

CS3245

# Information Retrieval

Lecture 12: Web Search

# 12

# Last Time

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## Chapter 11

1. Probabilistic Approach to Retrieval / Basic Probability Theory
2. Probability Ranking Principle
3. OKAPI BM25

## Chapter 12

1. Language Models for IR

# Today

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## Chapter 20

- Crawling

## Chapter 21

- Anchor Text
- PageRank



# CRAWLING

# Copyright violations



## Chilling Effects

Home | Weather Reports | Report Receiving  
Sending a C&D Notice | Search the Database | Topics

Topic Home | FAQs

Monitoring the legal climate for Internet activity

[Chilling Effects Clearinghouse](#) > [Piracy or Copyright Infringement](#) > Frequently Asked Questions

## Frequently Asked Questions (and Answers) about Piracy or Copyright Infringement

- [Q: What is the purpose of copyright law?](#)
- [Q: If I am accused of "piracy," what does this mean?](#)
- [Q: Is all copying piracy?](#)
- [Q: Why is "piracy" such a big issue now?](#)
- [Q: My website contains a disclaimer that clearly states that I do not support or promote copyright infringement. Will this protect me?](#)
- [Q: Why are copyright holders concerned about piracy?](#)
- [Q: What are the penalties for copyright infringement, such as making infringing copies of software?](#)
- [Q: I run a website but I never actually upload or download copyrighted materials. Could I be liable for what visitors to my site do?](#)
- [Q: What is vicarious liability?](#)
- [Q: What is contributory infringement?](#)
- [Q: So am I better off not monitoring my website if I want to avoid contributory infringement liability?](#)
- [Q: Am I protected by Digital Millennium Copyright Act's Safe Harbor?](#)
- [Q: Can I copy or distribute software that is out of print and has been abandoned for years?](#)
- [Q: Aren't I allowed to make a backup copy of my software?](#)
- [Q: Isn't sending my friend a music file from a CD I already own just like loaning her the physical CD?](#)
- [Q: Aren't I allowed to make a backup copy of my software?](#)

<http://chillingeffects.org/piracy/faq.cgi>

**R**  
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# What any crawler should do



- Be capable of **distributed** operation
- Be scalable: need to be able to increase crawl rate by adding more machines
- Fetch pages of higher quality first
- Continuous operation: get fresh version of already crawled pages

# How hard can crawling be?



- Web **search engines must crawl** their documents.
- Getting the content of the documents is easier for many other IR systems.
  - E.g., indexing all files on your hard disk: just do a recursive descent on your file system
- Ok: for web IR, getting the content of the documents takes longer . . .
  - . . . because of latency.
- But is that really a design/systems challenge?



# Basic crawler operation

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- Initialize queue with URLs of known seed pages
- Repeat
  - Take URL from queue
  - Fetch and parse page
  - Extract URLs from page
  - Add URLs to queue
- Fundamental assumption: The web is well linked.





# What's wrong with this crawler?

```
urlqueue := (some carefully selected set of seed urls)
while urlqueue is not empty:
  myurl := urlqueue.getlastanddelete()
  mypage := myurl.fetch()
  fetchedurls.add(myurl)
  newurls := mypage.extracturls()
  for myurl in newurls:
    if myurl not in fetchedurls and not in urlqueue:
      urlqueue.add(myurl)
      addtoinvertedindex(mypage)
```

# What's wrong with the simple crawler



- Scale: we need to **distribute**.
- We can't index everything: we need to **subselect**. How?
- Duplicates: need to integrate **duplicate detection**
- Spam and spider traps: need to integrate **spam detection**
- **Politeness**: we need to be “nice” and space out all requests for a site over a longer period (hours, days)
- **Freshness**: we need to recrawl periodically.
  - Because of the size of the web, we can do frequent recrawls only for a small subset.
  - Again, subselection problem or **prioritization**

# Magnitude of the crawling problem



- To fetch 20,000,000,000 pages in one month . . .  
... we need to fetch almost 8000 pages per second!
- Actually: many more since many of the pages we attempt to crawl will be duplicates, unfetchable, spam etc.

