

# **VE203**

## **Assignment 2**

*Jiang Yicheng*  
*515370910224*

September 27, 2016

---

# 1 $2 + 2 = 4$

## 1.1

We define the function  $\cdot + \cdot : \mathbb{N} \times \mathbb{N} \mapsto \mathbb{N}$  has the following properties:

1.  $n + 0 = n$
2. For  $m \neq 0$ , set  $m = \text{succ}(m')$ ,  $n + m = \text{succ}(n) + m'$

## 1.2

Since  $n + 1 = \text{succ}(n)$ ,  $2 = \text{succ}(1)$ ,  $3 = \text{succ}(2)$ ,  $4 = \text{succ}(3)$ , then

$$2 + 2 = 2 + \text{succ}(1) = \text{succ}(2) + 1 = 3 + 1 = \text{succ}(3) = 4$$

So we get that  $2 + 2 = 4$

## 1.3

We first prove that  $\forall m, n \in \mathbb{N}, \text{succ}(n + m) = \text{succ}(n) + m$   
 $\forall n \in \mathbb{N}$ ,

1. When  $m = 0$ ,  $\text{succ}(n + m) = \text{succ}(n + 0) = \text{succ}(n) = \text{succ}(n) + 0 = \text{succ}(n) + m$   
so the statement holds when  $m = 0$
2. Assume that the statement holds when  $m = m_0$ , i.e.  $\text{succ}(n + m_0) = \text{succ}(n) + m_0$ ,  
then according to definition of addition

$$\begin{aligned} \text{succ}(n + \text{succ}(m_0)) &= \text{succ}(\text{succ}(n) + m_0) \\ &= \text{succ}(\text{succ}(n)) + m_0 \\ &= \text{succ}(n) + \text{succ}(m_0) \end{aligned}$$

so the statement also holds when  $m = \text{succ}(m_0)$ .

To sum up,  $\forall m, n \in \mathbb{N}, \text{succ}(n + m) = \text{succ}(n) + m$ . Then we try to prove that  
 $\forall n \in \mathbb{N}, n + 0 = 0 + n$

1. When  $n = 0$ ,  $n + 0 = 0 + 0 = 0 + n$  so the statement holds when  $n = 0$
2. Assume that the statement holds when  $n = n_0$ , i.e.  $n_0 + 0 = 0 + n_0$ , then according  
to definition of addition and the statement we just prove

$$\begin{aligned} 0 + \text{succ}(n_0) &= \text{succ}(0) + n_0 = \text{succ}(0 + n_0) \\ &= \text{succ}(n_0 + 0) = \text{succ}(n_0) \\ &= \text{succ}(n_0) + 0 \end{aligned}$$

so the statement also holds when  $n = \text{succ}(n_0)$ .

To sum up,  $\forall n \in \mathbb{N}, n + 0 = 0 + n$ .

Now we start to prove that  $\forall m, n \in \mathbb{N}, n + m = m + n$ .  $\forall n \in \mathbb{N}$

1. When  $m = 0$ ,  $n + 0 = 0 + n$  has been proved, so the statement holds when  $m = 0$

- 
2. Assume that the statement holds when  $m = m_0$ , i.e.  $n + m_0 = m_0 + n$ , then according to definition of addition and the statement we just prove

$$\begin{aligned} n + \text{succ}(m_0) &= \text{succ}(n) + m_0 = \text{succ}(n + m_0) \\ &= \text{succ}(m_0 + n) \\ &= \text{succ}(m_0) + n \end{aligned}$$

so the statement also holds when  $m = \text{succ}(m_0)$ .

In conclusion,  $\forall m, n \in \mathbb{N}, n + m = m + n$ .

## 2 Straightforward Induction

**Proof:**

1. When  $n = 1, 2$ , since  $a_1 = 1 = 3 \cdot 2^{1-1} + 2 \cdot (-1)^1$ ,  $a_2 = 8 = 3 \cdot 2^{2-1} + 2 \cdot (-1)^2$ , so  $a_n = 3 \cdot 2^{n-1} + 2 \cdot (-1)^n$  holds for  $n = 1, 2$
2. Assume that  $\forall n \leq k (k \geq 2)$ ,  $a_n = 3 \cdot 2^{n-1} + 2 \cdot (-1)^n$  holds, then

$$\begin{aligned} a_{k+1} &= a_k + 2a_{k-1} = 3 \cdot 2^{k-1} + 2 \cdot (-1)^k + 2 \cdot (3 \cdot 2^{k-1-1} + 2 \cdot (-1)^{k-1}) \\ &= (3 + 3) \cdot 2^{k-1} + (2 - 4) \cdot (-1)^k \\ &= 3 \cdot 2^{k+1-1} + 2 \cdot (-1)^{k+1} \end{aligned}$$

so the equation also holds for  $n = k + 1$ .

