



What is **Relational Structure**?

The relational **Structure** uses a collection of tables to represent both data and the relationships among those data. Each table has multiple columns, and each column has a unique name. Tables are also known as relations. The relational Structure is an example of a record-based model. Recordbased models are so named because the database is structured in fixedformat records of several types. Each table contains records of a particular type. Each record type defines a fixed number of fields, or attributes. The columns of the table correspond to the attributes of the record type. The relational data model is the most widely used data model, and a vast majority of current database systems are based on the relational model.



What is Relational Structure?

What is the	Relational		
Structure?			

The relational Structure represents how data is stored in Relational Databases. A relational database consists of a collection of tables, each of which is assigned a unique name. Consider a relation STUDENT with attributes ROLL_NO, NAME, ADDRESS, PHONE, and AGE shown in the table.

ROLL_NO	NAME	ADDRESS	PHONE	AGE
1	RAM	DELHI	9455123451	18
2	RAMESH	GURGAON	9652431543	18
3	SUJIT	ROHTAK	9156253131	20
4	SURESH	DELHI		18



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important terminology of Relational Structure

TABLE:SQL Table is a collection of data which is organized in terms of rows and columns. In DBMS, the table is known as relation and row as a tuple. Table is a simple form of data storage. A table is also considered as a convenient representation of relations.

EMP_ID	EMP_NAME	CITY	PHONE_NO
1	Kristen	Washington	728920122 3
2	Anna	Franklin	937828288
3	Jackson	Bristol	926478383 8
4	Kellan	California	725472834 6
5	Ashley	Hawaii	963848267 8

EXAMPLE OF TABLE CREATION

SQL> CREATE TABLE EMPLOYEE (
EMP_ID INT NOT NULL,
EMP_NAME VARCHAR (25) NOT NULL,
PHONE_NO INT NOT NULL,
ADDRESS CHAR (30),
PRIMARY KEY (ID)
);

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important terminology of Relational Structure

Row: In the context of relational databases, a row represents a single record in a table. Each row contains values for each column defined in the table schema. In everyday database terminology, a row is the term most commonly used to refer to a horizontal entry in a table.

Working with Tuple in DBMS

In a relational database, a relation is defined by a set of attributes and a set of tuples that have values for those attributes.

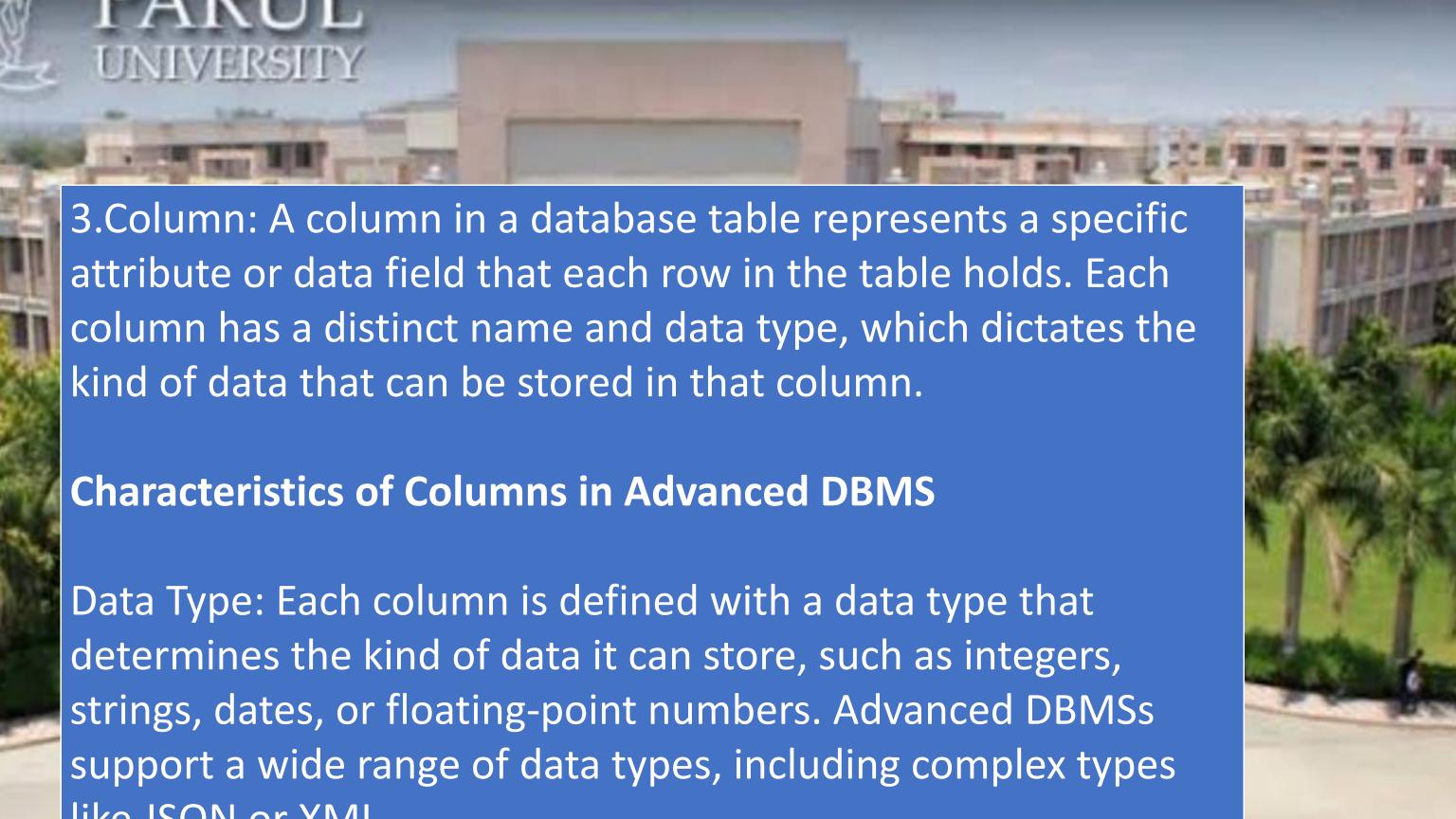
Example:

A relation called "CUSTOMER" might have attributes such as "customer_id", "first_name", "last_name", and "email". Each tuple in the relationship would have a unique value for the "customer_id" attribute and corresponding values for the other attributes, such as "John" for "first_name" and "Smith" for "last_name".



customer _id	first_na me	last_nam e	email
1	John	Smith	abc@gm ail.com
2	Abhishek	Bhosle	cde@gm ail.com
3	Natasha	Witch	fgh@gma il.com

Tuples are also used in the process of normalization in a relational database. Normalization is the process of organizing data in a database to minimize data redundancy and improve data integrity. In the process of normalization, a relation is broken down into multiple smaller relations, each with a specific purpose and containing a specific set of attributes and tuples.



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nts: Columns can have constraints that enforce certain rules on the mmon constraints include:

Key: Uniquely identifies each record in the table.

Key: Ensures referential integrity between tables by linking to a key in another table.

Ensures all values in the column are unique.

