

Computability Theory

Preliminary

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PRELIMINARIES

SET

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 - ▶ $\mathbb{N} = \{0, 1, 2, \dots\}$: natural numbers (nonnegative integers)
 - ▶ $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$: integers
 - ▶ \mathbb{R} : real numbers
 - ▶ \mathbb{E} : even numbers
 - ▶ \mathbb{O} : odd numbers

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- ▶ **Subset**: A set S is a subset of T , $S \subseteq T$, if every element of S is an element of T
- ▶ **Proper subset**: a subset of T is a subset other than the empty set \emptyset or T itself (Use of word proper, proper subsequence or proper substring)
- ▶ **Strict Subset**: S is a strict subset, $S \subset T$, if not equal to T

Set Operations

- ▶ **Union:** $S \cup T \rightarrow$ the set of elements that are either in S or in T .
 - ▶ $S \cup T = \{s | s \in S \text{ or } s \in T\}$
 - ▶ $\{a, b, c\} \cup \{c, d, e\} = \{a, b, c, d, e\}$
 - ▶ $|S \cup T| \leq |S| + |T|$
- ▶ **Intersection:** $S \cap T$
 - ▶ $S \cap T = \{s | s \in S \text{ and } s \in T\}$
 - ▶ $\{a, b, c\} \cap \{c, d, e\} = \{c\}$
- ▶ **Difference:** $S - T \rightarrow$ set of all elements in S not in T
 - ▶ $S - T = \{s | s \in S \text{ but not in } T\} = S \cap \overline{T}$
 - ▶ $\{1, 2, 3\} - \{1, 4, 5\} = \{2, 3\}$
- ▶ **Complement:**
 - ▶ Need universal set U
 - ▶ $\overline{S} = \{s | s \in U \text{ but not in } S\}$

Set Operations

► Cartesian Product

- $S \times T = \{(s, t) | s \in S, t \in T\}$
- In a graph $G = (V, E)$, the edge set E is the subset of Cartesian product of vertex set V . $E \subseteq V \times V$.

► Power Set

- 2^S set of all subsets of S
- Note: notation $|2^S| = 2^{|S|}$, meaning 2^S is a good representation for power set.
- $S = \{a, b, c\}$, then
 $2^S = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$
- Indicator Vector: We can use a zero/one vector to represent the elements in power set.

	a	b	c
\emptyset	0	0	0
$\{a\}$	1	0	0
$\{b\}$	0	1	0
$\{a, b, c\}$	1	1	1

Ordered Pair

- ▶ (x, y) : ordered pair of elements x and y ; $(x, y) \neq (y, x)$.
- ▶ (x_1, \dots, x_n) : ordered n -tuple \rightarrow as \tilde{x} .
- ▶ $A_1 \times A_2 \times \dots \times A_n = \{(x_1, \dots, x_n) : x_1 \in A_1, \dots, x_n \in A_n\}$.
- ▶ $A \times A \times \dots \times A = A^n$.
- ▶ $A^1 = A$.

FUNCTION

Function Notation

- ▶ f is a set of ordered pairs s.t. if $(x, y) \in f$ and $(x, z) \in f$, then $y = z$, and $f(x) = y$.
- ▶ $Dom(f)$: Domain of f , $\{x : f(x) \text{ is defined}\}$.
- ▶ $f(x)$ is undefined if $x \notin Dom(f)$.
- ▶ $Ran(f)$: Range of f , $\{f(x) : x \in Dom(f)\}$.
- ▶ f is a function from A to B : $Dom(f) \subseteq A$ and $Ran(f) \subseteq B$.
- ▶ $f : A \rightarrow B$: f is a function from A to B with $Dom(f) = A$.

Mapping

- ▶ **Injective**: if $x, y \in \text{Dom}(f)$, $x \neq y$, then $f(x) \neq f(y)$.
- ▶ **Inverse** f^{-1} : the unique function g s.t. $\text{Dom}(g) = \text{Ran}(f)$, and $g(f(x)) = x$.
- ▶ **Surjective**: if $\text{Ran}(f) = B$.
- ▶ **Bijjective**: both **injective** and **surjective**.

