第一题 本题考虑对称矩阵的 Gauß 消去法和 LU 分解

(a) 假设 A 是一个满足 $a_11 \neq 0$ 的对称矩阵,当 A 的第一列完成消去的时候我们得到

$$\begin{bmatrix} a_{11} & a_{12} \dots a_{1n} \\ 0 & & \\ \vdots & A^{(1)} \\ 0 & & \end{bmatrix}$$

证明 $A^{(1)}$ 是对称的。

证明: 设 $A = \{a_{ij}\}_{n \times n}$, $A^{(1)} = \{a'_{ij}\}_{(n-1)\times(n-1)}$, 由 $a_11 \neq 0$, 那么根据消去规则有

$$(0, a'_{11}, a'_{12}, \dots, a'_{1(n-1)}) = (a_{21}, a_{22}, \dots, a_{2n}) - \frac{a_{21}}{a_{11}} (a_{11}, a_{12}, \dots, a_{1n})$$
$$= (0, a_{22} - \frac{a_{21}}{a_{11}}) a_{12}, \dots, a_{2n} - \frac{a_{21}}{a_{11}} a_{1n})$$

同理有, 对 $\forall i, j \in N^*$,

$$a'_{ij} = a_{(i+1)(j+1)} - \frac{a_{(i+1)1}}{a_{11}} a_{1(j+1)}$$

 $\therefore A$ 为正定矩阵, $\forall i, j \in N^*, a_{ij} = a_{ji}$, 则对于 $A^{(1)}$ 有

$$a'_{ji} = a_{(j+1)(i+1)} - \frac{a_{(j+1)1}}{a_{11}} a_{1(i+1)}$$

$$= a_{(i+1)(j+1)} - \frac{a_{(i+1)1}}{a_{11}} a_{1(j+1)}$$

$$= a'_{ij}$$

· A⁽¹⁾ 是对称的

(b) 根据上一问的结论用伪代码的形式写出计算一个正定矩阵 *LU* 分解的算法,利用对称性节省计算量。

算法思路:

由于是对称正定矩阵,故其在每一轮高斯消元时都能保证 $a_{ii} \neq 0$ (见《Linear Algebra with Applications(9th Edition)》—Steven J.Leon). 由于每次消元的子矩阵都

是对称的,故只需算出其上三角部分,下三角部分直接对称赋值即可,节约计算量。

算法伪代码:

```
% 对正定对称矩阵A进行不做行交换的LU分解.
[sx,sy] = size(A);
U = A;
L = zeros(sx,sy);
for i = 1 : sy
    L(i,i) = 1;
    for j = i + 1 : sx
        L(j,i) = U(j,i) / U(i,i);
        U(j,i) = 0;
        for k = j : sy
            U(j,k) = U(j,k) - L(j,i) * U(i,k);
            U(k,j) = U(j,k);
        end
    end
end
```

(c) 编写程序, 用 Cholesky 分解解给定方程组 Ax = b。

代码如下

```
clear, clc

A = [4,-2,4,2;-2,10,-2,-7;4,-2,8,4;2,-7,4,7];
b = [8;2;16;6];
showAb(A,b);
% 对正定对称矩阵A进行不做行交换的LU分解.
[sx,sy] = size(A);
U = A;
L = zeros(sx,sy);
for i = 1 : sy
    L(i,i) = 1;
    for j = i + 1 : sx
    L(j,i) = U(j,i) / U(i,i);
    U(j,i) = 0;
```

```
for k = j : sy
           U(j,k) = U(j,k) - L(j,i) * U(i,k);
           U(k,j) = U(j,k);
        end
    end
end
showLU(L,U);
%对L进行行列数乘得到新的L,即Cholesky分解
for i = 1 : sx
    factor = sqrt(U(i,i));
   for j = i : sx
        L(j,i) = L(j,i) * factor;
    end
end
showLLt(L,L');
%根据L求解Y
Y = zeros(sy,1);
for i = 1 : sy
   Y(i) = b(i);
   for j = 1 : i - 1
        Y(i) = Y(i) - L(i,j) * Y(j);
    end
   Y(i) = Y(i) / L(i,i);
end
%根据L的转置求解X
X = zeros(sx,1);
for i = sy : -1 : 1
X(i) = Y(i);
for j = i + 1 : sx
X(i) = X(i) - L(j,i) * X(j);
X(i) = X(i) / L(i,i);
end
%输出解
fprintf("x = \n");
```

```
disp(X);
%%
%输出A和b
function showAb(A,b)
\% Display the matrix A and the vector b
   used in Gauss elimination
...此处省略
end
%%
%输出L和U
function showLU(L,U,varargin)
\% Display the matrix L and U
% got in LU decomposition
...此处省略
end
%%
%输出L和L'
function showLLt(L,Lt,varargin)
\% Display the matrix L and U
 got in LU decomposition
... 此处省略
end
```

