

Exit point of a matrix

Starting pt.

source pt - (0,0)

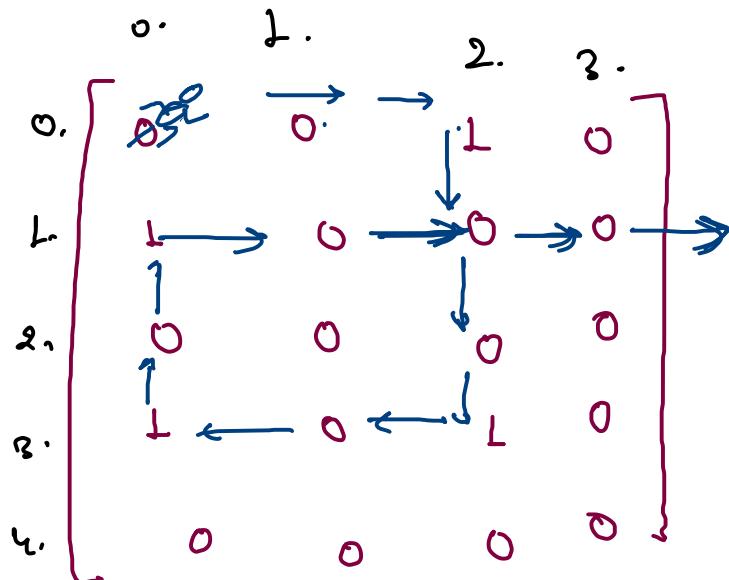
0 - Enterer \rightarrow

Direction controller.

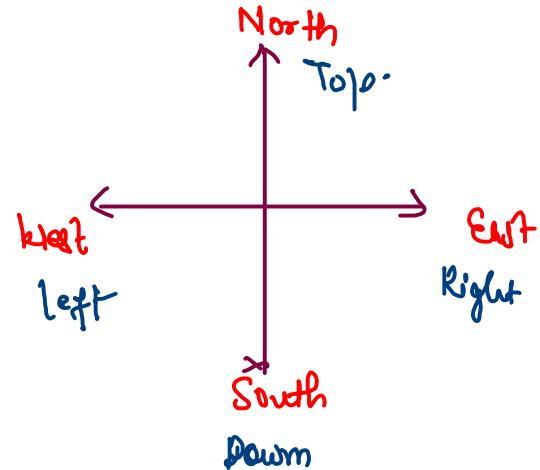
L - Encouter.

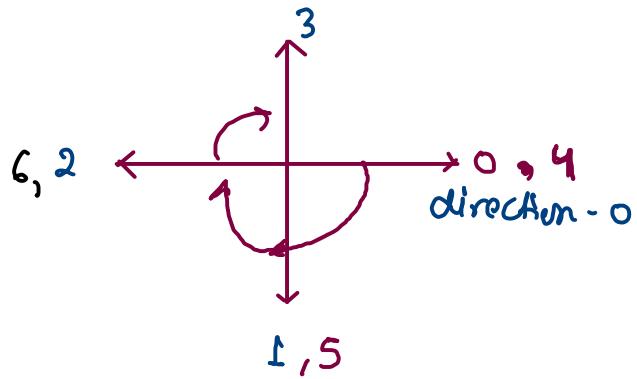
\nwarrow clockwise 90° turn.

\searrow go towards Right



Exit pt \rightarrow (1,3)

