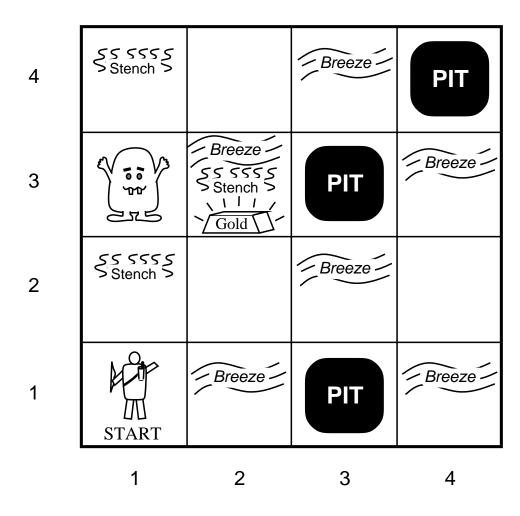
Artificial Intelligence Knowledge Representation and Reasoning (Logic)

Nilsson - Chapters 13, 14, 15

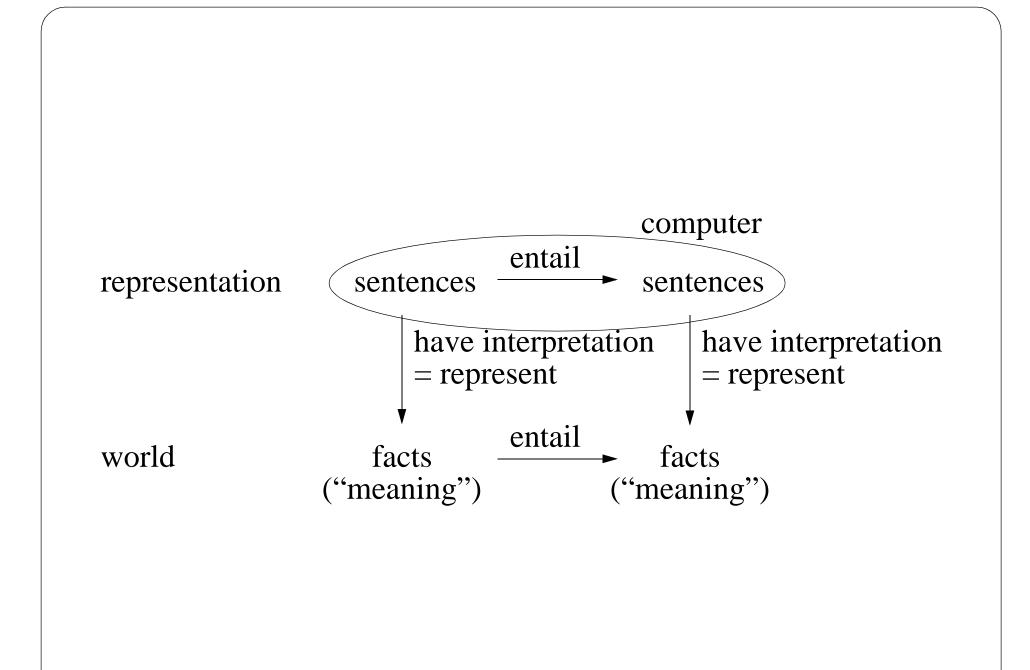
Russell and Norvig - Chapters 7, 8

"wumpus world"



knowledge representation languages should be expressive, concise, unambiguous, independent of context, effective

Agents are given knowledge about the world. Knowledge Representation: How can facts about the world be represented? Reasoning: How can the agent infer new facts from the given ones?



arithmetic

syntax

Is "1 > 2" a well-formed formula? Is "x + *y 5 >" a well-formed formula?

semantics

When is "x + 2 = 5" true?

propositional logic = sentences represent whole propositions

"2 is prime." P
"I ate breakfast today." Q

syntax = how a sentence looks like

Sentence -> AtomicSentence | ComplexSentence

AtomicSentence -> T(RUE) | F(ALSE) | Symbols

Symbols \rightarrow P | Q | R | ...

negation

ComplexSentence -> (Sentence) | NOT Sentence | Sentence Connective Sentence

Connective -> AND | OR | IMPLIES | EQUIV(ALENT) conjunction disjunction implication

equivalence





Precedence: NOT AND OR IMPLIES EQUIVALENT

syntax = how a sentence looks like

The constants T(RUE) and F(ALSE) are sentences.