# What a crawler must do



## Be polite

- Don't hit a site too often
- Only crawl pages you are allowed to crawl: robots.txt

## Be robust

- Be immune to spider traps, duplicates, very large pages, very large websites, dynamic pages etc

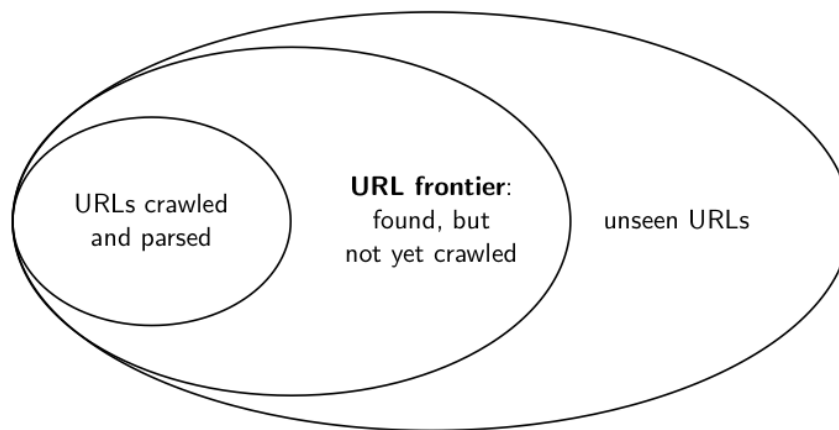
# Robots.txt



- Protocol for giving crawlers (“robots”) limited access to a website, originally from 1994
- Example:
  - User-agent: \*
  - Disallow: /yoursite/temp/
  - User-agent: searchengine
  - Disallow: /
- **Important:** cache the robots.txt file of each site we are crawling