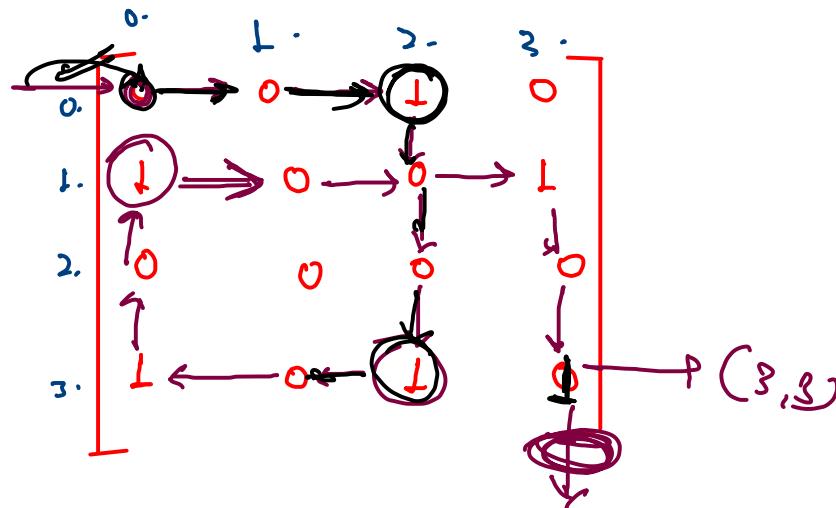




dir = 0

dir = 0 ✗ ✗ ✗ ✗ 5

r = 0 ✗ 3
(2, 0 & 1, 0)
c = 0 ✗ ✗ ✗ 0



direction (Management) d = 0 → Right

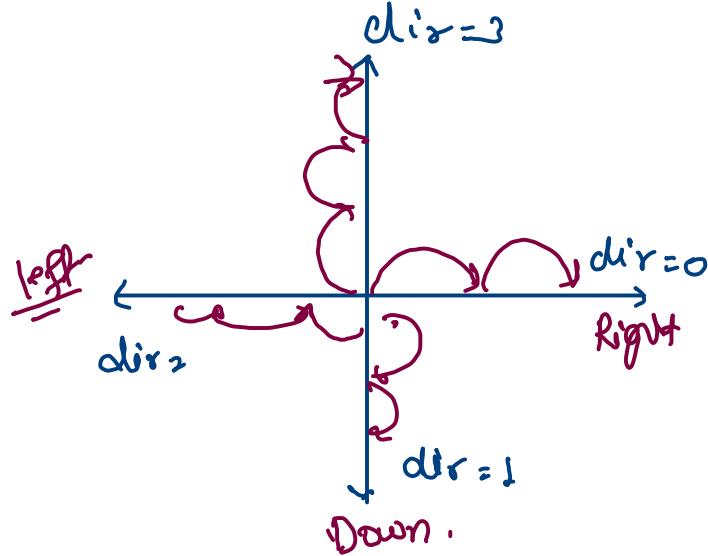
↓ → Down

2 → Left

3 → Top

dir = (dir + arr[i][c])
0, 1, 2, 3

Moves management



- $\text{dir} = 0$
fix \rightarrow Row
vary \rightarrow Column ++
- $\text{dir} = 2$
fix \rightarrow column
vary \rightarrow Row ++
- $\text{dir} = 1$
fix \rightarrow Row.
vary \rightarrow Column --;
- $\text{dir} = 3$
fix \rightarrow Colun
Move \times Row --

$r = \cancel{0} \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5} \cancel{6}$
 $c = \cancel{0} \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5} \cancel{6} \cancel{7} \cancel{8} \cancel{9} \cancel{10}$

$\text{dir} = \cancel{0} \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5} \cancel{6} \cancel{7} \cancel{8} \cancel{9} \cancel{10}$

$r = \cancel{0} \cancel{1} \cancel{2} \cancel{3} \cancel{4} \cancel{5} \cancel{6} -1$

