

## Agenda (TnS)

- Searching ↗
- Linear Search
- Binary Search

↗  
questions  
+  
partial  
+  
extra.

- ↗ Sorting-
- Bubble
- Selection
- Insertion
- Merge
- Quick
- Count
- Radix,

Complexity  
Analysis

Expectation: →

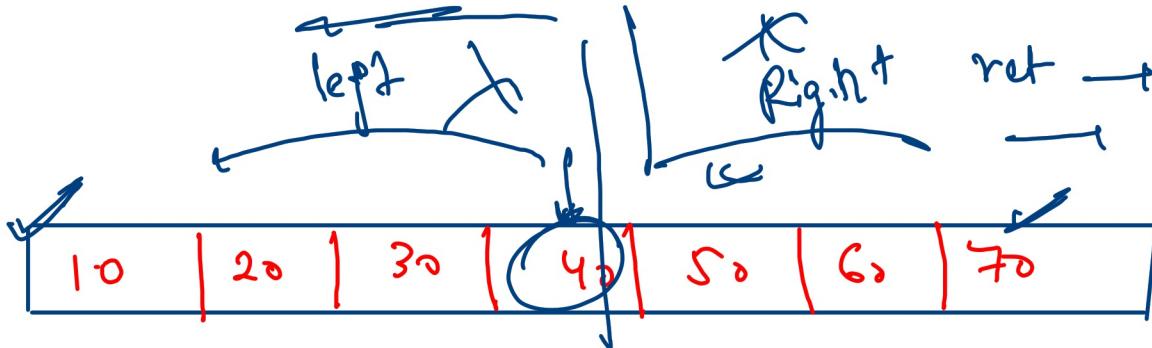
Searching (arr, lo, hi, data) → True  
False.

forth →

searching (arr, lo, mid, data)

searching (arr, mid+1, hi, data)

Merging →

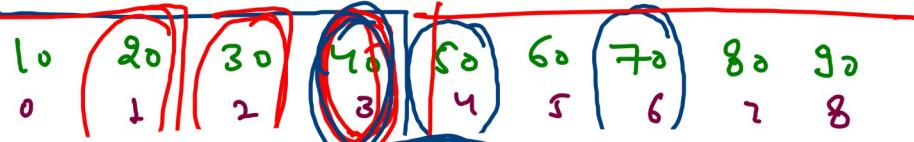


$$\underline{\text{mid}} = \underline{\underline{\text{lo}}} + \frac{\underline{\text{hi}} - \underline{\text{lo}}}{2}$$

data : if arr(mid) == data  
return;

$$\begin{aligned}
 & \Rightarrow \underline{\text{lo}} + \frac{\underline{\text{hi}}}{2} - \underline{\text{lo}} \\
 &= \frac{\underline{\text{lo}} + \underline{\text{hi}}}{2} \\
 &= \frac{\underline{\underline{\text{lo}}} + \underline{\underline{\text{hi}}}}{\underline{\underline{\underline{\text{lo}}} + \underline{\underline{\text{hi}}}}}
 \end{aligned}$$