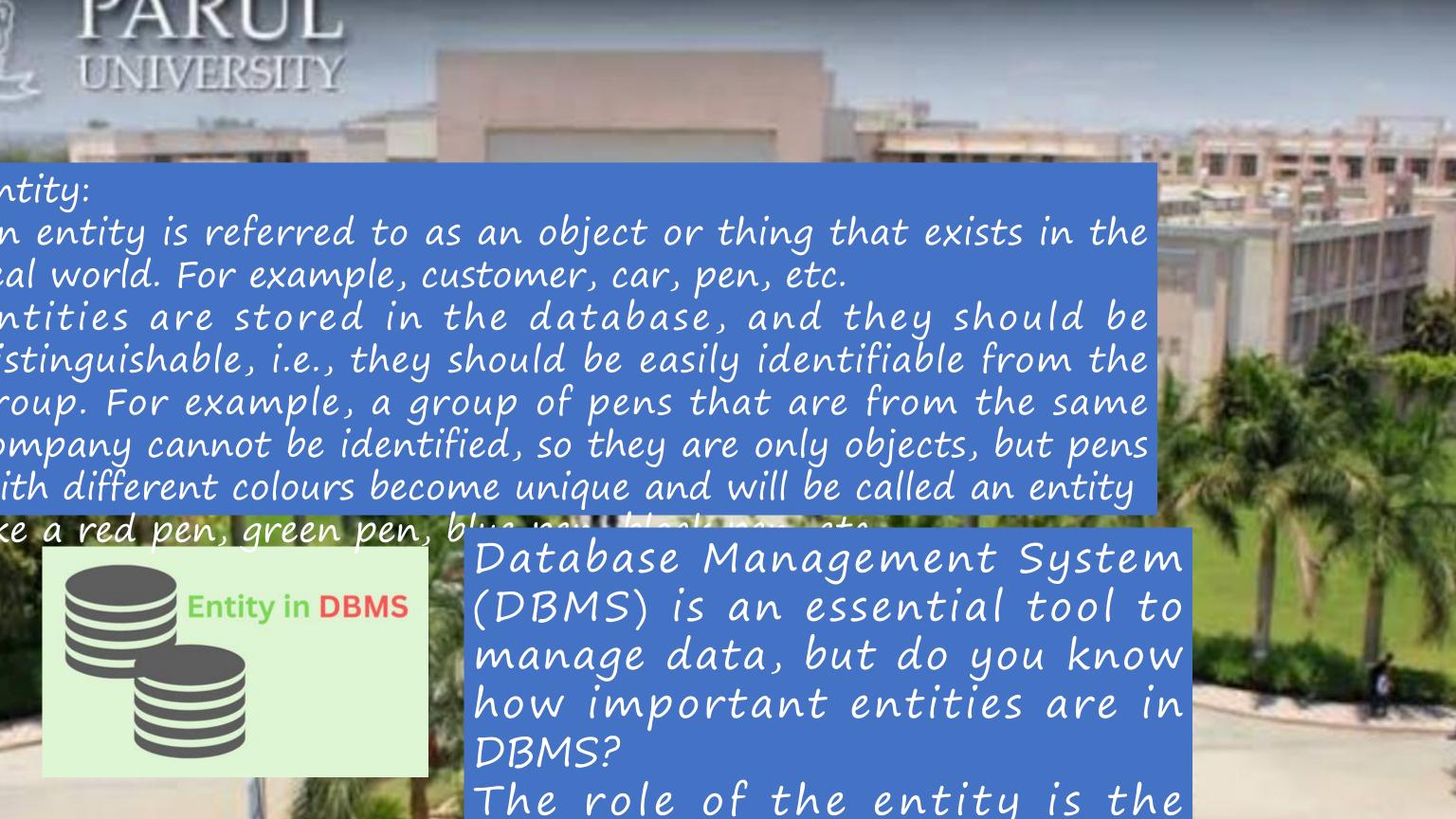
: Ensures that the column does not contain null values.

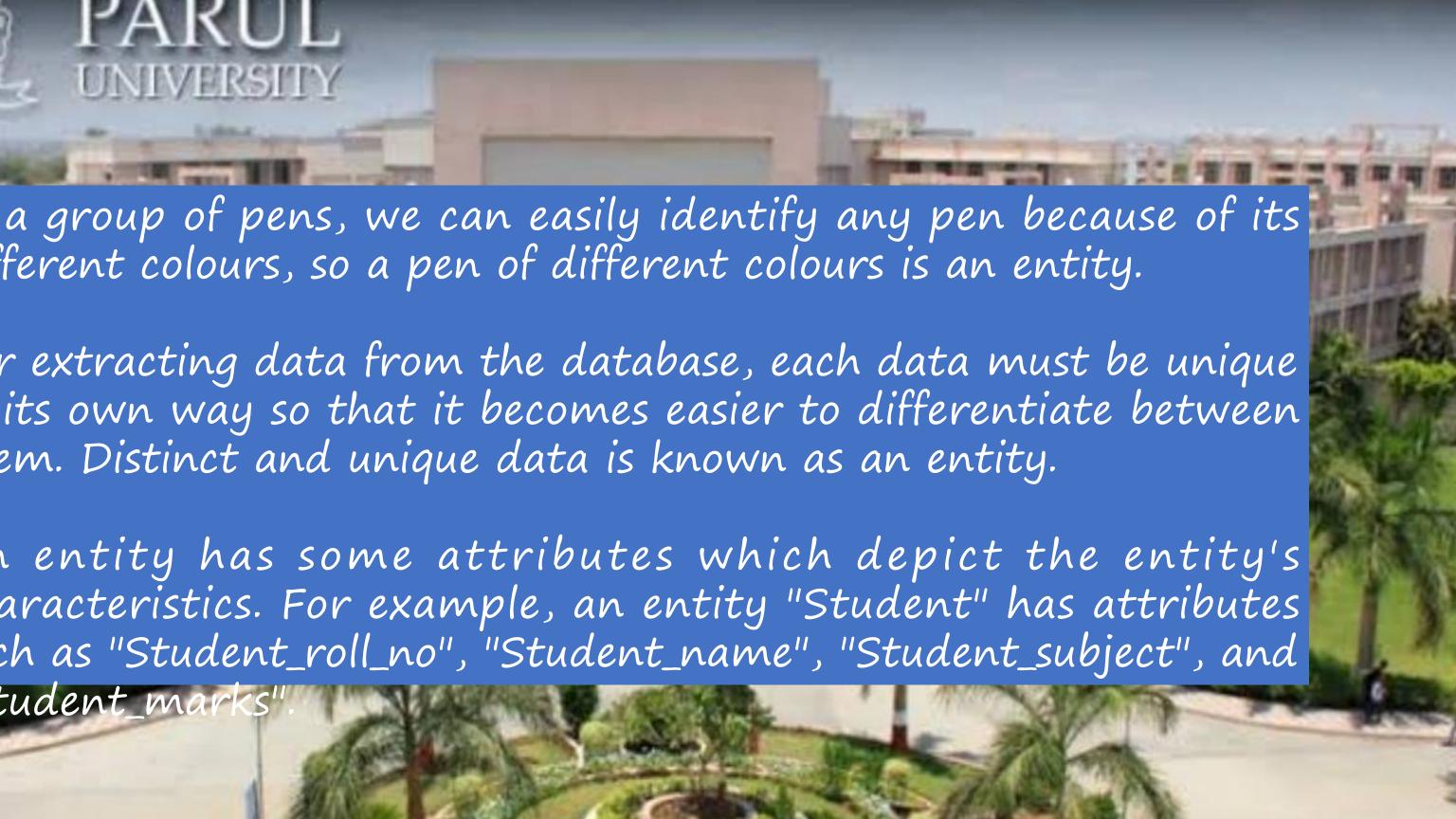
Provides a default value for the column if none is specified.

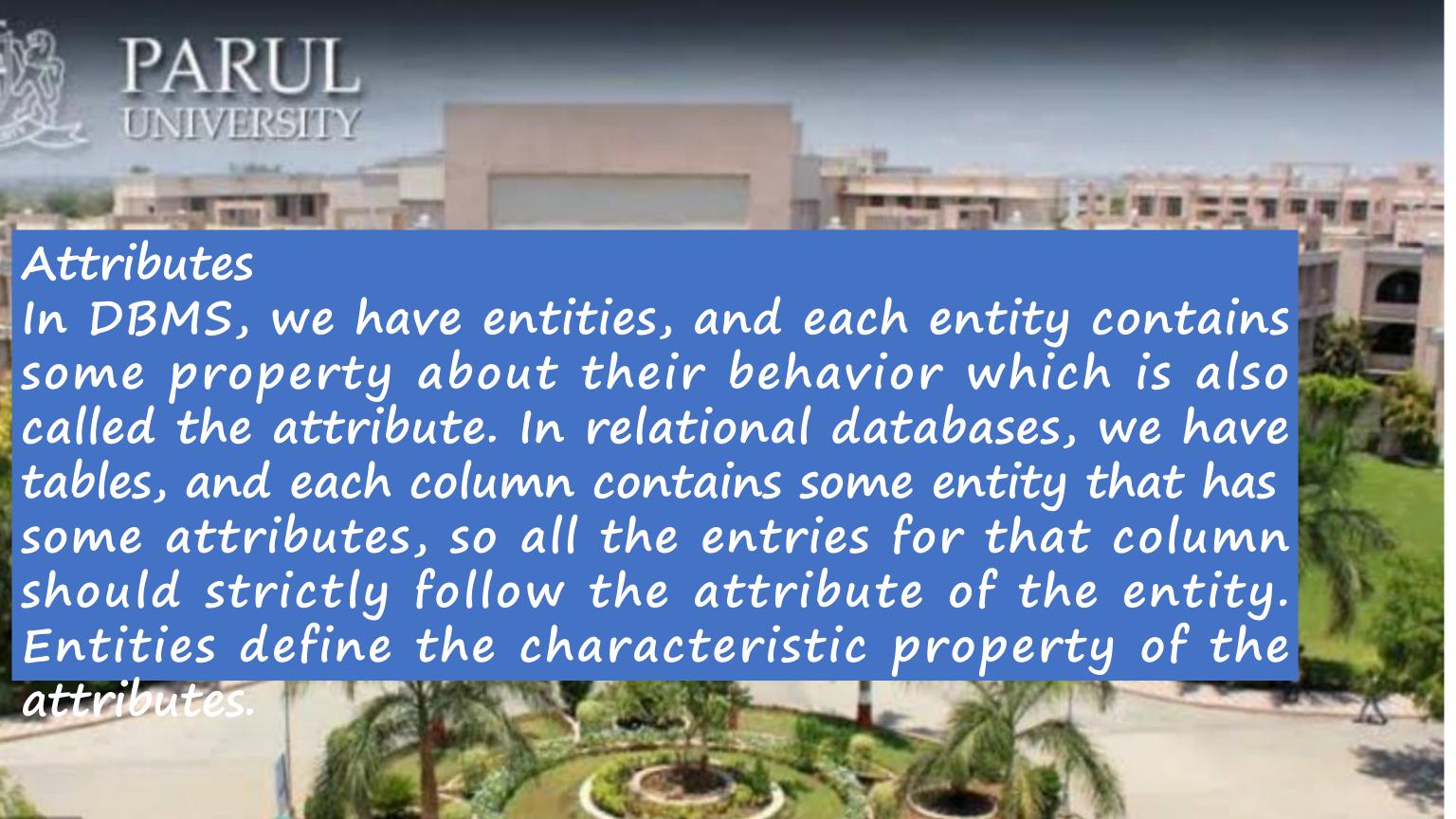
: Columns can be indexed to improve query performance. Advanced upport various indexing methods, including B-trees, hash indexes, e sophisticated types like bitmap indexes and full-text indexes.