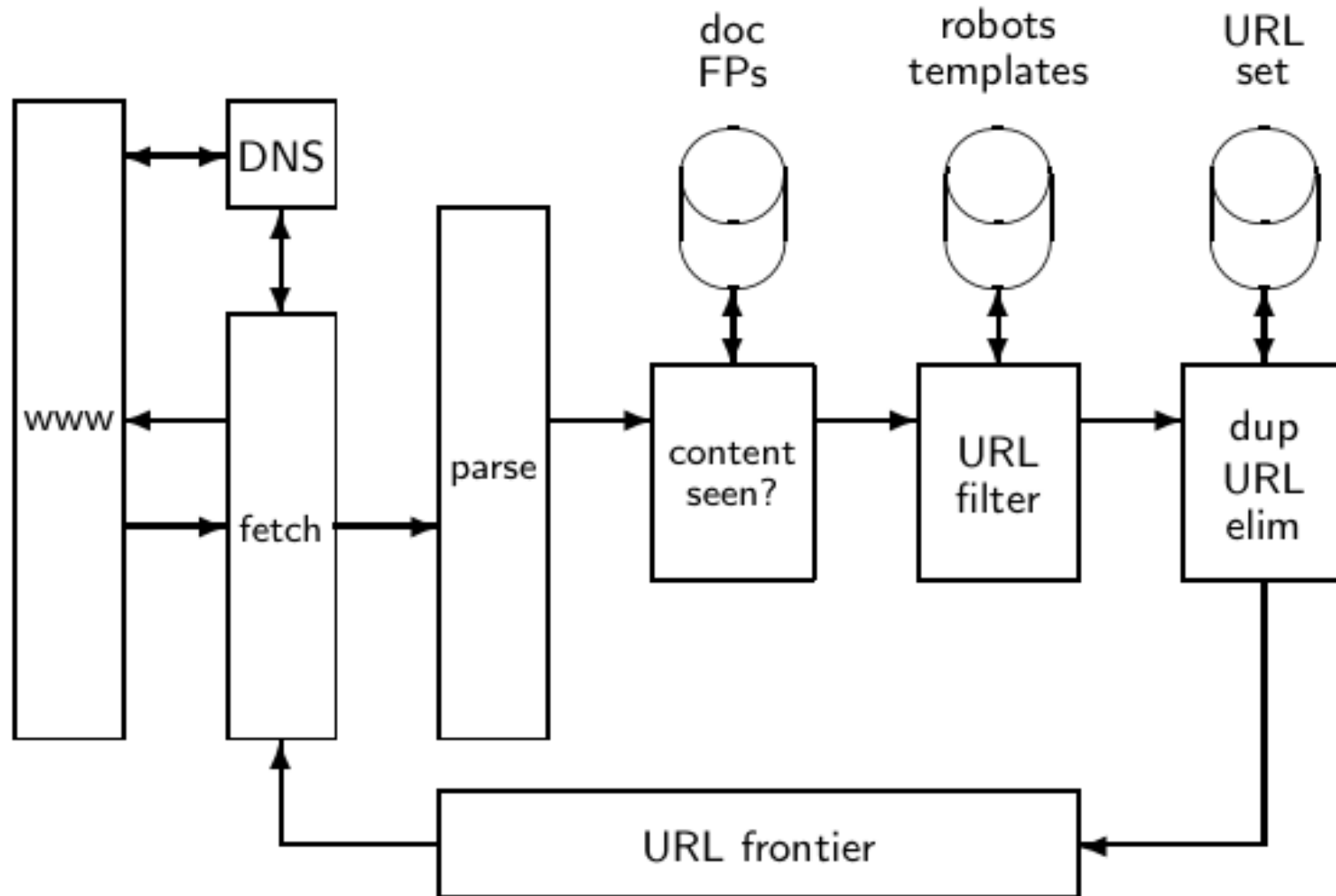


# URL Frontier



- The URL frontier is the data structure that holds and manages URLs we've seen, but that have not been crawled yet.
- Can include multiple pages from the same host
- Must avoid trying to fetch them all at the same time
- Must keep all crawling threads busy

# Basic Crawling Architecture





# URL normalization

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- Some URLs extracted from a document are **relative** URLs.
- E.g., at `http://mit.edu`, we may have `abouthome.html`
  - This is the same as: `http://mit.edu/abouthome.html`
- During parsing, we must normalize (expand) all relative URLs.





# Content seen

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- For each page fetched: check if the content is already in the index
- Check this using document fingerprints or shingles
- Skip documents whose content has already been indexed
- Still need to consider **Freshness**: Crawl some pages (e.g., news sites) more often than others