arr →



data = 80

right  
mid:2

2,3  
right  
mid:1

0,3

left

mid:4

0,8

\* Base Case

lo > hi

return false

lo > hi

mid+1

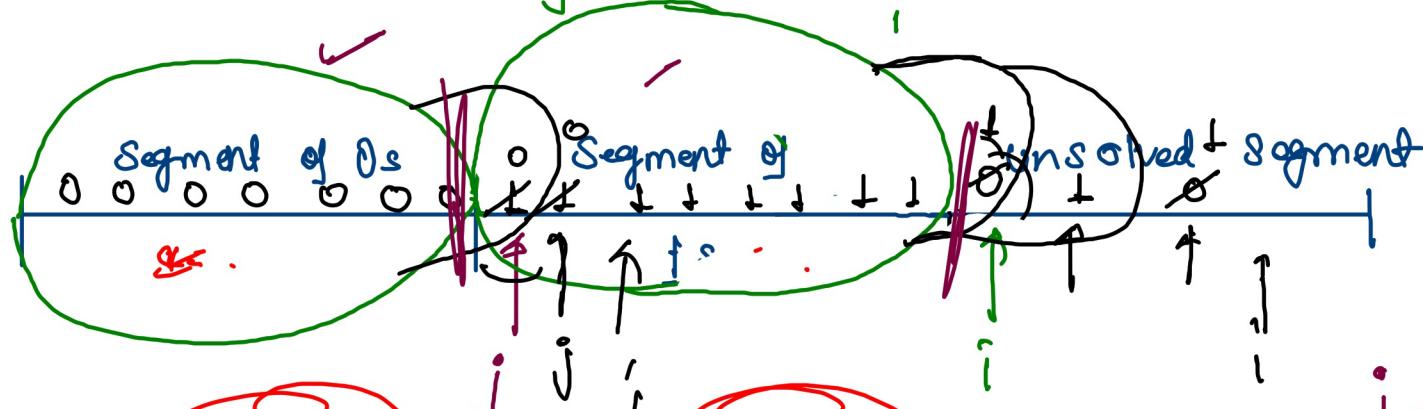
lo ↑

mid-1 ↓

for next  
level

```
public static boolean binarySearchRec(int[] arr, int lo, int hi, int data) {  
  
    int mid = lo + (hi - lo) / 2;  
  
    boolean res = false;  
    if(arr[mid] == data) {  
        res = true;  
    } else if(arr[mid] < data) {  
        // right part  
        res = binarySearchRec(arr, mid + 1, hi, data);  
    } else {  
        // left part  
        res = binarySearchRec(arr, lo, mid - 1, data);  
    }  
    return res;  
}
```

Sort 01



prime	Not prime
odd	even
$\text{arr}[i] = 1$	$\text{arr}[i] = 0$

// Increment in segment of 1.

i++;

swap( $\text{arr}$ ,  $i$ ,  $j$ );

$i++;$

$j++;$

$i \rightarrow$  first unsolved  
 $j \rightarrow$  first 1

Segregate Parity

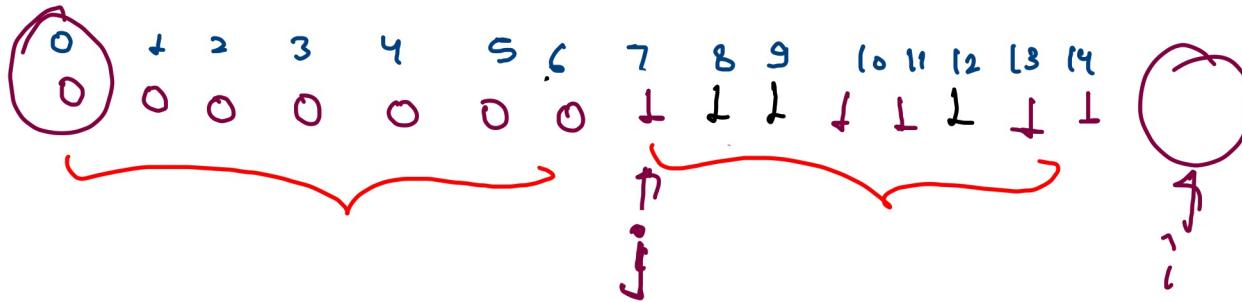
```

while( i < arr.length )
{
    if( arr[i] == 0 )
    {
        swap(arr, i, j)
        i++; j++;
    }
    else
    {
        i++;
    }
}

```

~~[ 000 | 111 | 010 ]~~ →  $i \Rightarrow$  first unresolved  
 $j =$  first 'one'

~~[ 0 1 0 | 1 0 0 | 1 1 0 ]~~  
~~( $i$ )  $\nearrow$  ( $j$ )~~  
~~( $j$ )  $\nearrow$~~   
~~0 0 0 | 1 1 1 1 0~~  
~~( $j$ )  $\nearrow$  ( $i$ )~~  
~~( $i$ )  $\nearrow$  ( $j$ )  $\nearrow$  ( $i$ )~~  
~~0 0 0 0 1 1 1 1 0~~  
~~( $i$ )~~  
[ 0 0 0 0 0 1 1 1 1 1 ]

