

# CLASS ROOM LOG -7

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# Topics covered

- ❑ Aliasing
- ❑ Tables as sets
  - ❑ Set operators (UNION, INTERSECT, EXCEPT)
- ❑ UNION ALL, INTERSECT ALL, EXCEPT ALL

## ❑ Substring Pattern matching Operator

- ❑ LIKE operator

- ❑ \_ operator

## ❑ Arithmetic operators

- ❑ + , - , \* , /

## ❑ BETWEEN, ORDER BY

## ❑ Nested queries

- ❑ IN Operator

- ❑ ALL Operator

- ❑ IS NULL Vs IS NOT NULL; EXISTS Vs NOT EXISTS

- ❑ JOIN

  - ❑ LEFT OUTER JOIN, RIGHT OUTER JOIN, NATURAL JOIN

- ❑ AGGREGATE FUNCTIONS

  - ❑ SUM, MIN, MAX, AVG, COUNT

- ❑ GROUP BY, HAVING clauses

- ❑ INSERT INTO, DELETE, UPDATE

# Aliasing

- ❑ It is used when we want to have multiple references to the same relation/column within a query

# Aliasing Syntax

## SQL Alias Syntax for Columns

```
SELECT column_name AS alias_name  
FROM table_name  
WHERE [condition];
```

## SQL Alias Syntax for Tables

```
SELECT column1, column2....  
FROM table_name AS alias_name  
WHERE [condition];
```

Consider the relation  
**EMPLOYEE**

**EMPLOYEE**(FNAME, LNAME,  
SSN, SUPER\_SSN, DNO)

SUPER\_SSN is FK of SSN

For each employee, retrieve the  
employee's last name and last name of  
his or her immediate supervisor

```
SELECT E.LNAME, S.LNAME
```

```
FROM EMPLOYEE AS E, EMPLOYEE  
AS S
```

```
WHERE E.SUPER_SSN=S.SSN;
```

Here E,S are called Aliases

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# Why SQL considers tables as Multisets?

- ❑ Duplicate elimination is expensive
- ❑ User may want to see the duplicates as the result of the query
- ❑ When an aggregate function is applied to tuples, in most cases we do not want to eliminate duplicates



# Tables as sets in SQL

- ❑ SQL treats tables( relations ) as Multisets
- ❑ **DISTINCT** : Used to eliminate duplicates in a relation

# Syntax for DISTINCT

```
SELECT DISTINCT column_name[s]  
FROM table_name  
WHERE [condition]
```

Consider the relation  
**PRODUCT**

**PRODUCT**(PNAME,PRICE,  
CATEGORY,PRODUCT-ID,  
MANUFACTURER)

Select the distinct categories from  
**PRODUCT** relation

**SELECT DISTINCT CATEGORY**

**FROM PRODUCT**

---

# Set operators

- ❑ UNION
- ❑ INTERSECT
- ❑ EXCEPT

These all operators remove duplicates after the query is performed.

# UNION

## Syntax :

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition]
```

## UNION

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition];
```

# Example for UNION

SELECT A

FROM R

UNION

SELECT A

FROM S

R UNION S

A
1
2
3
4
5

S

A	B
5	a1
4	a2
3	a3

R

A
1
2

# INTERSECT

## Syntax :

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition]
```

## INTERSECT

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition]
```

# Example for INTERSECT

SELECT A  
FROM R  
INTERSECT  
SELECT A  
FROM S

R INTERSECT S

a1
----

S

A	<u>B</u>
a1	1
a4	2
a3	3

R

<u>A</u>
a1
a2



# EXCEPT

The SQL **EXCEPT** clause/operator is used to combine two **SELECT** statements and returns rows from the first **SELECT** statement that are not returned by the second **SELECT** statement.