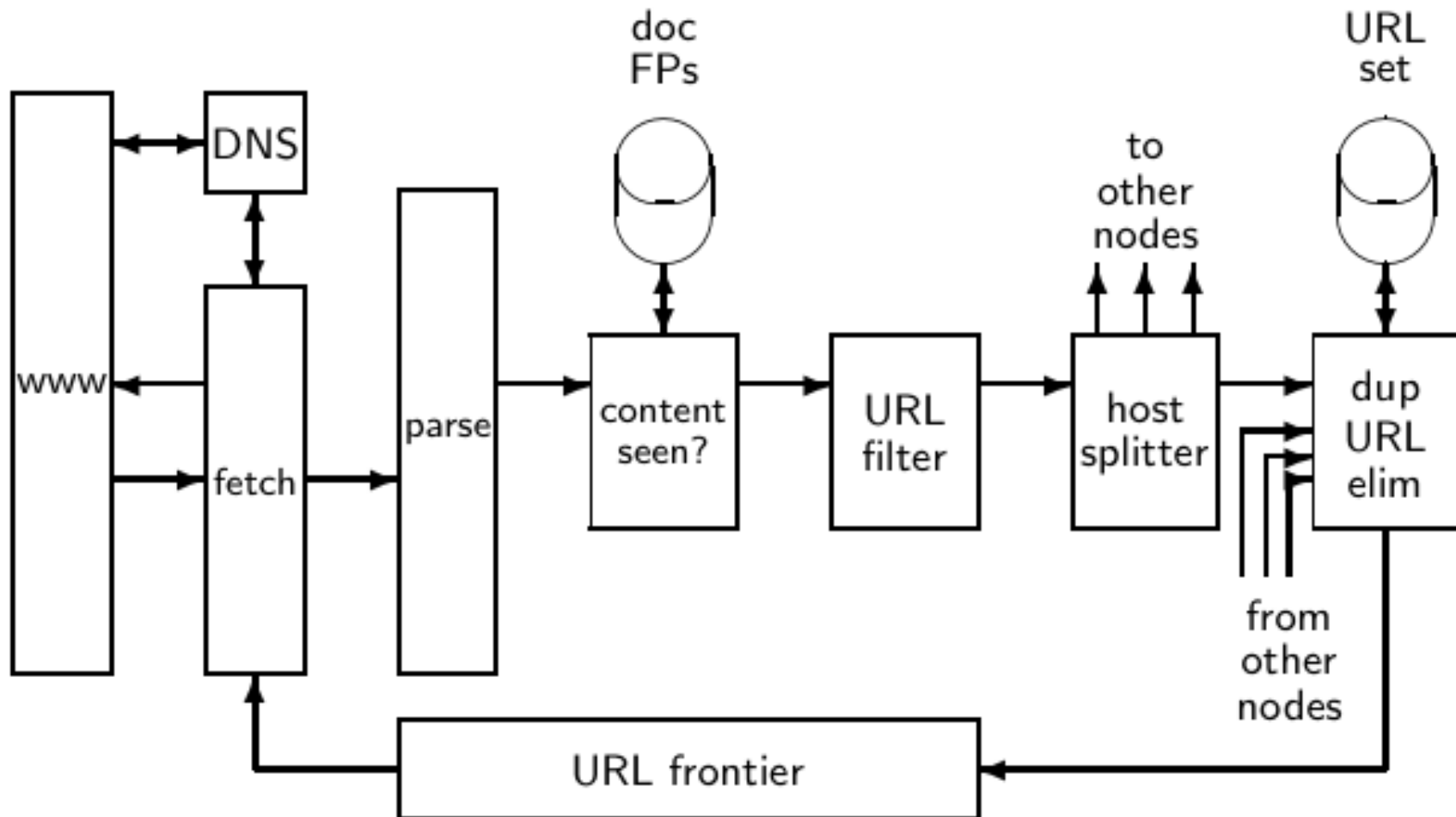
# Distributing the crawler



- Run multiple crawl threads, potentially at different nodes
  - Usually geographically distributed nodes
- Partition hosts being crawled into nodes



# Distributed crawling architecture



# A Crawler Issue: Spider traps

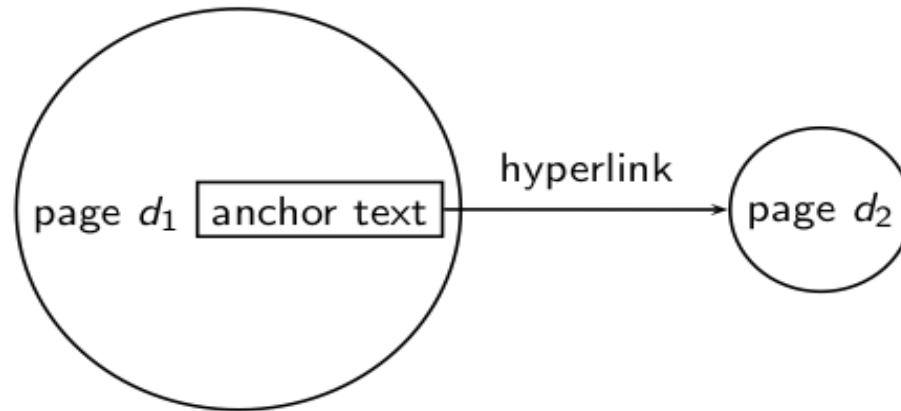


- Malicious server that generates an infinite sequence of linked pages
- Sophisticated spider traps generate pages that are not easily identified as dynamic.



# LINK ANALYSIS

# The web as a directed graph



- Assumption 1: A hyperlink is a quality signal.
  - The hyperlink  $d_1 \rightarrow d_2$  indicates that  $d_1$ 's author deems  $d_2$  high-quality and relevant.
- Assumption 2: The anchor text describes the content of  $d_2$ .
  - We use anchor text somewhat loosely here for: the text surrounding the hyperlink .
  - Example: “You can find cheap cars `<a href =http://...>here</a>`. ”
  - Anchor text: “You can find cheap cars here”

[text of  $d_2$ ] only vs.