To sum up,  $\forall n > 0, a_n = 3 \cdot 2^{n-1} + 2 \cdot (-1)^n$

## 3 The Fifth Peano Axiom

**Proof:** Assume that there exists some non-empty set  $S \subset \mathbb{N}$  doesn't have a least element, that is  $\forall m \in S, \exists m_0 \in S$  such that  $m_0 < m$ .

Set  $T = \mathbb{N} \setminus S$ . Since 0 is not the successor of any natural number, then if  $0 \in S$ , it would be the least number in  $S$ . So  $0 \in T$ .

Since 1 is the successor of 0, and since  $0 \notin S$ , then if  $1 \in S$ , 1 would be the least number in  $S$ . So  $1 \in T$ .

Since 2 is the successor of 1, and since  $0, 1 \notin S$ , then if  $2 \in S$ , 2 would be the least number in  $S$ . So  $2 \in T$ .

Repeat this way, we can prove that for each natural number  $n$ , its successor is in  $T$ . And since  $0 \in T$ , then according to induction axiom  $T = \mathbb{N}$ , and therefore  $S = \emptyset$ . This leads to contradiction.

So such kind of set doesn't exist, and *Well-Ordering Principle* holds.

## 4 Is a direct induction approach always successful?

First, we use induction to prove that  $(1 + x)^n \geq 1 + nx$  holds for any  $n \in \mathbb{N}$ , where  $x > -1$

---

**Proof:**

1. When  $n = 0$ ,  $(1+x)^n = (1+x)^0 = 1$ ,  $1+nx = 1+0 \cdot x = 1$ . So  $(1+x)^n \geq 1+nx$  holds when  $n=0$
2. Assume that  $(1+x)^n \geq 1+nx$  holds when  $n=k$ , where  $k \in \mathbb{N}$ , i.e.  $(1+x)^k \geq 1+kx$ . Then

$$\begin{aligned}(1+x)^{k+1} &\geq (1+x) \cdot (1+kx) \\ &= 1+kx+x+kx^2 \\ &= (1+(k+1)x+kx^2) \\ &\geq 1+(k+1)x\end{aligned}$$

This is because  $k \in \mathbb{N}, x > -1$

So  $(1+x)^n \geq 1+nx$  also holds when  $n = k+1$ .

According to 1, 2,  $\forall n \in \mathbb{N}$ ,  $(1+x)^n \geq 1+nx$ .

So  $\forall n \in \mathbb{N}$ ,  $(1+x)^n \geq nx$

## 5 Strong Induction

**Proof:**

1. When  $n = 1, 2$ , since  $1 = 2^0, 2 = 2^1$ , then we can see that the statement holds for  $n = 1, 2$ .
2. Assume that  $\forall n \in \mathbb{N}, 1 \leq n \leq k (k \geq 2)$ , it can be written as a sum of distinct powers of 2. Then for  $n = k+1$ :

(a) If  $k+1$  is even, then  $k \geq \frac{k+1}{2} \in \mathbb{N}^*$ , according to assumption we can set that  $\frac{k+1}{2} = p_0 2^0 + p_1 2^1 + \cdots + p_{(k+1)/2-1} 2^{(k+1)/2-1} (p_i \in \{0, 1\}, i = 0, 1, \dots, (k+1)/2-1)$ . This is practical since  $2^n = (1+1)^n > n \cdot 1 = n$ . Then  $k+1 = p_0 2^1 + p_1 2^2 + \cdots + p_{(k+1)/2-1} 2^{(k+1)/2}$  which is a sum of distinct powers of 2.

(b) If  $k+1$  is odd, then  $\frac{k}{2} \in \mathbb{N}^*, 1 \leq \frac{k}{2} \leq k$ , according to assumption we can set that  $\frac{k}{2} = p_0 2^0 + p_1 2^1 + \cdots + p_{k/2-1} 2^{k/2-1} (p_i \in \{0, 1\}, i = 0, 1, \dots, k/2-1)$ . Then  $k+1 = 2^0 + p_0 2^1 + p_1 2^2 + \cdots + p_{k/2-1} 2^{k/2}$  which is a sum of distinct powers of 2.

according to (a)(b),  $k+1$  always can be written as a sum of distinct powers of 2. So, the statement also holds for  $n = k+1$ .

From 1, 2, every  $n \in \mathbb{N} \setminus \{0\}$  can be written as a sum of distinct powers of 2.

---

## 6 Structural Induction

**Proof:** Use structural induction to prove

1. Since  $(0, 0) \in S$ , and  $5 \mid 0 + 0$ , the base case is established.
2. Assume that  $\forall a, b \in S, 5 \mid (a + b)$ , so we can set  $a + b = 5k, k \in \mathbb{N}$ . Then  $(a + 2, b + 3), (a + 3, b + 2) \in S$ , and  $(a + 2) + (b + 3) = (a + 3) + (b + 2) = a + b + 5 = 5(k + 1)$ . Since  $k + 1 \in \mathbb{N}$ ,  $5 \mid ((a + 2) + (b + 3)), 5 \mid ((a + 3) + (b + 2))$ .

