



DECEMBER 11-12, 2024
BRIEFINGS

How the Internet Dodged a Bullet: The KeyTrap Denial-of-Service Attacks against DNSSEC

Speaker(s):

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Contributors:

Haya Schulmann, Michael Waidner



Refresher: DNS and DNSSEC

Why is it always DNS?

DDoS attack that disrupted internet was largest of its kind in history, experts say

Dyn, the victim of last week's denial of service attack, said it was orchestrated using a weapon called the Mirai botnet as the 'primary source of malicious attack'

Understanding how Facebook disappeared from the Internet

2021-10-04

DNS poisoning slams web traffic from millions in China into the wrong hole

ISP blames unspecified attack for morning outage

[Home](#) > [News](#) > [Security](#) > Akamai DNS global outage takes down major websites, online services

Akamai DNS global outage takes down major websites, online services

By [Sergiu Gatlan](#)

July 22, 2021 12:39 PM 2

Salesforce cloud services go down worldwide

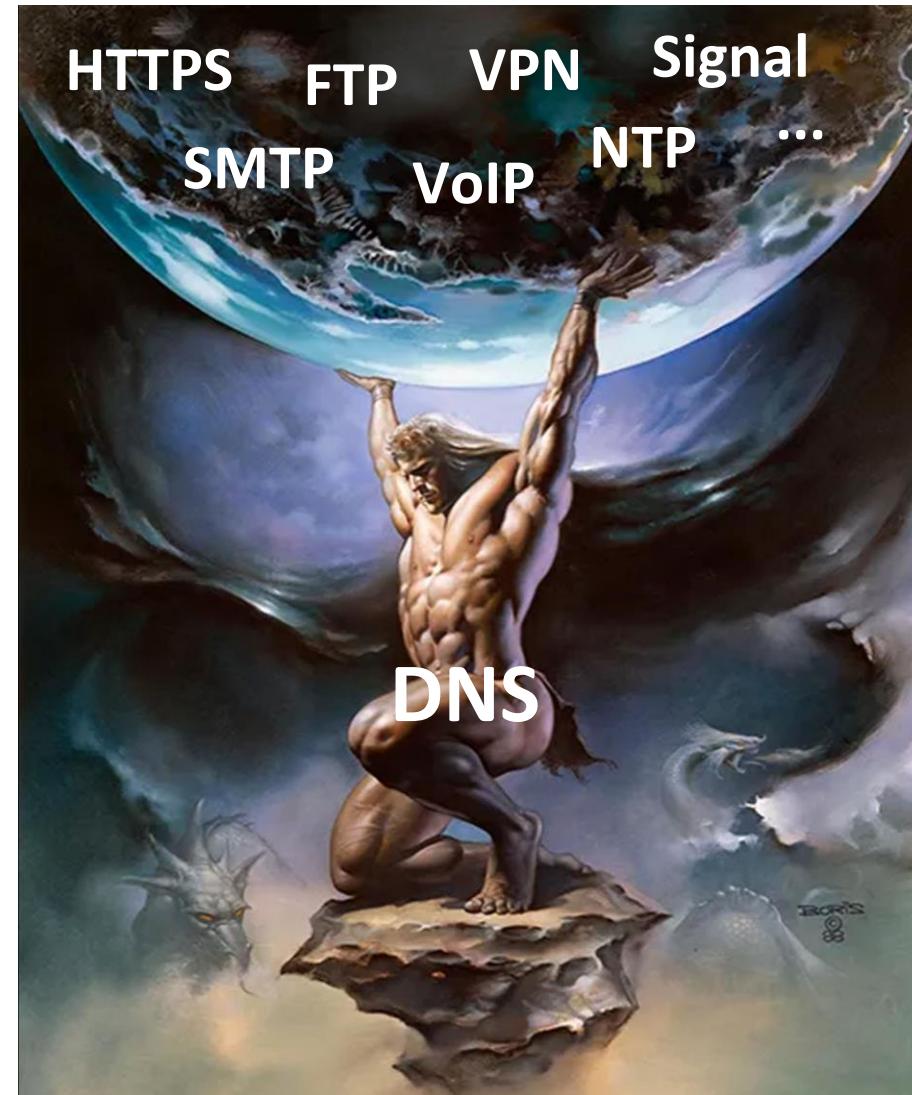
By [Juha Saarinen](#)

May 12 2021 9:06AM

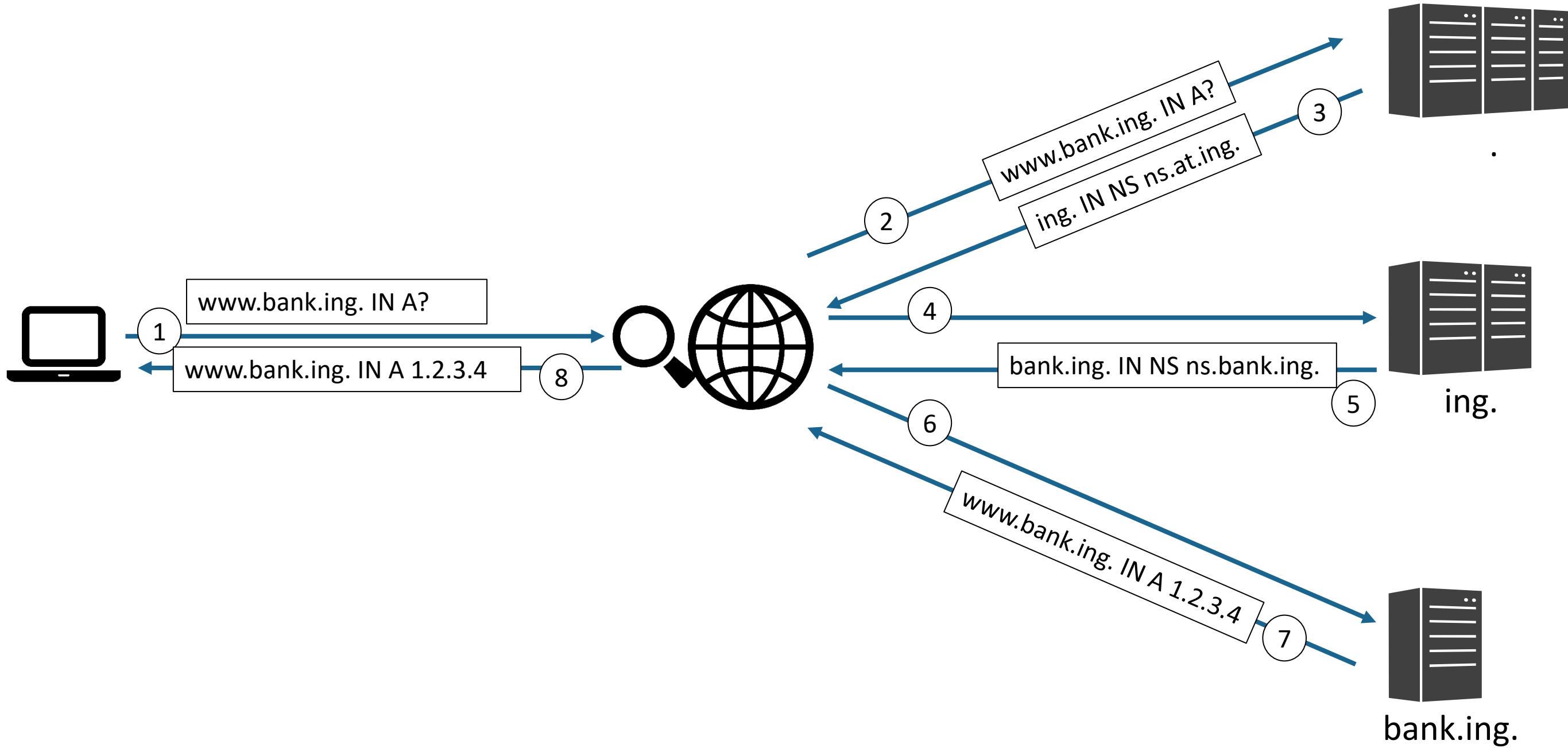


Caused by DNS issue.

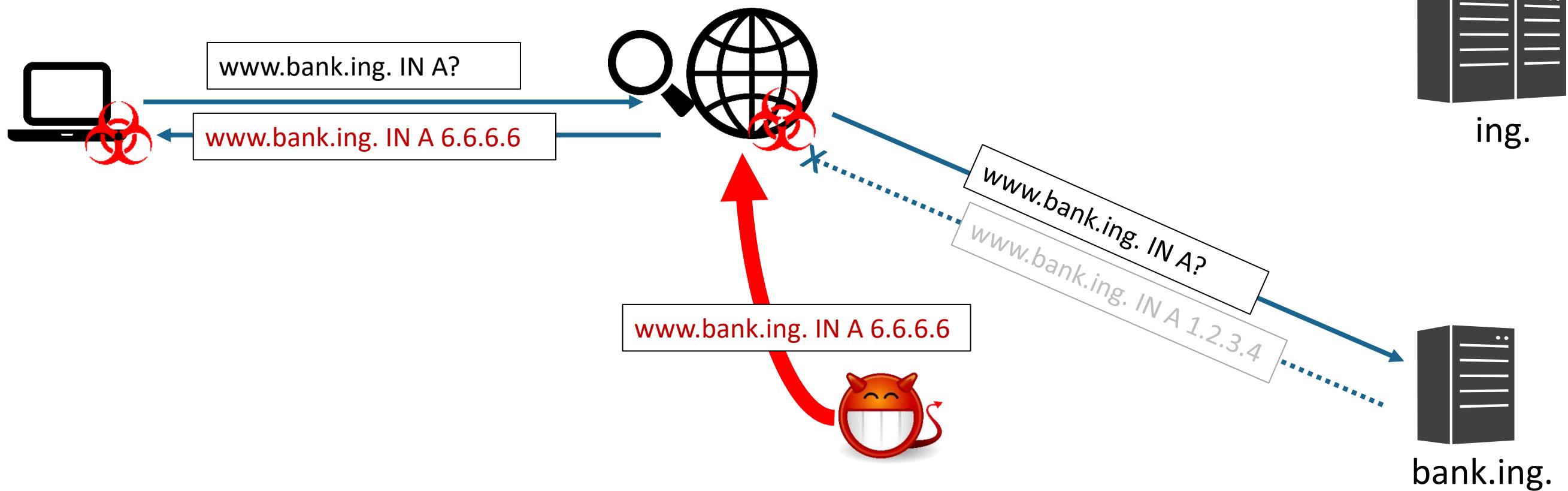
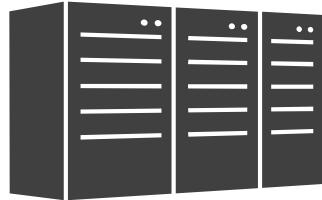
Why is it always DNS?



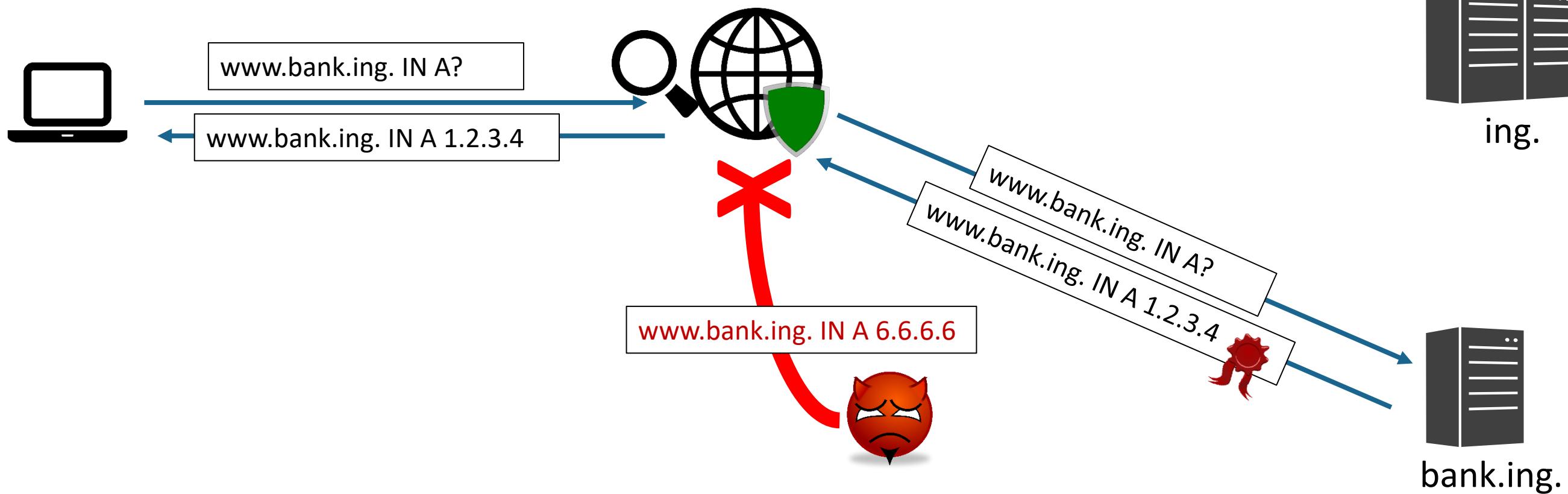
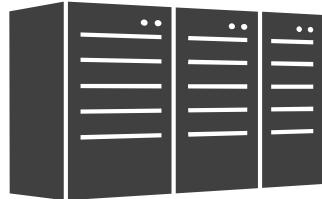
DNS Resolution



DNS Poisoning



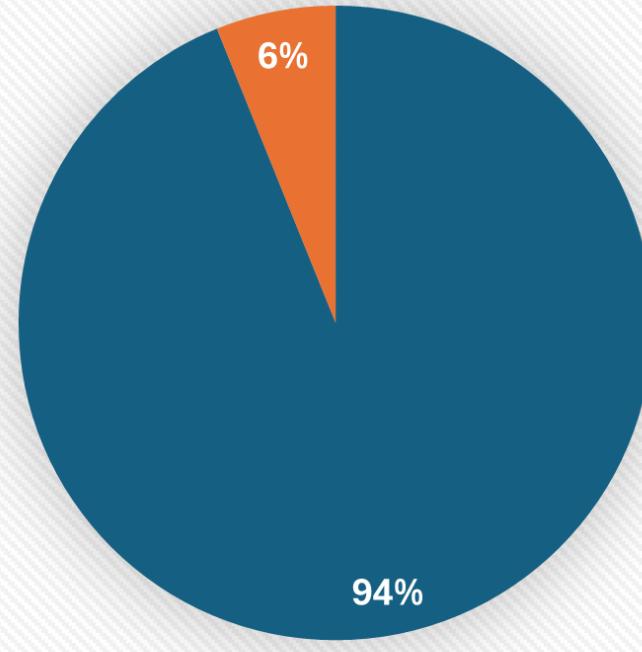
DNSSEC to the Rescue!