Optimization: Modern DBMSs employ techniques such as columnar







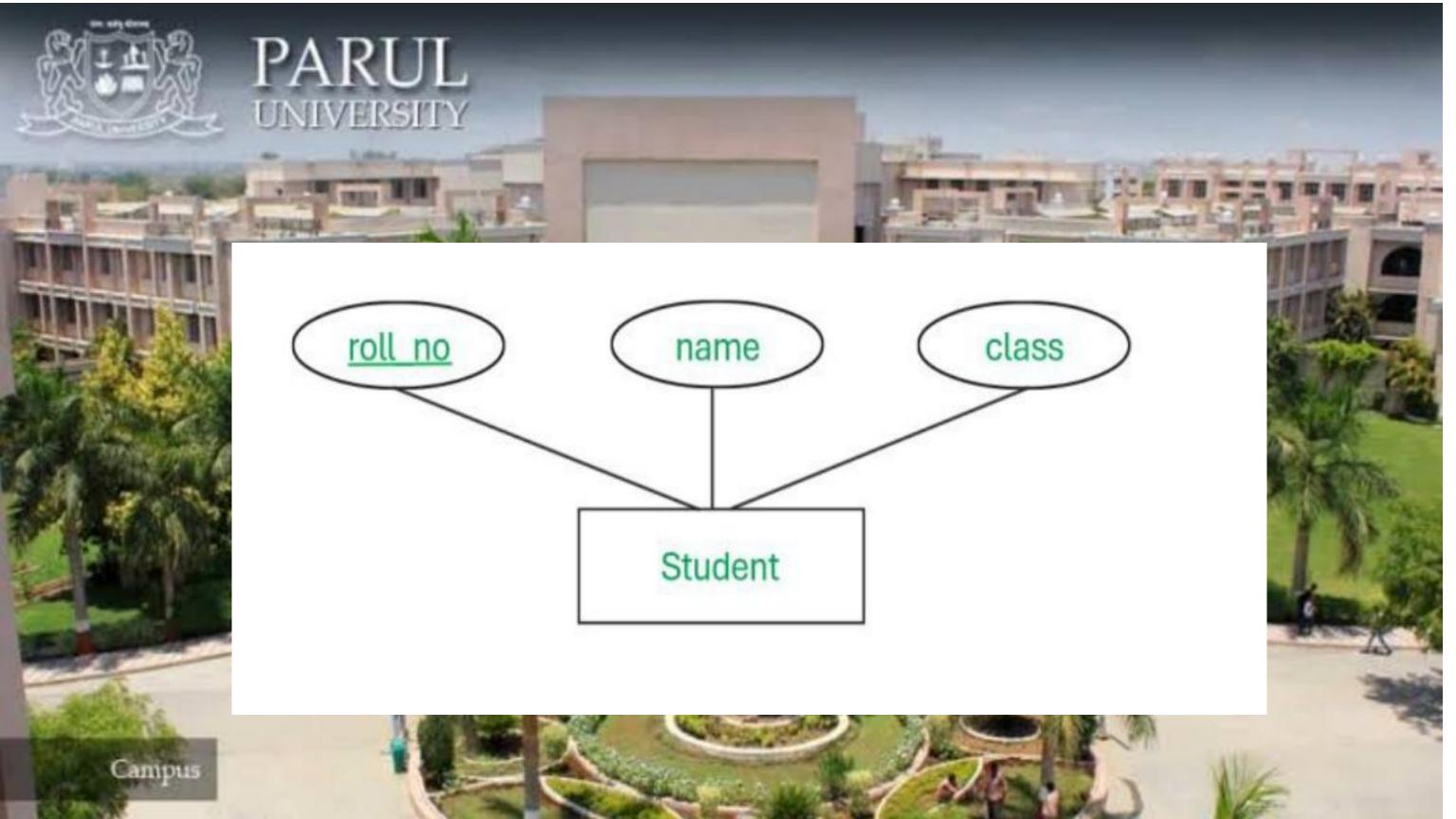


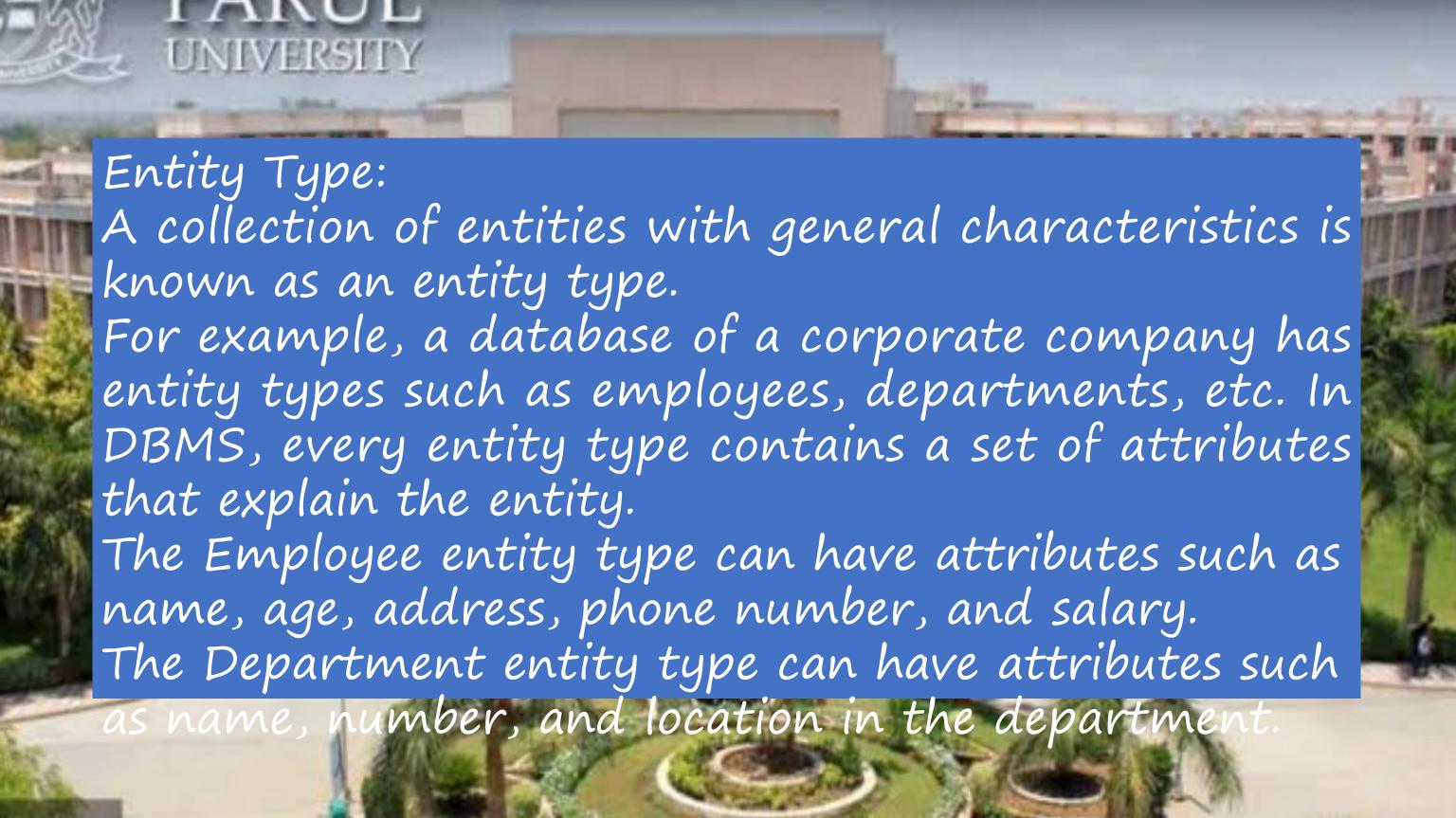
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ributes are properties or characteristics of an entity. ributes are used to describe the entity. The attribute is ring but a piece of data that gives more information about entity. Attributes are used to distinguish one entity from other entity. Attributes help to categorize the entity and entity can be easily retrieved and manipulate the entity. ributes can help the database to be more structural and archical. An entity with no attribute is of no use in the abase.

