

Principles of Distributed Systems

inft-3507

Dr. J.Burns

ADA University

Autumn 2025

Section 1: Introduction

This content is based on the following public resources: <https://www.distributed-systems.net/index.php/books/ds4/>

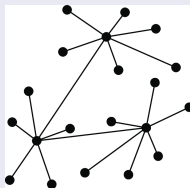
Distributed versus Decentralized systems

Distributed versus Decentralized

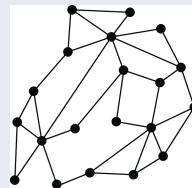
What many people state



Centralized



Decentralized



Distributed

When does a decentralized system become distributed?

- Adding 1 link between two nodes in a decentralized system?
- Adding 2 links between two other nodes?
- In general: adding $k > 0$ links....?

Alternative approach

Distributed system

- A distributed system consists of multiple computers (nodes) that work together to appear as a single system to the user.
- **Focus:** performance, scalability, fault tolerance
- **Control** typically/commonly centralized

Decentralized system

- A decentralized system distributes control and decision-making authority across multiple independent nodes rather than relying on a central authority.
- **Focus:** autonomy, resilience to control, trust minimization
- **Control:** No single node has ultimate authority

The Key Difference: Control & Authority

Distributed

Multiple nodes *compute together*, but control can still be centralized:

- Leader / coordinator often exists
- Decisions may flow from one organization or service

Decentralized

Multiple nodes *share control*:

- No single node (or org) has ultimate authority
- Governance/decision-making spread across participants

Key takeaway

A system can be **distributed yet centralized**; a **decentralized** system must distribute *control*.

Architecture & Coordination

Typical distributed coordination

- Client–server, master–worker, leader-based replication
- Specialized roles (leader, replicas, workers)

Typical decentralized coordination

- Peer-to-peer style participation
- Collective coordination (e.g., voting/consensus-like mechanisms)

Intuition

Distributed → “many machines, one system”

Decentralized → “many owners, shared authority”

Fault Tolerance & Failure Modes

Distributed systems

- Handle **node failures** (crashes, partitions)
- But may have **single point of control**

Decentralized systems

- Handle node failures *and* control failures
- Aim to avoid **single authority failure**

Practical distinction

Distributed: resilience to **hardware/service failure**

Decentralized: resilience to **control/censorship/authority failure**

Trust & Security Model

Distributed (often)

- Assumes trusted administrators / central policy
- Security enforced via access control and ops processes

Decentralized (often)

- Assumes some participants may be malicious
- Uses mechanisms to reduce reliance on trust in any one node

Rule of thumb

More decentralization usually means a **stronger adversarial model** (and more design complexity).

Performance & Efficiency Trade-offs

Distributed systems tend to optimize

- Latency, throughput, operational efficiency
- Coordination cost kept low via leaders / centralized control

Decentralized systems tend to trade performance for

- Independence, robustness, shared governance
- Higher coordination overhead (more parties must agree)

Trade-off

Decentralization often increases cost of coordination \Rightarrow lower peak performance.

Summary Comparison

Quick comparison

- **Goal:** Distributed → scalability/performance vs Decentralized → autonomy/resilience
- **Control:** Distributed → may be centralized vs Decentralized → shared control
- **Trust:** Distributed → trusted operators vs Decentralized → trust-minimized
- **Complexity:** Distributed → moderate vs Decentralized → higher

Design goals

What do we want to achieve?

Overall design goals

- Support sharing of resources
- Distribution transparency
- Openness
- Scalability

Sharing resources

Canonical examples

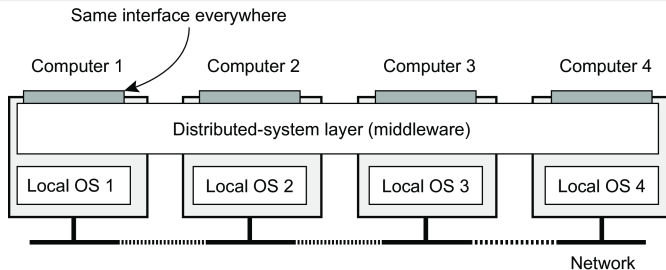
- Cloud-based shared storage and files
- Peer-to-peer assisted multimedia streaming
- Shared mail services (think of outsourced mail systems)
- Shared Web hosting (think of content distribution networks)

Observation

"The network is the computer"

(John Gage, Sun Microsystems)

Distribution transparency



What is transparency?