➤ Attack prevented

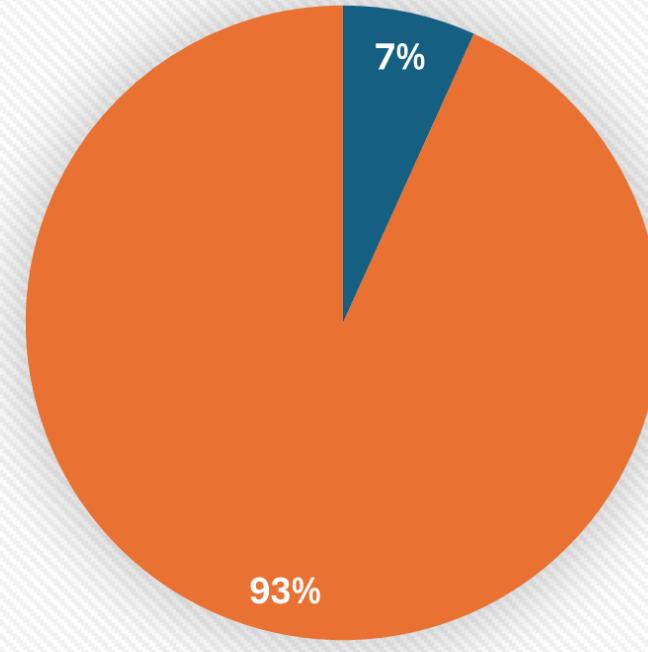
DNSSEC Adoption on the Internet

TLDs



■ signed ■ unsigned

Top 1M

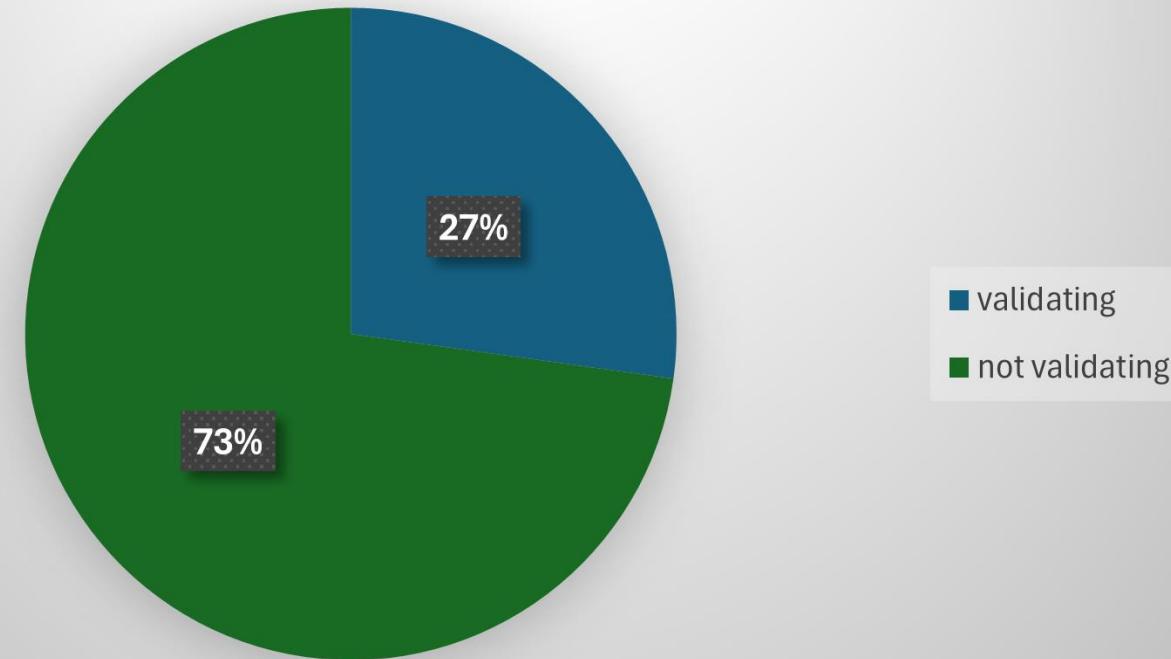


■ signed ■ unsigned

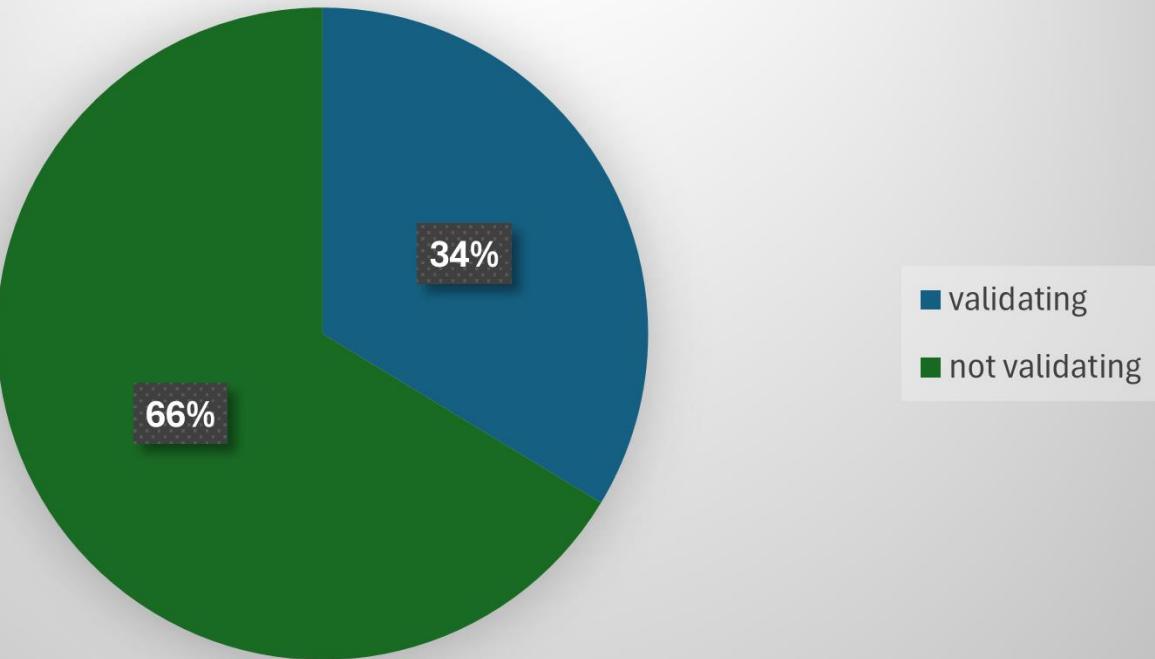
Adoption in domains is dragging

DNSSEC Adoption on the Internet

Open Resolvers



Web Clients



Better adoption in Resolvers

How DNSSEC Validation Works

```
~$ dig www.ietf.org -t A +dnssec
```

*What is the IP Address of
www.ietf.org?*

How DNSSEC Validation Works

```
~$ dig www.ietf.org -t A +dnssec
```

*What is the IP Address of
www.ietf.org?*

;; ANSWER SECTION:

www.ietf.org.	300	IN	A	104.16.45.99
www.ietf.org.	300	IN	A	104.16.44.99
www.ietf.org.	300	IN	RRSIG A 13 3 300 20241211 20241209	34505 ietf.org. eSTHK9ql[...]uvSgBA==

How DNSSEC Validation Works

```
~$ dig www.ietf.org -t A +dnssec
```

*What is the IP Address of
www.ietf.org?*

How to validate??

```
;; ANSWER SECTION:  
www.ietf.org.    300   IN    A    104.16.45.99  
www.ietf.org.    300   IN    A    104.16.44.99  
www.ietf.org.    300   IN    RRSIG A 13 3 300 20241211 20241209  
                           ↴ 34505 ietf.org. eSTHK9ql[...]uvSgBA==
```

Obtaining Public Keys

```
~$ dig ietf.org -t DNSKEY +dnssec
```

*What are the keys for
ietf.org?*

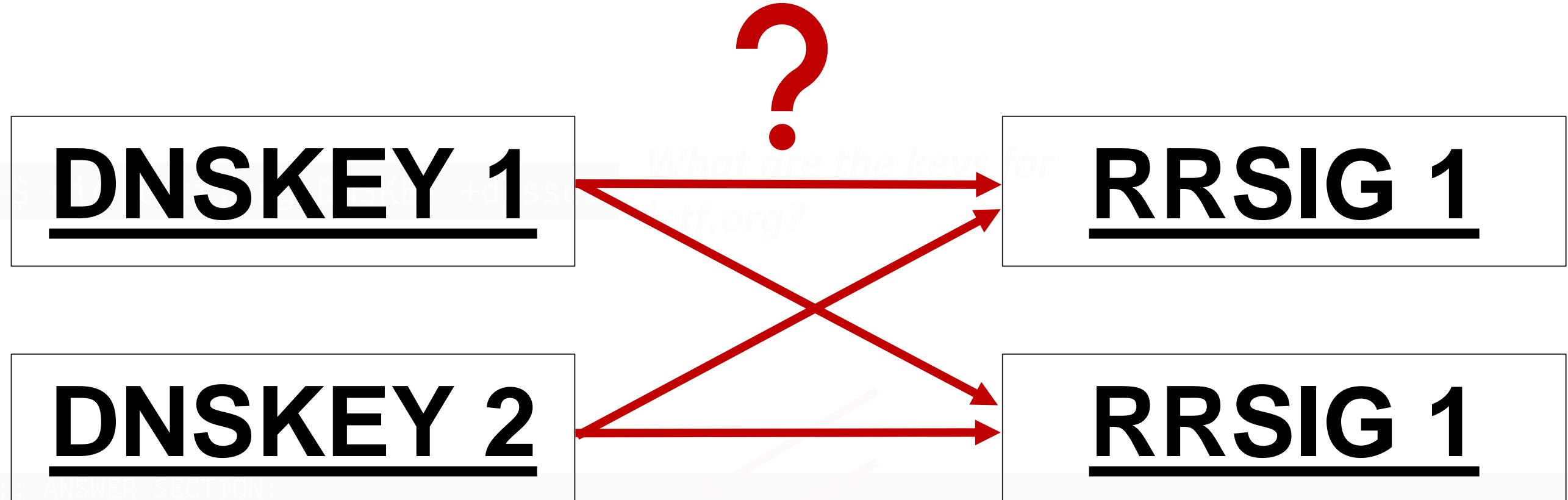
Obtaining Public Keys

```
~$ dig ietf.org -t DNSKEY +dnssec
```

*What are the keys for
ietf.org?*

```
; ANSWER SECTION:  
ietf.org.      3600   IN    DNSKEY 257 3 13 mdsswUyr[...]53eKGQ==  
ietf.org.      3600   IN    DNSKEY 256 3 13 oJMRESz5[...]a2XhSA==  
ietf.org.      3600   IN    RRSIG  DNSKEY 13 2 3600 20250130 20241130  
                           ↴ 2371 ietf.org. gdCgidVw[...]helodA==
```

Associating Signatures with Keys



Associating Signatures with Keys

;; ANSWER SECTION:

ietf.org. 3600 IN DNSKEY 257 3 13 mdsswUyr[...]53eKGQ==
ietf.org. 3600 IN DNSKEY 256 3 13 oJMRESz5[...]a2XhSA==

;; ANSWER SECTION:

www.ietf.org. 300 IN RRSIG A 13 3 300 20241211 20241209
↳ 34505 ietf.org. eSTHK9ql[...]uvSgBA==

Associating Signatures with Keys

;; ANSWER SECTION:

ietf.org.

ietf.org.

DNSKEY

13

DNSKEY

13

;; ANSWER SECTION:

www.ietf.org.

RRSIG

13

↳ 34505 ietf.org.

Associating Signatures with Keys



DNSKEY 1: ietf.org | Algo 13 | [KqXX...]



DNSKEY 2: ietf.org | Algo 13 | [KkxL...]



Signature: ietf.org | Algo 13 | Key Tag 34505

;; ANSWER SECTION:

ietf.org.

ietf.org.

[REDACTED]

DNSKEY
DNSKEY

13
13

;; ANSWER SECTION:

www.ietf.org.

