# Ve203 Discrete Mathematics (Fall 2016)

# **Assignment 6: Recurrence and Combinatrics**

Date Due: 4:00 PM, Thursday, the 10<sup>th</sup> of Noevmber 2016



This assignment has a total of (27 Marks).

#### Exercise 6.1

In this question, assume that f is an increasing function satisfying the recurrence relation  $f(n) = af(n/b) + cn^d$  with  $a \ge 1$ ,  $b \in \mathbb{N} \setminus \{0,1\}$ ,  $c,d \in \mathbb{R}_+$ . Our goal is to prove the Master Theorem 2.3.19 of the lecture.

- i) Show that if  $a = b^d$  and n is a power of b, then  $f(n) = f(1)n^d + cn^d \log_b n$ . (2 Marks)
- ii) Show that if  $a = b^d$ , then  $f(n) = O(n^d \log n)$ .

  (1 Mark)
- iii) Show that if  $a \neq b^d$  and n is a power of b, then

$$f(n) = C_1 n^d + C_2 n^{\log_b a},$$
  $C_1 = \frac{b^d c}{b^d - a},$   $C_2 = f(1) + \frac{b^d c}{a - b^d}.$ 

 $(2 \, \text{Marks})$ 

- iv) Show that if  $a < b^d$ , then  $f(n) = O(n^d)$ .

  (1 Mark)
- v) Show that if  $a > b^d$ , then  $f(n) = O(n^{\log_b a})$ . (1 Mark)

#### Exercise 6.2

A recursive algorithm for modular exponentiation is given in Example 3 of Section 4.4, page 312 of the textbook.

- i) Set up a divide-and-conquer recurrence relation for the number of modular multiplications required to compute a<sup>n</sup> mod m, where a, m, n ∈ Z<sub>+</sub>.
   (2 Marks)
- ii) Construct a big-O estimate for the number of modular multiplications required to compute  $a^n \mod m$ . (1 Mark)

## Exercise 6.3

Prove that in a bit string, the string 01 occurs at most one more time than the string 10. (2 Marks)

## Exercise 6.4

Consider the scheme for counting fractions shown below and let

$$\varphi \colon \mathbb{N}^* \times \mathbb{N}^* \to \mathbb{N}^*, \qquad \qquad \varphi(p,q) = \frac{(p+q-1)(p+q-2)}{2} + p,$$

where  $N^* = \mathbb{N} \setminus \{0\}$ . The goal of this question is to prove that, in the scheme shown, traversing successive diagonals from top to bottom, the fraction p/q is indeed the  $\varphi(p,q)$ th fraction encountered and that  $\varphi$  gives a bijection  $N^* \times \mathbb{N}^* \to \mathbb{N}^*$ .

i) Write the traversal of the fractions as a recursively defined sequence of pairs  $(p_n, q_n)$ .

(1 Mark)

ii) Prove that for any  $(p,q) \in \mathbb{N}^* \times \mathbb{N}^*$  there exists an  $n \in \mathbb{N}$  such that  $(p,q) = (p_n,q_n)$ .

(2 Marks)

- iii) Prove (by induction in n) that  $\varphi(p_n, q_n) = n$ . (2 Marks)
- iv) Deduce that  $\varphi$  is bijective.

 $(0.5 \, \text{Marks})$ 

v) Find the inverse of  $\varphi$ .

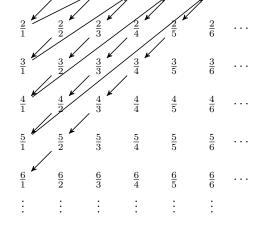
(2 Marks)

vi) How can  $\varphi$  be modified to give a bijective map  $\mathbb{N} \times \mathbb{N}^* \to \mathbb{N}$ ?

(1 Mark)

vii) How can  $\varphi$  be modified to give a bijective map  $\mathbb{Z} \times \mathbb{N}^* \to \mathbb{N}$ ?

 $(0.5 \, \text{Marks})$ 



#### Exercise 6.5

Show that the countable union of countable sets is countable, i.e., if  $\{A_i\}_{i=0}^{\infty}$  is an infinite family of sets, each of which is countable, then  $\bigcup_{i\in\mathbb{N}} A_i$  is countable. *Hint:* Let  $a_{ij}$  denote the *i*th element of  $A_j$ ... (2 Marks)

### Exercise 6.6

Let M, N be finite sets with card  $M = \operatorname{card} N$  and  $M \subset N$ . Prove that M = N. (2 Marks)

## Exercise 6.7

Use the Pigeonhole Principle or Theorem 2.4.21 of the lecture to prove the following theorem:

Let M, N be finite sets with card  $M > \operatorname{card} N$  and  $f : M \to N$ . Then f is not injective.

(2 Marks)