[text of  $d_2$ ] + [anchor text  $\rightarrow d_2$ ]



- Searching on [text of  $d_2$ ] + [anchor text  $\rightarrow d_2$ ] is often more effective than searching on [text of  $d_2$ ] only.
- Example: Query *IBM*
  - Matches IBM's copyright page
  - Matches many spam pages
  - Matches IBM wikipedia article
  - May not match IBM home page!
  - ... if IBM home page is mostly graphics
- Searching on [anchor text  $\rightarrow d_2$ ] is better for the query *IBM*.
  - In this representation, the page with most occurrences of *IBM* is [www.ibm.com](http://www.ibm.com)

# Anchor text containing *IBM* pointing to [www.ibm.com](http://www.ibm.com)



www.nytimes.com: "IBM acquires Webify"

www.slashdot.org: "New IBM optical chip"

www.stanford.edu: "IBM faculty award recipients"

[www.ibm.com](http://www.ibm.com)





# Indexing anchor text

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- Thus: Anchor text is often a better description of a page's content than the page itself.
- Anchor text can be weighted more highly than document text.

(based on Assumption 1 & 2)

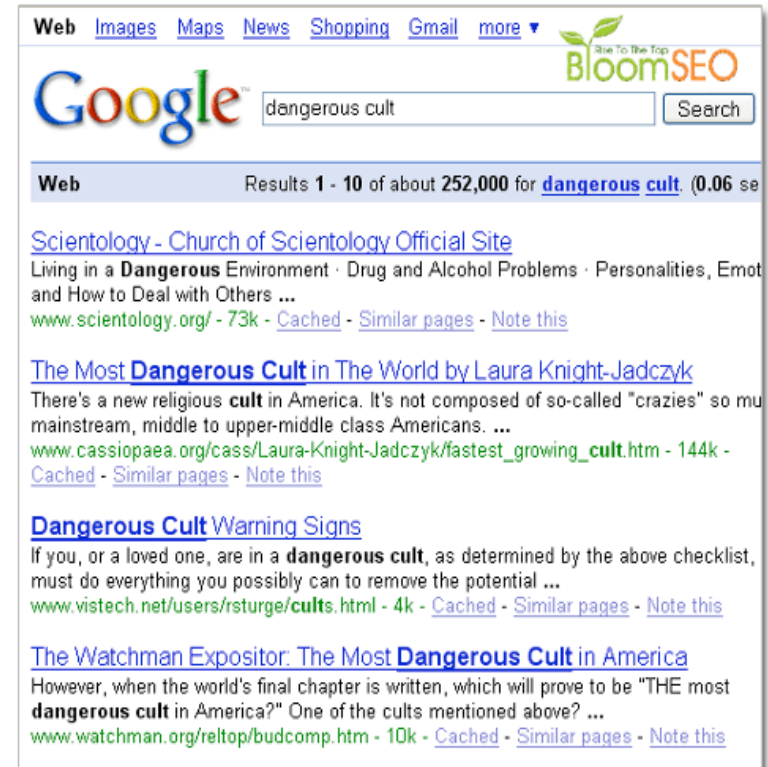


# Assumptions underlying PageRank

- Assumption 1: A link on the web is a quality signal – the author of the link thinks that the linked-to page is high-quality.
- Assumption 2: The anchor text describes the content of the linked-to page.
- Is Assumption 1 true in general?
- Is Assumption 2 true in general?

# Google bombs

- Is a search with “bad” results due to maliciously manipulated anchor text.
- E.g., [dangerous cult] on Google, Bing, Yahoo
  - Coordinated link creation by those who dislike the Church of Scientology
- Google introduced a new weighting function in January 2007 that fixed many Google bombs.
- Defused Google bombs: [who is a failure?], [evil empire]





# PAGERANK

# Origins of PageRank:

## Citation analysis (1)



- Citation analysis: analysis of citations in the scientific literature.
- Example citation: “[Miller \(2001\)](#) has shown that physical activity alters the metabolism of estrogens.”
- We can view “Miller (2001)” as a hyperlink linking two scientific articles.
- One application of these “hyperlinks” in the scientific literature:
  - Measure the similarity of two articles by the overlap of other articles citing them.
  - This is called [cocitation similarity](#).
  - Cocitation similarity on the web: Google’s “find pages like this” or “Similar” feature.



