

def 1: A preference is called continuous if:

Say we have 2 bundles  $x$  and  $y$  as functions of  $n$  (say  $(x_n, 0)$  &  $(y_n, 0)$ )

(the sequence (then we have 2 sequences  $\{x_n\}$  &  $\{y_n\}$ )  
must be converging)

then if  $x_n \geq y_n \forall n$

 it continues then  $\lim_{n \rightarrow \infty} x_n \geq \lim_{n \rightarrow \infty} y_n$

def 2: A preference is called continuous if (n)  
the upper contour set for any bundle is closed.

A set is closed if you take a sequence of points, then the limit of seq. is also included in the set.

def 1  $\Rightarrow$  def 2:

proof  
Say we have an element  $x^*$ ; to show that the if there is a sequence of points  $\{x^n\}$  such that  $x^n \geq x^* \forall n$  then  $\lim_{n \rightarrow \infty} x^n \geq x^*$ .

In def 1 take  $\{x^n\}$  as  $\{x_n\}$  and  $\{y^n\}$  as  $\{y_n\}$   
we see that if  $x_n \geq x^* \forall n$   
 $\Rightarrow \lim_{n \rightarrow \infty} x_n \geq x^* \Rightarrow$  Hence proved.

\* Connected set: the set which cannot be partitioned

To show all lexicographic sequences may be are discontinuous

e.g.

$$x^n = (2, 1) \quad y^n = (2-1, 2)$$

$$x^n \geq y^n \quad \forall n; \text{ as } \frac{1}{n} > 0 \forall n$$

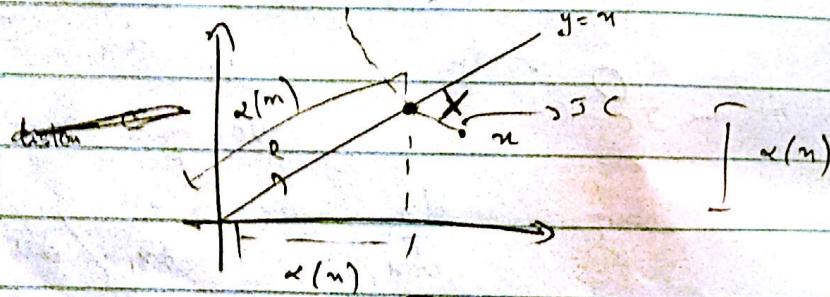
$$\text{but } \lim_{n \rightarrow \infty} x^n \not\leq \lim_{n \rightarrow \infty} y^n$$

### Representation Theorem

Suppose that  $\succeq$  is well behaved i.e. it is MON 2 convex. Additionally  $\succeq$  is continuous.

(i) Then  $\exists u$ , s.t.  $u$  represent  $\succeq$

(ii)  $u$  is continuous



The point  $x$  is  $\alpha(n)$

Clearly we have  $\alpha(n) \in \sim x$

Claim:  $\alpha(m) \in \sim x$  can be the utility function.

First we have to show that such an  $\alpha$  exists

Define a set  $\{x \mid \text{upper contour set of } m \cap \mathbb{Z}\}$

Lower contour set  $\{x \mid \text{lower contour set of } m \cap \mathbb{Z}\}$

The union of these two sets is entire line segment  $y = n$

Due to the connected property, there must be an element in their intersection. Hence  $\{x \mid \text{lower }\}$  is non-empty and that  $x \sim m$  (since it lies in both)

Proof of Repres.  
Maschler's  
Sours  
Microeconomics  
theory.

Homothetic pattern

Berk Paper

$$u(tn) = t u(n)$$

If  $n \rightarrow \infty \Rightarrow t \rightarrow \infty$

Quadrant

$$t \rightarrow \infty \Rightarrow u_n(t_0) \approx g_n(t_0)$$

$$u(n + t_0) = u(n) + t_0$$

for some indifference curve

$$u(y_1, n_2) = \text{constant}$$

$$\text{or } du = \frac{\partial u}{\partial y_1} dy_1 + \frac{\partial u}{\partial n_2} dn_2$$

$$\Rightarrow \frac{dn_2}{dy_1} = -\frac{\frac{\partial u}{\partial y_1}}{\frac{\partial u}{\partial n_2}}$$

$\frac{du}{dy_1}$  do this is not change in utility

$\frac{du}{dy_1} \rightarrow$  onto the change in  $y_1$

keeping the other constant

This is called marginal utility of  $y_1$ .