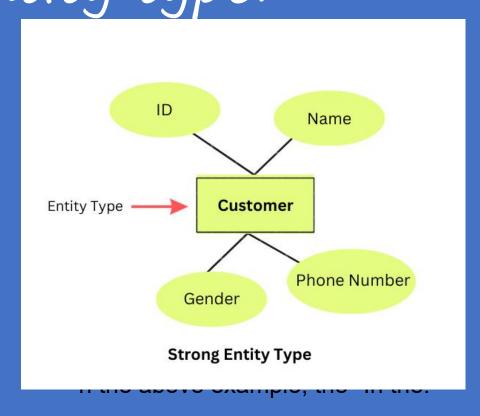
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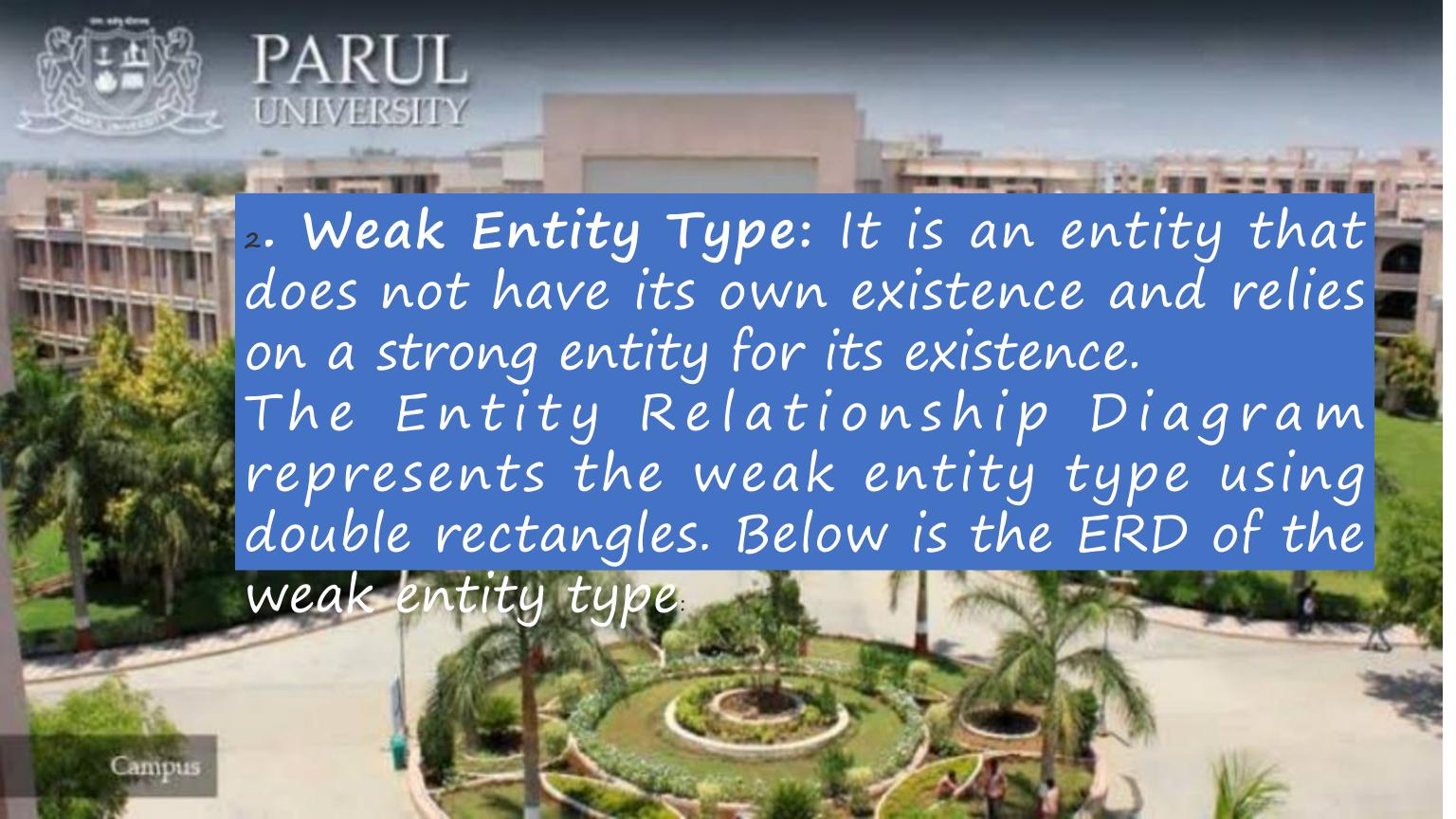
's take the student as an entity. Students will have tiple attributes such as roll number, name, and class. se attributes are used to describe the student in more





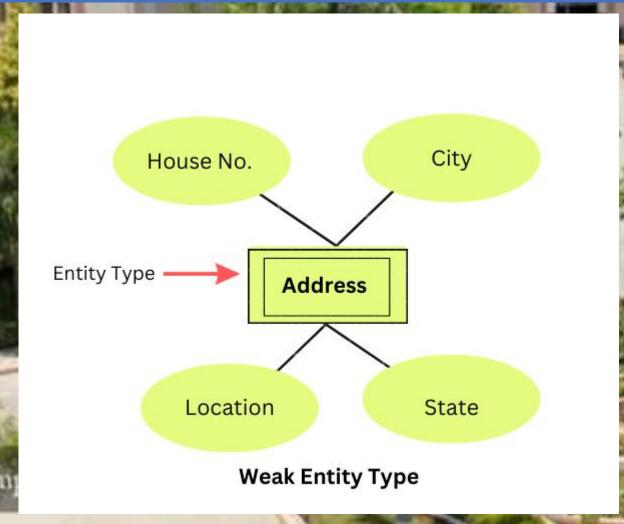
*Strong Entity Type: It is an entity that has its own existence and is independent. The entity relationship diagram represents a strong entity type with the help of a single rectangle. Below is the ERD of the strong entity type:





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In the above example, the "Customer" is the entity type with attributes such as ID, Name, Gender, and Phone Number. Customer is a strong entity type as it has a unique ID for each customer



In the above example, "Address" is a weak entity type with attributes such as House No., City, Location, and State.

The relationship between a strong and a weak entity type is known as an identifying relationship.

Using a double diamond, the Entity-Relationship Diagram represents a relationship between the strong and the weak entity type.

PARUL ER- MODEL IN DBMS? •ER model stands for an Entity-Relationship model. It is a high-level data model. This model is used to define the data elements and relationship for a specified system. ·It develops a conceptual design for the database. It also develops a very simple and easy to design view of data. ·In ER modeling, the database structure is portrayed as a diagram called an entity-relationship diagram. For example, Suppose we design a school database. In this database, the student will be an entity with attributes like address, name, id, age, etc. The address can be another entity with attributes like city, street name, pin code, etc and there will be a relationship between them.



Why Use ER Diagrams In DBMS?

ER diagrams represent the E-R model in a database, making them easy to convert into relations (tables).

ER diagrams provide the purpose of real-world modeling of objects which makes them intently useful.

ER diagrams require no technical knowledge and no hardware support.

These diagrams are very easy to understand and easy to create even for a naive user.

It gives a standard solution for visualizing the data logically.



Symbols Used in ER Model

ER Model is used to model the logical view of the system from a data perspective which consists of these symbols:

Rectangles: Rectangles represent Entities in the ER Model.

Ellipses: Ellipses represent Attributes in the ER Model.

Diamond: Diamonds represent Relationships among Entities.

Lines: Lines represent attributes to entities and entity sets with other

relationship types.

Double Ellipse: Double Ellipses represent Multi-Valued Attributes.

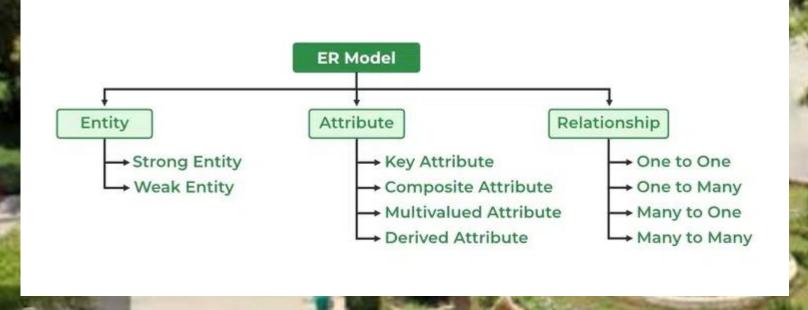
Double Rectangle: Double Rectangle represents a Weak Entity.



