

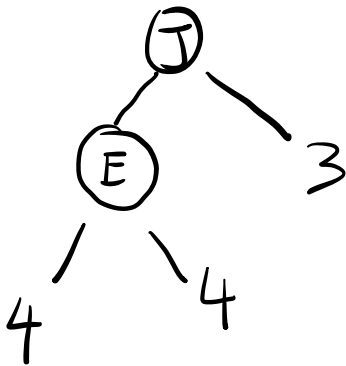
$$n_1 = 1 \quad n_2 = 2 \quad n_3 = 5 \quad n_4 = 5 \times 2 + 2 \times 2$$

$$\Rightarrow n_5 = 14 \times 2 + 2 \times 5 + 2 \times 2$$

$$= 28 + 10 + 4 = 42$$

2.

A B C D E F G H I J K L M



14 types of tree of 4

5 types of tree of 3

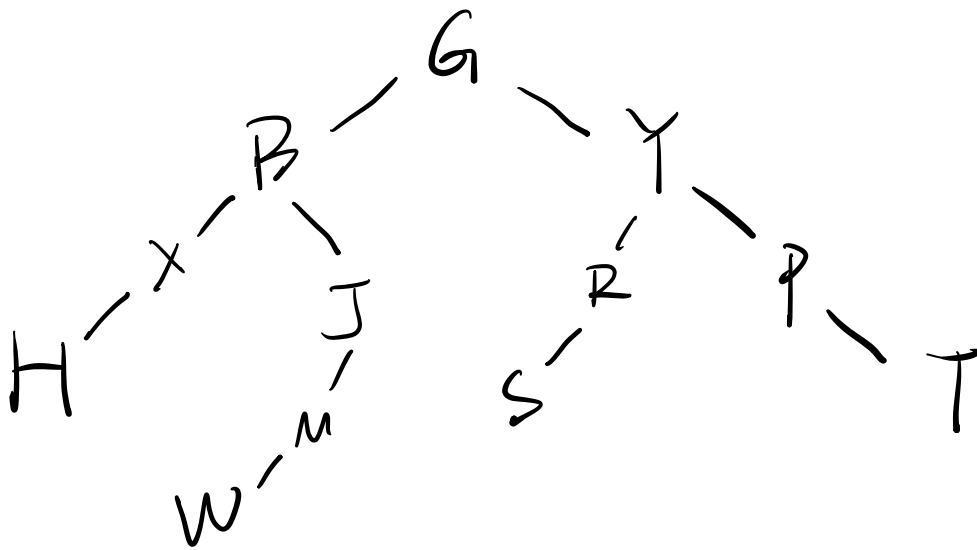
$$14 \times 14 \times 5 = 980$$

3. in order (left, Root, Right)

post order (left, Right, Root)

H X B W M J | G | S R Y P T

H X W M J B | S R T P Y G
L R



by pre order we know G is the root

Y is the Right child of root, and by

both we know H is the most left child

and T is the Right most child -

4.

(a) (1) extraMin()

follow the left pointers from the root until a node with no left pointers. delete that node and return the value.

It has $O(\log n)$ time

(2) Insert()

Insert the key in the AVL Tree and rotate to rebalance. It also takes $O(\log n)$ time.

(2) It is worse than a heap

running time for min & extraction of heap

is $O(1)$

insert has $O(\log n)$ which is the same for both,