## Citation analysis (2)

- Another application: Citation frequency can be used to measure the **impact** of an article .
  - Simplest measure: Each article gets one vote – not very accurate.
- On the web: citation frequency = **inlink count**
  - A high inlink count does not necessarily mean high quality ...
  - ... mainly because of link spam.
- Better measure: **weighted** citation frequency or citation rank
  - An article's vote is weighted according to its citation impact.
  - Circular? No: can be formalized in a well-defined way



## Citation analysis (3)

- Better measure: weighted citation frequency or citation rank, invented in the context of citation analysis by Pinski and Narin in the 1960s.
- This is basically PageRank.
- We can use the same formal representation for
  - citations in the scientific literature
  - hyperlinks on the web
- Appropriately weighted citation frequency is an excellent measure of quality ...
  - ... both for web pages and for scientific publications.



# Definition of PageRank

- The importance of a page is given by the importance of the pages that link to it.

$$x_i = \sum_{j \in B_i} \frac{1}{N_j} x_j$$

importance of page  $i$

pages  $j$  that link to page  $i$

importance of page  $j$

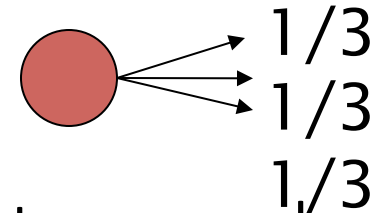
number of outlinks from page  $j$





# Pagerank scoring

- Imagine a browser doing a random walk on web pages:
  - Start at a random page
  - At each step, follow one of the  $n$  links on that page, each with  $1/n$  probability
- Do this repeatedly. Use the “long-term visit rate” as the page’s score
- This is a global score for the page, based on the topology of the network.
- Think of it as  $g(d)$  from Chapter 7

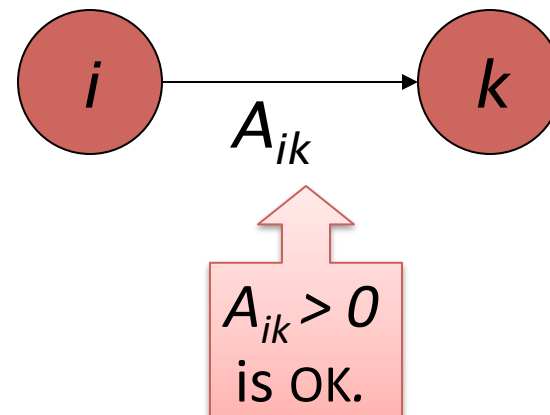


# Markov chains



A Markov chain consists of  $n$  states, plus an  $n \times n$  transition probability matrix  $A$ .

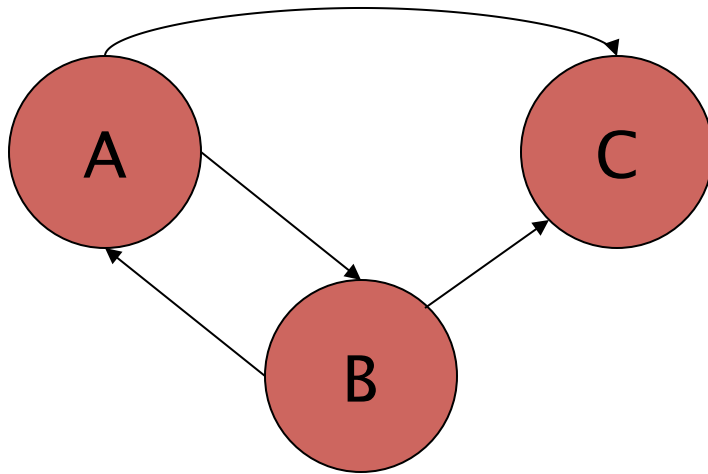
- At each step, we are in exactly one of the states.
- For  $1 \leq i, k \leq n$ , the matrix entry  $A_{ik}$  tells us the probability of  $k$  being the next state, given we are currently in state  $i$ .
- **Memorylessness property**: The next state depends only at the current state (first order Markov Chain)





