Problem 1

(a) Grammar:

```
O. exp' ::= exp $
  1. exp ::= id
  2. exp ::= ( exp )
  3. exp ::= ( type ) exp
  4. type ::= id
FIRST(exp) = \{ id, ( \} \}
FIRST(exp') = FIRST(exp) = \{ id, ( \} \}
FIRST(type) = \{ id \}
FOLLOW(exp) = \{ \$, \}, id \}
FOLLOW(exp') = \{ \}
FOLLOW(type) = \{ \ ) \ \}
Nullable(exp) = false
Nullable(exp') = false
Nullable(type) = false
```

(b) Grammar:

O. eqns' ::= eqns \$ (\$ is the end-of-file marker)

```
1. eqns ::= eq sup eqns
  2. eqns ::= eq
  3. eq ::= x
FIRST(eqns) = \{ x \}
FIRST(eqns') = FIRST(eqns) = \{ \mathbf{x} \}
FIRST(eq) = \{ x \}
FOLLOW(eqns) = \{ \}
FOLLOW(eqns') = \{ \$, x \}
FOLLOW(eq) = \{ \$, \sup, x \}
Nullable(eqns) = false
Nullable(eqns') = false
Nullable(eq) = false
```

Problem 2

(a) Symbol table at line 11:

Symbol	Type
$\overline{\mathbf{c}}$	integer
a	integer
y	integer
${f Z}$	integer
b	integer
\mathbf{m}	integer
n	integer

(b) Actions required for symbol table management:

When the parser enters a new procedure, it needs to create a new symbol table.

Assuming that there is a stack of symbol tables, it also needs to push the new table onto the stack.

When the parser exits the procedure, it needs to pop the created symbol table off of the stack of symbol tables.

Problem 3

```
VarDecl -> var IDList: typeID
    VarDecl.env = {IDList -> (typeID, var)}

IDList -> IDList, ID
    (type, kind) = IDList.env.lookup(IDList)
```

```
IDList.type = type
IDList.kind = kind

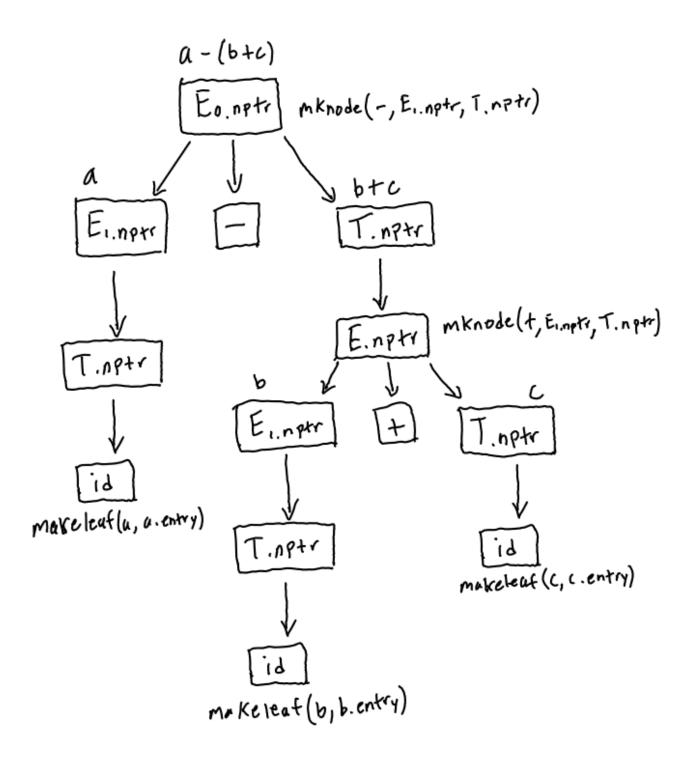
IDList -> ID

IDList.type = ID.type
IDList.kind = ID.kind
```

Yes, this scheme can operate in a single pass of the syntax tree.

Problem 4

Parse tree for a - (b + c)



Problem 5

In order to create distinct identifiers for each function I would change the

naming scheme of each identifier so that they take the form of:

```
{function name}_{return type}_{arg 1 type}_{arg 2
type}_...
```

This ensures that each overloaded version of the function will be associated with a unique identifier that can be easily constructed when looking up.