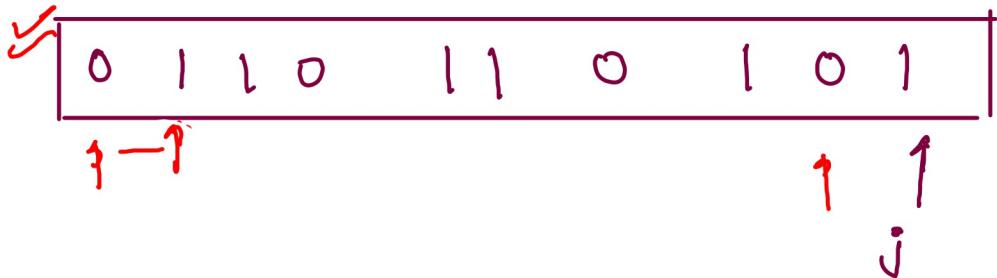


$j = 0$   
 ~~$i = 0$~~  *first Unsolved*

$i \rightarrow$  first Unsolved i

$j \rightarrow$  first 1

~~$arr[i] == 0$~~   $arr[i] == 1$   
 $i++$        $swap(i, j)$   
 $,$              $j--;$



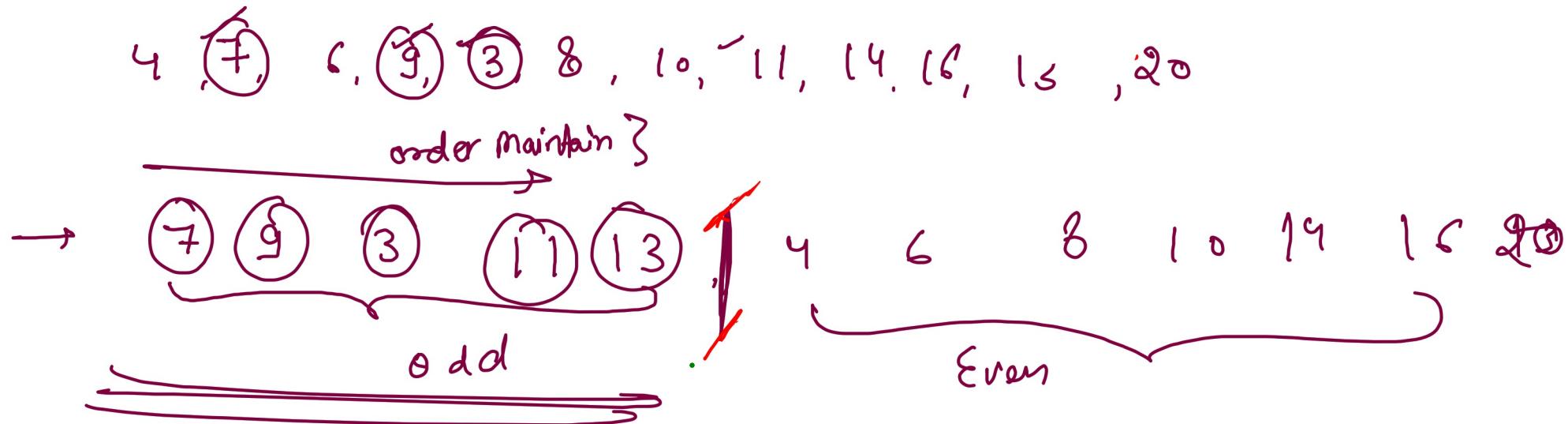
`if (arr[i] == 1) {  
 i++; // Increment in  
 segment of 1`

`} else {  
 swap(arr, i, j);  
 i++;  
 j++;`

$i \rightarrow$  first unsolved  
 $j = \{ \text{of} \}$  unsolved

Segregate Odd Even → Odd - Even but we can't maintain the order of even element.