The symbols P, Q, R, ... are sentences.

If S is a sentence, then (S) is a sentence.

If S is a sentence, then NOT S is a sentence.

If S1 and S2 are sentences, then S1 AND S2 is a sentence.

If S1 and S2 are sentences, then S1 OR S2 is a sentence.

If S1 and S2 are sentences, then S1 IMPLIES S2 is a sentence.

If S1 and S2 are sentences, then S1 EQUIV S2 is a sentence.

Examples: P, T, P IMPLIES NOT T, P AND (Q OR S)

semantics = what a sentence means

interpretation: assigns each symbol a truth value, either t(rue) or f(alse)

the truth value of T(RUE) is t(rue) the truth value of F(ALSE) is f(alse)

truth tables ("compositional semantics") the meaning of a sentence is a function of the meaning of its parts

A	В	NOT A	A AND B	A OR B	A IMPLIES B	A EQUIV B
t	t	f	t	t	t	t
t	f		f	t	f	\mathbf{f}
f	t	t	f	t	t	\mathbf{f}
f	f		f	f	t	t

examples

either I go to the movies or I go swimming (inclusive versus exclusive OR - our OR is inclusive)

A	В	A OR B	A XOR B
		(inclusive OR)	(exclusive OR)
t	t	t	f
t	f	t	t
f	t	t	t
f	f	\mathbf{f}	f

Read "A IMPLIES B" as "if A then B"

Read "A EQUIV B" as "if and only if A then B" or "A and B have the same truth value"

examples

2 is prime implies that 2 is even (implication does not imply causality)

2 is odd implies that 3 is even (false implies everything)

rewrite rules

A and NOT NOT A have identical truth tables.

A	A	A	NOT NOT A
t	t	t	t
f	f	f	f

Therefore, we can say A instead of NOT NOT A, and NOT NOT A instead of A without changing the truth value of any propositional ssentence.

For example, A AND B and A AND NOT NOT B have identical truth tables and are therefore equivalent.

A	В	A AND B	A AND NOT NOT B
t	t	t	t
t	f	f	f
f	t	\mathbf{f}	f
f	f	f	f

rewrite rules

NOT NOT A = A

NOT (A AND B) = NOT A OR NOT B

NOT (A OR B) = NOT A AND NOT B

A AND (B OR C) = (A AND B) OR (A AND C)

A IMPLIES B = NOT A OR B

A EQUIV B = (A IMPLIES B) AND (B IMPLIES A)

and many more

deriving formulas

```
A B A XOR B (= exclusive OR)

t t f A AND B

t f t A AND NOT B

f t t NOT A AND B

f f f NOT A AND NOT B
```

(A AND NOT B) OR (NOT A AND B)

problem with propositional logic

All cats are mammals. Bob is a cat.

Bob is a mammal.

first-order logic (= FOL)
= sentences refer to objects, their relationships, and their properties
= superset of propositional logic

constants objects Bob, CS3361, October 13 functions MotherOf, ColorOf predicates relations properties IsBlue

syntax = how a sentence looks like

Sentence -> AtomicSentence | Sentence Connective Sentence |

Ouantifier Variable Sentence | NOT Sentence

(Sentence)

AtomicSentence -> T(RUE) | F(ALSE) |

Predicate (Term, ..., Term) | **Predicate** |

Term = Term

Term -> Constant | Variable |

Function (Term, ..., Term)

Connective -> AND (\land) | OR (\lor) | IMPLIES (\Rightarrow) | EQUIV (\Leftrightarrow)

Quantifier -> FORALL (\forall) | EXISTS (\exists)

Constant -> $A \mid ... \mid Z \mid Bob \mid CS3361 \mid ...$

Variable -> $a \mid ... \mid z$

Predicate -> IsMotherOf | IsBlue | ...

Function -> MotherOf | ColorOf | ...