We have already said:

The most preferable

optimal solution is point X

where the budget

line is tangent to IC.

How to mathematically formulate it?

Utility man Problem.

$$\text{Max } u(n_1, n_2) \text{ s.t. } \begin{cases} p \cdot \vec{n} \leq M \\ M > 0 \end{cases}$$

Budget set is compact

We will use L.P.P.

Lagrangian Multiplier Technique

$$\mathcal{L}(n_1, n_2, \lambda) \quad [\lambda \geq 0]$$

$$= u(n_1, n_2) + \lambda [M - p \cdot n]$$

$\mathcal{L}$  is the function which you need to maximize

If you had another constraint, then add that also. (with another parameter)

The technique used to solve is Kuhn-Tucker

Say at  $n'$   $MRS \neq \frac{P_1}{P_2}$



$P_1 n_1 + P_2 n_2 = m$  budget line

From this statement can we say that this is not the point we're looking for?

Case 1: Slope of budget line is flatter than the slope of IC

Say we are at  $x$ , lets see can I get some more  $n_1$  (by giving away some  $n_2$ ).

Since the slope of budget line is less, according to our budget we need to lose lesser units of  $n_2$  than what the

MRS gives. Hence our utility increases since we would be left with more units of  $n_2$  (if we follow budget line)

W. Rudin Mathematical analysis

Bunk Pages

$$MU_1 = \frac{\partial U}{\partial x_1} \quad MU_2 = \frac{\partial U}{\partial x_2} \quad \text{Max } U(x_1, x_2) \\ \text{s.t. } p_1 x_1 + p_2 x_2 = M$$

$$L(n, \lambda) = u(n) + \lambda [M - p_1 n]$$

Conditions:  $\frac{\partial L}{\partial n_i} \leq 0 \quad \frac{\partial L}{\partial \lambda} \geq 0$

$$\Rightarrow \frac{\partial U}{\partial x_i} \leq \lambda p_i$$

Corner soln

$$M \geq p_n$$

Inequality might occur as corner soln, when consumption of atleast one good is 0  
If we get an interior soln

~~$$\frac{\partial L}{\partial n_i} = \frac{\partial U}{\partial x_i} = \lambda p_i$$~~

Equality

Optimal consumption bundle

It is called convex if

in ~~out~~ boundary  $\frac{\partial u}{\partial n} \leq 1 P_i$

$\Rightarrow$  Convexity of a set

for  $\forall$  any  $x \in X$  the upper  
contour is convex  $u(\cdot)$  is quasiconcave.

$x, x' \in X$

$$u'' \approx x_2 + (1-\alpha)x'$$

Now back to problem let  $n^*$  be the solution to the problem

$$\text{At } n^*, \frac{\partial u}{\partial n_1} \leq 1 P_1, \quad \frac{\partial u}{\partial n_2} \leq 1 P_2$$

good sd"

$$M \geq p_1 n_1 + p_2 n_2$$

if interior:  $\frac{\partial u}{\partial n_1} = 1 P_1, \quad \frac{\partial u}{\partial n_2} = 1 P_2, \quad M = p_2 n_2$

solutions

$$\frac{\partial u}{\partial n_1}, \quad \frac{\partial u}{\partial n_2}$$

$$= \frac{P_1}{P_2}$$

$$M = p_1 n_1$$

\* Take a monotonic transformation of  $u$  to help solve

Sufficient conditions for maximisation Buyer's problem

(both for interior and exterior)

$u(\cdot)$  is negative semi-definite

Cobb-Douglas utility function

Actual step in general  $u(x_1, x_2) = x_1^\alpha x_2^\beta$

a) compute utility at corner  $P_1x_1 + P_2x_2 = m$  Derive  $\gamma^*$

b) Solve by Lagrangian

$$c) \text{ If } \gamma^* \text{ satisfy I, solves } \max_{x_1, x_2} u(x_1, x_2)$$