$c = 0$

```

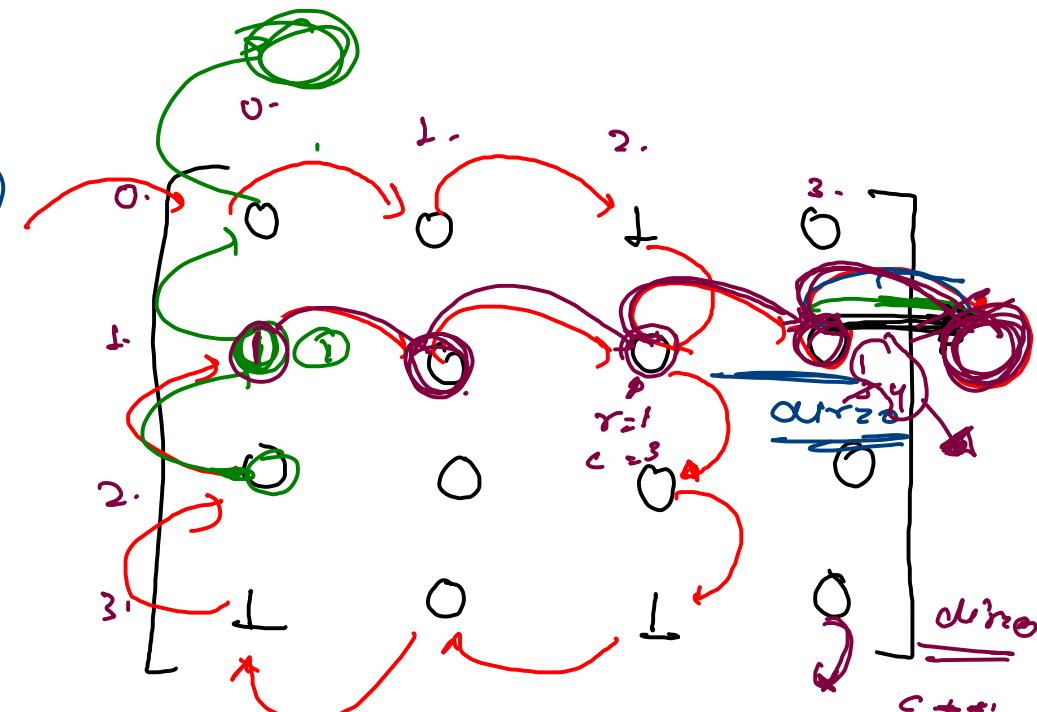
while(r >= 0 && r < arr.length && c >= 0 && c < arr[0].length) {
    // direction management
    dir = (dir + arr[r][c]) % 4;

    // moves management
    if(dir == 0) {
        c++;
    } else if(dir == 1) {
        r++;
    } else if(dir == 2) {
        c--;
    } else {
        r--;
    }
}

System.out.println(r + ", " + c);

```

Handling case of invalids
 $\text{dir} = 3 \rightarrow \text{invalid}$
 $r = -1$



point is invalid because of lost direction

$\text{dir} = 0$
 $\rightarrow c--$

Saddle Point \rightarrow least price
in Row best
max price in
column.

Row \rightarrow min
col \rightarrow max

(1)	12	13	14
21	22	23	24
33	32	31	34
40	42	43	44

* Is there any possibility of
No. saddle pt? - Yes

* How many saddle
point are possible
in a m.

Row min
Col max
Saddle point

① But why??

a	b	c	d
e	f	g	h
i	j	k	l
m	n	o	p

conclusion
 \rightarrow There is no possibility
 of more than 1
 Saddle point

Row \rightarrow Min

Col \rightarrow Max

$$l < j \rightarrow \textcircled{I}$$

$$l > h \rightarrow \textcircled{II}$$

$$\boxed{h < l < j} \xrightarrow{\text{combination of } \textcircled{I} \text{ & } \textcircled{II}}$$

if 'f' is saddle
 then.