Precedence: NOT AND OR IMPLIES EQUIVALENT

sentences refer to facts terms refer to objects

predicated can be understood as sets

= (equal) is a predicate with a predefined meaning

functions are just a convenient way of using predicates MotherOf(y) = x iff IsMotherOf(x,y) universal quantifier ("for all"): FORALL (∀)
existential quantifier ("there exists"): EXISTS (∃)
every variable must be introduced by a quantifier before it is used

Everything is blue. FORALL x IsBlue(x)

=

IsBlue(Sven) AND IsBlue(Tom's Mailbox) AND IsBlue(Sven's TV) AND ...

At least one thing is blue EXISTS x IsBlue(x)

=

IsBlue(Sven) OR IsBlue(Tom's Mailbox) OR IsBlue(Sven's TV) OR ...

universal quantifier ("for all"): FORALL (∀) existential quantifier ("there exists"): EXISTS (∃)

Not everything is blue. NOT FORALL x IsBlue(x)

=

NOT(IsBlue(Sven) AND IsBlue(Tom's Mailbox) AND ...)

=

NOT(IsBlue(Sven)) OR NOT(IsBlue(Tom's Mailbox)) OR ...

=

EXISTS x NOT IsBlue(x)

Everything is blue.

FORALL x IsBlue(x)

= NOT NOT FORALL x IsBlue(s)

= NOT EXISTS x NOT IsBlue(s)

universal quantifier ("for all"): FORALL (∀) existential quantifier ("there exists"): EXISTS (∃)

Nothing is blue. FORALL x NOT IsBlue(x)

_

NOT(IsBlue(Sven)) AND NOT(IsBlue(Tom's Mailbox)) AND ...

=

NOT(IsBlue(Sven) OR IsBlue(Tom's Mailbox) OR ...)

=

NOT EXISTS x IsBlue(x)

At least one thing is blue.

EXISTS x IsBlue(x)

= NOT NOT EXISTS x IsBlue(x)

= NOT FORALL x NOT IsBlue(x)

rewrite rules

Make sure that you convince yourself of the following rewrite rules (both that they make intuitively sense and that they are correct)

```
FORALL x P(x) = NOT EXISTS x NOT P(x)

EXISTS x P(x) = NOT FORALL x NOT P(x)

FORALL x FORALL y P(x,y) = FORALL y FORALL x P(x,y)

EXISTS x EXISTS y P(x,y) = EXISTS y EXISTS x P(x,y)

NOT FORALL x (P(x) IMPLIES Q(x)) = EXISTS x (P(x) AND NOT Q(x))

NOT EXISTS x (P(x) AND Q(x)) = FORALL x (P(x) IMPLIES NOT Q(x))
```

and many more

universal quantifier ("for all"): FORALL (∀) existential quantifier ("there exists"): EXISTS (∃)

Every mailbox is blue FORALL x (IsMailbox(x) IMPLIES IsBlue(x))

Everything is a blue mailbox (probably not what you mean) FORALL x (IsMailbox(x) AND IsBlue(x))

At least one mailbox is blue EXISTS x (IsMailbox(x) AND Blue(x))

At least one thing is not a mailbox or blue (probably not what you mean) EXISTS x (IsMailbox(x) IMPLIES IsBlue(x))

universal quantifier ("for all"): FORALL (∀) existential quantifier ("there exists"): EXISTS (∃)

Everyone has at least one mother FORALL y EXISTS x IsMotherOf(x,y)

Everyone shares at least one mother (= has the same mother). EXISTS x FORALL y IsMotherOf(x,y)

The translations above are actually not quite correct.

The correct translation is

Everyone has at least one mother

FORALL y (IsPerson(y) IMPLIES EXISTS x IsMotherOf(x,y))

It is unnecessary (but not incorrect) to say

DRALL y (IsPerson(y) IMPLIES EXISTS x (IsPerson(x) AND IsMotherOf(x,y)

examples

Sue is Bill's mother.

Sue is Bill's and Tom's mother.

Bill is smart, and so is Tom.

Either Bill or Tom is smart, but not both.

Sue has kids (= Sue has at least one kid).

Sue has no kids.

Sue has one kid.

Sue has at most one kid.

Sue has at least two kids.

All of Sue's kids are smart.

At least one of Sue's kids is smart.

None of Sue's kids are smart, neither is Sue.

Sue loves Bill.

Noone (= no person) loves Bill.

