19CSE205 – Program Reasoning

Jevitha KP Lecture 11, 12, 13 – Forward vs Backward Reasoning, Weakest Preconditions – Conditionals



Credits

- Adapted from :
 - Dr. Bharat Jayaraman, University of Buffalo, CSE449-459 Software verification course, Spring 2020.



Sample Code

- x = 17;
- y = 42;
- z = x + y;
- Write the facts that are true at every point in the program

- { true }
- x = 17;
- y = 42;
- z = x + y;

- { true }
- x = 17;
- $\{ x = 17 \}$
- y = 42;
- z = x + y;

```
{ true }
x = 17;
{ x = 17 }
y = 42;
{ x = 17 \( \Lambda \) y = 42 }
z = x+y;
```

```
• { true }
• x = 17;
• \{ x = 17 \}
• y = 42;
• \{ x = 17 \land y = 42 \}
• z = x + y;
• \{ x = 17 \land y = 42 \land z = 59 \}
```

- An assertion is a logical formula inserted at some point in a program.
- It is presumed to hold true at that point in the program.



Precondition and Postcondition as Assertions

- A precondition is an assertion inserted prior to execution
- A postcondition is an assertion inserted after execution.
- {true} is the precondition
- $\{x = 17 \land y = 42 \land z = 59 \}$ is the postcondition.
- All other assertions are called intermediate assertions.
- They serve as steps as you reason between precondition and postcondition, Eg: like the intermediate steps in a math problem to solution.



Forward reasoning

- Previous example forward reasoning.
- Simulated the execution of the program, considering each statement in the order they would be executed.
- Disadvantage assertions may accumulate a lot of irrelevant facts as you move through the program.
- Don't know which parts of the assertions will come in handy to prove something later and which parts won't.
- End up listing everything we know about the program.



Backward Reasoning

- When we write a block of code, we have a clear idea of what's supposed to be true after it executes (ie) we know the postcondition already,
- To prove that the expected postcondition will indeed hold true given the appropriate precondition.
- For this reason, backward reasoning is often more useful than forward reasoning - Weakest precondition



Backward Reasoning Process

- Push the postcondition up through the statements to determine the precondition.
- Start by writing down the postcondition you want at the end of the block.
- Then look at the last statement in the block and ask, "For the postcondition to be true after this statement, what must be true before it?"



Backward Reasoning Process

- Keep going until you've reached the top of the statement list.
- Whatever must be true before the first statement is the precondition.
- It is guaranteed that if this precondition is satisfied before the block of code is executed, then the postcondition will be satisfied afterward.



Example

- x = y;
- x = x + 1;
- $\{ x > 0 \}$

Example

- x = y;
- $\{x + 1 > 0\}$
- x = x + 1;
- $\{x > 0\}$

Example

- $\{y + 1 > 0\}$
- x = y;
- $\{x + 1 > 0\}$
- x = x + 1;
- $\{x > 0\}$

- False case:
- Eg: if y = -1 if precondition is false
- x = -1 (x = y)
- x = 0 (x = x + 1)
- x=0 post condition (x>0) is also false

- True case:
- Eg: if y = 2 if precondition is true
- x = 2 (x = y)
- x = 3 (x = x + 1)
- x= 3 post condition (x>0) is also true

Weakest Preconditions for conditional statements



- Statement : if (B) S1 else S2
- How do we define: WP(if (B) S1 else S2, O)
 - If-Part: _____
 - Else-Part: _____

- Statement : if (B) S1 else S2
- Defining WP: WP(if (B) S1 else S2, O)
 - If-Part: wp(S1,O)
 - Else-Part: wp(\$2,0)



- Statement : if (B) S1 else S2
- Defining WP: WP(if (B) S1 else S2, O)
 - If-Part: wp(S1,O)
 - Else-Part: wp(\$2,0)
- WP(if (B) S1 else S2, O) =



- Statement : if (B) S1 else S2
- Defining WP: WP(if (B) S1 else S2, O)
 - If-Part: wp(S1,O)
 - Else-Part: wp(S2,O)

```
    WP(if (B) S1 else S2, O) =
    B ==> wp(S1,O)
    &&
    not(B) ==> wp(S2,O)
```

- Statement : if (B) S1
- How do we define: WP(if (B) S1, O)
 - If block : _____
 - Else block : _____

- Statement : if (B) S1
- How do we define: WP(if (B) S1, O)
 - If block: wp(\$1,0)
 - Else block : O