## Syntax :

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition]
```

## EXCEPT

```
SELECT column1 [, column2 ]  
FROM table1 [, table2 ]  
[WHERE condition]
```

---

# Example for EXCEPT

SELECT A

FROM R

EXCEPT

SELECT A

FROM S

R EXCEPT S

A
2

S

A	B
1	a1
4	a2
3	a3

R

A
1
2

# UNION ALL, INTERSECT ALL, EXCEPT ALL

- ❑ They don't remove duplicates after these operations are performed

(a)

R	A
	a1
	a2
	a2
	a3

S	A
	a1
	a2
	a4
	a5

(b)

T	A
	a1
	a1
	a2
	a2
	a2
	a3
	a4
	a5

(c)

T	A
	a2
	a3

(d)

T	A
	a1
	a2

**FIGURE 8.5** The results of SQL multiset operations. (a) Two tables, R(A) and S(A). (b) R(A) UNION ALL S(A). (c) R(A) EXCEPT ALL S(A). (d) R(A) INTERSECT ALL S(A).

# LIKE Operator

- ❑ It is used as a Substring matching operator

## Syntax:

```
SELECT column_name(s)
```

```
FROM table_name
```

```
WHERE column_name LIKE pattern;
```

---

- ❑ Consider the relation PRODUCT

PRODUCT (PID, PNAME, PRICE,  
CATEGORY, MANUFACTURER)

Retrieve all PRODUCT details whose  
CATEGORY contains “ELEC”

SELECT \*

FROM PRODUCT

WHERE CATEGORY LIKE '%ELEC%';

# Wild card

- ❑ ‘\_’ replaces single character
- ❑ Consider relation  
**Customers**(CustID, CustName,  
Address, City, PostalCode, Country)

Select all customers with a City starting with any character, followed by "erlin"

```
SELECT *
```

```
FROM Customers
```

```
WHERE City LIKE '_erlin';
```

# ARITHMETIC OPERATORS

- ❑ ADDITION(+)
- ❑ SUBTRACTION(-)
- ❑ MULTIPLICATION(\*)
- ❑ DIVISION(/)
- ❑ These operators can be applied on numeric values or attributes with numeric domains

Consider relation

**PRODUCT** (PNAME, PRICE, PID, CATEGORY,  
MANUFACTURER)

Display price of all the products if  
each product price is given a 10  
percent hike.

```
SELECT 1.1*PRICE AS  
INCREASED_PRICE  
  
FROM PRODUCT
```

---



# BETWEEN

- ❑ The **BETWEEN** operator selects values within a range. The values can be numbers, dates etc...

## Syntax

```
SELECT column_name(s)
```

```
FROM table_name
```

```
WHERE column_name BETWEEN value1  
AND value2;
```

---

Consider relation

**EMPLOYEE**(ENAME, SALARY,  
DNO, SSN)

Retrieve all employees details in  
department 5 whose salary is between  
\$30,000 and \$40,000.

```
SELECT *  
FROM EMPLOYEE  
WHERE (SALARY BETWEEN 30000  
AND 40000) AND DNO =5;
```

---

# ORDER BY

- ❑ It is used to sort the result-set by one or more columns.

## Syntax

**SELECT** column-list

**FROM** table\_name

[**WHERE** condition]

[**ORDER BY** column1, column2, .. columnN] [**ASC** | **DESC**];

- ❑ **Order of selection :**

FROM -> Where -> SELECT -> ORDER BY

Consider relation

**PRODUCT** (PNAME, PID, PRICE,  
CATEGORY, MANUFACTURER)

Retrieve product-name, price, manufacturer details of PRODUCT relation ordered by price, name of the product