[REDACTED]

RRSIG

13

↳ 34505 ietf.org.

Associating Signatures with Keys



DNSKEY 1: ietf.org | Algo 13 | $f(KqXX...) = 2371$



DNSKEY 2: ietf.org | Algo 13 | $f(KkxL...) = 34505$



Signature: ietf.org | Algo 13 | Key Tag 34505

; ANSWER SECTION:

ietf.org.

ietf.org.

DNSKEY

13

DNSKEY

13

; ANSWER SECTION:

www.ietf.org.

RRSIG

13

↳ 34505 ietf.org.

Associating Signatures with Keys



DNSKEY 1: ietf.org | Algo 13 | $f(KqXX...) = 2371$



DNSKEY 2: ietf.org | Algo 13 | $f(KkxL...) = 34505$



Signature: ietf.org | Algo 13 | Key Tag **34505**

; ANSWER SECTION:

ietf.org.

ietf.org.

DNSKEY

13

DNSKEY

13

; ANSWER SECTION:

www.ietf.org.

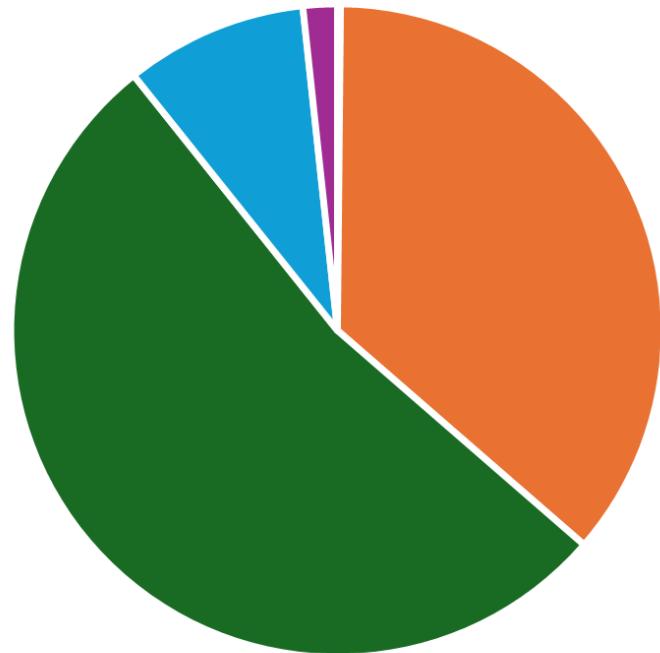
RRSIG

13

↳ 34505 ietf.org.

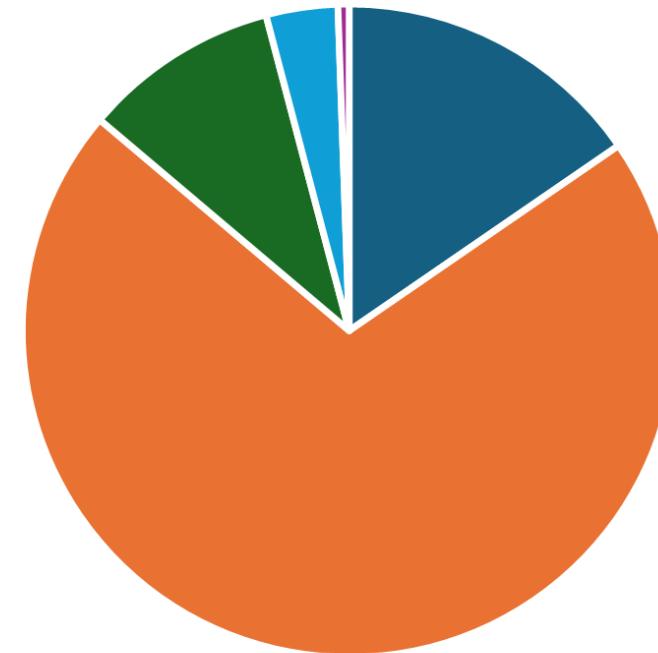
Frequencies of Keys in Domains

DNSKEYs in TLDs



■ 1 ■ 2 ■ 3 ■ 4 ■ >4

DNSKEYs in Top 1M



■ 1 ■ 2 ■ 3 ■ 4 ■ >4

Association of Keys and Signatures is efficient under normal conditions.



KeyTrap Attacks

Associating Signatures with Keys



DNSKEY 1: ietf.org | Algo 13 | $f(KqXX...) = 2371$



DNSKEY 2: ietf.org | Algo 13 | $f(KkxL...) = 34505$



Signature: ietf.org | Algo 13 | Key Tag **34505**

; ANSWER SECTION:

ietf.org.

ietf.org.

DNSKEY

13

DNSKEY

13

; ANSWER SECTION:

www.ietf.org.

RRSIG

13

↳ 34505 ietf.org.

Colliding Key Tags



DNSKEY 1: ietf.org | Algo 13 | $f(KqXX...) = 2371$



DNSKEY 2: ietf.org | Algo 13 | $f(KkxL...) = 34505$



Signature: ietf.org | Algo 13 | Key Tag **34505**

; ANSWER SECTION:

ietf.org.

ietf.org.

DNSKEY

13

DNSKEY

13

; ANSWER SECTION:

www.ietf.org.

RRSIG

13

↳ 34505 ietf.org.

Colliding Key Tags

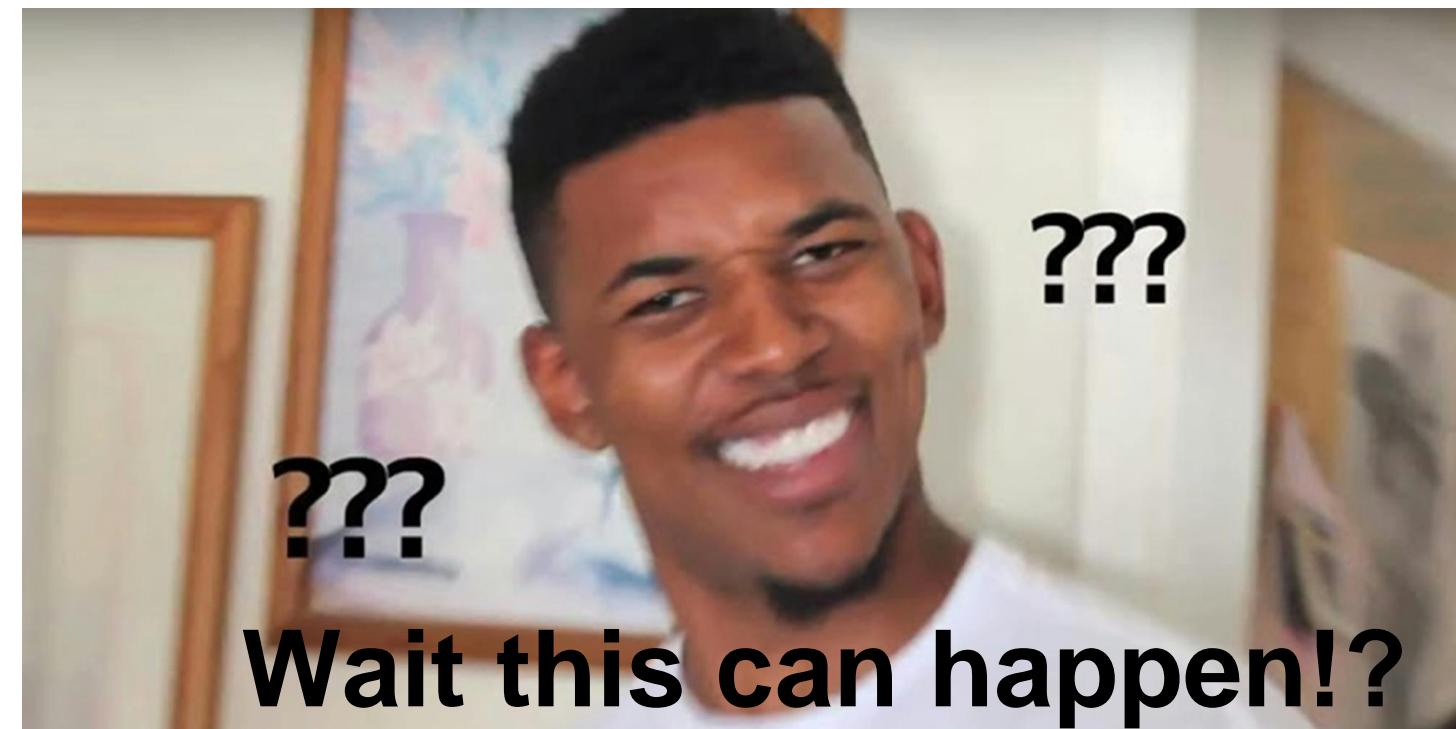


DNSKEY 1: ietf.org | Algo 13 | $f(KqXX...) = 34505$

DNSKEY 2: ietf.org | Algo 13 | $f(KkxL...) = 34505$



Signature: ietf.org | Algo 13 | Key Tag **34505**



A Closer Look at Signature Validation

However, it is essential to note that the key tag is not a unique identifier. It is theoretically possible for two distinct DNSKEY RRs to have the same owner name, the same algorithm, and the same key tag. The key tag is used to limit the possible candidate keys, but it does not uniquely identify a DNSKEY record. Implementations **MUST NOT** assume that the key tag uniquely identifies a DNSKEY RR.

RFC4034, Appendix B "Key Tag Calculation"

A Closer Look at Signature Validation

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RFC4034, Appendix B "Key Tag Calculation"

It is possible for more than one DNSKEY RR to match the conditions above. In this case, the validator cannot predetermine which DNSKEY RR to use to authenticate the signature, and it MUST try each matching DNSKEY RR until either the signature is validated or the validator has run out of matching public keys to try.

RFC4035, Section 5.3.1. "Checking the RRSIG RR Validity"

A Closer Look at Signature Validation

However, it is essential to note that there can be more than one unique identifier. It is theoretically possible for two or more distinct DNSKEY RRs to have the same key tag. This is because the same algorithm, and the same key tag, can be used to uniquely identify a DNSKEY record. Therefore, implementers MUST NOT assume that the key tag identifies a single DNSKEY RR.

RFC4034, Appendix B "Key Identifier"



more than one DNSKEY RR to match the signature above. In this case, the validator cannot determine which DNSKEY RR to use to authenticate the signature and it MUST try each matching DNSKEY RR until the signature is validated or the validator runs out of matching public keys to try.

Section 5.3.1. "Checking the RRSIG RR Validity"

A Closer Look at Signature Validation



- Linear time algorithm
- Resource-intensive public key crypto operations
- Resolver MUST try all keys?

That sounds like a bad idea...

A Closer Look at Signature Validation

Wait, wasn't there more like this?

- Linear time algorithm
- Resource intensive public key crypto operations
- Resolver MUST try all keys?

That sounds like a bad idea...

We can make it even worse

This document specifies that a resolver SHOULD accept any valid RRSIG as sufficient, and only determine that an RRset is Bogus if all RRSIGs fail validation.

If a resolver adopts a more restrictive policy, there's a danger that properly signed data might unnecessarily fail validation due to cache timing issues. Furthermore, certain zone management techniques, like the Double Signature Zone Signing Key Rollover method described in Section 4.2.1.2 of [RFC6781], will not work reliably. Such a resolver is also vulnerable to malicious insertion of gibberish signatures.

RFC6840 Section 5.4. "Caution about Local Policy and Multiple RRSIGs"

It's quadratic!

This document specifies that a resolver **SHOULD** accept any valid RRSIG as sufficient, and only determine that an RRset is Bogus if all RRSIGs fail validation.

If a resolver adopts a more restrictive policy, there's a danger that properly signed data might unnecessarily fail validation due to cache timing issues. Furthermore, certain zone management techniques, like the Double Signature Zone Signing Key Rollover method described in Section 4.2.1.2 of [RFC6781], will not work reliably. Such a resolver is also vulnerable to malicious insertion of gibberish signatures.

