

reject $H_{(i)}$,
go to the step (i+1) if pais = then else accept Heis- Hens Stepn: if pan = I then reject Han else accept Han. The procedure stops the first time when $P(t) = \Delta c = \frac{\Delta}{n-i+i}$ Note Holm's procedure is less conservative than
Bonferioni's procedure (which rejects all Ho;:

Pict Theorem 3. Holm's procedure controls the FNER.

Strongly (without assumption of independence) Let io = minti: Mo,; 15 true] $p_{(i)} \leq \dots \leq p_{(i_0-1)} \leq p_{(i_0)} \leq p_{(i_0+1)} \leq \dots \leq p_{(n)}$ $H_{(i)} = H_{(i_0-1)} + H_{(i_0)} + H_{(i_0+1)} + \dots + H_{(n)}$ $H_{3i} \text{ are true} = H_{0i} \text{ are true and probably a few } H_{3i}$ We have no hypotheses when Hoi are time. $=> n-i_0+1 \ge n_0 => i_0 \le n-n_0+1$

Global Testing vs. Multiple Testing global testing procedures: Bonferrom: reject H, if $p_{(i)} \leq \frac{1}{n}$. Simes reject H if $p_{(i)} \leq \frac{1}{n}$ or $p_{(i)} \leq \frac{1}{n}$... or (independence) or $p_{(i)} \leq \frac{1}{n}$ or $p_{(i)} \leq \frac{1}{n}$ or $p_{(n)} \leq \frac{1}{n}$. Fisher reject H if $\sum_{i=1}^{n} 2\log p_{i}^{-1} \geq \chi_{2n}^{2}(1-\lambda)$ (independence) Multiple testing

Our goal is to control FWER: we want

procedures for which FWER < 2.

Bonferoni: reject H; if p; = 1 controls

FWER multiple testing · Sims: if $f_j: p_{(j)} = \frac{jd}{n} \Rightarrow \text{reject } H_{(i)} \text{ for } (\text{Independence})$ does not control FWER $(i) = \max \{i : p(i) \leq \frac{i}{n}\} \Rightarrow \text{reject } H(i)$ $for \forall i \leq i_0$ Example: Paj = &; Paj = 2d, let no= h-1. Let pen is too small >> How, is rejected.
Then we have n-1 hypotheses Herr. Henry and the smalest p-value is compared with $\frac{2d}{n}$. But to control FWER P(2) should compare with $\frac{d}{n-1}$ fWER in this case will be n-2

Closure Principle. n hypotheses 1 Hi 3 i-1 Let us define:

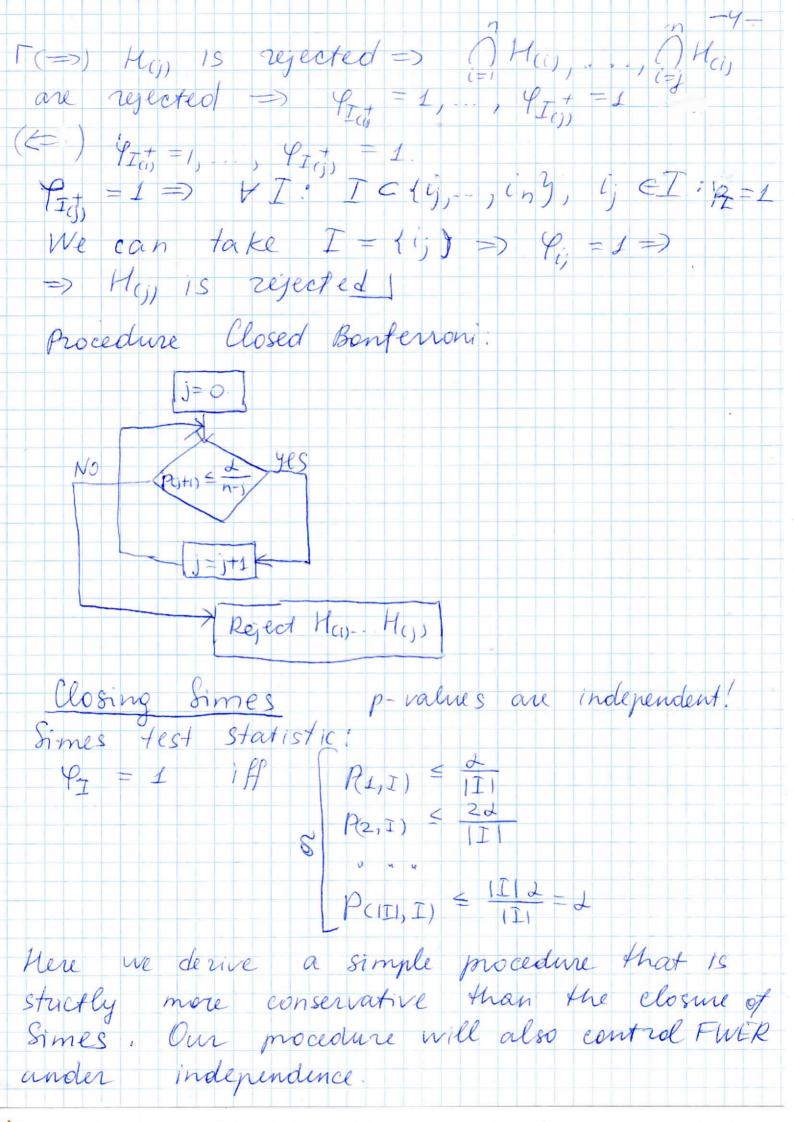
Hi; = Hi n H; = {pi, p; ~ U to, 173 $H_{I} = \bigcap_{i \in I} H_{i}$, $I \subset \{1, \dots, n\}$ Definition. the closure of the family of Hi 3i-1 is 1 HI: VIC 11- n3 3 Example: n=4. H123 H124 H134 H234 H12 H13 H14 H23 H24 M34 H1 H2 H3 H4 Let I is the test that is generated by a global testing proledure (for example, Bonferroni, Pisher, Simes, x etc.) We have set of hypotheses { Hi: i e I } we apply of for testing the global mull HI TH; (assume that 9, 15 & level test for testing Hi, HI is rejected, if PI=1.) $P(\varphi_T = 1/H_T) \leq 2$ The closure mocedure: If Hy is rejected for all Y: I'= I (at level 2) => HI is rejected