Operation

1. $f|X$: Restriction of f to X .
Domain $X \cap \text{Dom}(f)$. Write $f(X)$ for $\text{Ran}(f|X)$.
2. $f^{-1}(Y) = \{x : f(x) \in Y\}$: inverse image of Y under f .
3. $f \subseteq g$: g extends f , $f = g|_{\text{Dom}(f)}$.
 $\text{Dom}(f) \subseteq \text{Dom}(g)$ and $\forall x \in \text{Dom}(f)$, $f(x) = g(x)$.
4. $f \circ g$: composition of f and g . Domain
 $\{x : x \in \text{Dom}(g) \text{ and } g(x) \in \text{Dom}(f)\}$, value $f(g(x))$.
5. f_\emptyset : function defined nowhere. $\text{Dom}(f_\emptyset) = \text{Ran}(f_\emptyset) = \emptyset$.
 $f_\emptyset = g|_\emptyset$ for any function g .

\simeq : similar-or-equal-to

Suppose $\alpha(\tilde{x})$ and $\beta(\tilde{x})$ are expressions involving $\tilde{x} = (x_1, \dots, x_n)$, then $\alpha(\tilde{x}) \simeq \beta(\tilde{x})$ means $\forall \tilde{x}$, $\alpha(\tilde{x})$ and $\beta(\tilde{x})$ are either both defined, or both undefined, and if defined they are equal.

- ▶ $f(x) \simeq g(x)$ is a kind of $f = g$.
- ▶ $f(x) \simeq y$ means $f(x)$ is defined and $f(x) = y$.

Partial and Total Function

- ▶ **n -ary function:** $f(\tilde{x})$, $f(x_1, \dots, x_n)$, $f : \mathbb{N}^n \rightarrow \mathbb{N}$.
- ▶ **Partial function:** $Dom(f)$ is not necessarily the whole \mathbb{N}^n .
(In our class function means partial function)
- ▶ **Total function:** $Dom(f) = \mathbb{N}^n$.
- ▶ **Zero function:** $\mathbf{0} : \mathbb{N} \rightarrow \mathbb{N}$.

RELATIONS

Relation

If A is a set, a property $M(x_1, \dots, x_n)$ that holds for some n -tuple from A^n and does not hold for all other n -tuples from A^n is called an n -ary **relation** or **predicate** on A .

- ▶ Property $x < y$. $2 < 5$, $6 < 4$.
- ▶ f from \mathbb{N}^n to \mathbb{N} gives rise to predicate $M(\tilde{x}, y)$ by:
 $M(x_1, \dots, x_n, y)$ iff $f(x_1, \dots, x_n) \simeq y$.

Equivalence Relation

- ▶ A binary relation R on A is called **equivalence relation** if

$$\left. \begin{array}{ll} \text{reflexivity} & \forall x \in A \ R(x, x) \\ \text{symmetry} & R(x, y) \Rightarrow R(y, x) \\ \text{transitivity} & R(x, y), R(y, z) \Rightarrow R(x, z) \end{array} \right\} \text{equivalence}$$

- ▶ A binary relation R on A is called a **partial order** if

$$\left. \begin{array}{ll} \text{irreflexivity} & \neg R(x, x) \\ \text{transitivity} & R(x, y), R(y, z) \Rightarrow R(x, z) \end{array} \right\} \text{partial order}$$

Example

	reflexive	symmetric	transitive
$<$	No	No	Yes
\leq	Yes	No	Yes
Parent of	No	No	No
$=$	Yes	Yes	Yes

DIAGONAL METHOD

The First Theorem

Theorem

The set of reals is uncountable.

The Second Theorem

Theorem

The power set of a set is always of greater cardinality than the set itself.

Hand Writing

- ▶ Small letters for **elements** and **functions**.
 - ▶ a, b, c for elements,
 - ▶ f, g for functions,
 - ▶ i, j, k for integer indices,
 - ▶ x, y, z for variables,
- ▶ Capital letters for **sets**. A, B, S . $A = \{a_1, \dots, a_n\}$
- ▶ Small letters with tilde for **vectors**. \tilde{x}, \tilde{y} . $\tilde{v} = \{v_1, \dots, v_m\}$
- ▶ Bold capital letters for **collections**. **A, B**. **S** = $\{S_1, \dots, S_n\}$
- ▶ Blackboard bold capitals for **domains** (standard symbols). $\mathbb{N}, \mathbb{R}, \mathbb{Z}$.
- ▶ German script for **collection of functions**. $\mathcal{C}, \mathcal{I}, \mathcal{T}$.
- ▶ Greek letters for **parameters** or **coefficients**. α, β, γ .
- ▶ Double strike handwriting for bold letters.