RFC6840 Section 5.4. "Caution about Local Policy and Multiple RRSIGs"



Open Source is doing this...

```
// Get amount of signatures on the rrset
num = rrset_get_sigcount(rrset);

// Iterate all Signatures
for (i = 0; i < num; i++) ← Signature Loop
{
    sec = dnskeyset_verify_rrset_sig(env, ve, *env->now, rrset,
        dnskey, i, &sortree, reason, reason_bogus,
        section, qstate);
```

```
// Get amount of keys in keyset
size_t i, num = rrset_get_count(dnskey);

// Iterate over all keys for the set
for (i = 0; i < num; i++) ← Key Loop
{
    /* see if key matches keytag and algo */
    if (algo != dnskey_get_algo(dnskey, i) ||
        tag != dnskey_calc_keytag(dnskey, i))
        continue;

    sec = dnskey_verify_rrset_sig(env->scratch,
        env->scratch_buffer, ve, now, rrset, dnskey, i,
        sig_idx, sortree, &buf_canon, reason, reason_bogus,
        section, qstate);

    if (sec == sec_status_secure)
        return sec;
    else if (sec == sec_status_ineterminate)
        numindeterminate++;
}
```

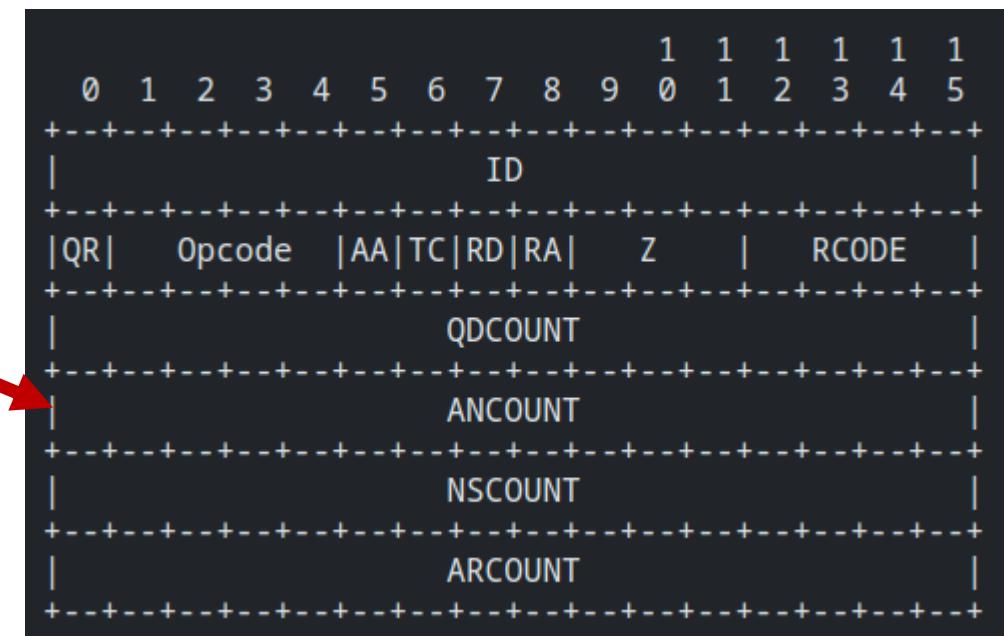
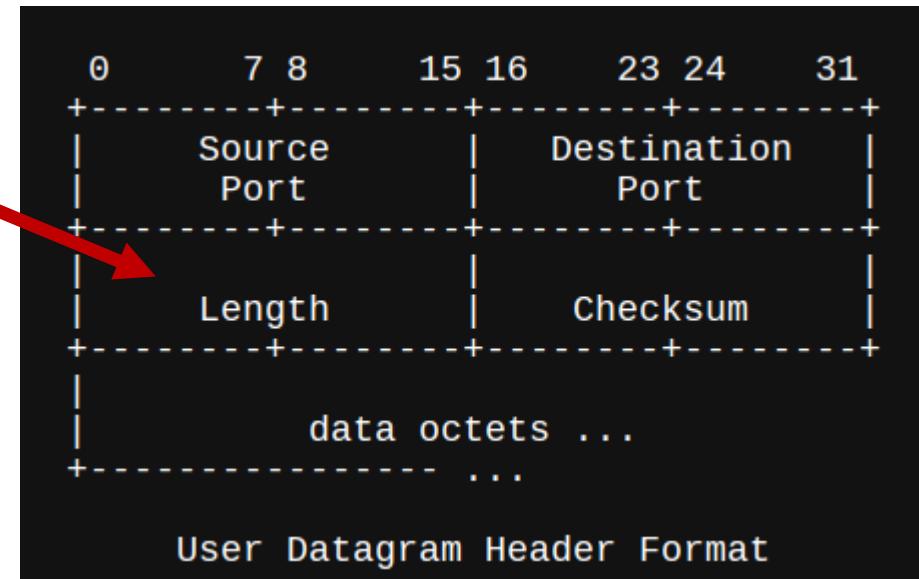
[https://github.com/NLnetLabs/unbound/
blob/master/validator/val_sigcrypt.c](https://github.com/NLnetLabs/unbound/blob/master/validator/val_sigcrypt.c) (line 704)

[https://github.com/NLnetLabs/unbound/
blob/master/validator/val_sigcrypt.c](https://github.com/NLnetLabs/unbound/blob/master/validator/val_sigcrypt.c) (line 641)

Nested Loops!

- Limited buffers?
- Other Mitigations in place?
→ there is only one way to find out

How much can “all” be?



Not many restrictions in the DNS protocol

Quick'n'Dirty Maths

Size DNS msg. 65 kB

Keys : 80 Byte $\Rightarrow 65.000 \text{ // } 80 \sim 820 \text{ p. msg.}$

Sigs: 110 Byte $\Rightarrow 65.000 \text{ // } 110 \sim 600 \text{ p. msg.}$

Validations: Keys \times Sigs $\Rightarrow 820 \times 600 \sim 490.000$

Time per Validation: 150 ns [1]

Stall from 1 req: $490.000 \times 150 \text{ ps} \sim 74 \text{ s}$

[1] <https://blog.cloudflare.com/go-crypto-bridging-the-performance-gap/>

Let's construct an attack...



How to obtain lots of DNSKEYs with the same tag?

- Generate many and keep only those conforming to a specific one

Zone File Contents: DNSKEYs

```
keytrap.org 60 IN DNSKEY 257 3 13 YpsCtFxj[...]0/G81g== ;{id=26539 (ksk)}
```

```
keytrap.org 60 IN DNSKEY 256 3 13 wwZzOFux[...]7s+jGg== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 AH18SDyd[...]vgk3gQ== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 8GP1H4eS[...]3X6dCA== ;{id=1337 (zsk)}  
...  
keytrap.org 60 IN DNSKEY 256 3 13 brtDSnWm[...]nz1K0w== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 WfKVRBtM[...]eC3A1w== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 6+GbtX4h[...]RacDLw== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 6KNR3F+Q[...]RcK2kQ== ;{id=1337 (zsk)}  
keytrap.org 60 IN DNSKEY 256 3 13 ijhRFa4f[...]L5cBtQ== ;{id=1337 (zsk)}
```

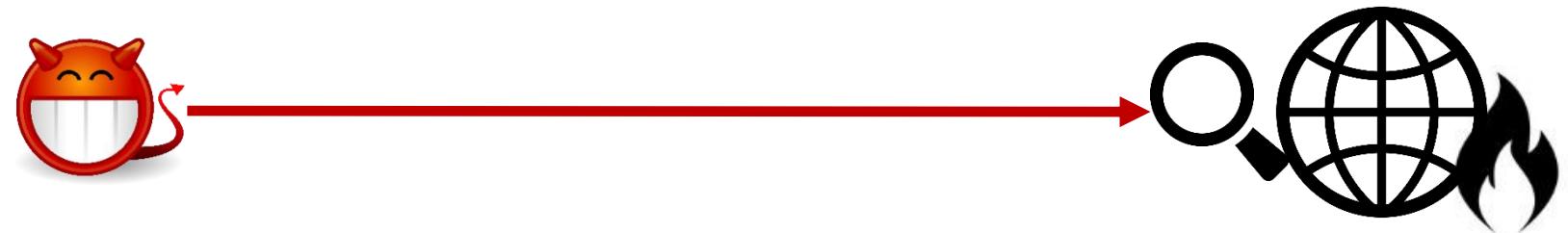
```
keytrap.org 60 IN RRSIG DNSKEY 13 4 60 20250117 20240119 26539 keytrap.org gTq+83Bx[...]hcLjQ==
```

Zone File Contents: RRSIGs

```
a0001.keytrap.org 0 IN A 10.2.3.4
```

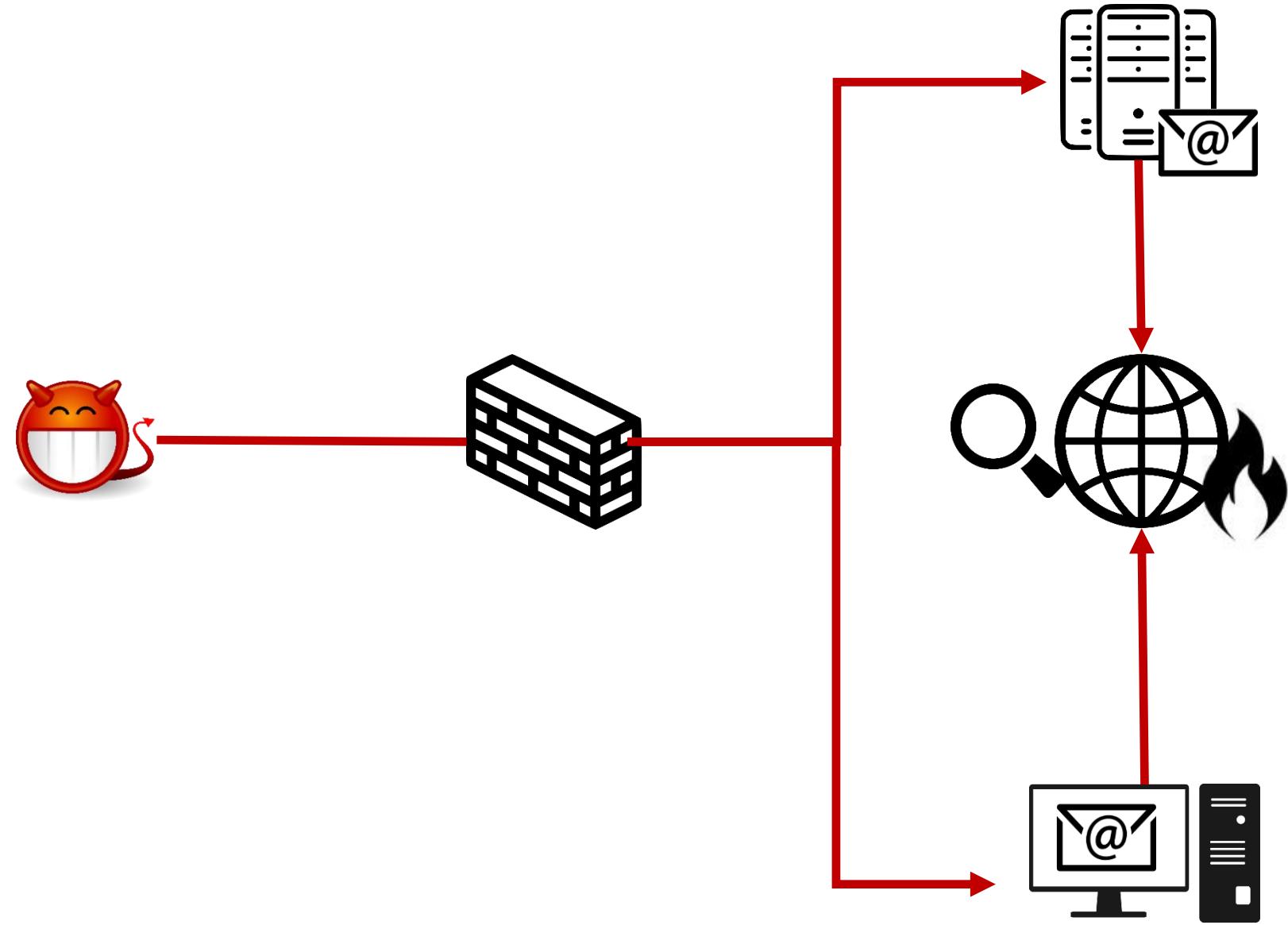
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
...												
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==
a0001.keytrap.org	0	IN	RRSIG	A	13	5	0	20250117	20240119	1337	keytrap.org	GVHrz5G0+[...]0V0bdw==

Attracting a Victim Resolver



Attack Vectors
- Querying directly

Attracting a Victim Resolver



Attack Vectors

- Querying directly
- SMTP Bounce
- HTML E-mail

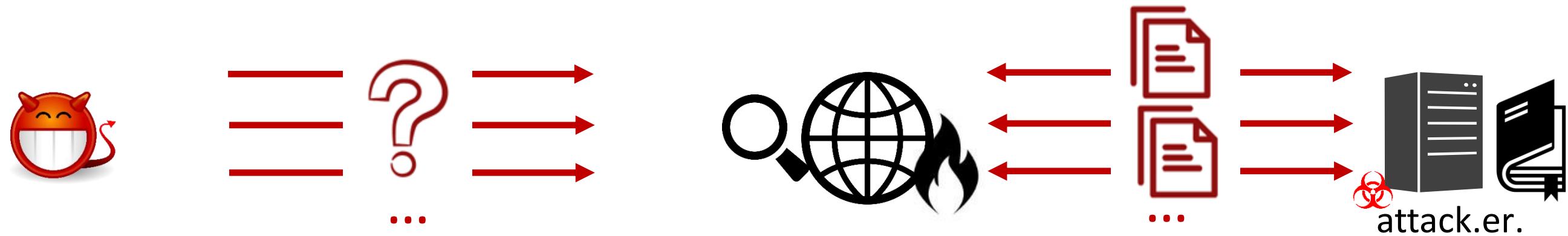
Scaling up the victim count

How to attract many resolvers at the same time?



- Online Ads
 - Internet Measurement Networks

Attack Procedure



- Flooding the victim resolver with malicious queries

Demo Time

(Demo)

KeyTrap Attacks



SigJam

Jamming the validator with
signatures only

LockCram

Collisions-only attack, using a
single RRSIG per response

KeySigTrap

Combined attack

HashTrap

KeySigTrap with DS records



Impact and Evaluations

Which algorithm to use?