Example n = 4

H₁₂₃ H₁₂₄ H₁₃₄ H₂₃₄ (M12) (M13) (M14) H23 H24 H34 · HI 18 rejected because & J C f1, 2, 3, 43 s. +. 113 € J Hg 13 rejected · H2 is not rejected, because H234, H23, H24
are not rejected. Theorem 1: The closure principle controls the Tho = 1 Hiv: mHoi is true. 3 - true nulls make false discourses make folse discoveries A = {V > 13 - we have at least one false disco-B = & reject Here y - we reject the set of time mulls For the closure procedure, A1B=A (see example, if we reject thoy, for example, then, for the cl. pr. we reject tho, 234). So $P(A) = P(A \cap B) = P(B) P(A/B) \leq 2$ = closure minciple marides us with a generic recipe for translating global testing proceedings into valid FIVER controlling mocedures. Only Issue Is that this is not applicable in machine For example, let n = 10 000 = we should test 2 10 000 hypotheses. Not prossible at all to

test all this products even for the simple tests like Bonferron. The guestion is it we can use the closure procedure to construct some exsy test which would be closure of some well-known global procedures. Closing Bronferroni's procedure to generate 97: $\varphi_{I} = 1$ iff min $p_{i} \leq \frac{\Delta}{|I|}$ 1. Pc1) = --- = pc1) = --- = Pcn) Mai Hij set I:

Paj = min 1 pi 1 e I 3, I e 1 ij - in 3, ij e I Let $I_{(j)} + = \{i_j ... i_n \}$ — the index set corresponding to the n-j+1 largest p-values: P(s), P(n). Then φ_{IG} = 1 implies that φ_{I} = 1. $|f| |\varphi_{I,j}| = 1 |\langle = \rangle | |p_{(j)}| \leq \frac{2}{n-j+1}$ Pc, is the smallest p- value with index in I $\Rightarrow |I| \le n - j + 1$ $\Rightarrow p(j) \le \frac{2}{n - j + 1} \le \frac{2}{|I|}$ 2. H_{cj}) is rejected $\Leftrightarrow \varphi_{I_{cij}}^+ = 1, \dots, \varphi_{I_{cij}}^+ = 1 \Leftrightarrow$ $\Rightarrow p_{(i)} \leq \frac{d}{n}, \dots, p_{(j)} \leq \frac{d}{n-j+1}$



-3-Lemma 2. < pcs, < - & Pg1) =_ Pc1) & Pc21 & ---E Plan Hen Hez Heji) Heji) Hig Hiz Hij Heji Hin) Min index set I Suppose that: (α) $i_j \in \mathcal{I}$ (6) $\exists j' \geq j$ such that $P(j') \leq \frac{d}{n-j'+j}$ Then Py = 1 for the Simes test Py Let $k:=\max\{\ell: i_{\ell}\in I; \ell \in J'\}$. By (a). k $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t \in Then$ $\ell \times i \times f \times s$ and $\ell \times i \times t$ and $\ell \times i \times t$ DI by definition ix & I. Let $a = |16j...6_{k-1}3 \cap I1|$, $b = |66_{k+1}...6_{n}3 \cap I1|$ then $a + b + 1 = |II| \Rightarrow |II| \le ab + a + b + 1 =$ $= \frac{1}{6t} = \frac{a+1}{1II}$ By definition of Simes procedure 47 = 1