- Statement : if (B) S1
- How do we define: WP(if (B) S1, O)
 - If block: wp(\$1,0)
 - Else block : O
- WP(if (B) S1 , O) =

&&

not(B) ==> O

Another definition of WP for Conditional Statements

```
• WP(if (B) S1 else S2, O) =
   B && wp($1,0)
   not(B) && wp(S2,O)
• WP(if (B) S1, O) =
   B && wp($1,0)
   not(B) && O
```

Checking Equivalence

- We can check equivalence between:
- B && wp(S1, O) | not(B) && wp(S2, O) and
- (B ==> wp(S1, O)) && (not(B) ==> wp(S2, O))
- Abbreviate:
- wp(S1, O) \rightarrow P
- wp(S2, O) \rightarrow Q

Checking Equivalence

- We can check equivalence between:
- B && wp(S1, O) || not(B) && wp(S2, O) and
- (B ==> wp(S1, O)) && (not(B) ==> wp(S2, O))
- Using Alt-Ergo, we can quickly check equivalence :

$$(B \&\& P) \parallel (not(B) \&\& Q) <==> (B ==>P) \&\& (not(B) ==> Q)$$

Checking Equivalence

```
1 logic p,q,b: prop
 3 goal a:
         (b and p) or (not(b) and q) <->
         (b \rightarrow p) and (not(b) \rightarrow q)
File "try-alt-ergo-file", line 4, characters 5-68: Valid (0.1060) (7 steps)
(goal a)
```

- Output condition : $x \ge 0$
- Input condition : i=0
- Program:

```
If(i \ge 0) then
```

$$x=i;$$

else

$$x = -i$$
;

Find WP and check I=>WP.

- O: $x \ge 0$; B: $i \ge 0$; S1: x = i; S2: x = -i;
- WP(if(B) S1 else S2, O)
- b ==> wp(S1,O) && not(b) -> wp(S2,O)
- WP(S1,O) = (x>=0)(x=i) = (i>=0)
- WP(S2,O) = $(x>=0) \{x=-i\} = (-i>=0) = (i<=0)$
- (i>=0) ==> (i>=0) && not(i>=0) ==> (i<=0)
- (i>=0) ==> (i>=0) && (i<0) ==> (i<=0)
- Alternative WP:
- (i>=0) && (i>=0) || (i<0)&&(i<=0)



- WP Ans:
- i > = 0 = > (i > = 0) & (i < 0) = > (i < = 0)

I==>WP

• i=0 ==> (i>=0) ==> (i>=0) && (i<0) ==> (i<=0)

Alternative I==> WP

• i=0 ==> (i>=0) && (i>=0) || (i<0)&&(i<=0)



Alt-ergo

```
goal a:

forall i: int.

(i = 0) ->

(i>=0 -> i>=0)

and

(i<0 -> i<=0)
```

```
2 goal a:
 3 forall i: int.
 5 (i = 0) ->
         ((i>=0) \rightarrow (i>=0)) and ((i<0) \rightarrow (i<=0))
File "try-alt-ergo-file", line 4, characters 1-72: Valid (0.1060) (1 steps)
(goal a)
```

Alt-ergo – Alternative WP

```
goal a:
forall i: int.
(i = 0) ->
   (i>=0 and i>=0)
   or
   (i<0) and i<=0)</pre>
```

```
1 goal a:
 2 forall i: int.
 4 (i = 0)
        ((i)=0) and (i)=0)
         ((i<0) \text{ and } (i<=0))
File "try-alt-ergo-file", line 2, characters 1-89: Valid (0.0960) (1 steps)
(goal a)
```

- If(x<5)
- x=x*x
- else
- x=x+1
- 0: x > = 9

- B => wp(S1,O) && not(B) => wp(S2,O)
- B: (x<5)
- WP(S1,O) = $x > = 9 \{x = x^*x\} = x^*x > = 9$
- WP(S2,O) = $x > = 9\{x = x + 1\} = x + 1 > = 9 = x > = 8$
- (x<5) => (x*x>=9) && not(x<5) => (x>=8)
- (x<5)=>(x*x>=9) && (x>=5) => (x>=8)

