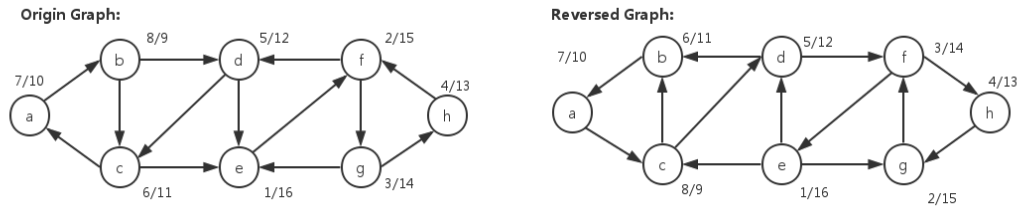


Algorithmn HW6

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Problem 9.16



Firstly, apply DFS and update the *post* value;

Secondly, reverse graph G to G^T ;

Thirdly, apply DFS to G^T starting at vertex with maximum *post* value, e this time;

A single DFS of this G^T can approach all the vertexes, so the whole graph G is strongly connected.

Problem 9.32

Require: a given graph $G = (V, E)$

Ensure: whether it is a bipartite graph

```

1: function DFS( $v$ )
2:   mark  $v$  visited
3:    $v.mark \leftarrow (color \leftarrow !color)$  ▷ mark neighbour vertexes with different colors
4:   for each edge  $(v, w) \in E$  do
5:     if  $w$  is marked unvisited then
6:       DFS( $w$ )
7:     return True
8:   else
9:     return  $(w.mark == v.mark?)$  False : True
10:  end if
11: end for
12: end function
13:  $color \leftarrow 1$  ▷ main function
14: mark each vertex  $v \in V$  unvisited
15: for each vertex  $v \in V$  marked unvisited do ▷ start dfs
16:   if DFS( $v$ ) == False then
17:     return False
18:   end if
19: end for
20: return True

```

This is an $O(V + E)$ algorithm

Reverse graph

Require: a directed graph $G = (V, E)$

Ensure: the reversed graph $G^T = (V, E^R)$

- 1: **for** each edge $(v, u) \in E$ **do**
- 2: append (u, v) to E^R
- 3: **end for**
- 4: replace E with E^R

Find a cycle

Require: an undirected graph $G = (V, E)$, $e = (u, v) \in E$

Ensure: whether G has a cycle containing e

- 1: delete e from E
- 2: apply DFS to G with *pre/post* signature
- 3: **if** $v.pre > u.post$ or $u.pre > v.post$ **then** \triangleright exist such cycle iff u and v are still in same component
- 4: **return** *False* \triangleright which means there shouldn't have any cross edge between them
- 5: **else**
- 6: **return** *True*
- 7: **end if**

Hamiltonian path in a DAG

Require: a DAG $G = (V, E)$

Ensure: whether it has a Hamiltonian path

- 1: $topo[] \leftarrow$ the topological sorted vertex sequence of G
- 2: **for** $i \leftarrow 1$ to $topo.size$ **do** \triangleright whether every consecutive pairs are connected
- 3: find the edge between $topo[i - 1]$ and $topo[i]$ in E
- 4: **if** cannot found **then** \triangleright once unconnected, there won't exist
- 5: **return** "*can not find*"
- 6: **end if**
- 7: **end for**
- 8: **return** $topo$ \triangleright the topological sort sequence is actually the Hamiltonian Path