$$\frac{\partial x_1^{\alpha-1} x_2^\beta}{\partial x_1} - \lambda P_1 = 0$$

$$U = \alpha x_1 + \beta x_2$$

$$\beta x_1^{\alpha-1} x_2^{\beta-1} - \lambda P_2 = 0$$

$$\gamma^* = \lambda P_1$$

$$\gamma^*$$

$$\beta = \lambda P_2$$

$$\gamma^*$$

$$x_1 = \gamma^* \quad | \quad x_2 = \beta$$

$$\lambda P_1$$

$$\lambda P_2$$

$$\beta x_1^{\alpha-1} x_2^{\beta-1} (\gamma^*) + m$$

$$x_1^{\alpha} x_2^{\beta} (\alpha + \beta) = m$$

$$x_1^{\alpha} x_2^{\beta} = \frac{m}{\alpha + \beta}$$

$$\alpha + \beta$$

$$P_1 \gamma^* + P_2 \beta = m$$

$$\lambda P_1$$

$$\lambda P_2$$

$$\Rightarrow \lambda = \alpha + \beta$$

$$\boxed{x_1 = \frac{\gamma^* (m)}{P_1 (\alpha + \beta)} \quad | \quad x_2 = \frac{\beta (m)}{P_2 (\alpha + \beta)}} \quad |$$

$\Rightarrow$  homogeneous

This  $(x_1^*, x_2^*)$  is such a bundle which I can buy and want to buy. This  $\gamma^*$  is an ordinary

demand function ( $\gamma^*$  is a function of  $P_1, P_2$  and  $m$ )

maximized level of utility when the prices/income are given by  $P_1, P_2, m$ :

$$= u(n_1^*(P, m), n_2^*(P, m)) \\ = v(P_1, P_2, m)$$

$\hookrightarrow$  for any  $P_1, P_2$  and  $m$ ,  $v$  gives the maximum level of utility

$v$  is an indirect utility function.

$\frac{\partial v}{\partial m} \rightarrow$  is known as marginal utility of money

$$\frac{\partial v}{\partial m} = \frac{\partial u}{\partial m}(n_1^*(P, m), n_2^*(P, m))$$

$$= \frac{\partial u}{\partial n_1} \left|_{\substack{n_1=n_1^* \\ n_2=n_2^*}} \right. \frac{\partial n_1}{\partial m} + \frac{\partial u}{\partial n_2} \left|_{\substack{n_2=n_2^* \\ n_1=n_1^*}} \right. \frac{\partial n_2}{\partial m}$$

$$= \lambda P_1 \frac{\partial n_1}{\partial m}, \lambda P_2 \frac{\partial n_2}{\partial m} = \lambda \left( P_1 \frac{\partial n_1}{\partial m} + P_2 \frac{\partial n_2}{\partial m} \right)$$

but since  $P_1 n_1 + P_2 n_2 = m \rightarrow \frac{\partial v}{\partial m}$

$$= \lambda = \frac{\partial v}{\partial m}$$

Properties of ordinary demand function

(i)  $n_i(P, m)$  is homogeneous of degree 0 in prices and income.

$$(ii) P n^* = M$$

$\hookrightarrow$  Walras' Law  $\rightarrow$  market clearance; i.e., in equilibrium, we spend total money to get (buy) the best product

(iii)  $n^*(P, m)$  is a convex set. (Set as there can be multiple answers)

$$\text{lit } n, n' \in n^*$$

$\text{lit } n'' = \alpha n + (1-\alpha)n'$  using convexity, we can say utility is concave and it follows diminishing returns

- Task: • find demand functions for standard utility functions
- if ~~strict~~ preference is strictly convex, prove that demand function is a singleton set.

Another type of question:

your answer must be such that its utility  $\geq$  some given  $u^*$ . You have to minimize the expenditure ( $n$ ) [ $p_1$  and  $p_2$  are fixed]

This is the Expenditure Minimization Problem

The answer  $h$  is a function of  $(p_1, p_2, u^*)$

it is called compensated demand function

Now to solve

$$\begin{aligned} \text{Minimize } & (p \cdot n) \text{ s.t. } u(n) \geq u^* \\ \geq & (n_1, n_2, \mu) \\ = & p \cdot n + \mu (u(n) - u^*) \end{aligned}$$

~~p + u~~ First order condition:

Let  $h$  be the final answer

$$u(h) = u^* \quad \rightarrow \text{Sir wrote MP,}$$

$$\text{tangency} \Rightarrow \frac{\partial u}{\partial h_1} = -1 p_1 \quad \frac{\partial u}{\partial h_2} = -1 p_2$$

$$\frac{\frac{\partial u}{\partial h_1}}{\frac{\partial u}{\partial h_2}} = MRS_{h_1=h^*} = \frac{p_1}{p_2}$$

Can you link MP and GMP?