$$f > \{ b, j, n \}$$

$$f < \{ e, g, h \}$$

$$f < h \rightarrow \textcircled{III}$$

$$f > j \rightarrow \textcircled{IV}$$

$$\boxed{j < f < h}$$

Combination

$$l > \{ d, h, p \}$$

$$l < \{ i, j, k \}$$

$$j < h$$

$$h < j$$

Hence proved

~~Contradiction~~

a	b	c	d
e	f^*	g	h
i	j	k	l^*
m	n	o	p

if f^* is saddle pt

$$f < \{e, g, h\} \rightarrow f < h \rightarrow \textcircled{I}$$

$$f > \{b, j, n\} \rightarrow f > j \rightarrow \textcircled{II}$$

if l^* is saddle pt -

$$l < \{i, j, k\} \rightarrow l < j \rightarrow \textcircled{III}$$

$$l > \{d, e, p\} \rightarrow l > h \rightarrow \textcircled{IV}$$

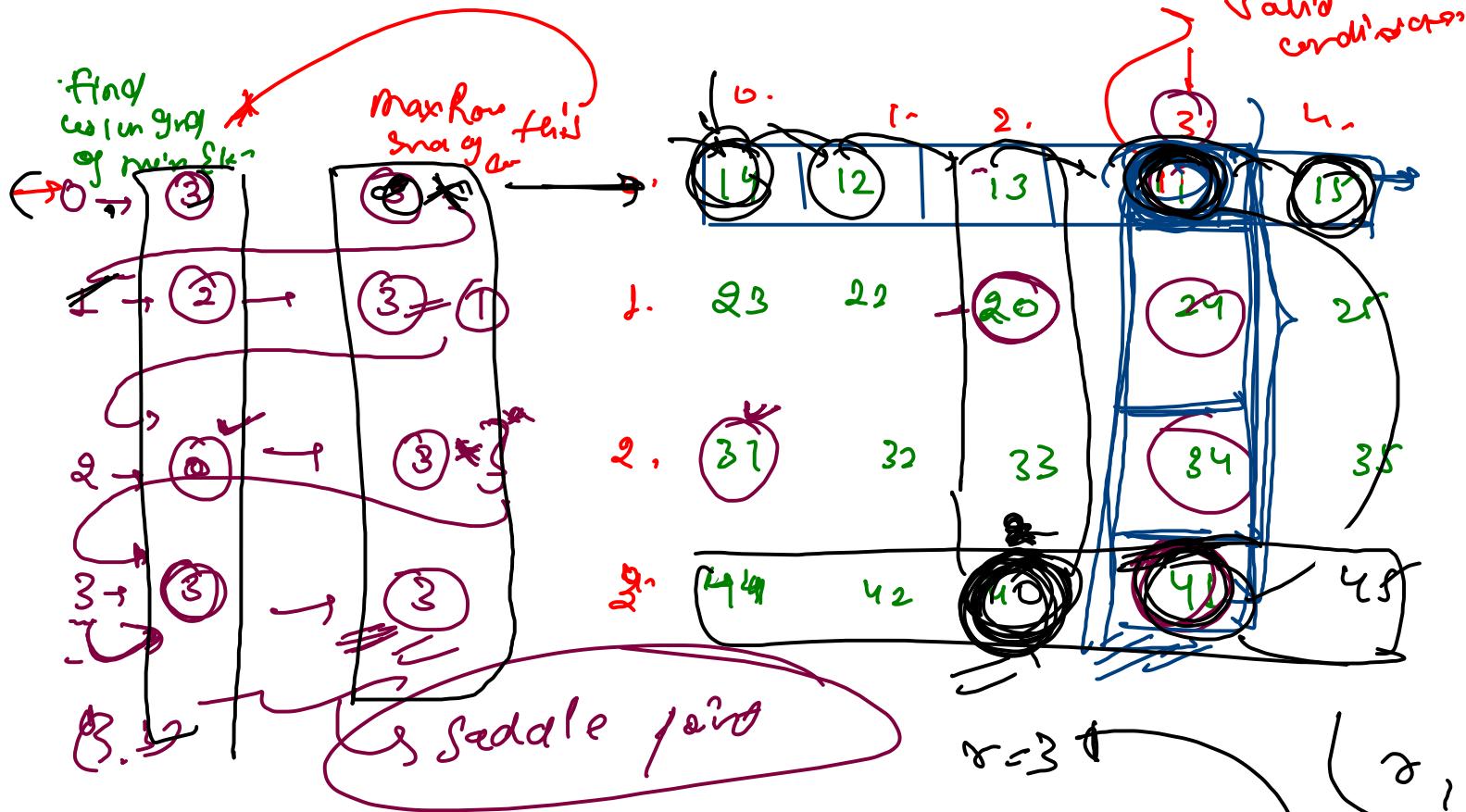
combination of \textcircled{III} & \textcircled{IV}