Figures	Symbols	Represents
Rectangle		Entities in ER Model
Ellipse		Attributes in ER Model
Diamond	\Diamond	Relationships among Entities
Line		Attributes to Entities and Entity Sets with Other Relationship Types
Double Ellipse		Multi-Valued Attributes
Double Rectangle		Weak Entity



Components of ER Diagram
ER Model consists of Entities, Attributes, and Relationships among Entities in a
Database System.



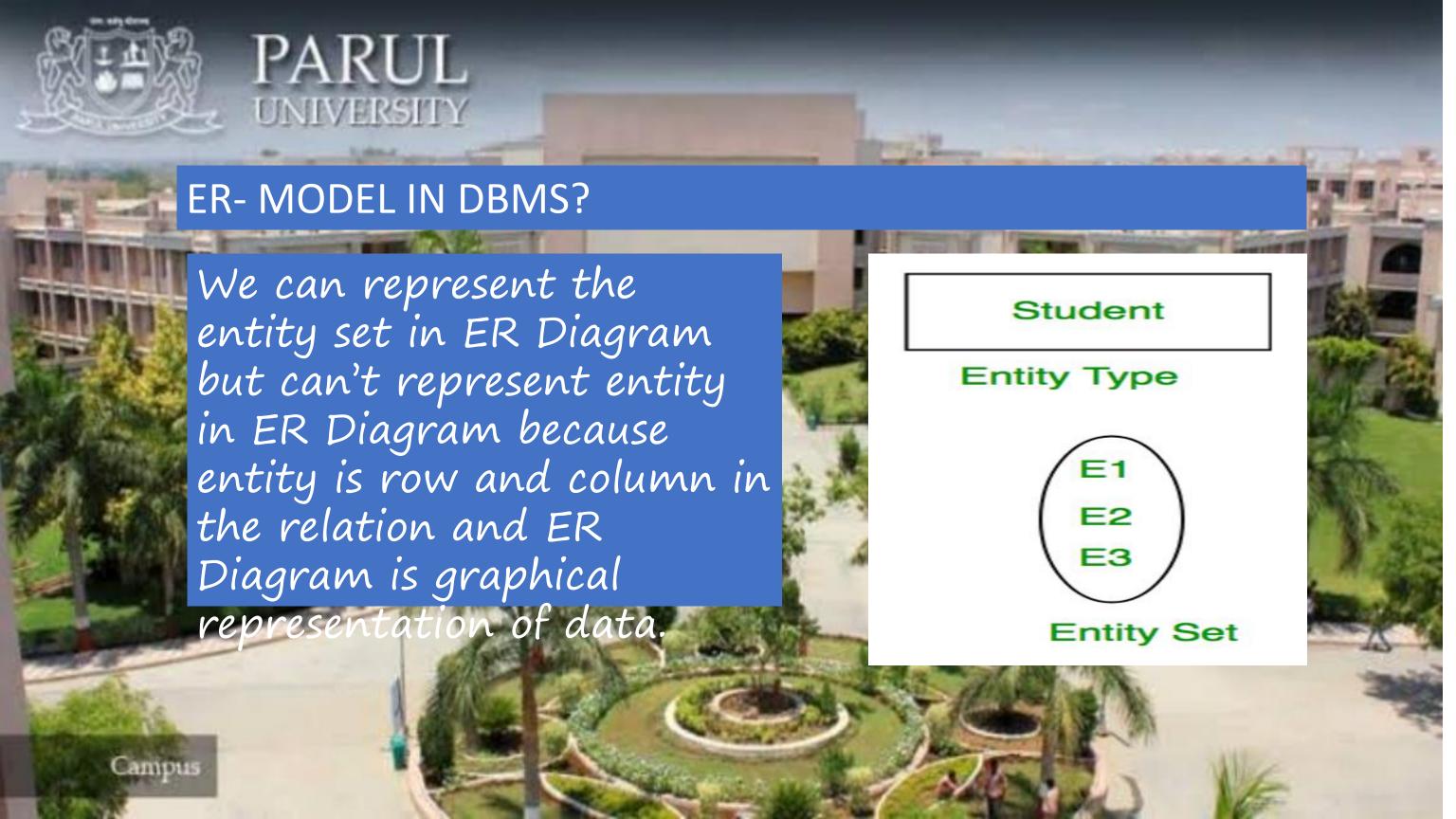


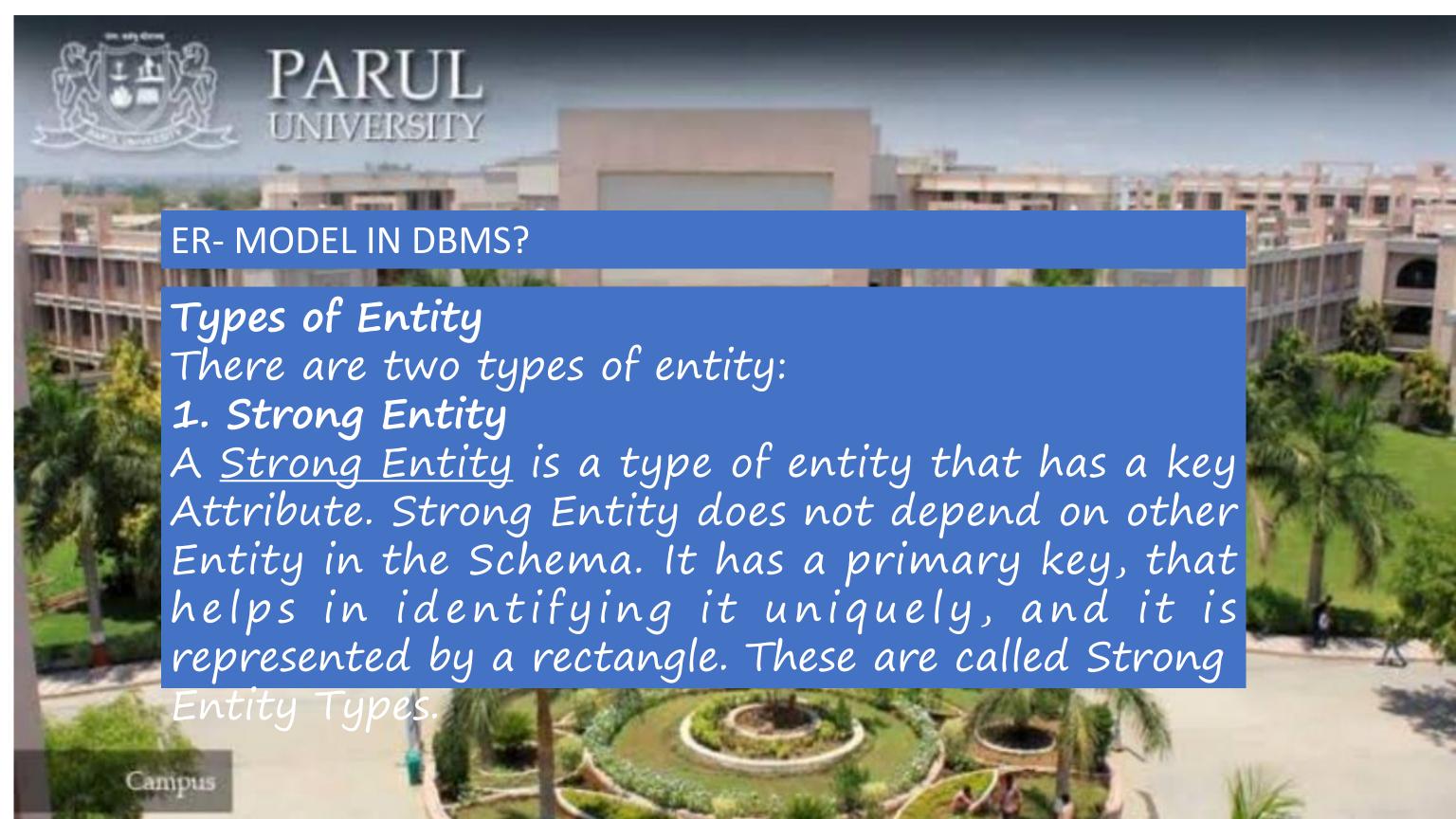
What is Entity?

An Entity may be an object with a physical existence – a particular person, car, house, or employee – or it may be an object with a conceptual existence – a company, a job, or a university course.

What is Entity Set?