```
SELECT PNAME, PRICE, MANUFACTURER  
FROM    PRODUCT  
ORDER BY PRICE, PNAME
```

# NESTED QUERIES

- ❑ Dynamically created tables from **SELECT** queries can be used within other **SELECT** queries:
- ❑ In  $(v \text{ **IN** } V)$ , Operator **IN**, compares a value  $v$  with a set (or multiset) of values  $V$  and evaluates to **TRUE** if  $v$  is one of the elements in  $V$
- ❑  $(v > \text{ALL } V)$  returns **TRUE** if the value  $v$  is greater than all the values in the set (or multiset)  $V$ .

# NESTED QUERIES USING IN

## CONSIDER THE FOLLOWING RELATIONS

**Company**(name, city)

**Product**(pname, maker)

**Purchase**(id, product, buyer)

Where **Purchase.product** is FK to  
**Product.pname** and

**Product.maker** is FK to **Company.name**

Find city names of the companies that  
manufacture products bought by Jay

```
SELECT Company.city
```

```
FROM      Company
```

```
WHERE Company.name IN
```

```
( SELECT Product.maker
```

```
FROM Purchase, Product
```

```
WHERE
```

```
Product.pname=Purchase.product AND Purchase  
.buyer = 'Jay');
```

---

# EXAMPLE OF NESTED QUERIES USING ALL

Consider relation **EMPLOYEE**(ENAME, SALARY, DNO, SSN) and find the names of employees whose salary is greater than the salary of all the employees in department 5

```
SELECT E.ENAME
FROM EMPLOYEE AS E
WHERE E.SALARY > ALL (
SELECT S.SALARY
FROM EMPLOYEE AS S
WHERE S.DNO=5);
```

# Correlated Nested Queries

- ❑ Whenever a condition in the WHERE clause of a nested query references some attribute of a relation declared in the outer query, the two queries are said to be correlated



# EXISTS, NOT EXISTS, NULL, NOT NULL

- ❑ EXISTS is true if result of the select operation is not empty
- ❑ EXISTS is false if result of the selection operation is empty
- ❑ NOT EXISTS is true if result of the selection operation is empty
- ❑ NOT EXISTS is false if result of the selection operation is not empty
- ❑ NOT NULL is true when result of the selection operation is not null
- ❑ NOT NULL is false when result of the selection operation is null
- ❑ NULL is true when result of the selection operation is null
- ❑ NULL is false when result of the selection operation is not null

# Example of correlated Nested queries

1. Consider relation **EMPLOYEE**(ENAME, SALARY, DNO, SSN),  
**DEPARTMENT**(DEPT\_ID, DNAME, MANAGED-BY, LOCATIONS)

Write a query to find the departments which do not have employees at all?

```
SELECT DEPT_ID,DNAME  
FROM DEPARTMENT AS D  
WHERE NOT EXISTS(  
SELECT *  
FROM EMPLOYEES AS E  
WHERE E.DNO = D.DEPT_ID)
```

# EXISTS

1) Consider relations `customers(customer_id, name, favourite-website)`, `orders(order_id, customer_id, order_date)`. Find all of the records from the customers table where there is at least one record in the orders table with the same `customer_id`.

```
SELECT *  
FROM customers  
WHERE EXISTS  
  (SELECT *  
   FROM orders  
   WHERE customers.customer_id = orders.customer_id);
```

# JOIN

- ❑ It performs join operation in the **FROM** clause of the query and outputs join table

Example: Consider relations

**EMPLOYEE**(ENAME, SALARY,  
ADDRESS, DNO, SSN),

**DEPARTMENT**(DNUMBER, DNAME,  
MANAGED-BY, LOCATIONS)

Here DNO is FK to DNUMBER

Retrieve the name and address of every employee who works for the 'Research' department

```
SELECT NAME, ADDRESS  
FROM (EMPLOYEE JOIN DEPARTMENT  
ON DNO=DNUMBER)  
WHERE DNAME='Research';
```

---

# NATURAL JOIN

CONSIDER RELATIONS

**foods**(ITEM\_ID, ITEM\_NAME,  
ITEM\_UNIT, COMPANY\_ID)

**company**( COMPANY\_ID,  
COMPANY\_NAME, COMPANY\_CITY)

**NATURAL JOIN ON foods,company**