Cipher	Keys	Signatures	Validations
ED448	907	454	411 778
ED25519	1 391	696	968 136
ECDSAP384SHA384	589	519	305 691
ECDSAP256SHA256	828	696	576 288
RSA-512	788	696	548 448
RSA-1024	444	413	183 372
RSA-2048	237	228	54 036
RSA-4096	122	119	14 518

Max validations

Min validations

Up to ~1M signature validations per DNS Request

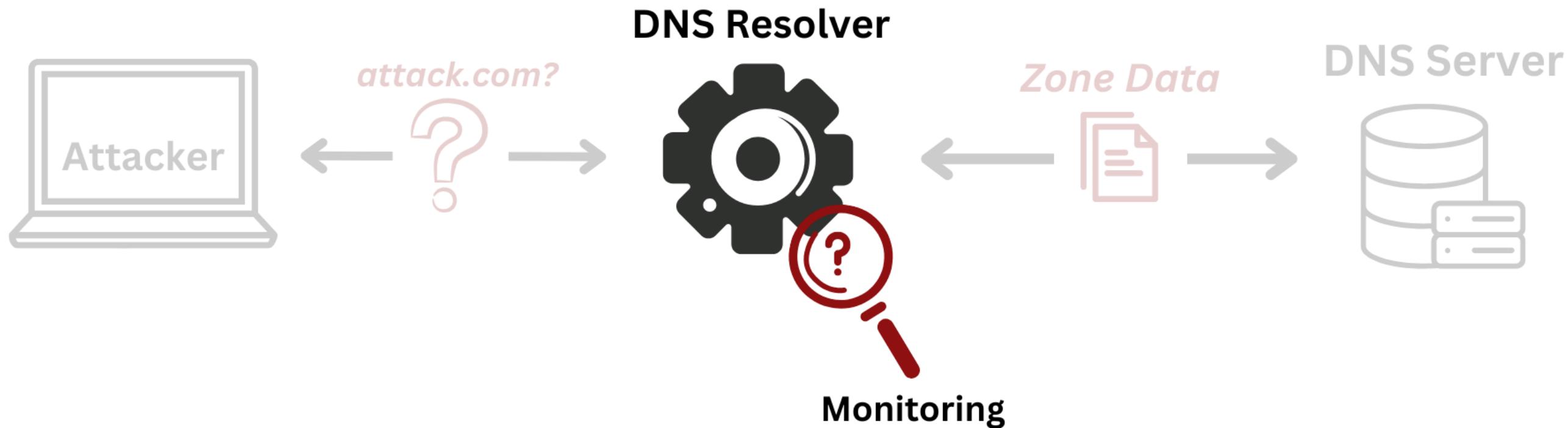
Test Setup with Unbound.



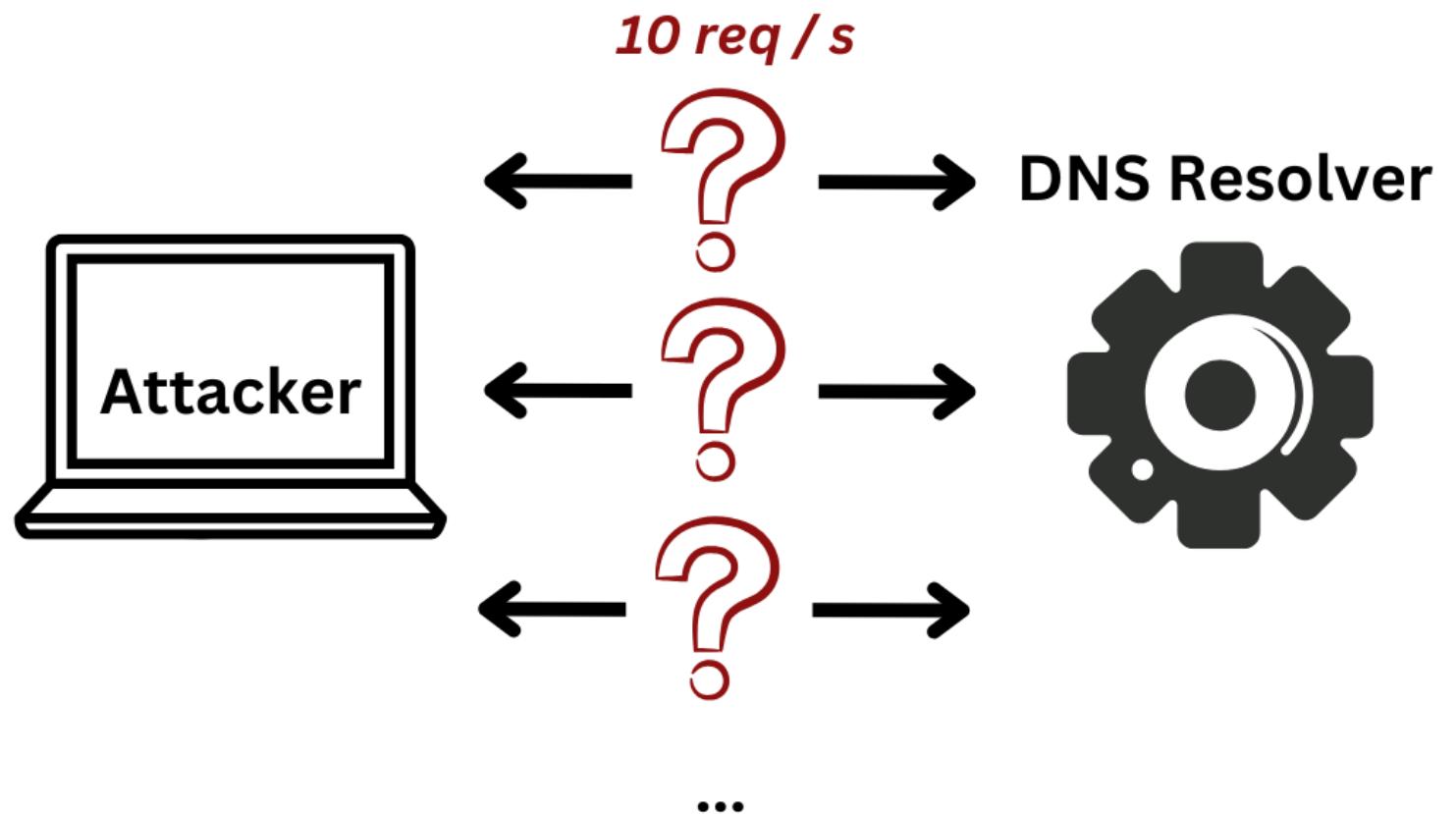
Test Setup with Unbound.



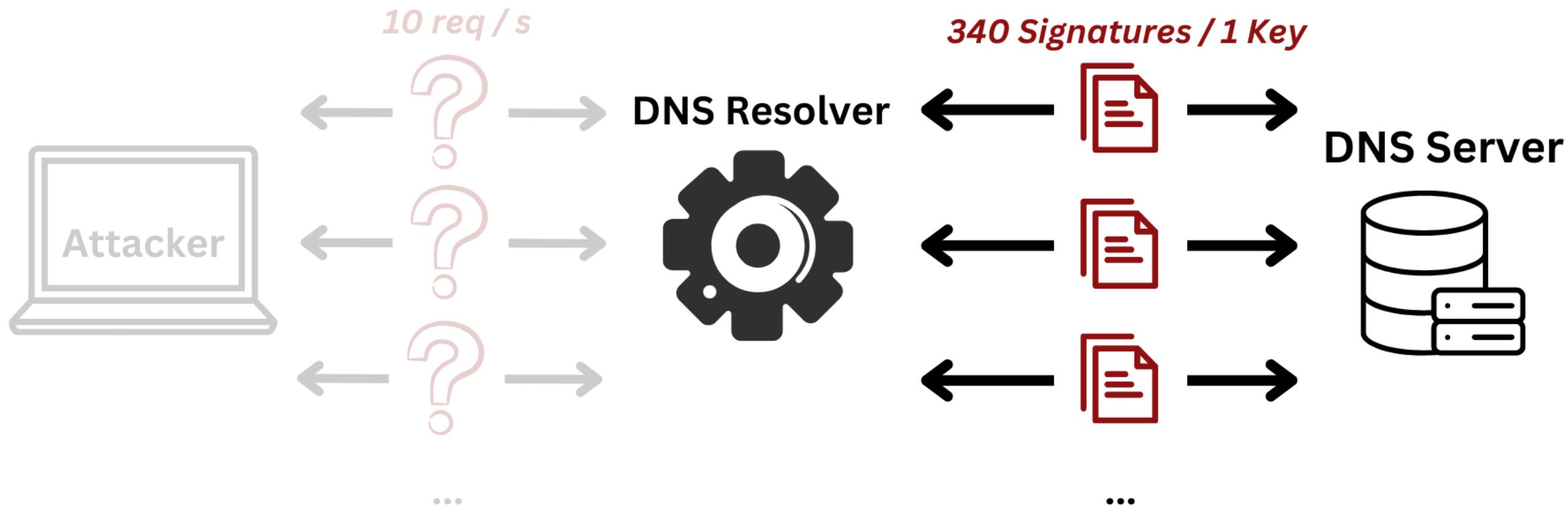
Test Setup with Unbound.



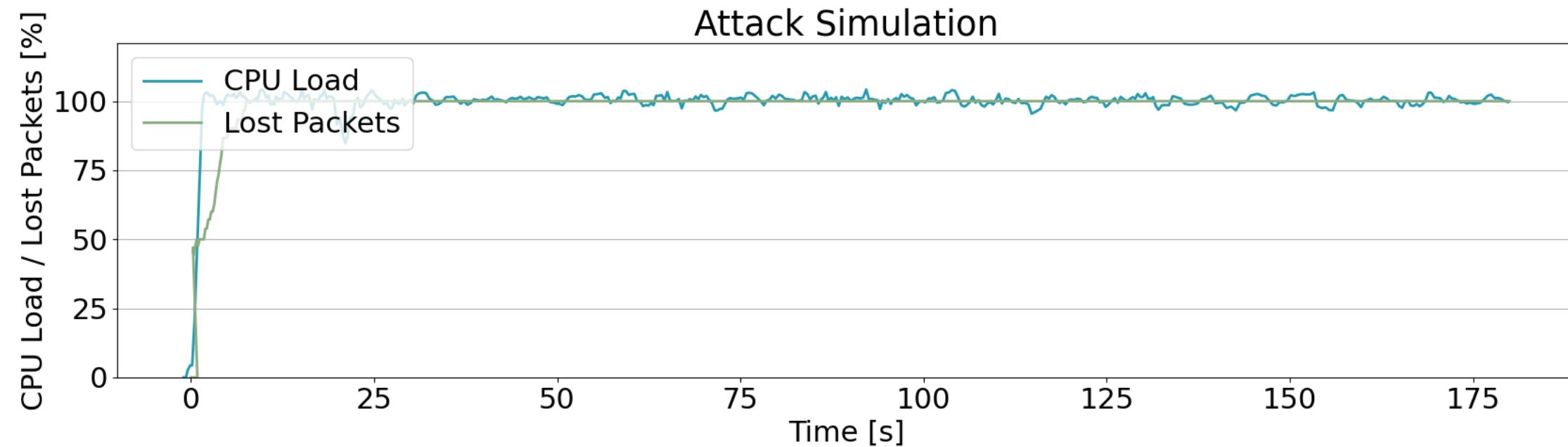
Test 1: SigJam (Many Signatures / 1 Key)



Test 1: SigJam (Many Signatures / 1 Key)



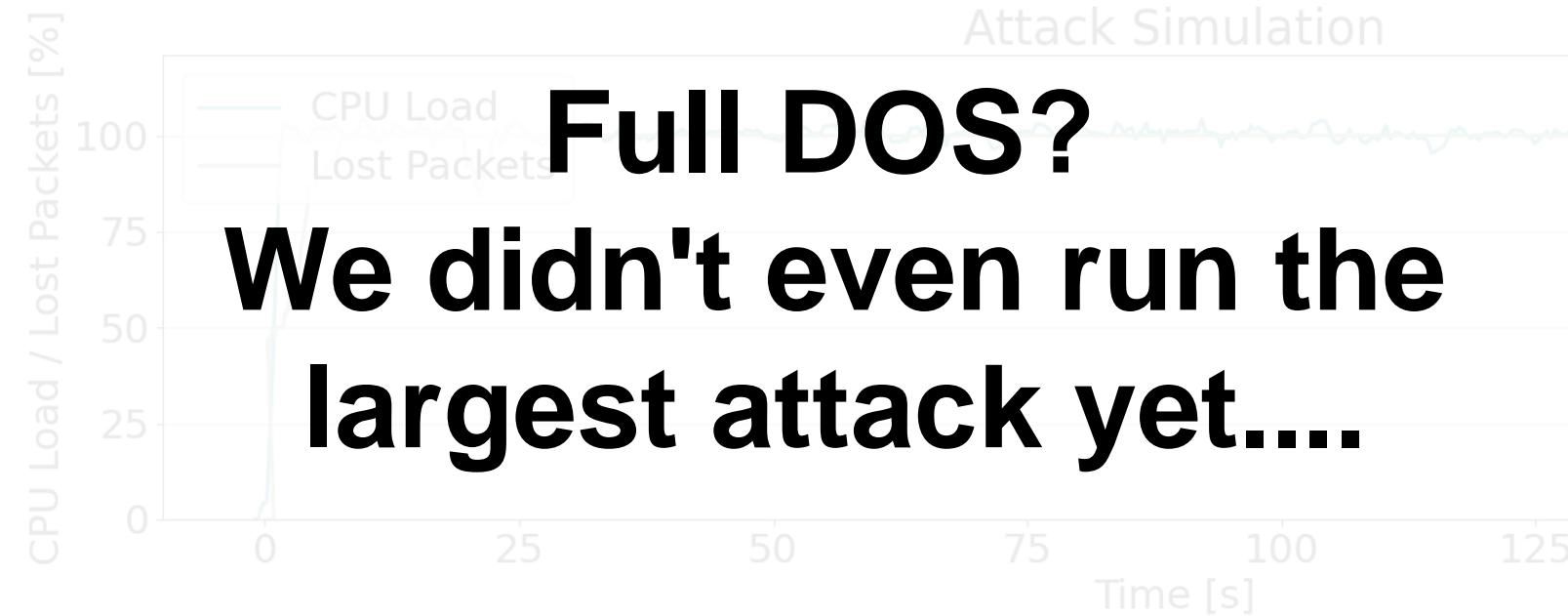
340 Signature Validations per Request
10 req/s