According to 1,2, we can see that  $\forall (a, b) \in S, 5 \mid (a + b)$

## 7 Some easy practice of relation properties

### 7.1 $x+y=0$

1. Since  $1 \in \mathbb{Z}, 1 + 1 = 2 \neq 0$ , this shows that the relation is not reflexive.
2.  $\forall (x, y) \in R, y + x = x + y = 0$ , so  $(y, x) \in R$ . So the relation is symmetric.
3. Since  $1 + (-1) = 0, (-1) + 1 = 0$ , then  $(1, -1), (-1, 1) \in R$ , while  $1 + 1 = 2 \neq 0$ , so the relation is not transitivity.

### 7.2 $2 \mid (x-y)$

1.  $\forall x \in \mathbb{Z}, (x, x) \in R$  since  $2 \mid 0 = x - x$ , which shows that the relation is reflexive.
2.  $\forall (x, y) \in R$ , set  $x - y = 2k, k \in \mathbb{Z}$ . Then  $y - x = -2k = 2(-k)$ . Since  $-k \in \mathbb{Z}, 2 \mid (y - x)$ . So  $(y, x) \in R$ . So the relation is symmetric.
3.  $\forall (x, y), (y, z) \in R$ , set  $x - y = 2k_1, y - z = 2k_2 (k_1, k_2 \in \mathbb{Z})$ . Then  $x - z = x - y + y - z = 2(k_1 + k_2)$ . Since  $k_1 + k_2 \in \mathbb{Z}, 2 \mid (x - z)$ . So  $(x, z) \in R$ . So the relation is transitivity.

### 7.3 $xy=0$

1. Since  $1 \in \mathbb{Z}, 1 \cdot 1 = 1 \neq 0$ , this shows that the relation is not reflexive.
2.  $\forall (x, y) \in R, yx = xy = 0$ , so  $(y, x) \in R$ . So the relation is symmetric.
3. Since  $1 \cdot 0 = 0, 0 \cdot 2 = 0$ , then  $(1, 0), (0, 2) \in R$ , while  $1 \cdot 2 = 2 \neq 0$ , so the relation is not transitivity.

### 7.4 $x=1$ or $y=1$

1. Since  $2 \in \mathbb{Z}, 2 \neq 1$ , then  $(2, 2) \notin R$ . This shows that the relation is not reflexive.
2.  $\forall (x, y) \in R, x = 1 \vee y = 1$ , so  $(y, x) \in R$ . So the relation is symmetric.
3. Since  $(3, 1), (1, 2) \in R$ , while  $(3, 2) \notin R$ , then the relation is not transitivity.

---

### 7.5 $x = \pm y$

1.  $\forall x \in \mathbb{Z}, (x, x) \in R$  since  $x = x$ , which shows that the relation is reflexive.
2.  $\forall (x, y) \in R, x = y \vee x = -y$ . So  $y = x \vee y = -x$ . So  $(y, x) \in R$ . So the relation is symmetric.
3.  $\forall (x, y), (y, z) \in R, x = y \vee x = -y, y = z \vee y = -z$ . Then  $x = z \vee x = -z$ . So  $(x, z) \in R$ . So the relation is transitivity.

### 7.6 $x = 2y$

1. Since  $1 \in \mathbb{Z}, 1 \neq 2 = 2 \cdot 1$ , then  $(1, 1) \notin R$ . This shows that the relation is not reflexive.
2.  $(2, 1) \in R$ , while  $(1, 2) \notin R$ . So the relation is not symmetric.
3. Since  $(4, 2), (2, 1) \in R$ , while  $(4, 1) \notin R$ , then the relation is not transitivity.

### 7.7 $xy \geq 0$

1.  $\forall x \in \mathbb{Z}, (x, x) \in R$  since  $x \cdot x = x^2 \geq 0$ , which shows that the relation is reflexive.
2.  $\forall (x, y) \in R, xy = xy \geq 0$ . So  $(y, x) \in R$ . So the relation is symmetric.
3. Since  $1 \cdot 0 = 0 \geq 0, 0 \cdot (-1) = 0 \geq 0$ , then  $(1, 0), (0, -1) \in R$ , while  $1 \cdot (-1) = -1 < 0$ , so the relation is not transitivity.

### 7.8 $x = 1$

1. Since  $2 \in \mathbb{Z}, 2 \neq 1$ , then  $(2, 2) \notin R$ . This shows that the relation is not reflexive.
2.  $(1, 2) \in R$ , while  $(2, 1) \notin R$ , so the relation is not symmetric.
3.  $\forall x, y, z \in \mathbb{Z}, (x, y), (y, z) \in R$ , then  $x = 1$  and therefore  $(x, z) \in R$ , then the relation is transitivity.