order of odd → closest n'th change sequence of odd Element does not change



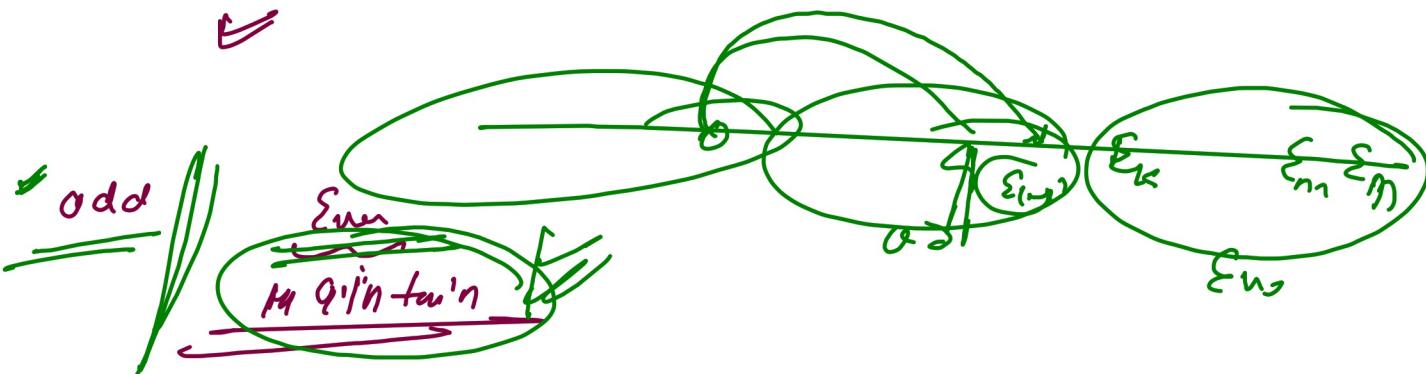
$\epsilon_1 \quad 0, \quad \epsilon_2 \quad \epsilon_3 \quad 0_2 \quad 0_3 \quad \epsilon_4 \quad 0_{4-}$   
 ↓      ↓      14      8      ↓      3      18      ↓  
 10      5      7

Sort → ~~odd + Even~~

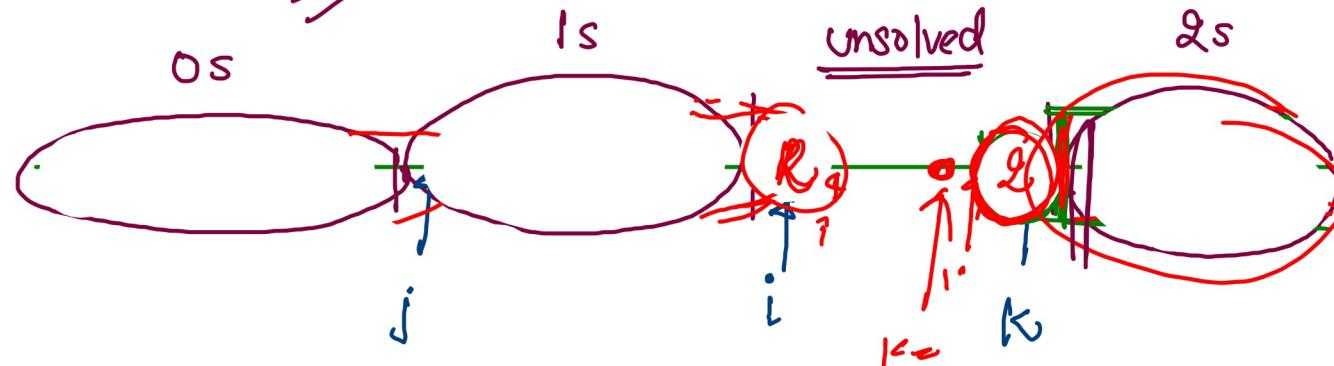
even      may      be      arrange      in  
 random      order