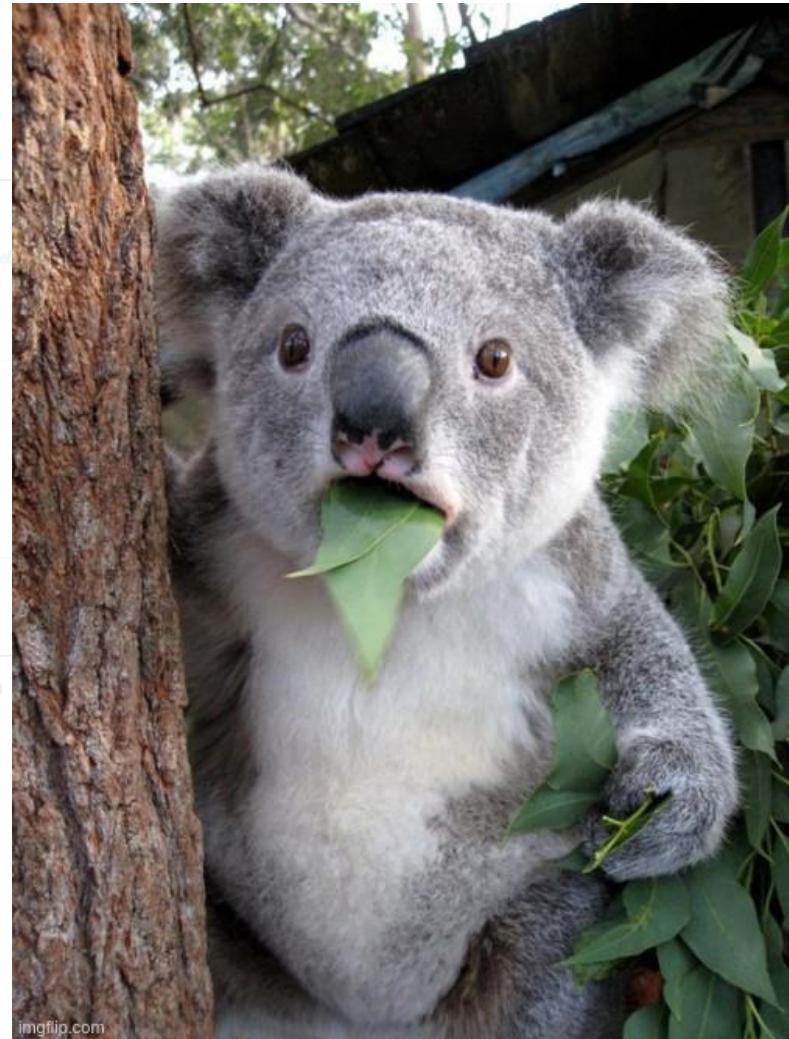
Measured on Intel(R) Xeon(R) Gold 6242 CPU @ 2.80GHz with Unbound

Evaluations

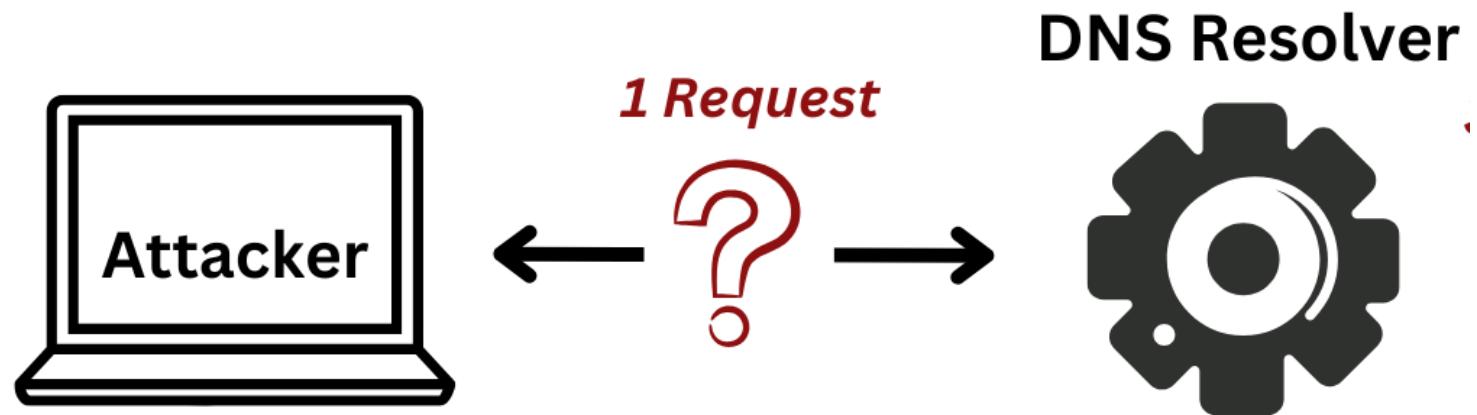
340 Signature Validations per Request
10 req/s



Measured on Intel(R) Xeon(R) Gold 6242 CPU @ 2.80GHz



Test 2: KeySigTrap (Many Signatures / Many Keys)



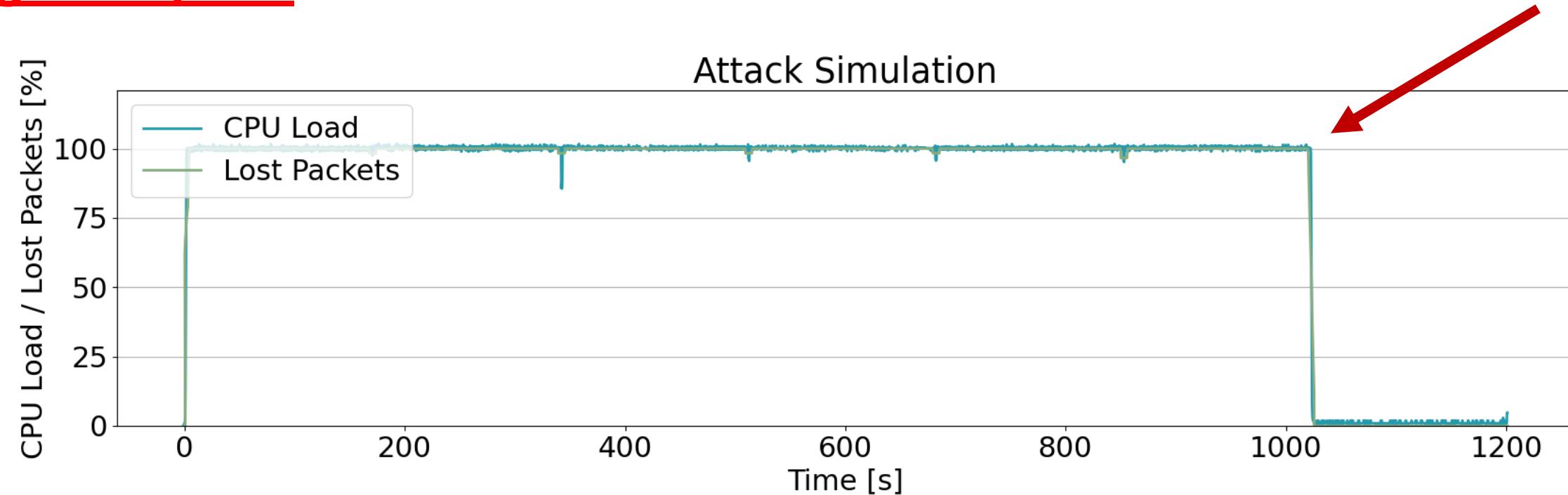
Test 2: KeySigTrap (Many Signatures / Many Keys)



Evaluations

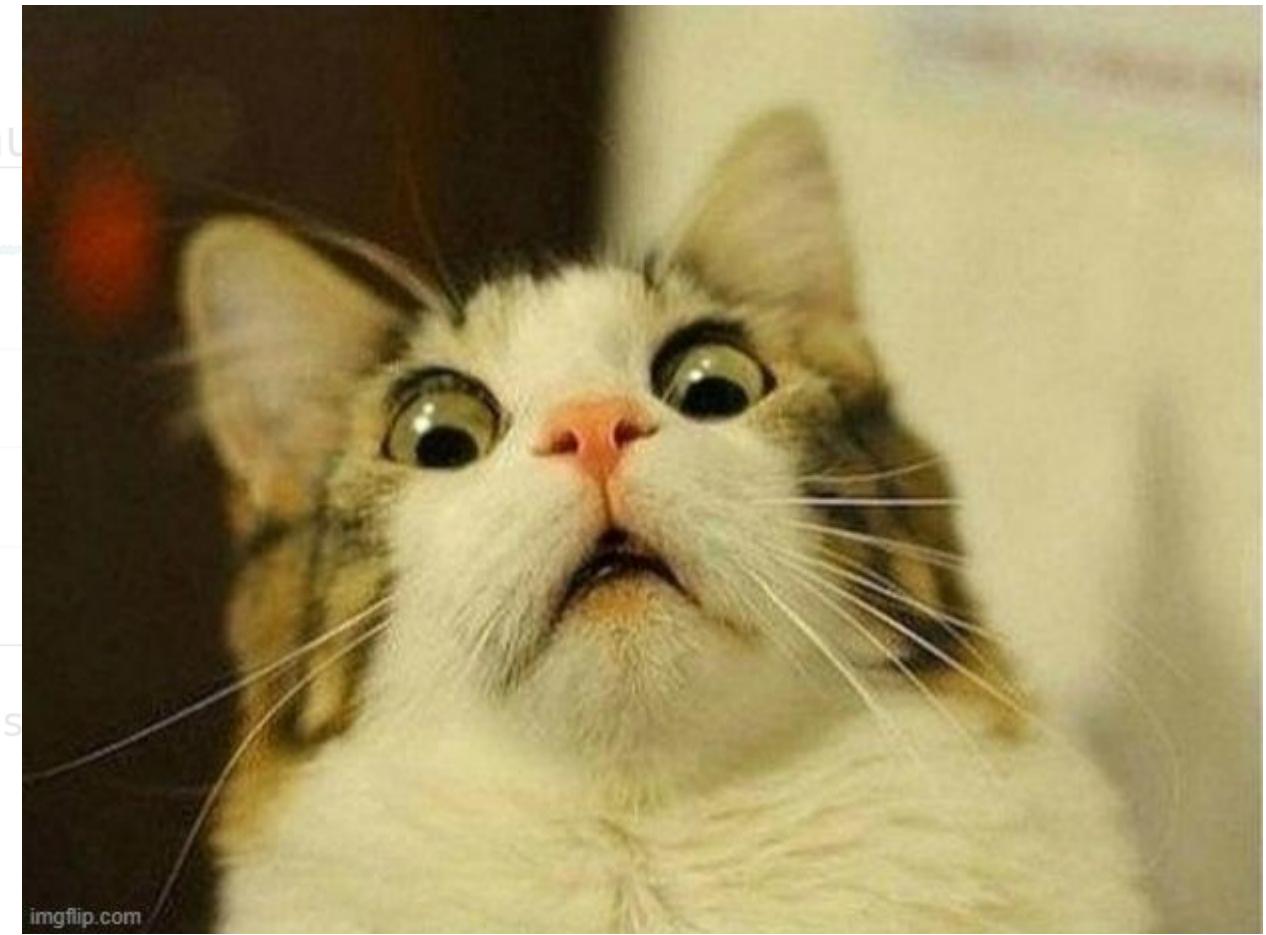
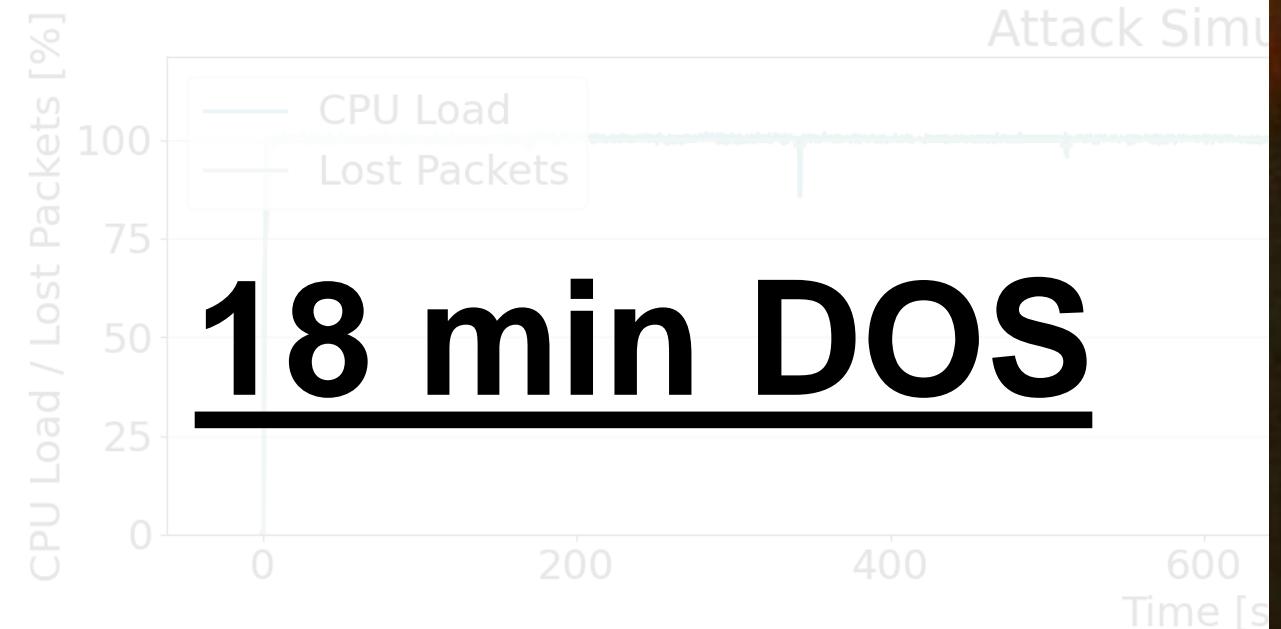
340 Signatures x 582 Keys = 221160 Validations
Single Request

340 Signatures x 582 Keys = 221160 Validations
Single Request



Measured on Intel(R) Xeon(R) Gold 6242 CPU @ 2.80GHz

340 Signatures x 582 Keys = 221160 Validations
Single Request



Measured on Intel(R) Xeon(R) Gold 6242 CPU @ 2.80GHz

Evaluations

Resolver	Stall Duration (Single Request)
Unbound	1014s
Bind9	58632s
Knot	51s
Akamai	186s
PowerDNS	170s
Windows Server	132s
Stubby	184s
Cloudflare 1.1.1.1	Vulnerable
Amazon Route 53	Vulnerable
Google DNS	Vulnerable

All tested DNSSEC implementing resolvers vulnerable!