输出结果如图 1

第二题 Richardson 迭代方法, 对于通用迭代格式

$$x^{(k+1)} = M^{-1}Nx^{(k)} + M^{-1}b$$

Richardson 迭代的 $M=\frac{1}{\omega}I, N=\frac{1}{\omega}I-A$, 此处 $\omega>0$ 。考虑 A 为正定的情况,并设 A 的最小和最大特征值分别为 λ_1 和 λ_n

(a) 证明 Richardson 迭代方法在 $\omega < \frac{2}{\lambda_n}$ 的情况下收敛

证明:由

$$x^{(k+1)} = M^{-1}Nx^{(k)} + M^{-1}b$$

有

$$Mx^{(k+1)} = Nx^{(k)} + b$$



图 1: 第一题 (c) 小题结果

代入精确值 (代入 $k \to \infty$) 后与原式相减得到

$$M(x^{(k+1)} - x) = N(x^{(k)} - x)$$

令误差 $e^k = |x^k - x|$ 有

$$Me^{(k+1)} = Ne^{(k)}$$

即

$$e^{(k+1)} = (I - \omega A)e^{(k)}$$

如此,考察矩阵 $G=I-\omega A$ 的特征值,设由 A 的特征值构成的对角矩阵为 D_A ,于是

$$I - \omega A = I - \omega P^{-1} D_A P$$

$$= P^{-1} I P - P^{-1} \omega D_A P$$

$$= P^{-1} (I - \omega D_A) P$$

$$= P^{-1} D_{I - \omega A} P$$

所以 G 的谱半径 $\rho(G_{\omega})$,即 G 的特征值的绝对值最大值为 $\max |1-\omega\lambda_i|$,而由题意,在 $\omega<\frac{2}{\lambda_n}$ 的时候

$$\rho(G_{\omega}) = \max |1 - \omega \lambda_i|$$

$$= \max\{|1 - \omega \lambda_1|, |1 - \omega \lambda_n|\}$$

$$< 1$$

所以 Richardson 迭代方法在该条件下收敛

(b) 证明其谱半径以及取最佳值时候的 ω 值由等式

$$1 - \omega \lambda_1 = \omega \lambda_n - 1$$

解得 $\omega = \frac{2}{\lambda_1 + \lambda_2}$, 记作 ω_b 引用上一小题结论有

$$\rho(G_{\omega}) = \max\{|1 - \omega \lambda_1|, |1 - \omega \lambda_n|\}$$

$$= \begin{cases} 1 - \omega \lambda_1 &, \omega \leq \omega_b \\ \frac{\lambda_n - \lambda_1}{\lambda_n + \lambda_1} &, \omega = \omega_b \\ \omega \lambda_n - 1 &, \omega \geq \omega_b \end{cases}$$

易证其取到最小值的时候, 即最优情况下, $\omega = \omega_b = \frac{2}{\lambda_1 + \lambda_n}$.

(c) 设计一个方法用 Matlab 随机数生成函数 rand 构造出 2 一个 5×5 的特征值为 1, 2, 3, 4, 5 的正定矩阵作为 A, 再用 rand 构造出一个 5×1 的向量作为 b。然后用上述 Richardson 迭代方法解 Ax = b,作图验证收敛半径。

