

Overview

- (a very brief) Introduction to Viper
- Predicates
- Research Question
- Automatic Predicate Inlining
- The Story So Far...
- Questions?

Viper

Verification Infrastructure for Permission-based Reasoning

- architecture for verification tools/prototypes
- intermediate verification language
- common use cases
 - sequential/concurrent programs
 - mutable state
 - separation logic (think Rust)
- provided verifiers
 - Boogie verification condition generation
 - Z3 symbolic execution
- compiles down to Z3

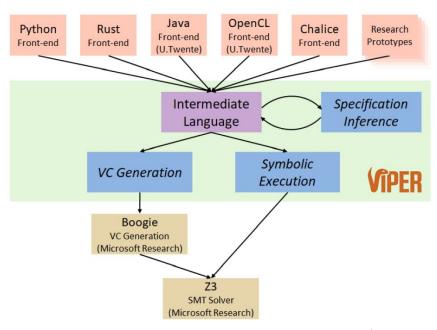


image from Programming Methodology Group | ETH Zurich

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
   x.f := x.f + i
```

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
                                         field declaration
    x.f := x.f + i
```

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
                                       signature
    x.f := x.f + i
```

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
                                       precondition
    x.f := x.f + i
```

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
                                        postcondition
    x.f := x.f + i
```

```
field f: Int
method increment(x: Ref, i: Int)
    requires acc(x.f)
    ensures true
                                      body
    x.f := x.f + i
```

Predicates

Working with Tuples

```
method makeTuple(this: Ref, a: Int, b: Int)
    requires acc(this.left) && acc(this.right)
    this.left := a
    this.right := b
method incrementTuple(this: Ref, i: Int, j: Int)
    requires acc(this.left) && acc(this.right)
    this.left := this.left + i
    this.right := this.right + j
```

Working with Tuples

```
method makeTuple(this: Ref, a: Int, b: Int)
   this.left := a
   this.right := b
                                                   code duplication
method incrementTuple(this: Ref, i: Int, j: Int)
   requires acc(this.left) && acc(this.right) ←
   this.left := this.left + i
   this.right := this.right + j
```

Predicates

```
predicate allowAccess(this: Ref)
   acc(this.left) && acc(this.right)
                       method incrementTuple(this: Ref, i: Int, j: Int)
                           requires allowAccess(this)
                           unfold allowAccess(this)
                           this.left := this.left + i
                           this.right := this.right + j
                           fold allowAccess(this)
```

Recursive Predicates: Lists

```
predicate list(this: Ref) {
  acc(this.elem) & acc(this.next) &
  (this.next \neq null \Longrightarrow list(this.next))
                         method append(this: Ref, e: Int)
                           requires list(this)
                           ensures list(this)
                           unfold list(this);
                           if (this.next == null) {...
                           fold list(this);
```

Research Question

Problem(s)

- A lot of Viper code is automatically generated
 - can lead to redundancies that a human developer would detect and remove
- An unfold is an operation that changes the program state, replacing the predicate resource with the assertions specified by its body. [1]
 - this is a lot of work performed at *runtime*

17

How do inlined predicates impact the performance of Viper?

Automatic Predicate Inlining

Automatic Predicate Inlining

```
inline predicate allowAccess(this: Ref)
{
    acc(this.left) && acc(this.right)
}
```

predicates quantified with the **inline** keyword will be automatically expanded when they're used

```
inline predicate allowAccess(this: Ref)
   acc(this.left) && acc(this.right)
                                                      rewrite methods
method makeTuple(this: Ref, a: Int, b: Int)
   requires allowAccess(this)
   unfold allowAccess(this)
   this.left := a
   this.right := b
   fold allowAccess(this)
                                   method makeTuple(this: Ref, a: Int, b: Int)
                                        requires acc(this.left) && acc(this.right)
                                        this.left := a
                                        this.right := b
```

Dealing with (Mutually) Recursive Predicates

```
predicate list(this: Ref) {
   acc(this.elem) & acc(this.next) &
   (this.next \neq null \Longrightarrow list(this.next))
infinite inline!
·list(this)
acc(this.elem) & acc(this.next) &
(this.next \neq null \implies list(this.next))
acc(this.elem) & acc(this.next) &
(this.next \neq null \Longrightarrow
  acc(this.elem) & acc(this.next) &
  (this.next \neq null \implies list(this.next)))
acc(this.elem) & acc(this.next) &
(this.next \neq null \Longrightarrow
  acc(this.elem) & acc(this.next) &
  (this.next \neq null \Longrightarrow
    acc(this.elem) & acc(this.next) &
```

options:

- 1. trust the programmer is sus
- 2. you may inline... once
- 3. detect cycles & error out

Eventually: Automatic Automatic Predicate Inlining

- Insert the inline keyword wherever possible
 - requires the detection of recursive/mutually-recursive cycles

Validation (Future Goals)

- Industrial Benchmarks
 - Provided by Prof. Alex Summers (SPL)
- Run two different versions of Viper
 - one with inlining
 - one without

The Story So Far...



Milestones

Decide on a language construct/keyword for inlining



Implement PoC that inlines all simple predicates



Recursive cycle detection V



Extend Viper parser with the inline keyword

Extend PoC to inline only predicates marked with inline

Benchmark performance of Viper programs with inlined predicates

Merge with upstream Viper project 5



Questions?