Evaluations

Resolver	Stall Duration (Single Request)
Unbound	1014s
Bind9	58632s
Knot	51s
Akamai	186s
PowerDNS	170s
Windows Server	132s
Stubby	184s
Cloudflare 1.1.1.1	Vulnerable
Amazon Route 53	Vulnerable
Google DNS	Vulnerable

All tested DNSSEC implementing resolvers vulnerable!

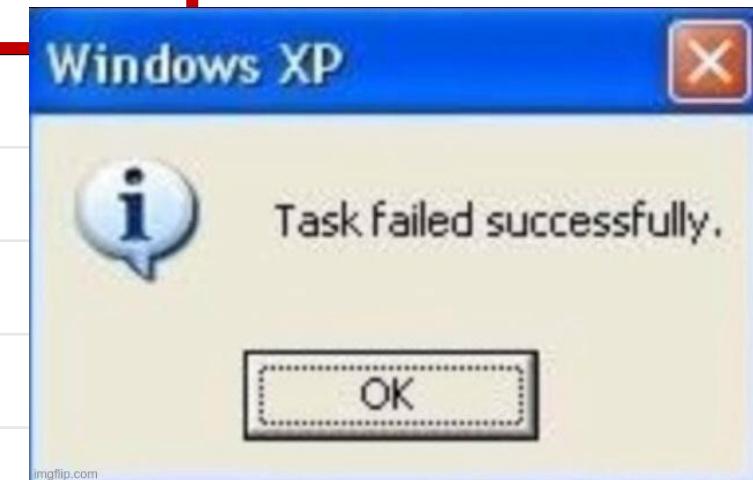
Evaluations

Resolver	Stall Duration (Single Request)
Unbound	1014s
Bind9	58632s
Knot	51s
Akamai	186s
PowerDNS	170s
Windows Server	132s
Stubby	184s
Cloudflare 1.1.1.1	Vulnerable
Amazon Route 53	Vulnerable
Google DNS	Vulnerable

All tested DNSSEC implementing resolvers vulnerable!

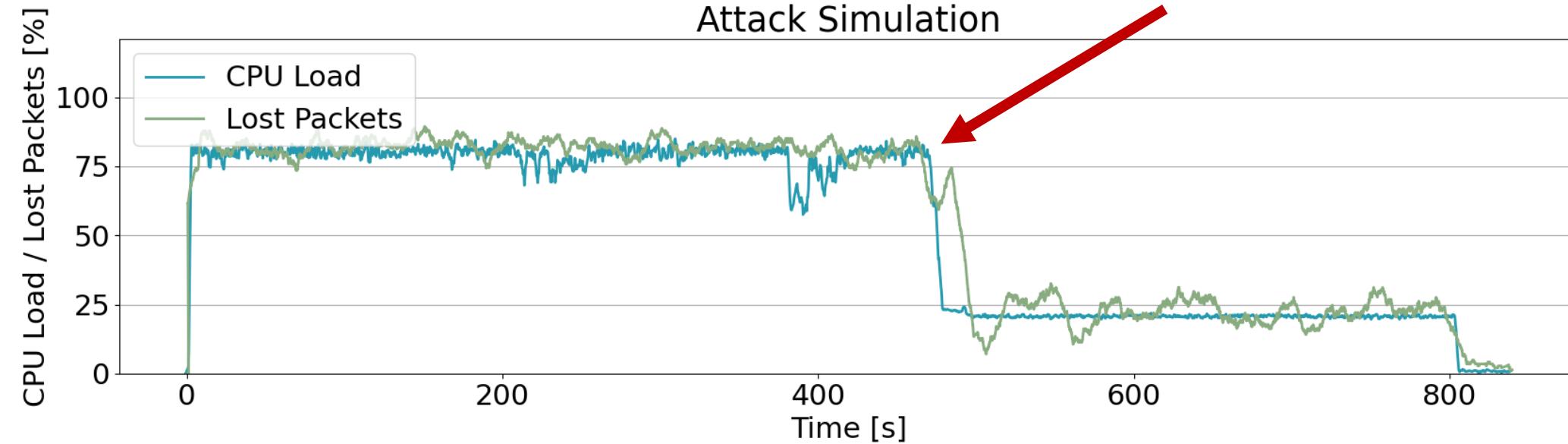
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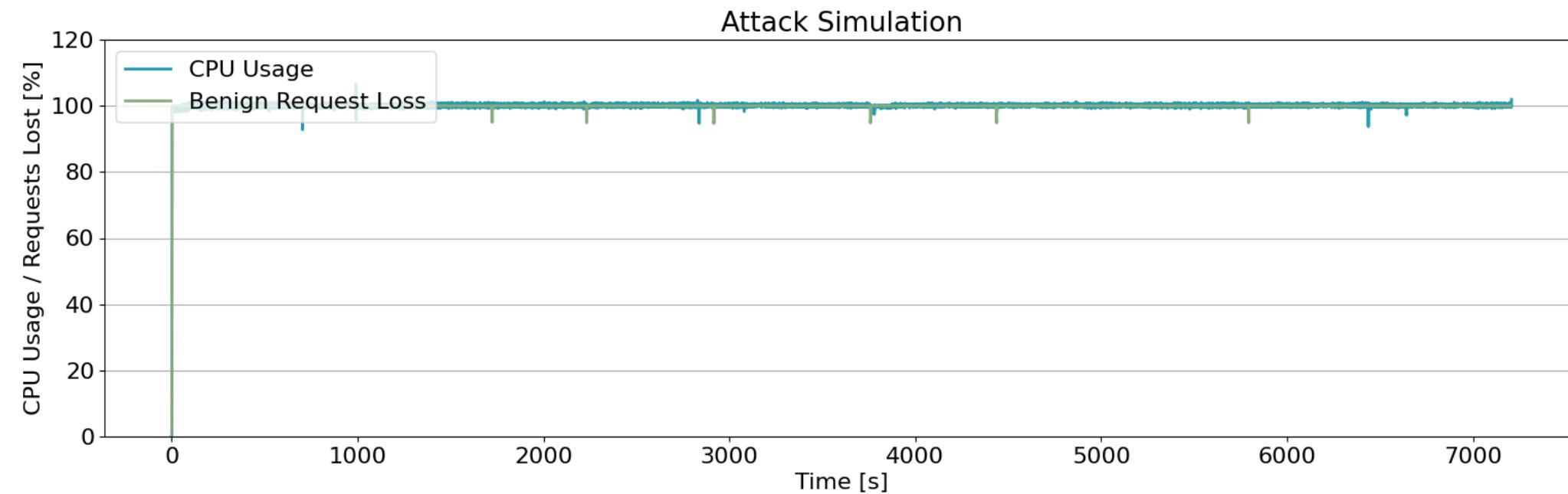
All tested DNSSEC implementing resolvers vulnerable!

Does Multi-Threading help?



No, but attacker needs to send a few more requests....

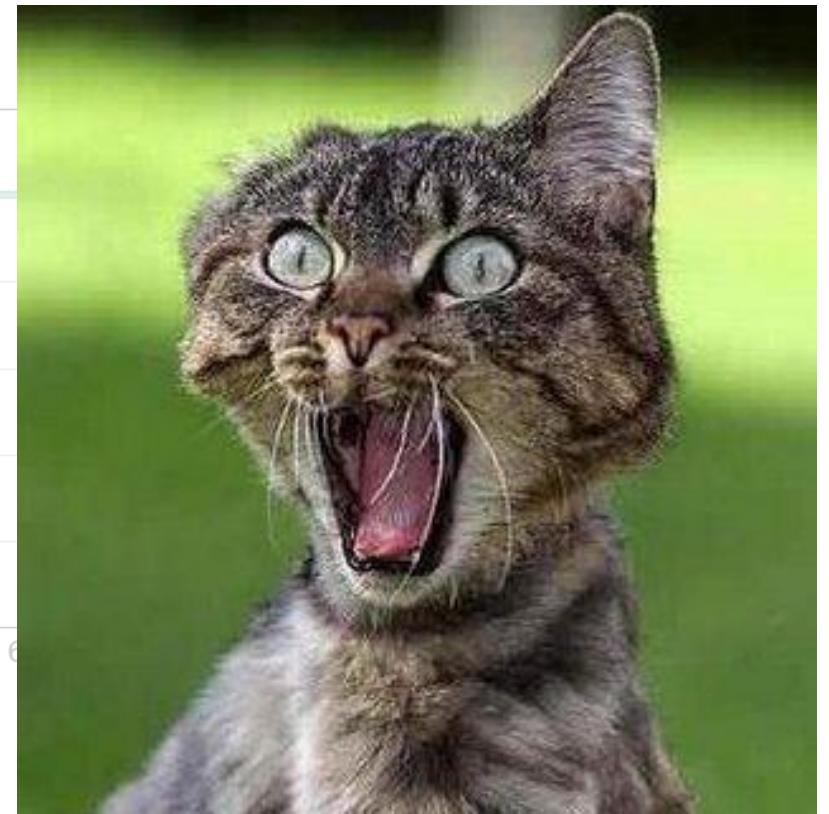
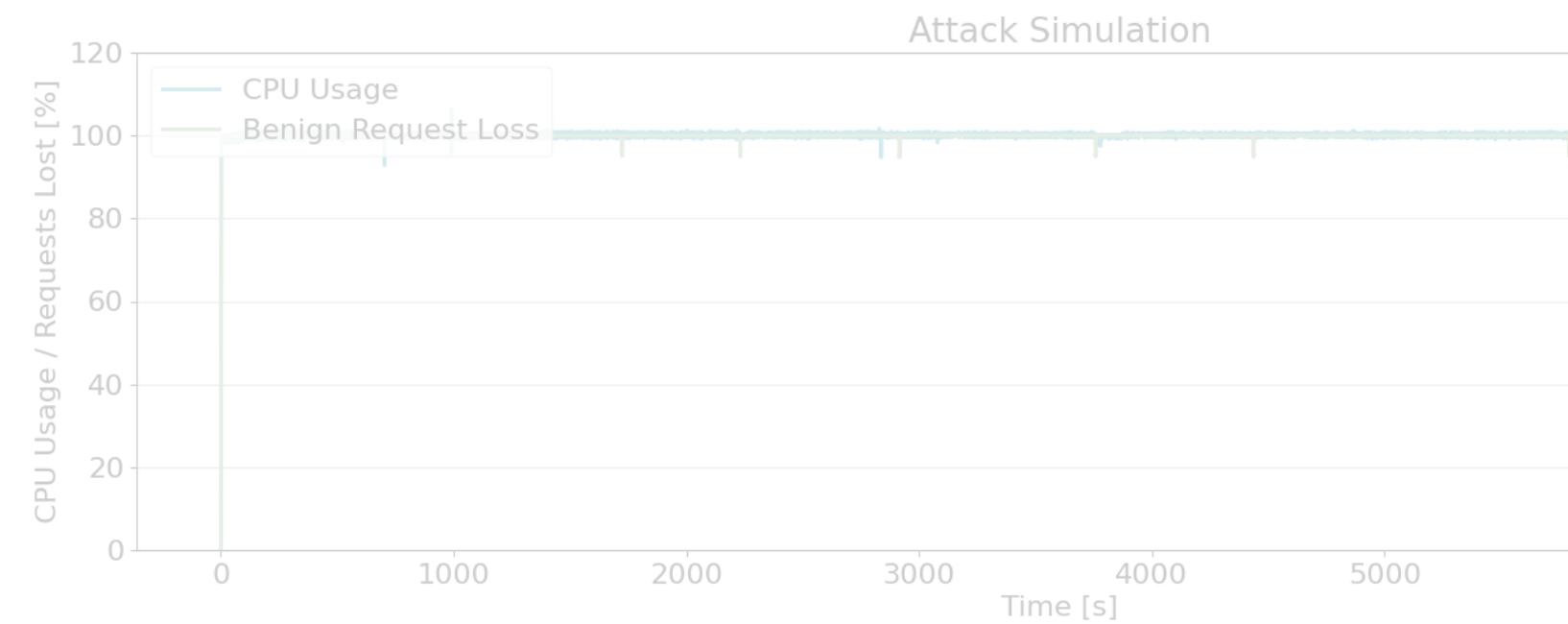
Viable attack: Comprehensive persistent DoS