代码如下

```
clear, clc
P = orth(rand(5,5));
B = diag([1,2,3,4,5]);
A = P \backslash B * P;
b = rand(5,1);
F = Q(x) (x > 2/(1+5)) .* (5 * x - 1) + (x <= 2/(1+5)) .* (1 - x);
%%
%遍历
wstep = 1e-3;
wnum = 2/5 /wstep;
wSet = linspace(wstep,2/5-wstep,wnum-1);
rhoReal = F(wSet);
rhoSet = zeros(wnum-1,1);
for i = 1 : wnum-1
    G = eye(5) - wSet(i) * A;
    x next=zeros(5,1);
    e_next = 1;
```

```
while(e_next>1e-11)
        x_curr = x_next;
        x_next = G*x_curr+wSet(i)*b;
        e_curr = e_next;
        e_next = norm(x_next-x_curr);
    end
    rhoSet(i) = e_next/e_curr;
end
%%
%计算rho
wBest = 2/(1+5);
G = eye(5) - wBest * A;
while(norm(x_next-x_curr)>1e-11)
    x_curr = x_next;
    x_next = G*x_curr+wBest*b;
end
fprintf("x=\n");
disp(x_next);
%%
%绘图验证
figure(1);
p1 = plot(wSet,rhoSet,'k'); hold on
p2 = plot(wSet,rhoReal,'b'); hold on
legend ([p1 ,p2], 'approx ', 'real');
xlabel('omega');
ylabel('rho');
%%
%输出A和b
function showAb(A,b)
\% Display the matrix A and the vector b
%
  used in Gauss elimination
```

```
... 此处省略
end
%%
%输出L和U
function showLU(L,U,varargin)
\% Display the matrix L and U
  got in LU decomposition
...此处省略
end
%%
%输出L和L'
function showLLt(L,Lt,varargin)
\% Display the matrix L and Lt
   got in LU decomposition
...此处省略
end
```

运行输出绘图结果如图 2, 可见和表达式基本吻合。计算得到方程解如图 3.

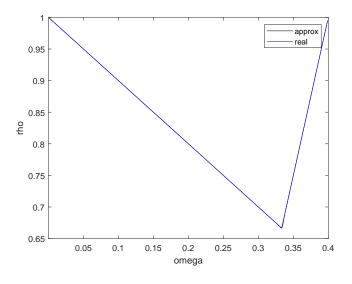


图 2: 第二题 (c) 小题作图

第三题 从另一个角度考虑 Gau 积分

(a) 由于高斯积分可以用 n 个采样点上的采样值精确求得一个 2n-1 阶多项式的定积分. 现取 n=6, 那么可以精确求得任意一个 11 阶多项式的定积分, 定义一组

命令行窗口

不熟悉 MATLAB? 请参阅有关快速入门的资源。

X=

-0.1902

0.2529

0.2514

0.4437

图 3: 第二题 (c) 小题方程组解

基函数 $\{1, x, x^2, \dots, x^{11}\}$, 设积分结点为 $\{x_i | 1 \le i \le 12, x_1 < x_2 < \dots < x_{12}\}$, 有

$$\begin{cases} \sum_{i=1}^{6} \omega_i x_i^0 - 2 &= 0\\ \sum_{i=1}^{6} \omega_i x_i^1 - \frac{0}{2} &= 0\\ \sum_{i=1}^{6} \omega_i x_i^2 - \frac{2}{3} &= 0\\ & \dots\\ \sum_{i=1}^{6} \omega_i x_i^{11} - \frac{0}{12} &= 0 \end{cases}$$

又由积分结点和积分权重关于原点的对称性,原方程组简化为(此处只需在意结

点的排布顺序, 哪边更大并不影响)

$$\begin{cases} \sum_{i=1}^{3} \omega_{i} x_{i}^{0} - 1 & = 0 \\ \sum_{i=1}^{3} \omega_{i} x_{i}^{2} - \frac{1}{3} & = 0 \\ \dots & \\ \sum_{i=1}^{3} \omega_{i} x_{i}^{10} - \frac{1}{11} & = 0 \\ x_{1} + x_{6} & = 0 \\ x_{2} + x_{5} & = 0 \\ x_{3} + x_{4} & = 0 \\ \omega_{1} - \omega_{6} & = 0 \\ \omega_{2} - \omega_{5} & = 0 \\ \omega_{3} - \omega_{4} & = 0 \\ x_{1}^{2}, x_{2}^{2}, x_{3}^{2} & > 0 \end{cases}$$

将 x_i^2 视作一个变量 s_i , 方程式变为

$$\begin{cases} \sum_{i=1}^{3} \omega_{i} s_{i}^{0} - 1 &= 0\\ \sum_{i=1}^{3} \omega_{i} s_{i}^{1} - \frac{1}{3} &= 0\\ \dots\\ \sum_{i=1}^{3} \omega_{i} s_{i}^{5} - \frac{1}{11} &= 0\\ s_{1}, s_{2}, s_{3} &> 0 \end{cases}$$

(b) 直接求导写出其雅可比矩阵

$$\begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 \\ \omega_1 & \omega_2 & \omega_3 & s_1 & s_2 & s_3 \\ 2\omega_1 s_1 & 2\omega_2 s_2 & 2\omega_3 s_3 & s_1^2 & s_2^2 & s_3^2 \\ 3\omega_1 s_1^2 & 3\omega_2 s_2^2 & 3\omega_3 s_3^2 & s_1^3 & s_2^3 & s_3^3 \\ 4\omega_1 s_1^3 & 4\omega_2 s_2^3 & 4\omega_3 s_3^3 & s_1^4 & s_2^4 & s_3^4 \\ 5\omega_1 s_1^4 & 5\omega_2 s_2^4 & 5\omega_3 s_3^4 & s_1^5 & s_2^5 & s_3^5 \end{bmatrix}$$