An Entity is an object of Entity Type and a set of all entities is called an entity set. For Example, E1 is an entity having Entity Type Student and the set of all students is called Entity Set. In ER diagram, Entity Type is represented as:

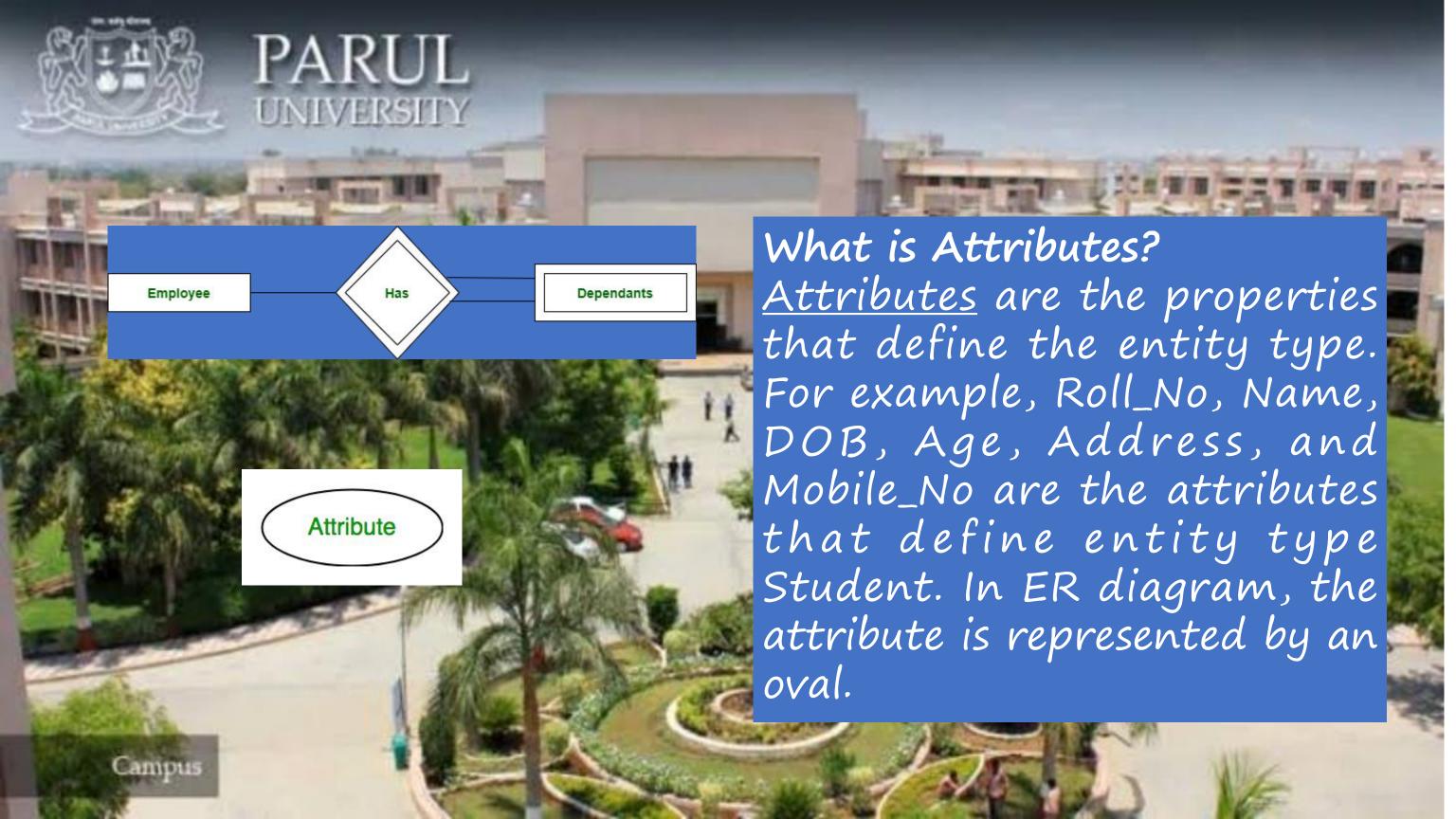




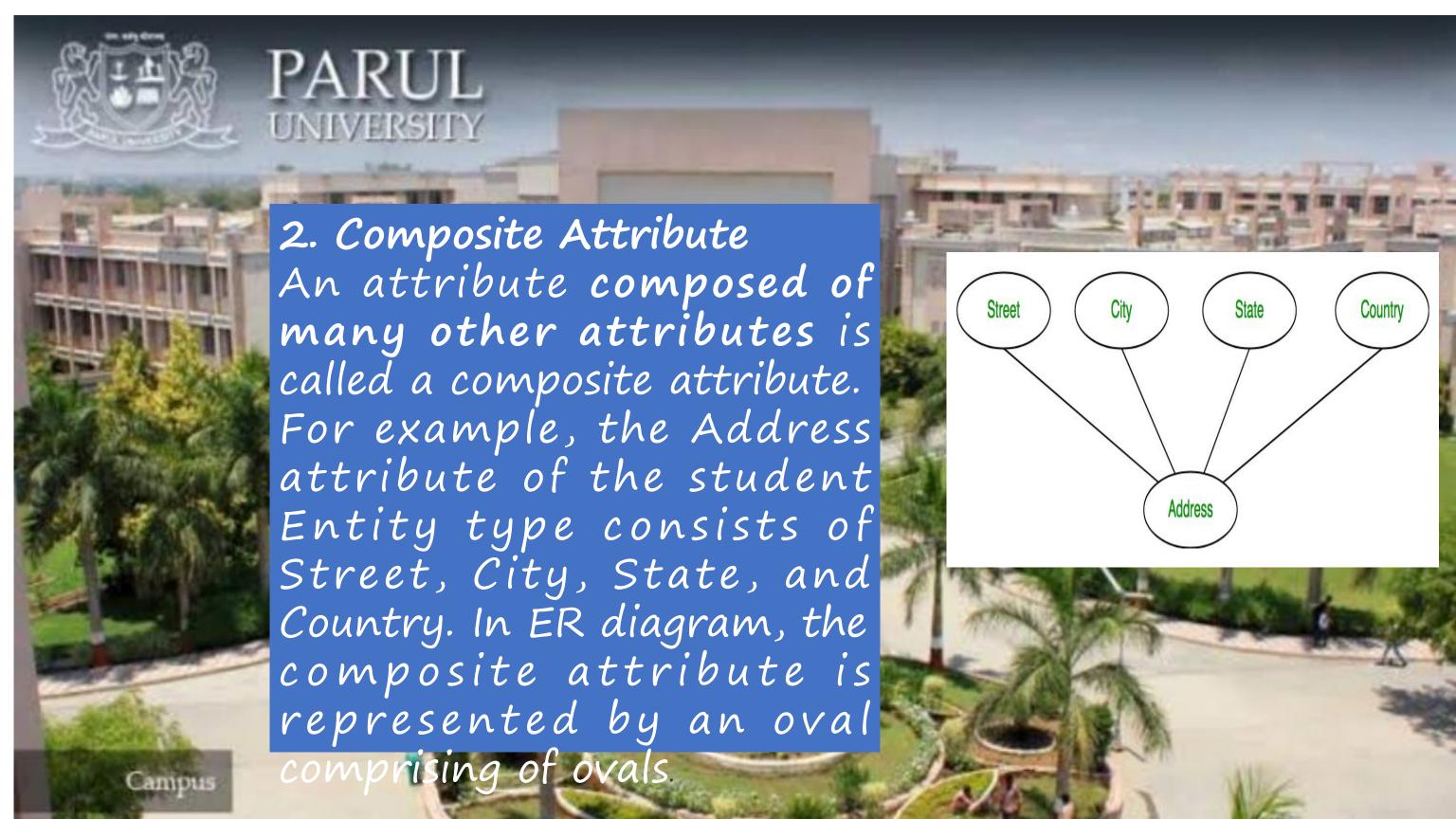


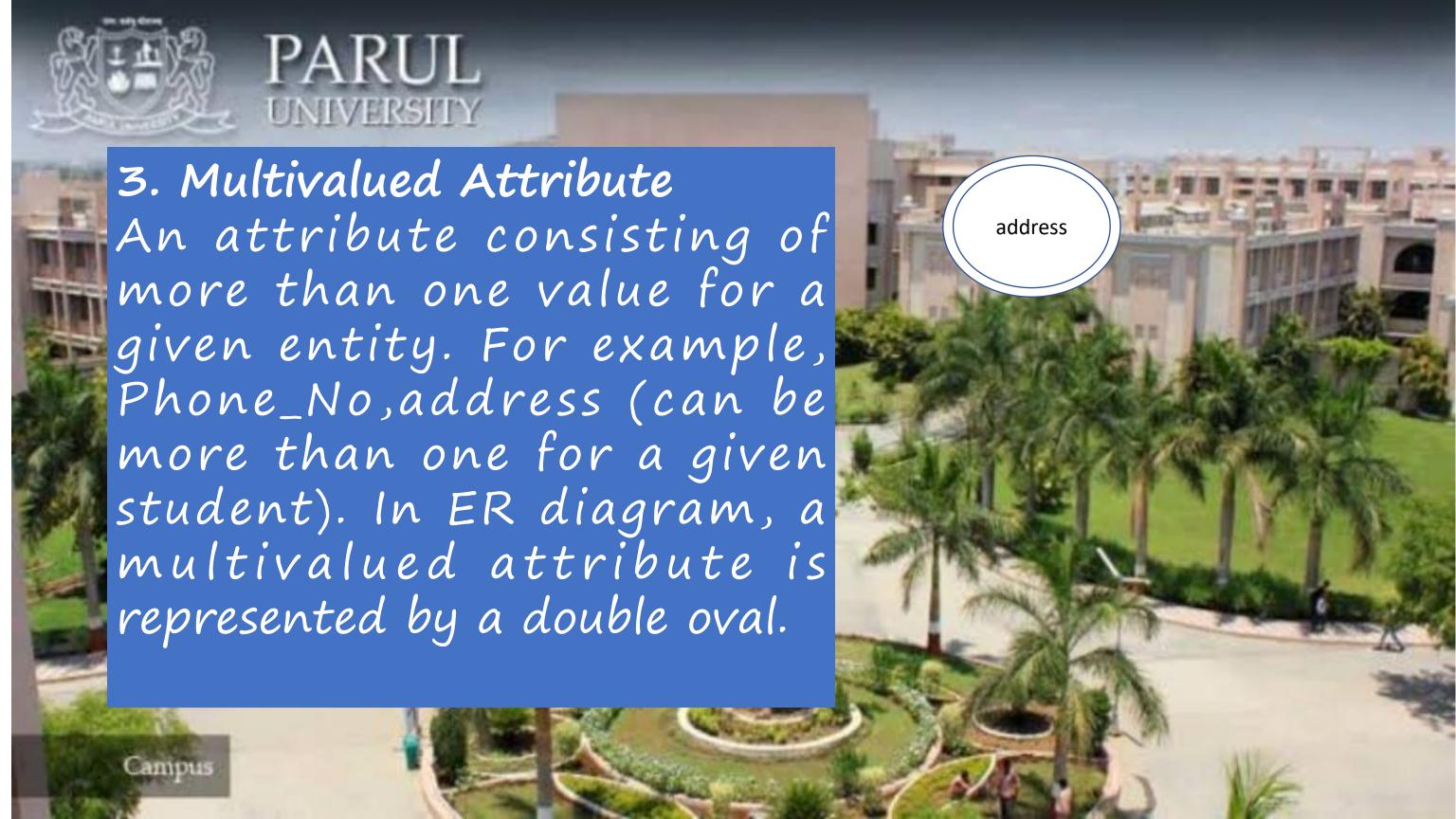
2. Weak Entity

An Entity type has a key attribute that uniquely identifies each entity in the entity set. But some entity type exists for which key attributes can't be defined. These are called Weak Entity types For Example, A company may store the information of dependents (Parents, Children, Spouse) of an Employee. But the dependents can't exist without the employee. So Dependent will be a Weak Entity Type and Employee will be Identifying Entity type for Dependent, which means it is Strong Entity Type. A weak entity type is represented by a Double Rectangle. The participation of weak entity types is always total. The relationship between the weak entity type and its identifying strong entity type is called identifying relationship and it is represented by a double diamond.











4. Derived Attribute An attribute that can be derived from other attributes of the entity type is known as a

derived attribute. e.g.;