{ 1      5      3      7 }      { 10      14      8      18 }

order Maintain



## Sort 0 1 2 (Dutch Flag)



Algorithm



$i \Rightarrow$  first unsolved

$j \Rightarrow$  first 1

$k \Rightarrow$  last unsolved

} decision factor  $\rightarrow i$

Quick Sort

$arr[i] := 1$

$arr[i] := 0$

$\Rightarrow arr[i] := 2$

$i++;$

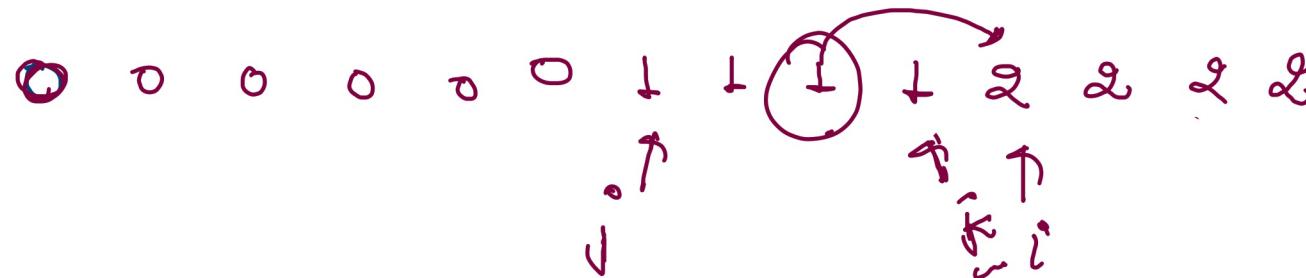
$swap(i, j);$

$swap(i, k); \rightarrow$  ~~increment~~  
~~count[i] = > 2~~

// Incremented the  
Region 1

$i++;$   
 $j++;$

$k--;$



$i = 0$

$j = 0$

$k = \text{arr.length} - 1$

$k > i$

→ Stopping.

$i \leq k$

Both are Unuseful

$\text{arr}[i] == 0$

$\text{swap}(i, j);$

$i++;$

$j++;$

$\text{arr}[i] == 1$

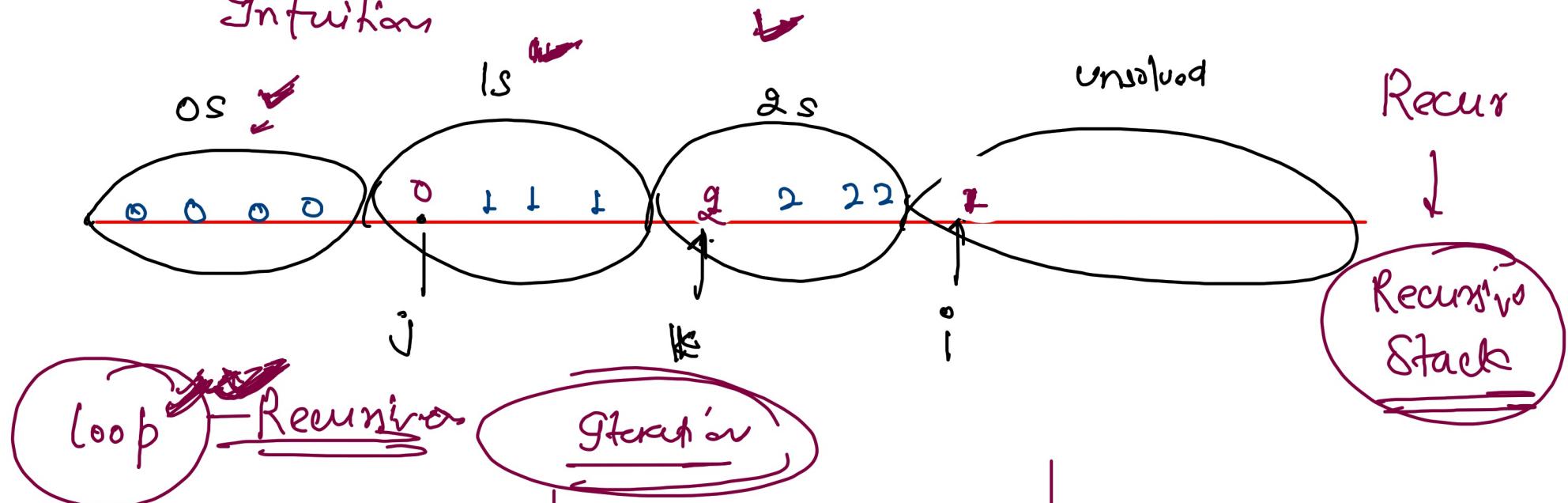
$i++;$

$\text{arr}[i] == 2$ .