(c) 选取等距节点 $\{x_i\} = \{-1, -0.6, -0.2, 0.2, 0.6, 1\}$, 转化为 $\{s_i\} = \{0.04, 0.36, 1\}$. 观察积分权重公式

$$\int_{-1}^{1} \frac{\prod_{k \neq i} (x - x_k)}{\prod_{k \neq i} (x_i - x_k)} dx$$

其被积函数在除了 x_i 之外的所有插值点其值都为 0,由于其本身是一个多项式函数,所以对其可以粗略看成仅在 x_i 附近的两个点之间对积分做了主要贡献,由于插值点均匀分布,那么其对于在每个插值点上的权重积分便具有粗糙的对称性,换言之所有权重趋向于值相近,而贡献总和由方程组的第一个方程可以粗略估计为 2,故最终确定权值为各个点均分 2。

```
clear, clc
n = 6;
X \text{ curr} = [0,0,0,0,0,0]';
X \text{ next} = [0.04, 0.36, 1, 0.33, 0.33, 0.33]';
err past = 3;
err curr = 2;
err_next = 1;
DF = zeros(6,6);
n = 1;
while(norm(X_next-X_curr)>1e-7)
    X_curr = X_next;
    %构造DF
    fprintf("n=");
    disp(n);
    n = n + 1;
    DF(1,:) = [0,0,0,1,1,1];
    DF(2,:) = [X curr(4:6)', X curr(1:3)'];
    DF(3,:) = DF(2,:) .* [X_curr(1:3)', X_curr(1:3)'];
    DF(4,:) = DF(3,:) .* [X curr(1:3)', X curr(1:3)'];
    DF(5,:) = DF(4,:) .* [X curr(1:3)', X curr(1:3)'];
    DF(6,:) = DF(5,:) .* [X curr(1:3)', X curr(1:3)'];
    DF(3,1:3) = DF(3,1:3) * 2;
    DF(4,1:3) = DF(4,1:3) * 3;
    DF(5,1:3) = DF(5,1:3) * 4;
```

```
DF(6,1:3) = DF(6,1:3) * 5;
    %求解F(x)
   L1 = zeros(6,3);
    U1 = X_curr(4:6)';
    Fx = zeros(6,1);
    L1(1,:) = [1,1,1];
    for i = 2 : 6
    L1(i,:) = L1(i-1,:) .* X_curr(1:3)';
end
Fx = L1 * U1' - [1,1/3,1/5,1/7,1/9,1/11]';
% 迭代求解
X_next = X_curr - DF\Fx;
err_past = err_curr;
err_curr = err_next;
err_next = norm(X_next-X_curr);
p = log(err_curr/err_next)/log(err_past/err_curr);
fprintf("p = ");
disp(p);
end
```

运行输出结果如图 4 和图 5

```
p = 1.521018179088277

n= 2

p = 0.760162536516021

n= 3

p = 2.081693478112024

n= 4

p = 1.107914148944804

n= 5

p = 2.501462037360084

n= 6

p = 2.040779175369055
```

图 4: 第三题 (c) 小题迭代次数与收敛阶数

```
-0.932469514203152
-0.66129386466263
-0.238619186083195
0.238619186083195
0.66129386466263
0.932469514203152

W =
0.171324492379171
0.360761573048140
0.467913934572688
0.467913934572688
0.360761573048140
```

图 5: 第三题 (c) 小题解

(d) 同理先列出其表达式

$$\omega_{i} = \int_{-1}^{1} \frac{\prod_{k \neq i} (\cos \frac{x\pi}{n} - \cos \frac{k\pi}{n})}{\prod_{k \neq i} (\cos \frac{i\pi}{n} - \cos \frac{k\pi}{n})} dx$$
$$= \int_{-1}^{1} \frac{\prod_{k \neq i} (\sin \frac{(x+k)\pi}{2n} \sin \frac{(x-k)\pi}{2n})}{\prod_{k \neq i} (\sin \frac{(i+k)\pi}{2n} \sin \frac{(i-k)\pi}{2n})} dx$$

