

Lecture Five

1) About handle

$$G: S \Rightarrow aAcBe$$

$$A \rightarrow b | Ab$$

$$B \rightarrow d$$

for a right-sentential form

$$abbede$$

← the last right most derivation.

(1) b is a handle, although we have $B \rightarrow d$, d is not a handle

(2) reduction 1.

$$\underline{a}bbede$$

↑ handle

(3) reduction 2

$$aAcde$$

now, d is a handle.

$$\underline{a}Acde$$

↑ handle

(4) reduction 3

$$\underline{aAcBe}$$

↑ handle.

(5) reduction 4.

$$S \quad \square$$

the leftmost reductions are

$$abbede, aAbcde, aAcde, aAcBe, S$$

$\xleftarrow{\text{reduction } A \rightarrow b}$
 $\xleftarrow{\text{reduction } A \rightarrow Ab}$
 $\xleftarrow{\text{reduction } B \rightarrow d}$
 $\xleftarrow{\text{reduction } S \rightarrow aAcBe}$

2. Viable Prefixes

For example, suppose $E \xRightarrow{rm} F * id \xRightarrow{rm} (E) * id$

at various times during the parse, the stack will hold $($, $(E$, and (E) , but it must not hold $(E)*$, since (E) is a handle, which the parser must reduce to F before shifting $*$.

$$A \rightarrow \beta_1 \beta_2 \quad \textcircled{1} \text{ if } \beta_2 \neq \epsilon \text{ shift}$$

$\textcircled{2}$ if $\beta_2 = \epsilon$ if $A \rightarrow \beta_1$ is the handle, we should reduce by this production.

two valid items may tell us to do different things for the same viable prefix.

3. Canonical LR(1) $S \Rightarrow BB$
 $B \Rightarrow aB \mid b$

(1) rightmost derivation $S \xRightarrow{*}_{rm} aaBab \Rightarrow aaaSab$

$[B \rightarrow a.B, a]$ is valid for viable prefix $\gamma = aaa$

Let $aaa.Bab$
 $\begin{array}{ccc} \overline{\uparrow} & \overline{\uparrow} & \overline{\uparrow} \\ \delta & A & W \end{array}$

(2) rightmost derivation $S \xRightarrow{*}_{rm} BaB \Rightarrow BaaB$

$[B \rightarrow a.B, \$]$ is valid for viable prefix Baa .

$Baa.B\$$
 $\begin{array}{ccc} \overline{\uparrow} & \overline{\uparrow} & \overline{\uparrow} \\ \delta & A & W \end{array}$

4. $I_0: S \rightarrow \cdot S, \$$
 $S \rightarrow \cdot CC, \$$
 $C \rightarrow \cdot cC, c/d$
 $C \rightarrow \cdot d, c/d$

$I_1: S' \rightarrow S \cdot, \$$

$I_2: S \rightarrow C \cdot C, \$$
 $C \rightarrow \cdot cC, \$$
 $C \rightarrow \cdot d, \$$

$I_3: C \rightarrow c \cdot C, dd$
 $C \rightarrow \cdot cC, dd$
 $C \rightarrow \cdot d, c/d$

$I_4: C \rightarrow d \cdot, c/d$
 $I_5: S \rightarrow CC \cdot, \$$

$I_6: C \rightarrow c \cdot C, \$$
 $C \rightarrow \cdot cC, \$$
 $C \rightarrow \cdot d, \$$

$I_7: C \rightarrow d \cdot, \$$

$I_8: C \rightarrow cC \cdot, c/d$

$I_9: C \rightarrow cC \cdot, \$$