KeyTrap allows to disable DNSSEC resolver with max. a few packets / min

Evaluations

Viable attack: Comprehensive persistent DoS

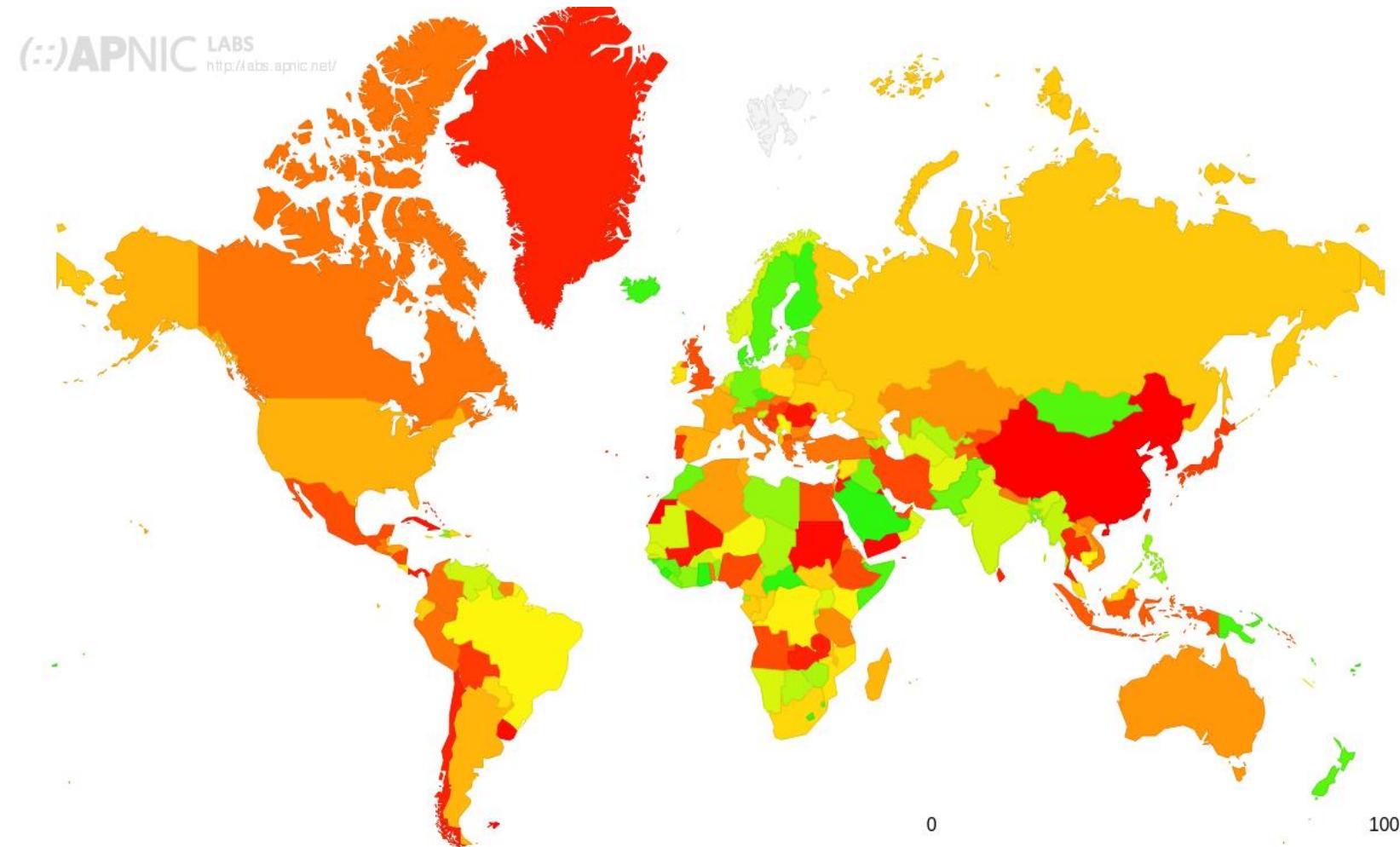


KeyTrap allows to disable DNSSEC resolver with max. a few packets / min



Impact

Global Impact

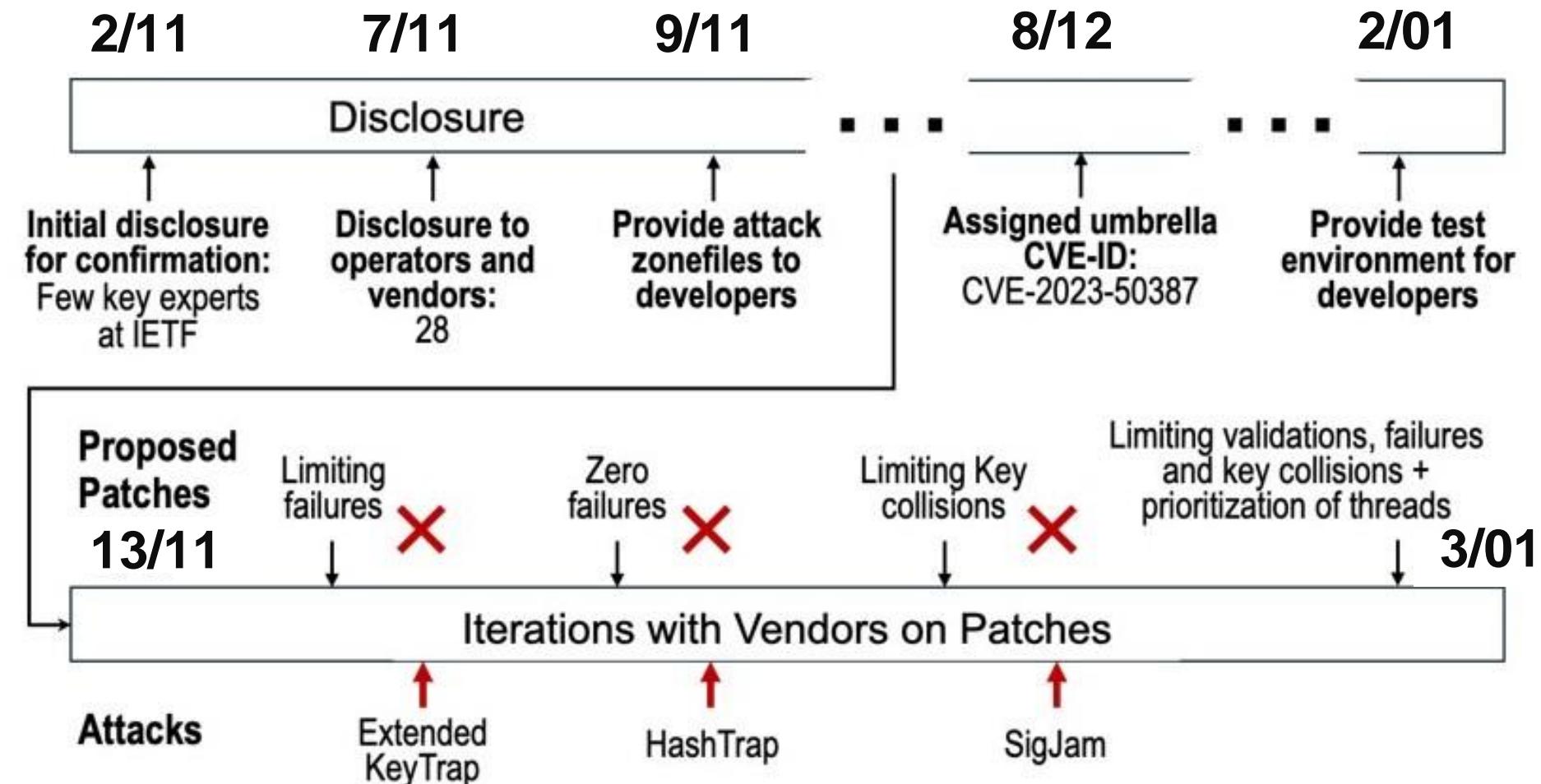


Measurements by APNIC Labs



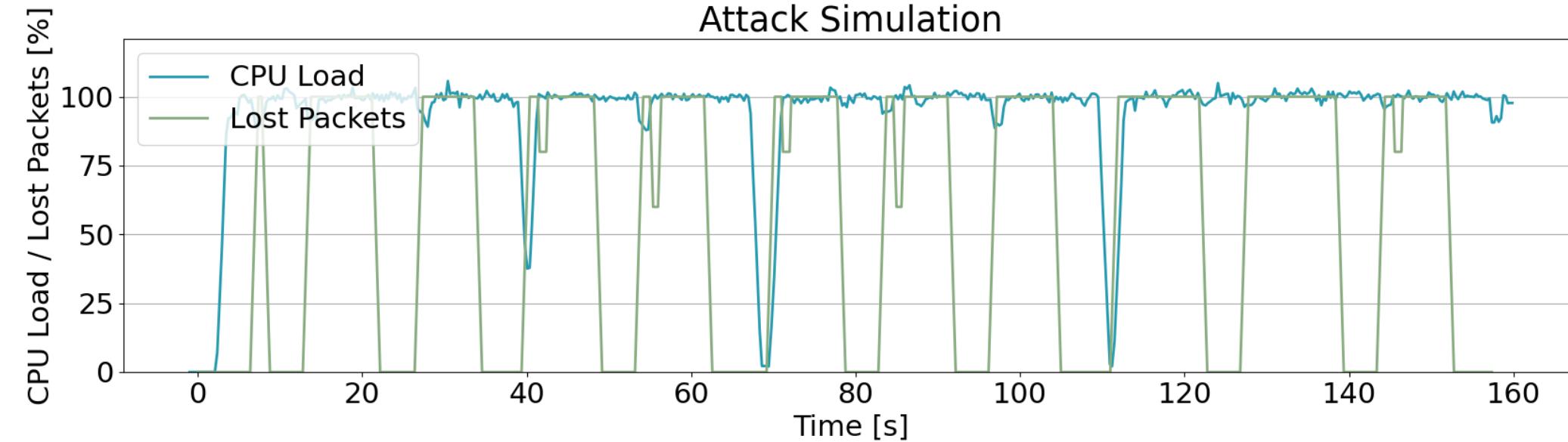
Disclosure: Patch-Break-Fix Process

Disclosure

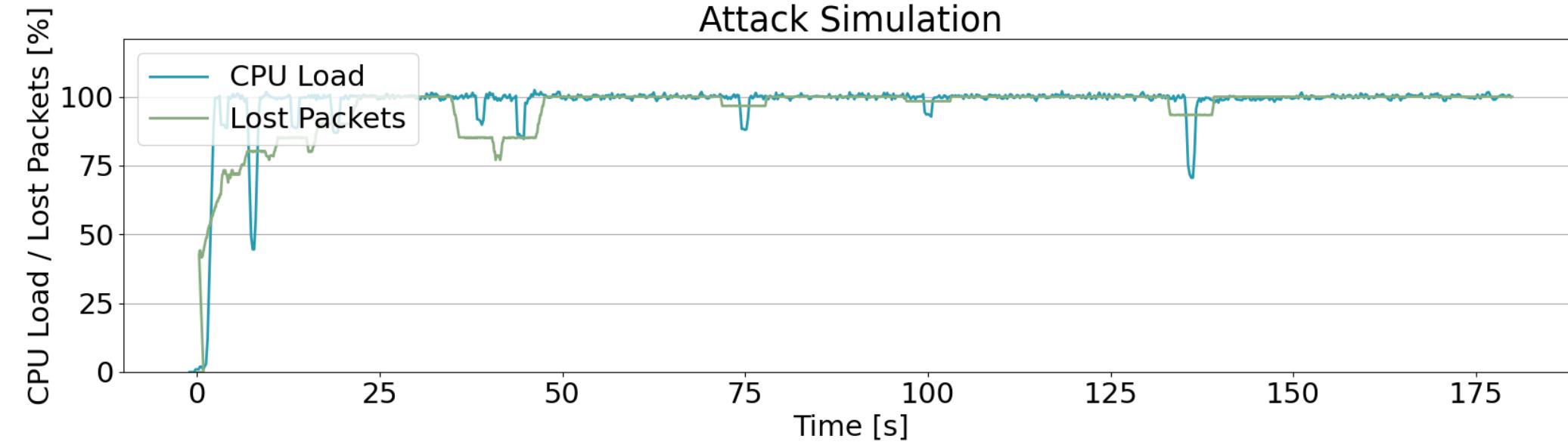


Fixing KeyTrap included > 30 people and took 3 months

(Patched) Resolver under continuous attack



(Patched) Resolver under Hash attack





Lessons Learned

- The KeyTrap Attacks allow for comprehensive DoS on a plethora of DNSSEC-validating resolvers
- Vulnerability stems from "eager validation" in RFCs
Desire to ensure availability compromised availability
- Specification needs to consider impact on resource consumption by implementations



- Read our paper here -
<https://doi.org/10.1145/3658644.3670389>