Everybody (= every person) loves somebody. - ambiguos

Somebody (= some person) loves everybody. - ambiguous

Everyone (= every person) has exactly one mother.

If two people have the same mother

then they are either both smart or neither of them is smart.

semantics = what a sentence means

propositional logic

interpretation (= "world" = model)
assigns
each propositional symbol a truth value
(every row in a truth table corresponds to one interpretation)

first-order logic

interpretation (= "world" = "model"):
 assigns
each constant a meaning (= an object),
 each function a meaning,
 each predicate a meaning.

sentence is

valid (= a tautology)
if it holds for all interpretations
(e.g. all rows of the truth table are true for a propositional sentence)

satisfiable
if it holds for at least one interpretation
(e.g. at least one row of the truth table is true for a propositional sentence)

unsatisfiable (= a contradiction)
if it holds for no interpretation
(e.g. all rows of the truth table are false for a propositional sentence)

valid? satisfiable? unsatisfiable?

2 is prime.

2 is prime or 2 is not prime.

FORALL x Eats(Lion,x) "the lion eats everything"

AND

FORALL x FORALL y (EATS(x,y) IMPLIES Hunts(x,y)) if x eats y then x hunts y

AND

FORALL x FORALL y (Hunts(x,y) IMPLIES NOT Hunts(y,x)) if x hunts y, then y doesn't hunt x

(1) FORALL x Eats(Lion,x) "the lion eats everything"

- (2) FORALL x FORALL y (EATS(x,y) IMPLIES Hunts(x,y)) if x eats y then x hunts y
- (3) FORALL x FORALL y (Hunts(x,y) IMPLIES NOT Hunts(y,x)) if x hunts y, then y doesn't hunt x

This knowledge base is unsatisfiable.

- (1) entails: (4) Eats(Lion,Lion) "the lion eats itself"
- (2) and (4) entail: (5) Hunts(Lion, Lion) "the lion hunts itself"
- (3) and (5) entail: (6) NOT Hunts(Lion, Lion) "the lion does not hunt itself

However, (5) and (6) cannot be true at the same time.

sentence 4
sentence 2
sentence 3

sentence 1 AND sentence 2 AND sentence 3 AND sentence 4

knowledge base = sentence (read: entails)

= whenever an interpretation makes the knowledge base true,
 it also makes the sentence true as well
 (advantage: the computer does not need to know the interpretation)

example:

does "heads, I win; tails, you lose" entail "I win"

(note: background knowledge is important!)

IN THE FOLLOWING, WE ONLY DEAL WITH PROPOSITIONAL LOGIC

for first order logic, truth tables do not work since there are infinitely many interpretations for first order logic, resolution might not terminate if a sentence is not entailed by the knowledge base (= resolution is semi-decidable for first order logic)

method 1: truth table

HE = heads comes up, TA = tails comes up, I = IW win, UL = you lose

heads, I win tails, you lose TA IMPLIES IW either heads or tails (background knowledge) HE EQUIV NOT TA IWPLIES UL I win or you lose (background knowledge) IW EQUIV UL

HE	TA	IW	UL	HE IMPLIES IW	TA IMPLIES UL	HE EQUIV NOT TA	IW EQUIV UI
t	t	t	t	t	t	f	t
t	t	t	f	t	f	f	f
t	t	\mathbf{f}	t	f	t	f	f
t	t	<u>f</u>	f	f	f	f	t
t	f	t	t	t	t	t	t
t	f	t	f	t	t	t	$\frac{}{}$
t	f	f	t	f	t	t	f
t	f	<u>f</u>	f	f	t	t	t
f	t	t	t	t	t	t	t
f	t	t	f	t	f	t	f
f	t	f	t	t	t	t	$f \mid$
f	t	f	f	t	f	t	t
f	f	t	t	t	t	f	t
f	f	t	f	t	t	f	$f \mid$
f	f	f	t	t	t	f	$f \mid$
f	f	f	f	t	t	f	t
each row corresponds to one interpretation							

inference procedures

knowledge base ⊢ sentence = the inference procedure can infer the sentence from the knowledge base

