

Compilation 2024

# Dolphin: phase 5

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# LLVM -- for translating aggregates

- Structure types
- Fixed-size arrays
- Named types
- Global Variables
- String variables
- Casting
- Pointer to integer conversion
- Computing physical size of types
- `getelementptr` (Gep)

# Translating Types

- Start translation by translating types to LLVM --

```
let trans_type tp =  
  match tp with  
  | TAst.Determined (TAst.Void _) -> Ll.Void  
  | TAst.Determined (TAst.Int _) -> Ll.I64  
  | TAst.Determined (TAst.Bool _) -> Ll.I1  
  | TAst.Determined (TAst.Byte _) -> Ll.I8  
  ...
```

# Flexible Array Elements

- In C (as of C99) we can pack an array together with its length into a single struct by having an array field of **unspecified length** as the **last** field.

```
struct array { int64_t len; char contents[]; };
```

- On the LLVM side, this is represented as:

```
%array_type = type {i64, [0 x i8]}
```

- We use pointers to these types to represent both Dolphin's arrays and strings

# Translating Arrays & Strings

- Start translation by translating types to LLVM --

```
let trans_type tp =  
  match tp with  
  | TAst.Determined (TAst.Void _) -> Ll.Void  
  | TAst.Determined (TAst.Int _) -> Ll.I64  
  | TAst.Determined (TAst.Bool _) -> Ll.I1  
  | TAst.Determined (TAst.Byte _) -> Ll.I8  
  | TAst.Determined (TAst.Str _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  | TAst.Determined (TAst.Array _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  ...
```

# Translating String Literals

```
int main () {  
    var x = "Hello World\n";  
    output_string(x, get_stdout());  
    return 0;  
}
```

The Dolphin program above is translated to the following LLVM -- program

```
%dolphin_record_stream = type { }  
%array_type = type { i64, [0 x i8] }  
  
@string_literal$1 = global { i64, [12 x i8] } {i64 12, [12 x i8] c"Hello World\0A"}  
  
declare void @output_string(%array_type*, %dolphin_record_stream*)  
declare %dolphin_record_stream* @get_stdout()  
  
define i64 @dolphin_fun_main () {  
    %x$0 = alloca %array_type*  
    %conv_string_literal_packed$2 = bitcast { i64, [12 x i8] }* @string_literal$1 to %array_type*  
    store %array_type* %conv_string_literal_packed$2, %array_type** %x$0  
    %load_local_var$3 = load %array_type*, %array_type** %x$0  
    %call$4 = call %dolphin_record_stream* @get_stdout ()  
    call void @output_string (%array_type* %load_local_var$3, %dolphin_record_stream* %call$4)  
    ret i64 0  
after_return$5:  
    unreachable  
}
```

# Translating Array Creation

- Use the runtime function

```
struct array *allocate_array(int32_t elem_size, int64_t numelems, void* contents)
```

- Compute the size of elements in LLVM
- Determine the default value based on the type
- Allocate space for the default value using LLVM's `alloca`
- Store the default value into the allocated space
- Pass the pointer to the default value along with size of array's elements and its length to the `allocate_array`
- Note how `allocate_array` returns the correct type!

# Translating `length_of`

- `length_of` acts on arrays and strings
- These are both represented as pointers to `%array_type*` in LLVM
- Use Gep to read the length field of `%array_type`



# Translating the Lval $a[i]$

- Translate  $a$ , it should produce an LLVM operand of type `%array_type*`
- Use Gep to get a pointer to the filed contents
  - Treat this as a point type `i8*`
  - Bitcast it to a pointer of type `t*` where  $t$  is the translation of the Dolphin type  $A$  such that array  $a$  has type  $[A]$  in Dolphin (use `trans_type`)
- Use Gep to index into the obtained pointer of type `t*`
- **NOTE:** The `trans_lval` function should **return a pointer**; the `trans_expr` function should it appropriately

# Translating `nil`

- Translate `nil` to LLVM's `null`
- We update `trans_type` to support `TAst.Undetermined _`

```
let trans_type tp =  
  match tp with  
  | TAst.Determined (TAst.Void _) -> Ll.Void  
  | TAst.Determined (TAst.Int _) -> Ll.I64  
  | TAst.Determined (TAst.Bool _) -> Ll.I1  
  | TAst.Determined (TAst.Byte _) -> Ll.I8  
  | TAst.Determined (TAst.Str _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  | TAst.Determined (TAst.Array _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  | TAst.Undetermined _ -> Ll.Ptr Ll.I8  
  ...
```

# Translating Record Declarations

- We translate record declarations into a **named struct types** in LLVM --
- We translate lists

```
record list {head : int; tail : list;}
```

into

```
%dolphin_record_list = type { i64, %dolphin_record_list* }
```

- Use for `trans_type` translating types of fields

# Translating Types Including Records

Sometimes we need the type of records not wrapped in `Ll.Ptr`



```
let trans_type raw_records tp =  
  match tp with  
  | TAst.Determined (TAst.Void _) -> Ll.Void  
  | TAst.Determined (TAst.Int _) -> Ll.I64  
  | TAst.Determined (TAst.Bool _) -> Ll.I1  
  | TAst.Determined (TAst.Byte _) -> Ll.I8  
  | TAst.Determined (TAst.Str _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  | TAst.Determined (TAst.Array _) ->  
    Ll.Ptr (Ll.Namedt (Sym.symbol "array_type"))  
  | TAst.Undetermined _ -> Ll.Ptr Ll.I8  
  | TAst.Determined (TAst.Record {recordname = RecordName {sym; _}}) ->  
    let raw_res = Ll.Namedt (encode_record_name sym) in  
    if raw_records then raw_res else Ll.Ptr raw_res  
  | TAst.Determined TAst.ErrorType -> assert false
```

# Translating Field Access Lvals

- To translate  $a.f$ , first translate  $a$
- We know the Dolphin type  $A$  of  $a$  from semantic analysis
- Translate  $A$  into an LLVM type (raw record type)  $t$
- Figure out the index of field  $f$  in the LLVM type  $t$
- Use Gep to obtain a pointer to field  $f$

# Translating Record Creation

- Use the runtime function

```
void *allocate_record(int32_t size)
```

- Compute the size of the record type in LLVM
- Call the `allocate_record` function to allocate space for the array
- The returned pointer has type `i8*` on the LLVM side
- Bitcast the obtained pointer into a pointer of the record
- For each field
  - Translate the initialization expression for that field
  - Compute the address of the field using Gep (as on the previous slide)
  - Write the computed initialization value into the field

# Translating comparison

- Check the underlying type
  - if it is two strings being compared, use the runtime function for comparing strings
  - If it is not strings
    - If the comparison is `<`, `<=`, `>`, or `>=`
      - The objects being compared are integers; use LLVM's `icmp` instruction
    - Otherwise, the comparison is `==` or `!=`
      - Use LLVM's `icmp` instruction
      - The type of objects being compared (needed by `icmp`) can be obtained by calling `trans_type`

# Stdlib Types and their Translation

- The standard library features types that should be represented in Dolphin programs, e.g., the `stream` type
- These types are represented (hardcoded) as **empty records** in Dolphin, and translated into LLVM as such  
`%dolphin_record_stream = type {}`
- **Note:** Semantic analysis should prohibit creating instance of these records, i.e., the following should be rejected with an error:  
`var z = new stream {};`