# Markov chains

- Clearly, for all  $i$ ,  $\sum_{k=1}^n A_{ik} = 1$ .
- Markov chains are abstractions of random walks



Try this: Calculate the matrix  $A_{ik}$  using  $1/n$  possibility

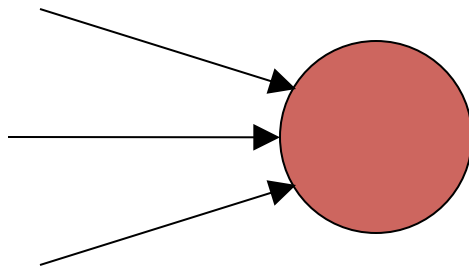
$A_{ik}$ :

	A	B	C
A			
B			
C			

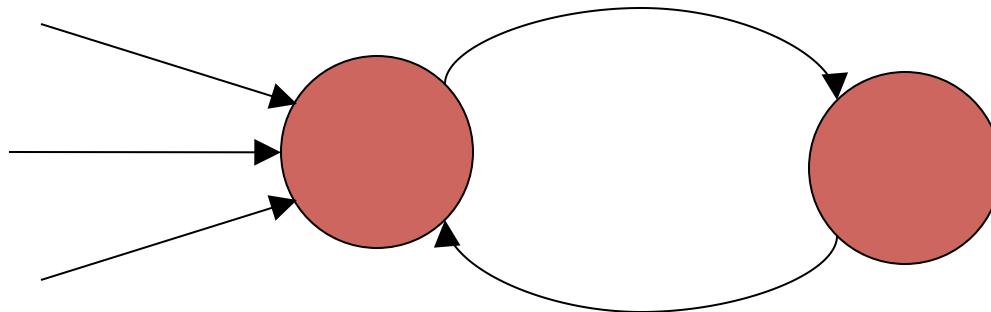


# Not quite enough

- The web is full of dead ends.
  - What sites have dead ends?
  - Our random walk can get stuck.



Dead End



Spider Trap



# Teleporting

- At each step, with probability 10%, teleport to a random web page
- With remaining probability (90%), follow a random link on the page
  - If a dead-end, stay put in this case

Follow!

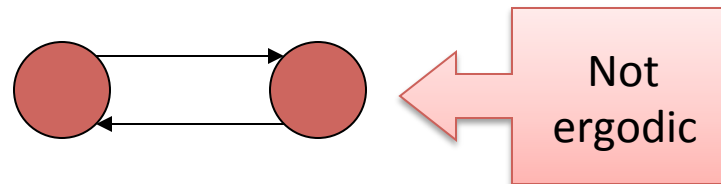
Teleport!

$$\vec{rank} = (1 - \alpha) A \times \vec{rank} + \alpha \left[ \frac{1}{N} \right] N \times 1$$



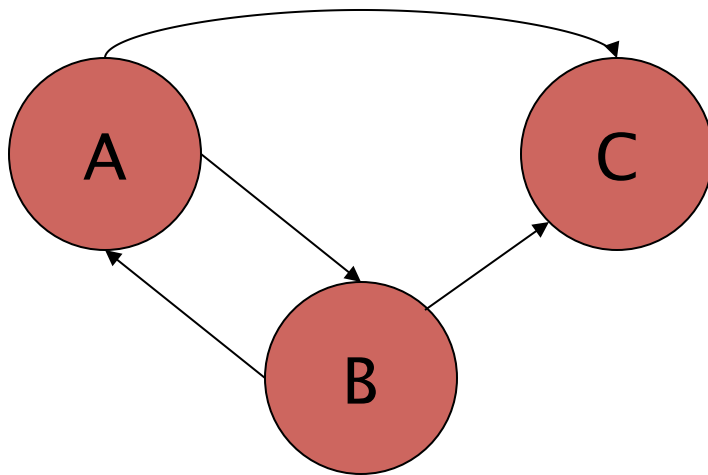
# Ergodic Markov chains

- A Markov chain is ergodic if
  - you have a path from any state to any other
  - you can be in any state at every time step, with non-zero probability



- With teleportation, our Markov chain is ergodic
- Theorem: With an ergodic Markov chain, there is a stable long term visit rate.