*The phenomenon by which a distributed system attempts to **hide** the fact that its processes and resources are **physically distributed across multiple computers**, possibly **separated by large distances**.*

Observation

Distribution transparency is handled through many different techniques in a layer between applications and operating systems: a **middleware layer**

Distribution transparency

Types

Transparency	Description
Access	Hide differences in data representation and how an object is accessed
Location	Hide where an object is located
Migration	Hide that an object may move to another location
Replication	Hide that an object is replicated
Concurrency	Hide that an object may be shared by several independent users
Failure	Hide the failure and recovery of an object

Openness of distributed systems

Open distributed system

A system that *offers components* that can easily be used by, or *integrated into other systems*. An open distributed system itself will often consist of components that originate from elsewhere.

What are we talking about?

Be able to interact with services from other open systems, irrespective of the underlying environment:

- Systems should conform to well-defined *interfaces*
- Systems should easily *interoperate*
- This means publicly documented protocols and APIs
- Vendor-neutral standards
- Interoperability and extensibility

Dependability

Basics

A **component** provides **services** to **clients**. To provide services, the component may require the services from other components \Rightarrow a component may **depend** on some other component.

Specifically

A component C depends on C^* if the **correctness** of C 's behavior depends on the correctness of C^* 's behavior. (Components are processes or channels.)

Dependability

Requirements related to dependability

Requirement	Description
Availability	Readiness for usage
Reliability	Continuity of service delivery
Safety	Very low probability of catastrophes
Maintainability	How easy can a failed system be repaired

Traditional reliability - MTTF/MTTR

The following *traditional* metrics were derived from shop-floor machine reliability modelling over many years.

Traditional metrics

- **Mean Time To Failure** (*MTTF*): The average time until a component fails.
- **Mean Time To Repair** (*MTTR*): The average time needed to repair a component.
- **Mean Time Between Failures** (*MTBF*): Simply $MTTF + MTTR$.

You will often see *MTTF*, *MTTR*, *MTBF* used for modelling system reliability over time.

Reliability v Availability

- *Availability* is about how much of the time the system is usable, allowing multiple failures and repairs.
- *Availability* is measured in *Uptime percentage* or *SLA "nines"* (eg, 5 9s)

$$A = \lim_{T \rightarrow \infty} \frac{\text{time system is up in } [0, T]}{T}$$

- *Reliability*: time to first failure
- *Reliability* is about whether failure occurs at all during an interval.

$$R(t) = \Pr\{\text{system survives without failure for time } t\}$$

- *Reliability* is measured Failure rate λ or MTTF

Availability v Reliability: Numerical Example

System 1

- Fails every 10 minutes
- Repair time = 1 second
- *Availability* $\approx 99.83\%$
- *Reliability over 1 hour* \approx almost zero

System 2

- Fails every 6 months
- Repair time = 6 hours
- *Availability* $\approx 99.83\%$
- *Reliability over 1 hour* \approx almost 1

Conclusion: Two systems, with the same *availability*, but completely different *reliability*.

A Note on Security

Observation

A distributed system that is not secure, is not dependable

What we need

- **Confidentiality**: information is disclosed only to authorized parties
- **Integrity**: Ensure that alterations to assets of a system can be made only in an authorized way

Authorization, Authentication, Trust

- **Authentication**: verifying the correctness of a claimed identity
- **Authorization**: does an identified entity has proper access rights?
- **Trust**: one entity can be assured that another will perform particular actions according to a specific expectation

Security mechanisms

Keeping it simple

It's all about **encrypting** and **decrypting** data using **security keys**.

Notation

$K(data)$ denotes that we **use key K** to **encrypt/decrypt $data$** .

Security mechanisms

Symmetric cryptosystem

With **encryption key** $E_K(data)$ and **decryption key** $D_K(data)$:

if $data = D_K(E_K(data))$ then $D_K = E_K$. Note: encryption and decryption key are the same and should be kept **secret**.

Asymmetric cryptosystem

Distinguish a **public key** $PK(data)$ and a **private (secret) key** $SK(data)$.