```
SELECT *  
FROM foods NATURAL JOIN  
company;
```

---

## LEFT OUTER JOIN

It returns all rows from the left table, even if there are no matches in the right table

### Syntax

```
SELECT column_name(s)
```

```
FROM table1 LEFT OUTER JOIN table2 ON  
table1.column_name=table2.column_name;
```

## RIGHT OUTER JOIN

It returns all rows from the right table, even if there are no matches in the left table

### Syntax

```
SELECT column_name(s)
```

```
FROM table1 RIGHT OUTER JOIN table2 ON  
table1.column_name=table2.column_name;
```

**Example: Employee  $\Leftarrow$  Department**

**SELECT \***

**FROM Employee RIGHT OUTER**

**JOIN Department ON EmployeeID= Manager**

INPUT RELATIONS

Employee

Employee ID	Name	Department
E1	A	D1
E2	B	D2
E3	C	D1

Department

Dept ID	Dept Name	Manager
D1	X	E3
D2	Y	E2
D3	Z	null

# AGGREGATE OPERATORS

- ❑ COUNT, SUM, MAX, MIN, AVG
- ❑ COUNT returns the number of tuples or values as specified in a query.
- ❑ The functions SUM, MAX, MIN, and AVG are applied to a set or multiset of numeric values and return, respectively, the sum, maximum value, minimum value, and average (mean) of those values.



For example consider relation

EMPLOYEE							
FNAM	MINI	LNAM	<u>SSN</u>	BDATE	ADDRESS	SEX	SALARY

Find the sum of the salaries of all employees, the maximum salary, the minimum salary, and the average salary.

```
SELECT SUM (SALARY),  
MAX (SALARY), MIN (SALARY),  
AVG(SALARY)  
FROM EMPLOYEE;
```

---

# Grouping: GROUP BY and HAVING Clauses

- ❑ **GROUP BY :** When we want aggregate functions to be applied to groups.
  - ❑ So, to mention grouping attributes we use GROUP BY
- ❑ **HAVING :** Used when we want to put condition on aggregation operation

## GROUP BY and HAVING Clauses

Consider relation

**EMPLOYEE**(ENAME, SALARY, DNO, SSN) and compute each department, count the number of employees and their average salaries

```
SELECT DNO, COUNT(*), AVG(SALARY)
FROM EMPLOYEE GROUP BY DNO;
```

Consider relation

**EMPLOYEE**(ENAME, SALARY, DNO, SSN) and return all departments having more than twenty employees, and show the number of employees and their average salary.

```
SELECT DNO, COUNT(*), AVG(SALARY)
FROM EMPLOYEE
GROUP BY DNO
HAVING COUNT(*) > 20;
```

# INSERT INTO, DELETE, AND UPDATE

```
CREATE TABLE EMPLOYEE
( FNAME          VARCHAR(15)      NOT NULL ,
  MINIT          CHAR            ,
  LNAME          VARCHAR(15)      NOT NULL ,
  SSN            CHAR(9)         NOT NULL ,
  BDATE         DATE             ,
  ADDRESS        VARCHAR(30)     ,
  SEX            CHAR            ,
  SALARY         DECIMAL(10,2)   ,
  SUPERSSN       CHAR(9)        ,
  DNO            INT             NOT NULL ,
  PRIMARY KEY (SSN) ,
  FOREIGN KEY (SUPERSSN) REFERENCES EMPLOYEE(SSN) ,
  FOREIGN KEY (DNO) REFERENCES DEPARTMENT(DNUMBER) ) ;

CREATE TABLE DEPARTMENT
( DNAME          VARCHAR(15)      NOT NULL ,
  DNUMBER        INT             NOT NULL ,
  MGRSSN         CHAR(9)         NOT NULL ,
  MGRSTARTDATE   DATE            ,
  PRIMARY KEY (DNUMBER) ,
  UNIQUE (DNAME) ,
  FOREIGN KEY (MGRSSN) REFERENCES EMPLOYEE(SSN) ) ;