$\text{swap}(i, k);$

$k--;$

## Intuitions



$$\text{arr}[i] = 0$$

~~swap(i, j)~~

~~swap(i, k)~~

$i++;$   
 $j++;$   
 $k \rightarrow r,$

$$\checkmark \text{arr}[i] = 1$$

~~swap(i, k);~~

$i++;$

~~| C F |~~

$$\text{arr}[i] = 2.$$

$i++;$

## Solve Polynomial

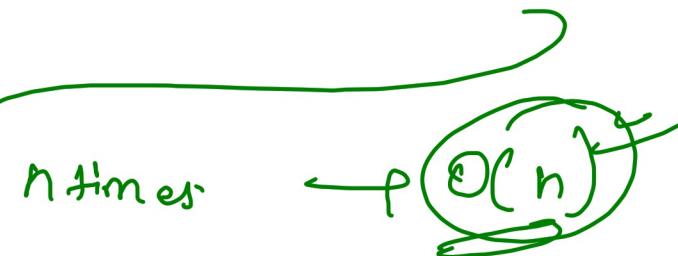
$$f(x, n) \rightarrow 1 \cdot x^n + 2 \cdot x^{n-1} + 3 \cdot x^{n-2} + \dots + n \cdot x^1$$

$$f(x, n) \rightarrow \cancel{n \cdot x^1} + (n-1) \cdot x^2 + \cancel{(n-2) \cdot x^3} + \dots + \cancel{1 \cdot x^n}$$



$O(1)$

$n$  times



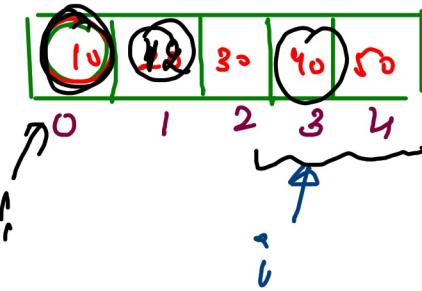
$n \log n$

## Merge Two Sorted Arrays

if (arr1[i] > arr2[j]) {  
 arr1↑

    res[k] = arr2[j];

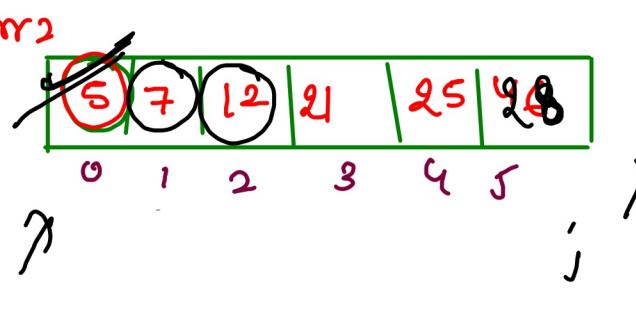
    j++;



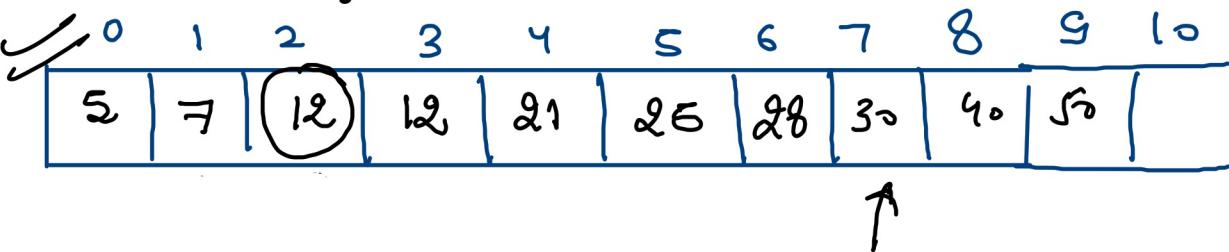
} else {

    res[k] = arr1[i];  
    arr1↑

    k++



res →



Largest subarray having eq no. of 0's and 1's

```
int findSubArray(int arr[], int n)
{
    int sum = 0;
    int maxsize = -1, startindex = 0;
    int endindex = 0;

    // Pick a starting point as i
    for (int i = 0; i < n - 1; i++) {
        sum = (arr[i] == 0) ? -1 : 1;

        // Consider all subarrays starting from i
        for (int j = i + 1; j < n; j++) {
            if (arr[j] == 0)
                sum += -1;
            else
                sum += 1;

            // If this is a 0 sum subarray, then
            // compare it with maximum size subarray
            // calculated so far

            if (sum == 0 && maxsize < j - i + 1) {
                maxsize = j - i + 1;
                startindex = i;
            }
        }
    }
}
```