- Encrypt message from *Alice* to *Bob*: $data = \underbrace{SK_{bob}(\overbrace{PK_{bob}(data)}^{\text{Sent by Alice}}))}_{\text{Action by Bob}}$
- Sign message for *Bob* by *Alice*: $[data, \underbrace{data \stackrel{?}{=} PK_{alice}(SK_{alice}(data))}_{\text{Check by Bob}}] = [data, \underbrace{SK_{alice}(data)}_{\text{Sent by Alice}}]$

Security mechanisms

Secure hashing

In practice, we use **secure hash functions**: $H(data)$ returns a **fixed-length string**.

- Any change from $data$ to $data^*$ will lead to a **completely different string** $H(data^*)$.
- Given a hash value, it is computationally impossible to find a $data$ with $h = H(data)$

Practical digital signatures

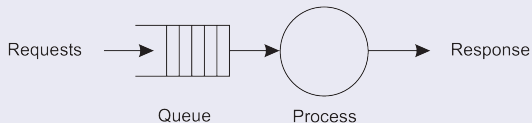
Sign message for *Bob* by *Alice*:

$$[data, \underbrace{H(data) \stackrel{?}{=} PK_{alice}(sgn)}_{\text{Check by Bob}}] = [data, H, \underbrace{sgn = SK_{alice}(H(data))}_{\text{Sent by Alice}}]$$

Performance Models

Performance Models

A centralized service can be modeled as a simple queuing system



Assumptions and notations

- The queue has infinite capacity \Rightarrow arrival rate of requests is not influenced by current queue length or what is being processed.
- Arrival rate requests: λ
- Processing capacity service: μ requests per second

Performance models

Utilization U of a service is the fraction of time that it is busy

$$U = \frac{\lambda}{\mu}$$

Average number of requests in the system

$$\bar{N} = \frac{U}{1 - U}$$

Average throughput

$$X = \frac{\lambda}{\mu} \cdot \mu = \lambda$$

via the Principle of Equilibrium (or conservation)

Performance models

Response time: total time take to process a request after submission

From Little's Law:

$$\begin{aligned}\bar{N} &= X \cdot R \Rightarrow R = \frac{\bar{N}}{X} \\ \Rightarrow R &= \frac{1}{\mu \cdot (1 - U)}\end{aligned}$$

Observations

- If U is small, response-to-service time is close to 1: a request is immediately processed
- If U goes up to 1, the system comes to a grinding halt.
Solution: increase μ .

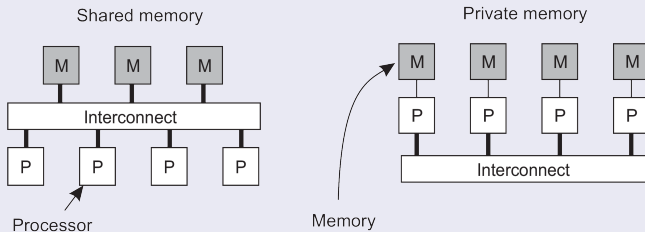
A simple classification of distributed systems

Parallel computing

Observation

High-performance distributed computing started with parallel computing

Multiprocessor and multicore versus multicomputer



Distributed shared memory systems

Observation

Multiprocessors are relatively easy to program in comparison to multicomputers, yet have problems when increasing the number of processors (or cores). **Solution:** Try to implement a **shared-memory model** on top of a multicomputer.

Example through virtual-memory techniques

Map all main-memory pages (from different processors) into one **single virtual address space**. If a process at processor *A* addresses a page *P* located at processor *B*, the OS at *A* **traps and fetches** *P* from *B*, just as it would if *P* had been located on local disk.

Problem

Performance of distributed shared memory could never compete with that of multiprocessors, and failed to meet the expectations of programmers. It has been widely abandoned by now.

Pitfalls

Developing distributed systems: Pitfalls

Observation

Many distributed systems are needlessly complex, caused by mistakes that required patching later on. Many **false assumptions** are often made.

False (and often hidden) assumptions

- The network is reliable
- The network is secure
- The network is homogeneous
- The topology does not change
- Latency is zero
- Bandwidth is infinite
- Transport cost is zero
- There is one administrator

Summary

Summary and Conclusions

We have discussed some important principles in Distributed Systems, namely:

- Distributed and Decentralized Systems
- Resource Sharing
- Distribution Transparency
- Openness and Security
- Performance Models
- Classification of Distributed Systems
- Pitfalls of Distributed Systems Development