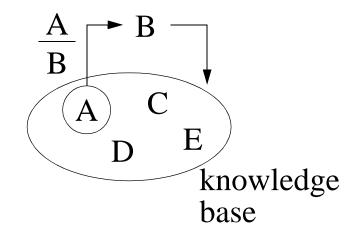
the inference procedure is **sound**if knowledge base ⊨ sentence
whenever knowledge base ⊢ sentence

logical inference (deduction)

the inference procedure is **complete**if knowledge base ⊢ sentence
whenever knowledge base ⊨ sentence

inference procedures = the repeated application of inference rules

$$\begin{array}{c} A \vdash B \\ \hline A \\ \hline B \end{array}$$



local reasoning why does it work? monotonicity!

in propositional logic it holds that $\frac{\text{knowledge base that includes A}}{B}$ if $\frac{A}{B}$

method 2: resolution

proof by contradiction

we want to show:

"heads, I win; tails, you lose" entails "I win;" in other words, whenever an interpretation makes "heads, I win; tails, you lose" true, it also makes "I win" true

instead, we show:

there is no interpretation that makes "heads, I win; tails, you lose" and "NOT(I win)" true at the same time; in other words "heads, I win; tails, you lose" and "NOT(I win)" entails FALSE.

resolution

P OR Q, NOT Q OR R P OR R

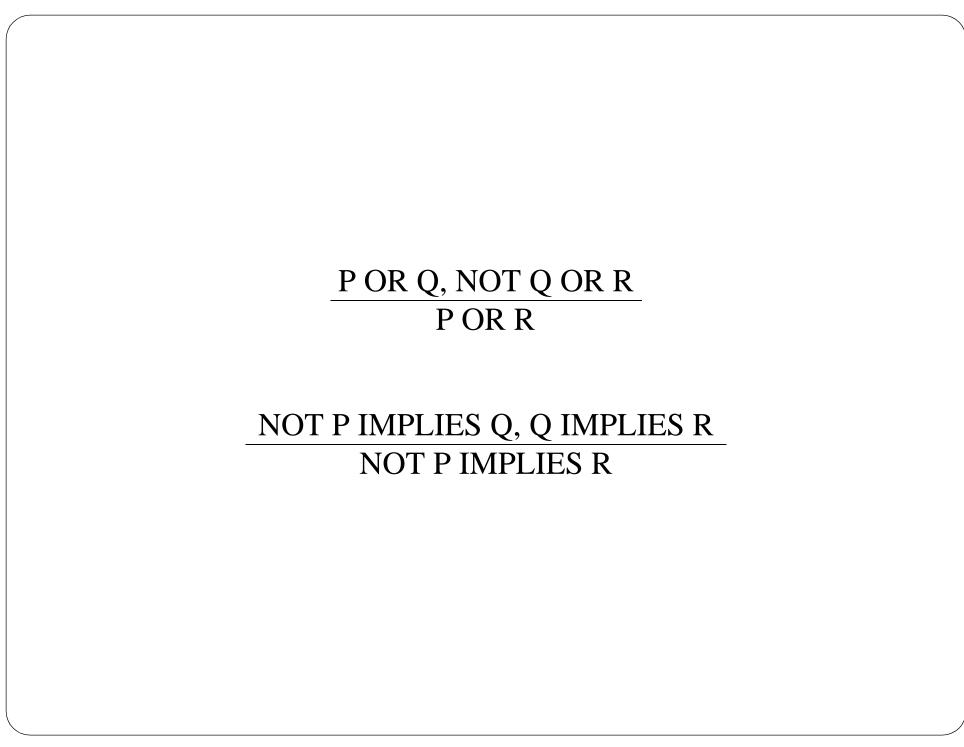
(read: "(P OR Q) AND (NOT Q OR R)" entails "P OR R")

P and R can be arbitrary sentences afterwards we remove duplicates from the resulting disjunct

examples:

A OR C OR NOT E, B OR NOT C OR NOT D OR F A OR B OR NOT D OR NOT E OR F

A OR B, NOT A OR B



proof by contradiction

we want to show:

"heads, I win; tails, you lose" entails "I win;" in other words, whenever an interpretation makes "heads, I win; tails, you lose" true, it also makes "I win" true

instead, we show:

there is no interpretation that makes "heads, I win; tails, you lose" and "NOT(I win)" true at the same time; in other words "heads, I win; tails, you lose" and "NOT(I win)" entails FALSE.

to do that, we add the negated sentence "NOT(I win)" to the knowledge base "heads, I win; tails, you lose" and use resolution to derive false ("empty") - if we can do that, "heads, I win; tails you lose" entails "I win."