Example 2 – Alt-Ergo

- WP: (x<5)=>(x*x>=9) && (x>=5) => (x>=8)
- 1: x = 3
- I => WP
- (x=3) => (x<5) => (x*x>=9) && (x>=5) => (x>=8)
- Alt-ergo:
- logic x:int
- goal a:
- (x=3) ->
- (x<5 -> x*x>=9)
- ullet and
- (x>=5 -> x>=8)

```
1 logic x:int
    goal a:
    (x=3) ->
           (x<5) \rightarrow (x*x>=9)
           (x>=5) \rightarrow (x>=8)
File "try-alt-ergo-file", line 4, characters 1-69: Valid (0.1230) (2 steps)
(goal a)
```

Example 2 – Alt-Ergo

- Alt-ergo: Invalid
- logic x:int
- goal a:
- (x=2) ->
- (x<5 -> x*x>=9)
- and
- (x>=5 -> x>=8)

```
logic x:int
    goal a:
        (x=2) ->
              (x<5 -> x*x>=9)
              (x>=5 -> x>=8)
File "try-alt-ergo-file", line 4, characters 5-81: I don't know (0.1270) (3
steps) (goal a)
```

Example - 3

```
if(x!=0)
  z=x
else
z=x+1
• O: z>0
```

- B=> WP(s1,o) && not(B) => wp(s2,o)
- $Wp(s1,o) = z>0 \{z=x\} = x>0$
- Wp(s2,o) = z>0 {z=x+1} = x+1>0
- (x!=0)=> x>0 && not(x!=0) => x+1>0
- WP: (x!=0) => x>0 && (x=0) => x+1>0
- WP: (x!=0) && x>0 || x=0 && x>-1

Example 3 – Alt-ergo

- WP: (x!=0) => x>0 && (x=0)=> x+1>0
- I : x= 2
- Alt-ergo : (!= should be written as <> in Alt-Ergo)
- x=2 ->
- (x <> 0 -> x > 0)
- And
- (x=0 -> x+1>0)

```
logic x:int
     goal a:
          (x=2) \rightarrow
                      (x \diamond 0 \rightarrow x \diamond 0)
                      (x=0 \rightarrow x+1>0)
File "try—alt—ergo—file", line 4, characters 5—88: Valid (0.1180) (2 steps)
(goal a)
```

Example 3 – Alt-ergo

- WP: (x!=0) => x>0 && (x=0)=> x+1>0
- | : x= -1
- Alt-ergo : (!= should be written as <> in Alt-Ergo)
- x=-1 ->
- (x <> 0 -> x > 0)
- And
- (x=0 -> x+1>0)

```
1 logic x:int
    goal a:
    x=-1 ->
             (x < 0 \rightarrow x > 0)
             and
             (x=0 \rightarrow x+1>0)
File "try-alt-ergo-file", line 4, characters 1-63: I don't know (0.1160) (3
steps) (goal a)
```

Example 4

```
if(a==b)
b = 2*a+1
else
b = 2*a
O:b>1
```

- B=> WP(s1,o) && not(B) => wp(s2,o)
- Wp(s1,o) = b>1 {b=2*a+1} = 2*a+1>1
- $Wp(s2,o) = b>1 \{b=2*a\} = 2*a>1$
- (a=b)=> 2*a+1>1 && not(a=b) => 2*a>1
- WP: (a=b)=> 2*a>0 && not(a=b) => 2*a>1
- WP: (a=b) && 2*a>0 || not(a=b) && 2*a>1

Example 4 – Alt-Ergo

- I : a = 2 and b = -1
- (a=2 and b=-1) => ((a=b) && 2*a>0) || (not(a=b) && 2*a>1)
- Alt-ergo:
- logic a,b:int
- goal g1:
- (a=2 and b=-1) ->
- ((a=b) and 2*a>0)
- or
- ((a<>b) and 2*a>1)

```
1 logic a,b:int
  3 goal g1:
    (a=2 \text{ and } b=-1) \rightarrow
             ((a=b) \text{ and } 2*a>0)
             ((a > b) and 2*a > 1)
File "try-alt-ergo-file", line 4, characters 1-82: Valid (0.1320) (3 steps)
(goal g1)
```

Example 4 – Alt-Ergo

- I : a = -1 and b = 2 Invalid
- (a=-1 and b=2) => (a=b && 2*a>0) || (not(a=b) && 2*a>1)
- Alt-ergo:
- logic a,b:int
- goal g1:
- (a=2 and b=-1) ->
- ((a=b) and 2*a>0)
- or
- ((a<>b) and 2*a>1)

```
logic a,b:int
    goal g1:
    (a=-1 \text{ and } b=2) ->
             ((a=b) \text{ and } 2*a>0)
             ((a \rightarrow b) and 2*a > 1)
File "try-alt-ergo-file", line 4, characters 1-82: I don't know (0.1210) (3
steps) (goal g1)
```