由其乘积式估得, 其赋予均值的估计是合理的。

```
clear,clc
for n = 4 : 30
[realx,realw] = gauss(n);
ns = n-round(n/2);
ChewbyX = linspace(0, n-1, n);
ChewbyPoint = cos(ChewbyX * pi / (n-1));
X_{\text{next}} = [\text{ChewbyPoint}(1:ns), \text{ones}(1, \text{round}(n/2))/\text{round}(n/2)];
s = ones(n,1);
while (norm(s) > 1e-7)
    X curr = X next;
    DF = zeros(n);
    Fx = zeros(n,1);
    if (rem(n,2) == 1)
         [Fx,DF] = SingleDF(ns,X_curr);
    else
         [Fx,DF] = DoubleDF(ns,X_curr);
```

```
end
    s = -inv(DF)*Fx';
    X_next = X_curr + s';
end
if(norm(X next-
            [abs(realx(1:n-round(n/2))'),
            realw(1:round(n/2))])>1e-7)
    break;
end
disp(n);
end
%%
function [Fx,DF] = DoubleDF(ns,X)
    Fx = zeros(2*ns,1);
    DF(1,:) = [zeros(1,ns), ones(1,ns)];
    Fx(1) = sum(X(ns+1:2*ns))-1;
    DF(2,1:ns) = 2*[X(1:ns) .* X(ns+1:2*ns)];
    DF(2,ns+1:2*ns) = X(1:ns) .* X(1:ns);
    Fx(2) = sum(X(ns+1:2*ns).*DF(2,ns+1:2*ns)) - 1/3;
    for i = 1 : 2 * ns - 2
        DF(i+2,1:ns) = (i+1)/i * DF(i+1,1:ns)
                        .* DF(2,ns+1:2*ns);
        DF(i+2,ns+1:2*ns) = DF(i+1,ns+1:2*ns)
                             .* DF(2,ns+1:2*ns);
        Fx(i+2) = sum(X(ns+1:2*ns))
                .*DF(i+2,ns+1:2*ns))-1/(2*i+3);
    end
    Fx = Fx';
end
%%
function [Fx,DF] = SingleDF(ns,X)
    DF(1,:) = [zeros(1,ns), ones(1,ns), 1/2];
    Fx(1) = sum(X(ns+1:2*ns))+X(2*ns+1)/2-1;
    DF(2,1:ns) = 2*[X(1:ns) .* X(ns+1:2*ns)];
```

```
DF(2,ns+1:2*ns) = X(1:ns) .* X(1:ns);
   Fx(2) = sum(X(ns+1:2*ns).*DF(2,ns+1:2*ns))-1/3;
    for i = 1 : 2 * ns - 1
       DF(i+2,1:ns) = (i+1)/i * DF(i+1,1:ns)
                            * DF(2,ns+1:2*ns);
        DF(i+2,ns+1:2*ns) = DF(i+1,ns+1:2*ns)
                            .* DF(2,ns+1:2*ns);
       Fx(i+2) = sum(X(ns+1:2*ns))
                            *DF(i+2,ns+1:2*ns))-1/(2*i+3);
    end
end
%%
function [X] = LUC(A,b)
%对正定对称矩阵A进行不做行交换的LU分解.
[sx,sy] = size(A);
   U = A;
   L = zeros(sx, sy);
    for i = 1 : sy
       L(i,i) = 1;
        for j = i + 1 : sx
           L(j,i) = U(j,i) / U(i,i);
           U(j,i) = 0;
            for k = i+1 : sy
               U(j,k) = U(j,k) - L(j,i) * U(i,k);
                U(k,j) = U(j,k);
            end
        end
    end
   %根据L求解Y
   Y = zeros(sy,1);
   for i = 1 : sy
       Y(i) = b(i);
        for j = 1 : i - 1
```

```
Y(i) = Y(i) - L(i,j) * Y(j);
        end
    end
    %根据U求解X
    X = zeros(sx,1);
    for i = sy : -1 : 1
            X(i) = Y(i);
        for j = i + 1 : sx
            X(i) = X(i) - U(i,j) * X(j);
        end
        X(i) = X(i) / U(i,i);
    end
end
%%
function [x,w] = gauss(N)
    beta = .5./sqrt(1-(2*(1:N-1)).^(-2));
    T = diag(beta,1) + diag(beta,-1);
    [V,D] = eig(T);
    x = diag(D); [x,i] = sort(x);
    w = 2*V(1,i).^2;
end
```

运行程序输出结果如图 2, 可见在该初值选取下最大的 n 是 10.

图 6: 第三题 (d) 小题解