CREATE TABLE DEPT_LOCATIONS
( DNUMBER        INT             NOT NULL ,
  DLOCATION       VARCHAR(15)     NOT NULL ,
  PRIMARY KEY (DNUMBER, DLOCATION) ,
  FOREIGN KEY (DNUMBER) REFERENCES DEPARTMENT(DNUMBER) ) ;

CREATE TABLE PROJECT
( PNAME          VARCHAR(15)      NOT NULL ,
  PNUMBER        INT             NOT NULL ,
  PLOCATION       VARCHAR(15)     ,
  DNUM           INT             NOT NULL ,
  PRIMARY KEY (PNUMBER) ,
  UNIQUE (PNAME) ,
  FOREIGN KEY (DNUM) REFERENCES DEPARTMENT(DNUMBER) ) ;

CREATE TABLE WORKS_ON
( ESSN           CHAR(9)         NOT NULL ,
  PNO            INT             NOT NULL ,
  HOURS          DECIMAL(3,1)    NOT NULL ,
  PRIMARY KEY (ESSN, PNO) ,
  FOREIGN KEY (ESSN) REFERENCES EMPLOYEE(SSN) ,
  FOREIGN KEY (PNO) REFERENCES PROJECT(PNUMBER) ) ;

CREATE TABLE DEPENDENT
( ESSN           CHAR(9)         NOT NULL ,
  DEPENDENT_NAME VARCHAR(15)     NOT NULL ,
  SEX            CHAR            ,
  BDATE         DATE            ,
  RELATIONSHIP   VARCHAR(8)     ,
  PRIMARY KEY (ESSN, DEPENDENT_NAME) ,
  FOREIGN KEY (ESSN) REFERENCES EMPLOYEE(SSN) ) ;
```

# INSERT INTO

- ❑ By using Insert data tuples can be added to tables
- ❑ Constraints that can be violated
  - ❑ Domain constraint
  - ❑ Key constraint
  - ❑ Referential integrity constraint
  - ❑ Entity integrity constraint
- ❑ Insertion fails if any of the above constraints is violated

1) Insert a tuple into EMPLOYEE table

```
INSERT INTO EMPLOYEE  
VALUES ('1002', 002, 'Bharath Kumar', 'M',  
'9-5-1973', 'Rajajinagar Bangalore 10',  
300000, 007, 5);
```

2) It allows the user to specify explicit attribute names that correspond to the values provided in the INSERT command. Remaining attribute values are set to their default/NULL

```
INSERT INTO EMPLOYEE (FNAME,  
LNAME, DNO, SSN)  
VALUES ('Richard', 'Marini', 4, '653298653');
```

# DELETE

- ❑ The DELETE command removes tuples from a relation. It includes a WHERE to select the tuples to be deleted
- ❑ Referential Integrity may get violated  
-delete operation is rejected
- ❑ In table2 put this constraint  
foreign-key attr-name references Table1name  
on delete{cascade/restrict}

DELETE FROM EMPLOYEE

WHERE LNAME = 'BROWN'

---

# UPDATE

- ❑ By using Insert data tuples can be added to tables
- ❑ Constraints that can be violated
  - ❑ Domain constraint
  - ❑ Key constraint
  - ❑ Referential integrity constraint
  - ❑ Entity integrity constraint
- ❑ Updation fails if any of the above constraints is violated

Change the location and controlling department number of project number 10 to 'BANGALORE' and 5, respectively

```
UPDATE PROJECT
```

```
SET PLOCATION = 'BANGALORE', DNUM = 5
```

```
WHERE PNUMBER=10;
```

---

# REFERENCES

1. SQL Tutorial | W3Schools. (n.d.). from <http://www.w3schools.com/SQL/default.asp>
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3. Elmasri, R., & Navathe, S. B. (2000). Fundamentals of database systems. Reading: Wesley.



**Thank you**