Comparison of Multiple-if vs Else-if

```
@requires marks = 75
                               @requires marks = 75
@ensures grade = B
                               @ensures grade = B
@program {
                               @program {
                                  if (marks > 90)
  grade = F;
                                        grade = A;
                                 else if (marks > 70)
  if (marks > 50) grade = C;
                                            grade = B;
  if (marks > 70) grade = B;
                                        else if (marks > 50)
  if (marks > 90) grade = A;
                                                  grade = C;
                                             else grade = F;
```

Which one is easy to verify???



```
If (a == b)
S1;
If (b == c)
S2;
If( c == a)
S3;
```

How many execution paths??

```
If (a == b)
S1;
If (b == c)
S2;
If( c == a)
S3;
```

Work out here...

$$A=1,b=2,c=3 --> None$$

$$A=1,b=1,c=3 --> S1$$

None

S1

S2

S3

ΑII

```
If (a)
                          None
   S1;
                          S1
If (b)
                          S2
  S2;
                          S3
If(c)
                          S1,S2
  S3;
                          S1,S3
                          S2,S3
                          S1,S2,S3
```

```
If (a == b)
S1;
If (b == c)
S2;
If( c == a)
S3;
```

2 ^ n

```
@requires marks = 75
@ensures grade = B
@program {

   grade = F;

   if (marks > 50) grade = C;
   if (marks > 70) grade = B;
   if (marks > 90) grade = A;
}
```

How many execution paths??

```
@requires marks = 75
@ensures grade = B
@program {

    grade = F;

    if (marks > 50) grade = C;
    if (marks > 70) grade = B;
    if (marks > 90) grade = A;
}
.
```

- How many execution paths??
- S1
- S1,S2
- S1,S2,S3
- \$1,\$2,\$3,\$4

```
@requires marks = 75
@ensures grade = B
@program {

    grade = F;

    if (marks > 50) grade = C;
    if (marks > 70) grade = B;
    if (marks > 90) grade = A;
}
```

- How many execution paths??
- S1
- S1,S2
- S1,S2,S3
- \$1,\$2,\$3,\$4

```
if (B1)
s1;
```

if (B2)

s2;

{O}

Try to derive the weakest precondition for this program



```
if (B1)
  s1;
{(B2 && wp(s2,O)) || (not B2 && O)}
if (B2)
  s2;
{O}
```



```
\{(B1 \&\& wp(s1,P)) || (not B1 \&\& P)\}
if (B1)
 s1;
P = \{(B2 \&\& wp(s2,O)) || (not B2 \&\& O)\}
if (B2)
 s2;
{O}
```



```
{(B1 && wp(s1,{
(B2 \&\& wp(s2,O))
(not B2 && O)
(not B1 && {
(B2 \&\& wp(s2,O)) \parallel (not B2)
&& O)})
```

- Larger the number of execution paths, greater is the size of verification conditions.
- The size of verification condition generated increases exponentially (eg. For n cases, its 2ⁿ

Else-if

```
If(B1)
S1;
else if(B2)
S2;
...
Sn;
```

How many execution paths for n cases?

Else-if

```
If(B1)
 S1;
else if(B2)
 S2;
Sn;
```

How many execution paths for n cases?

ONLY n paths!

s1 or s2 or s3 orsn

Else-if

```
If(B1)
                   B1 && wp(S1,O)
 S1;
                   || not(B1) && B2 && wp(S2,O)
else if(B2)
                   || not(B1) && not(B2) && B3 && wp(S3,O)
 S2;
else if(B3)
 S3;
                   not(B1 || B2 || ... || Bn-1) && wp(Sn,O)
else
Sn;
                   Increase in size of VC generated is linear
```



Comparison

```
@requires marks = 75
                               @requires marks = 75
@ensures grade = B
                               @ensures grade = B
@program {
                               @program {
                                 if (marks > 90)
  grade = F;
                                       grade = A;
                                 else if (marks > 70)
  if (marks > 50) grade = C;
  if (marks > 70) grade = B;
                                       else if (marks > 50)
  if (marks > 90) grade = A;
```

Else-if is easier to verify: smaller VC and less complex (grade is assigned exactly once)



grade = B;

grade = C;

else grade = F;