Add the negated sentence to the knowledge base.

For each sentence in the knowledge base, do:

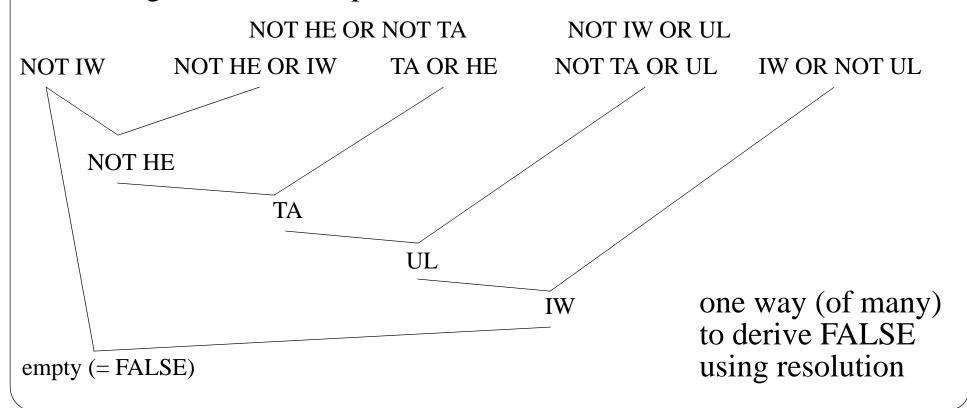
- 1. Eliminate EQUIV rewrite A EQUIV B as A IMPLIES B AND B IMPLIES A
- 2. Eliminate IMPLIES rewrite A IMPLIES B as NOT A OR B
- 3. Move NOT inward rewrite NOT(A AND B) as NOT A OR NOT B rewrite NOT(A OR B) as NOT A AND NOT B rewrite NOT(NOT A) as A
- 3. Distribute AND over OR rewrite A AND B OR C as (A OR C) AND (B OR C) rewrite A OR B AND C as (A OR B) AND (A OR C)
- 4. Flatten nested conjuncts and disjuncts rewrite A AND (B AND C) as A AND B AND C rewrite (A AND B) AND C as A AND B AND C rewrite A OR (B OR C) as A OR B OR C rewrite (A OR B) OR C as A OR B OR C
- 5. replace the original sentence with all of the obtained disjuncts example: for A OR B AND C, add both A OR B and C

HE = heads comes up, TA = tails comes up, IW = I win, UL = you lose

heads, I win tails, you lose either heads or tails (background knowledge) I win or you lose (background knowledge) I don't win (negated sentence)

HE IMPLIES IW
TA IMPLIES UL
HE EQUIV NOT TA
IW EQUIV UL
NOT IW

knowledge base after steps 1-5:



resolution is sound and complete

If a sentence is entailed by the knowledge base, resolution will eventually produce FALSE.

This requires a good search strategy.

There are several heuristics for good search strategies out there, including resolving with unit clauses whenever possible (to reduce the length of existing clauses) or resolving with clauses derived from the sentence to be shown to be entailed by the knowledge base (because FALSE cannot be produced from a satisfiable knowledge base alone).

If a sentence is not entailed by the knowledge base, then resolution will eventually not be able to produce new clauses, without having produced FALSE at that time.

(This statement is true only for propositional logic.)

This requires an exhaustive enumeration of all producable clauses and thus can take a long time.

resolution is sound and complete

think about this:

why can't we just use resolution to derive the sentence we are interested in (instead of adding its negation and derive false)?

answer

the empty knowledge base (KB) entails "P OR NOT P" but resolution cannot be used to derive "P OR NOT P" from the empty KB

resolution is sound and complete

think about this:

why can't we just use modus ponens instead of resolution to derive the sentence we are interested in?

modus ponens

P, P IMPLIES Q

Q

it rains, it rains implies the grass is wet the grass is wet

answer
"P IMPLIES Q" and "NOT P IMPLIES Q" entails "Q"
but modus ponens cannot show that