# Markov chains (2<sup>nd</sup> Try)



Try this: Calculate the matrix  $A_{ik}$  using a 10% chance of teleportation

$A_{ik}$ :

	A	B	C
A			
B			
C			



# Probability vectors

- A probability (row) vector  $\mathbf{x} = (x_1, \dots, x_n)$  tells us where the walk is at any point
- E.g.,  $(\underset{1}{000}\dots\underset{i}{1}\dots\underset{n}{000})$  means we're in state  $i$ .

More generally, the vector  $\mathbf{x} = (x_1, \dots, x_n)$  means  
The walk is in state  $i$  with probability  $x_i$ .

$$\sum_{i=1}^n x_i = 1.$$



# Change in probability vector



- If the probability vector is  $\mathbf{x} = (x_1, \dots, x_n)$  at this step, what is it at the next step?
- Recall that row  $i$  of the transition prob. Matrix  $\mathbf{A}$  tells us where we go next from state  $i$ .
- So from  $\mathbf{x}$ , our next state is distributed as  $\mathbf{x}\mathbf{A}$ .



# Pagerank algorithm

- Regardless of where we start, we eventually reach the steady state  $a$ 
  - Start with any distribution (say  $x=(10\dots0)$ )
  - After one step, we're at  $xA$
  - After two steps at  $xA^2$ , then  $xA^3$  and so on.
  - “Eventually” means for “large”  $k$ ,  $xA^k = a$
- Algorithm: multiply  $x$  by increasing powers of  $A$  until the product looks stable

# Steady State



- For any ergodic Markov chain, there is a unique long-term visit rate for each state
  - Over a long period, we'll visit each state in proportion to this rate
  - It doesn't matter where we start

# Eigenvector formulation



- The flow equations can be written

$$\mathbf{r} = \mathbf{A}\mathbf{r}$$

- So the rank vector is an eigenvector of the adjacency matrix
  - In fact, it's the first or principal eigenvector, with corresponding eigenvalue 1



# Pagerank summary

- Pre-processing:
  - Given graph of links, build matrix **A**
  - From it compute **a**
  - The pagerank  $a_i$  is a scaled number between 0 and 1
- Query processing:
  - Retrieve pages meeting query
  - Rank them by their pagerank
  - Order is *query-independent*



# PageRank issues

- Real surfers are not random surfers.
  - Examples of nonrandom surfing: back button, short vs. long paths, bookmarks, directories – and search!
  - → Markov model is not a good model of surfing.
  - But it's good enough as a model for our purposes.
- Simple PageRank ranking (as described on previous slide) produces bad results for many pages.
  - Consider the query [video service].
  - The Yahoo home page (i) has a very high PageRank and (ii) contains both *video* and *service*.
  - If we rank all Boolean hits according to PageRank, then the Yahoo home page would be top-ranked.
  - Clearly not desirable.

# How important is PageRank?



- Frequent claim: PageRank is the most important component of web ranking.
- The reality:
  - There are several components that are at least as important: e.g., anchor text, phrases, proximity, tiered indexes ...
  - Rumor has it that PageRank in his original form (as presented here) now has a negligible impact on ranking!
  - However, variants of a page's PageRank are still an essential part of ranking.
  - Addressing link spam is difficult and crucial.



# Summary

- Crawling – Obtaining documents for indexing
  - Need to be polite
- PageRank – A  $G(d)$  for asymmetrically linked documents
- Chapters 20 and 21 of IIR
- Resources
  - Paper on [Mercator crawler](#) by Heydon et al.
  - [Robot exclusion standard](#)

“PageRank reflects our view of the importance of web pages by considering more than 500 million variables and 2 billion terms. Pages that believe are important pages receive a higher PageRank and are more likely to appear at the top of the search results”