(maxsize < i+1)

$$2 < 3-0+1=4 \checkmark$$

$$4 < 5-0+1=6$$

$$6 < 8$$



$$\text{maxsize} = 1 \cancel{2} 4 \cancel{6} \underline{8} \rightarrow$$

$$i = 0 \leq 7$$

~~start~~

$$st = 0$$

$$j = 1 \leq 9$$



$$(i)(j) \cancel{x} \cancel{x} \cancel{x}(j)$$

$$\text{sum} \Rightarrow 1 \cancel{2} \cancel{x} \cancel{1} \cancel{1} \quad i = 1 \text{ no eq of sub}$$

$$\text{sum} = -1 \cancel{x} \cancel{x} \phi -1$$

$$\text{max} = 8 < (2, 4) \times$$

```
int largestSubarray( int arr) {
```

Complexity - O(n)

```
    int n = arr.length;
```

```
    int sum=0, maxsize = -1, stindx = -1;
```

```
    for( int i=0; i<n-1; i++) {
```

```
        sum = (arr[i]==0) ? -1 : 1; // initialize
```

```
        for( int j=i+1; j<n; j++) {
```

```
            if (arr[j]==0) sum = sum -1;
```

```
            else sum = sum + 1;
```

```
            if( sum==0 && maxsize < j-i+1) {
```

```
                maxsize = j-i+1;
```

```
            } stindx = i;
```

endIndex = stindx + maxsize - 1;

```
}
```

~

## Method - 2 use hash Map

rlen = -1 2 4 5 6  
cd = 4 4 5 6

-1	0
0	1
1	2

$$4 < 4 - 0 \times$$

$$4 < 5 - 0 \checkmark$$

$$5 - 0 = 5$$

$$5 < 6 - 0$$

```
// Initialize sum of elements
int sum = 0;

// Initialize result
int max_len = 0;
int ending_index = -1;
int start_index = 0;

for (int i = 0; i < n; i++) {
    arr[i] = (arr[i] == 0) ? -1 : 1;
}

// Traverse through the given array
for (int i = 0; i < n; i++) {
    // Add current element to sum

    sum += arr[i];

    // To handle sum=0 at last index

    if (sum == 0) {
        max_len = i + 1;
        ending_index = i;
    }

    // If this sum is seen before,
    // then update max_len if required
    if (hM.containsKey(sum)) {
        if (max_len < i - hM.get(sum)) {
            max_len = i - hM.get(sum);
            ending_index = i;
        }
    }
    else // Else put this sum in hash table
        hM.put(sum, i);
}

for (int i = 0; i < n; i++) {
    arr[i] = (arr[i] == -1) ? 0 : 1;
}

int end = ending_index - max_len + 1;
↓ starting index = 6 - 6 + 1 = 1
```

0	1	2	3	4	5	6
0	1	1	0	0	1	0

-1	1	1	-1	-1	1	-1
----	---	---	----	----	---	----

original arr change

$$\text{sum} = -1 0 1 0 -1 -2 -1$$

$$i = \emptyset 1 2 3 4 5 6$$

1	1	0	0	1	0
---	---	---	---	---	---

$$\text{len} = 6$$

$$1 \rightarrow 3 \quad 0 \rightarrow 3$$

$$\downarrow \text{starting index} = 6 - 6 + 1 = 1$$

```
int maxlen = -1, st = 0, end = 0;  
HashMap<I,I> map = new HashMap<>();
```

Step 1 : change the original array

```
for (int i=0; i < n; i++)  
    arr[i] = (arr[i] == 0) ? -1 : 1;
```

Step 2 :

```
for (int i=0; i < n; i++)  
{  
    sum = sum + arr[i];  
  
    if (sum == 0)  
    {  
        maxlen = i+1;  
    }  
    end = i;
```

```
    }  
    if (map.containsKey(sum))  
    {  
        if (maxlen < i - map.get(sum))  
        {  
            maxlen = i - map.get(sum);  
            end = i;  
        }  
    }  
    else  
    {  
        map.put(sum, i);  
    }  
    st = end - maxlen + 1;
```