$$l < j, l > h.$$

$$\frac{j < f < h}{\therefore j < h}$$

using contradiction

$$\frac{h < i}{h < j}$$



$$\begin{array}{c}
 \boxed{\text{MnC} = 2} \\
 \text{Max f} = 3 \\
 (\alpha, \text{mnC}) \rightarrow \text{Saddle point}
 \end{array}$$

Search In Sorted 2D matrix

data ≥ 29

- Row wise sorted
- Columnwise sorted
- Data \rightarrow 27 \rightarrow 29

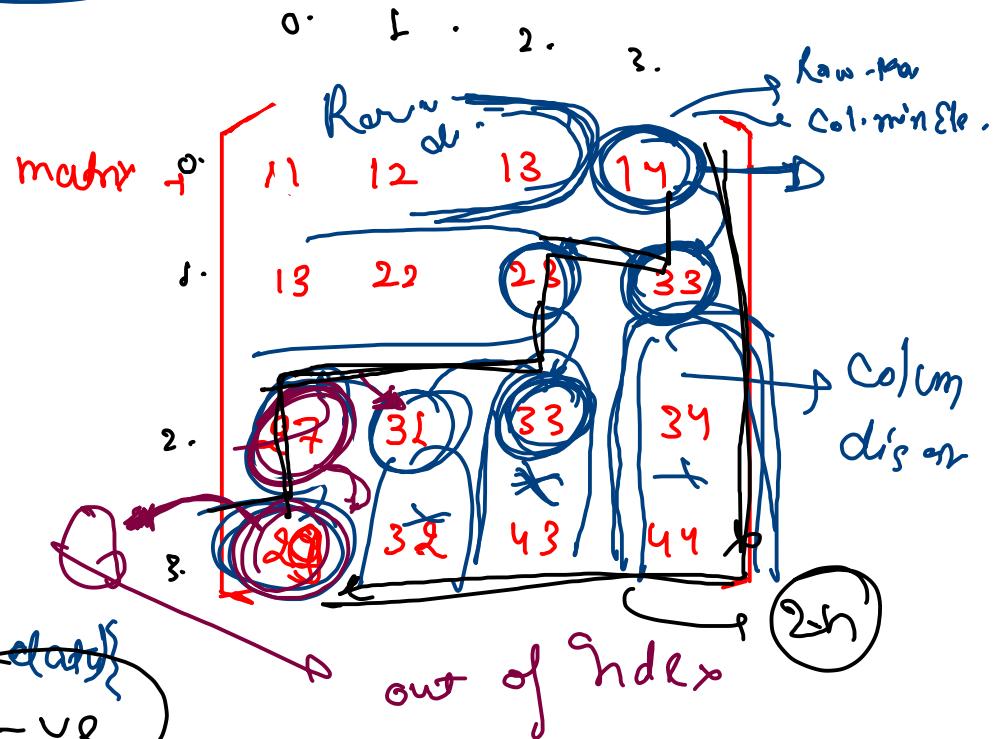
Iteration:

```

if (arr[r][c] == data) {
    print(r, c); return; } const
} else if (arr[r][c] > data) {
    c--;
} else if (arr[r][c] < data) {
    r++;
}

```

* for i.e.



3

Exit Condition from function / function terminations

→ Return keyword —

← all lines are Executed —