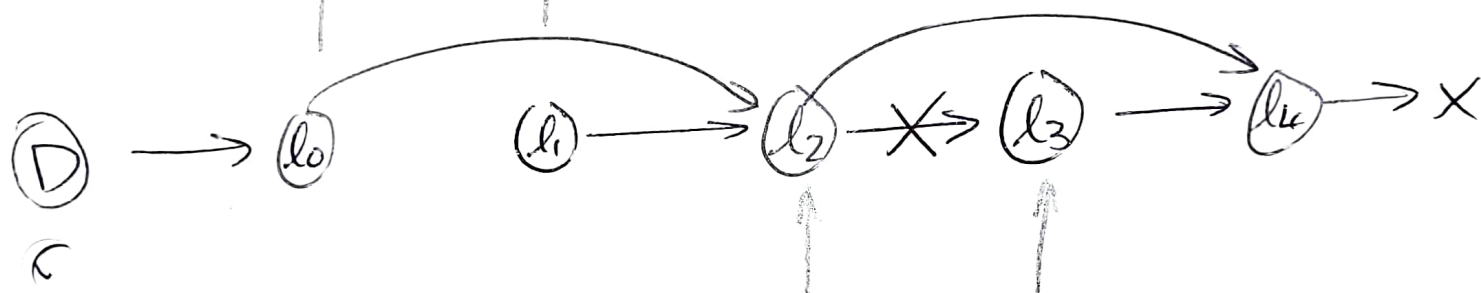
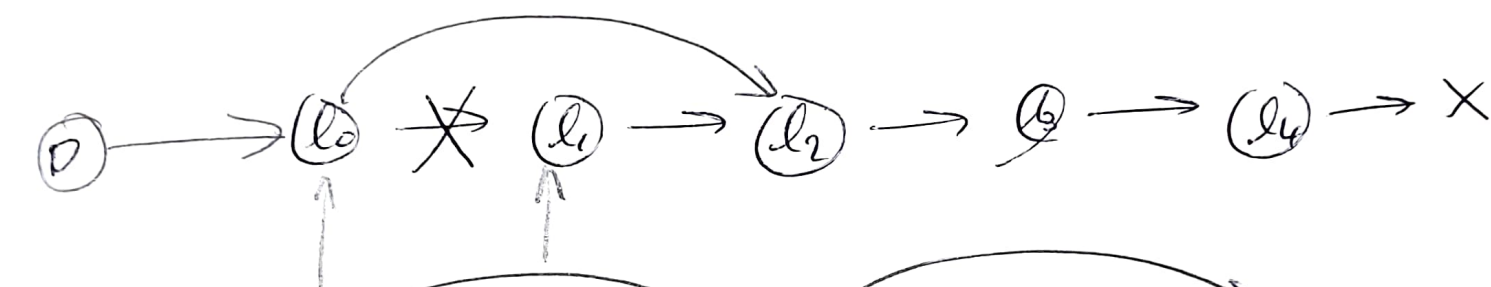
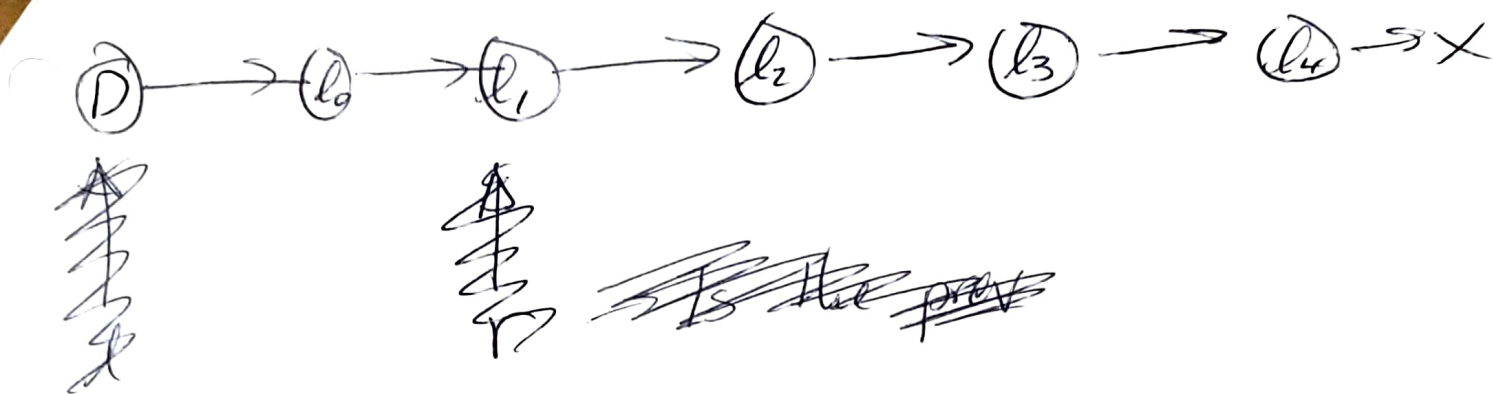
→ No,

↑  
r

↑  
l

→ I need to check the next node or prev node






There are two steps to this:


- ① ~~##~~ If the previous node is in even position, make it point ~~##~~ to the next even node.
- ② When all even nodes are found, make the last even node point to the first odd node.

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① If the curr. node is odd, make the prev. node point to curr.next

↳ ~~take~~ two steps increment?

↳ when curr == None, 

→ I might have to know the length of the list. → 

→ I should keep track of the first odd node and traverse from there.

◦ When we reach the last even node,  
make it point to the first odd node.

◦ Then make it point to  $x.next.next$

↳ What happen when I reach the end?

↳ Either the last odd point to  
Null or the last even.

→ 32 min Quit



Are we working w/ arrays, linked lists?

So the input is LL and output is the max value w/in that list.

This might mean for every values that are popped, keep track of the max by taking an item and comparing w/ all others.

This will take  $O(n)$  time  $O(1)$  space

```
def f(l):  
    max = 0  
    while l:  
        l.pop()  
        if l.pop() > max:  
            max = l.pop()  
    l.pop()  
    return max
```

6 min

→ Had to use OOP for this problem