Something to do with  $v$  and  $u$

$$p_1 h_1^* + p_2 h_2^* = e(p_1, p_2, u^*)$$

$\hookrightarrow$  Expenditure function.

Given, the indirect utility function, ( $v$ ), can we derive the ordinary demand function? Similarly with ( $e$ )

Ques 1) Perfect substitutes

$$u(x_1, x_2) = \alpha x_1 + \beta x_2$$

$$p_1 x_1 + p_2 x_2 = m$$

Find the ordinary demand function:  
(logically)

Consider  $\frac{x}{p_i}$ : what is this

If  $\frac{x}{p_i}$  is more, that means  $x_1$  is better as it gives higher utility  $\rightarrow$  lower

price

$$\therefore x_1^* = \begin{cases} \frac{m}{p_1} & \text{if } \frac{x}{p_1} > \frac{\beta}{\alpha} \\ 0 & \text{if } \frac{x}{p_1} \leq \frac{\beta}{\alpha} \end{cases}$$

$$\left[ 0, \frac{m}{p_1} \right] \text{ if } \frac{x}{p_1} = \frac{\beta}{\alpha}$$

(any value)

(\*)

The value demand function we get from the solution of EMP is called compensated demand function. It is homogenous of degree 1 in prices.

$h^*$   $\rightarrow$  Hicksian

$n^*$   $\rightarrow$  Martinez / Walrasian

Do something like  $p \rightarrow h$   $h' \rightarrow h''$   $h'' = \alpha p + (1-\alpha)p'$  Bunk Pages  
 to show  $p''h'' \geq \alpha ph + (1-\alpha)p'h'$

## Recoverability

$\sum \frac{u_i(p)}{m} \Rightarrow u(p, m)$  Indirect utility function

$$\frac{\partial u}{\partial p_1} = \frac{\partial u}{\partial p_1} (u^*(p_1, p_2, m), u_2^*(p_1, p_2, m))$$

$$\frac{du}{dm} = -1$$

$$= \frac{\partial u}{\partial m_1} + \frac{\partial u}{\partial m_2}$$

$$\frac{\partial u}{\partial m_i} = \lambda p_i$$

$$= \lambda \left[ \frac{p_1 \partial m_1}{\partial p_1} + \frac{p_2 \partial m_2}{\partial p_2} \right]$$

$$(p_1 m_1 + p_2 m_2 = m)$$

$$m_1 + p_2 \frac{\partial m_2}{\partial p_1} = 0 \quad m_1 + p_1 \frac{\partial m_1}{\partial p_1} + p_2 \frac{\partial m_2}{\partial p_1} = 0$$

$$\therefore p_1 \frac{\partial m_1}{\partial p_1} + p_2 \frac{\partial m_2}{\partial p_2} = -m_1$$

$$\therefore \frac{\partial u}{\partial p_1} / \frac{\partial u}{\partial m} = -m_1$$

Roy's Identity

$$C(p_1, p_2, u^*) \rightarrow h_1(p_1, p_2, u^*)$$

$$C(p, m) = p_1 h_1(p, u^*) + p_2 h_2(p, u^*)$$

$$\frac{\partial e}{\partial p_1} = h_1$$

Shephard's Lemma

Slack equation ← on your own

CES utility function

$$U = (m_1^p + m_2^p)^{1/p}$$

How to express  $e$  in terms of  $u$ ?

$u \rightarrow f(p, m) \Rightarrow$  express  $m$  in terms of  $u$  and  $p$

quasiconcavity etc.

↳ Mathematical economical  $h_1, h_2$

How to convert from  $v$  to  $u$ ?

Given an indirect utility function  $u^*(P_1, P_2, M)$ , we can find  $U(n_1, n_2)$  by:

$$U(n_1, n_2) = \min_{P_1, P_2} u^*(P_1, P_2, M) \text{ such that } P_1 n_1 + P_2 n_2 = M$$

How? Let  $n = (n_1, n_2)$  be the demanded bundle when prices are  $(P_1, P_2)$  and income is  $M$ . Then, by definition,  $u^*(P, M) = U(n)$ . Let  $P'$  be another such vector with  $P' n = M$ . Since  $n$  is still affordable  $u^*(P', M) \geq U(n) + P' n = M$ .