Age (can be derived from

DOB). In ER diagram, the

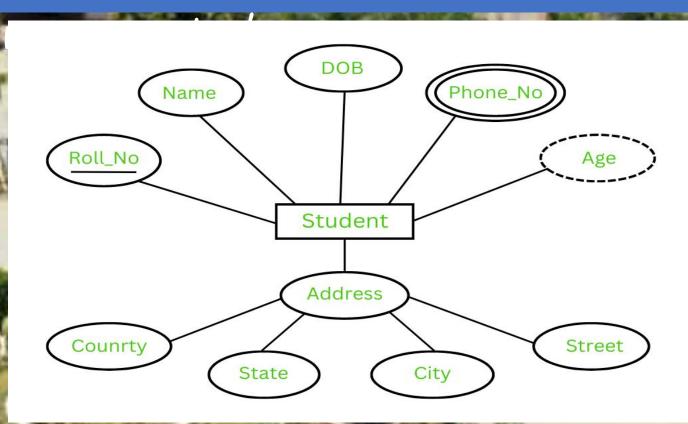
derived attribute is

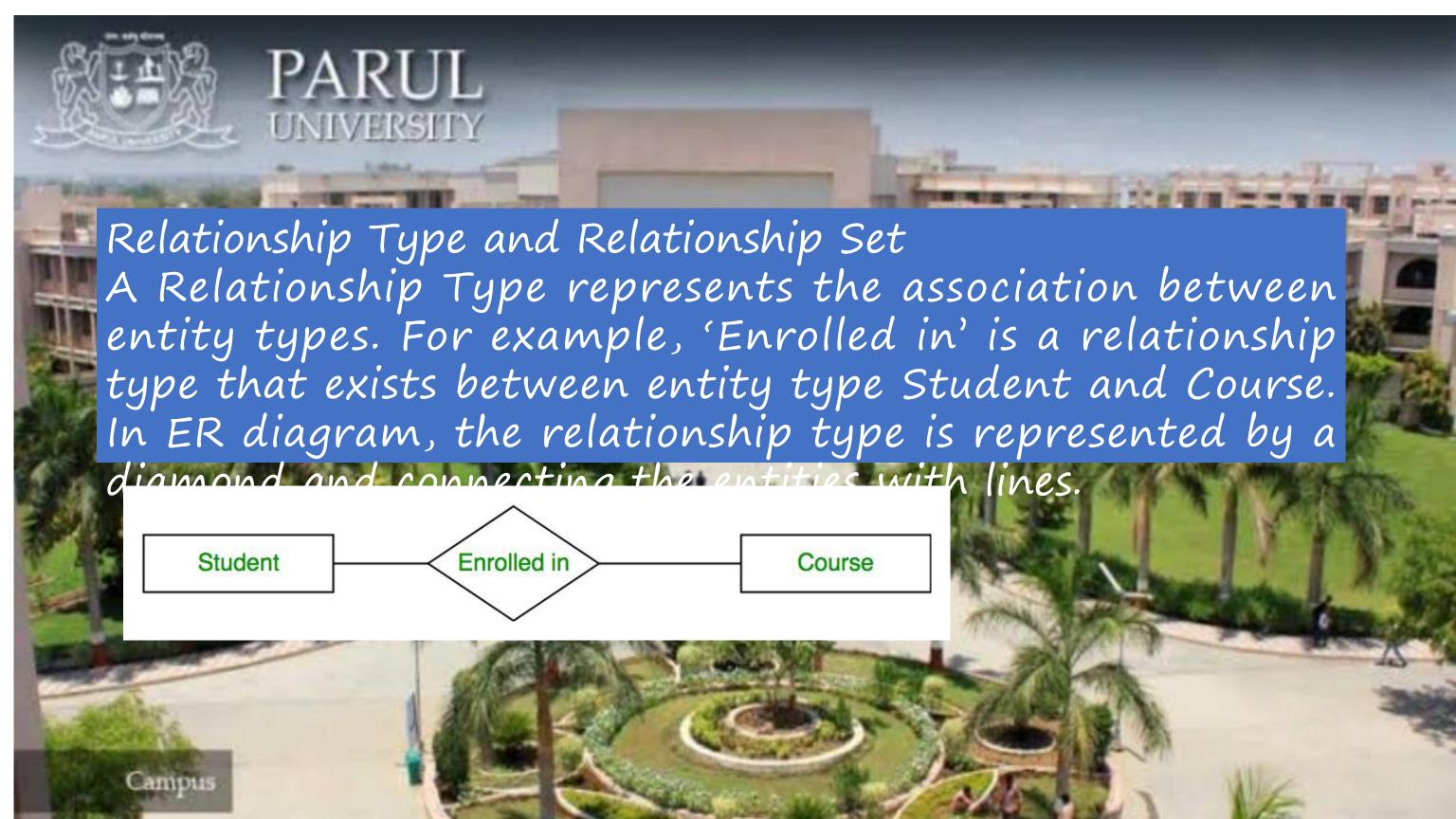
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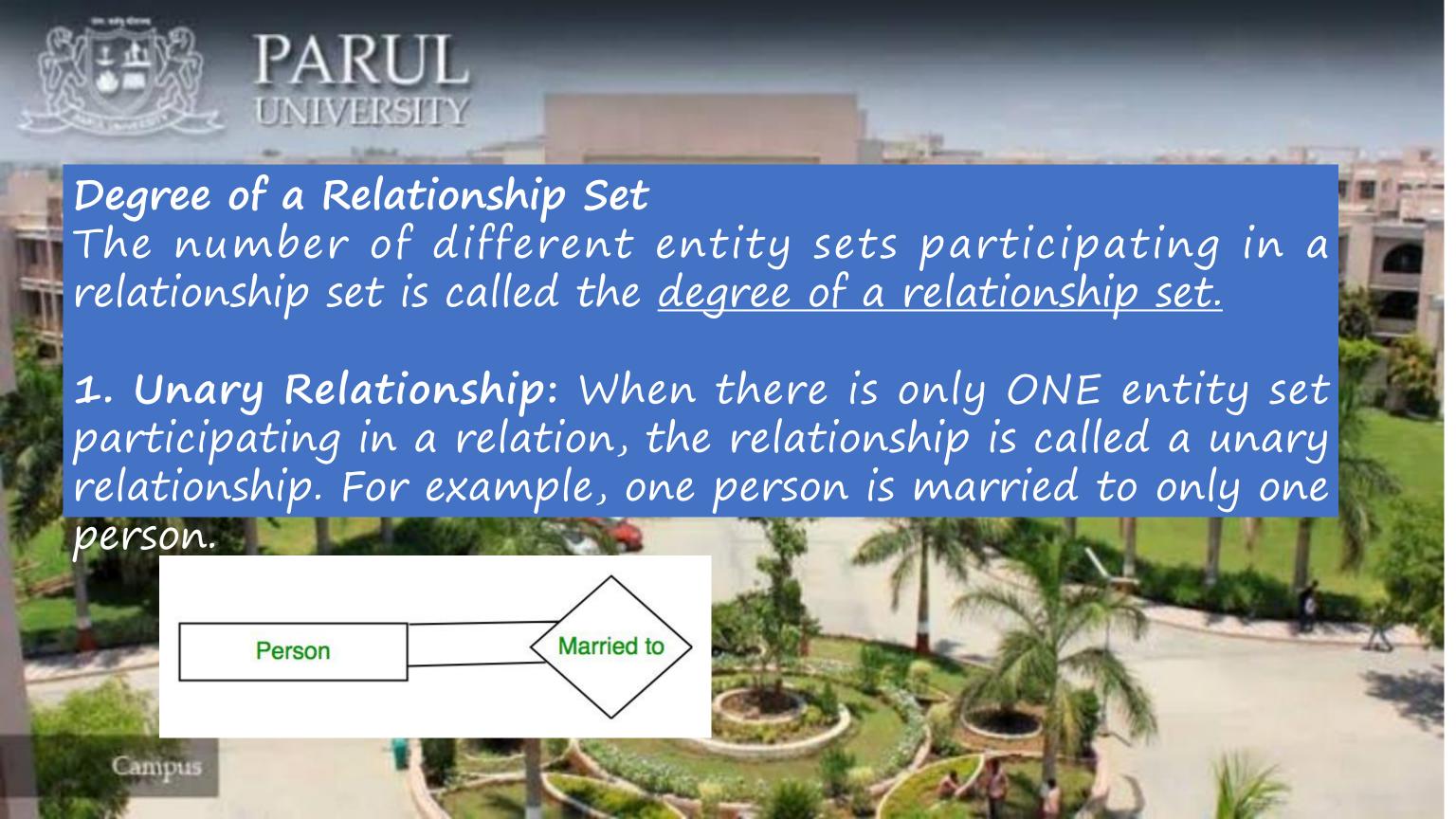
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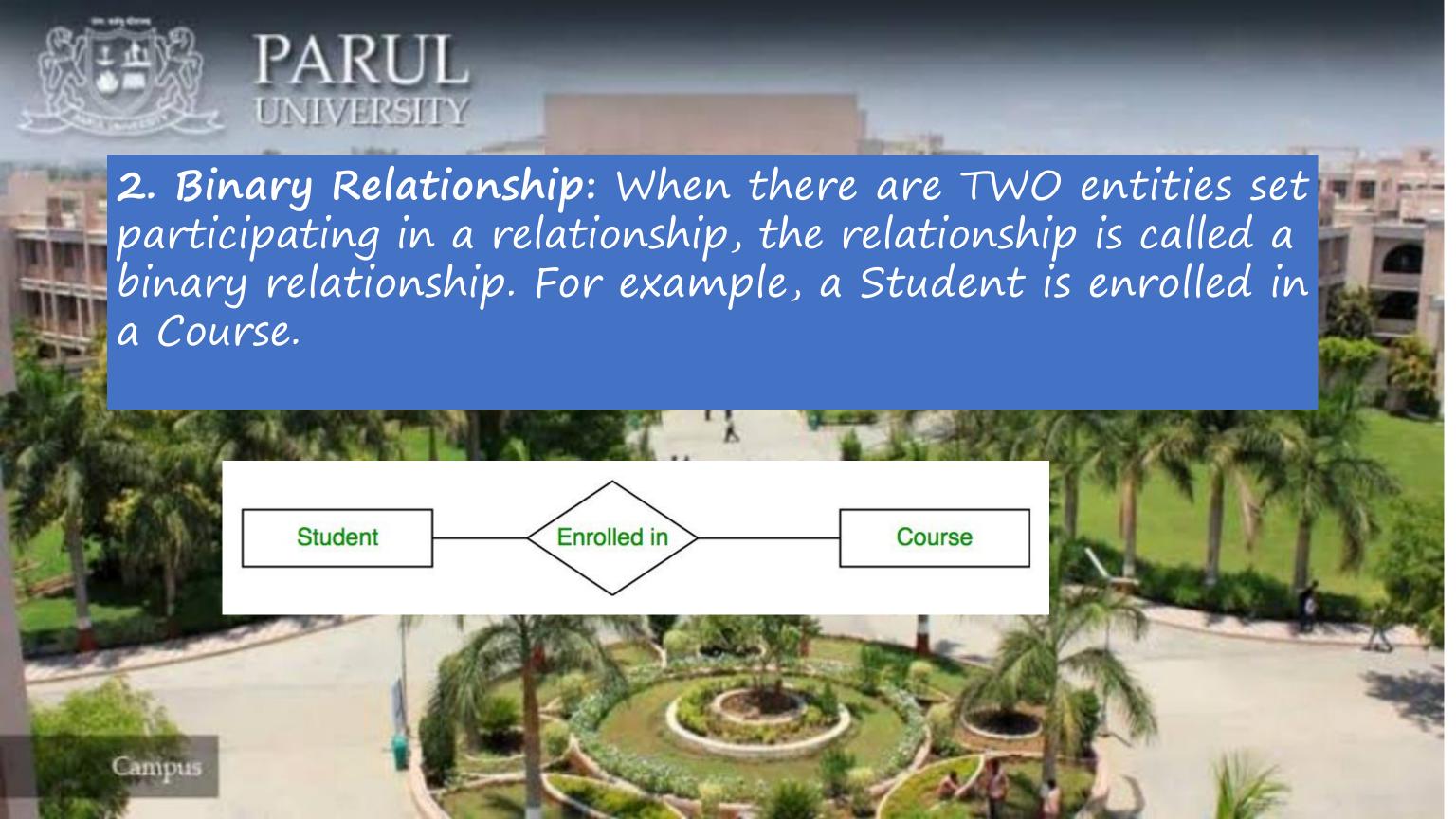
The Complete Entity Type

Student with its Attributes can be





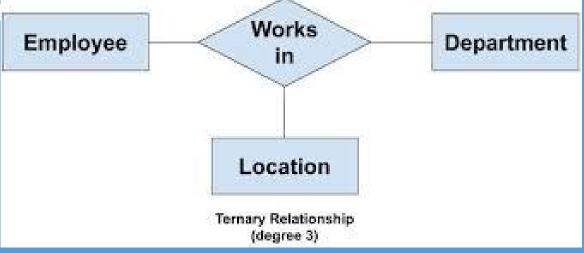




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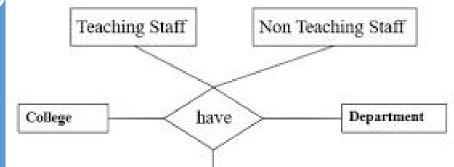
3. Ternary Relationship: When there are three entity sets participating in a relationship, the relationship is called a

ternary relation



4. N-ary Relationship: When there are n entities set participating in a relationship, the relationship is called an

n-ary relationship



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What is Cardinality?

The number of times an entity of an entity set participates in a relationship set is known as <u>cardinality</u> Cardinality can be of different types:

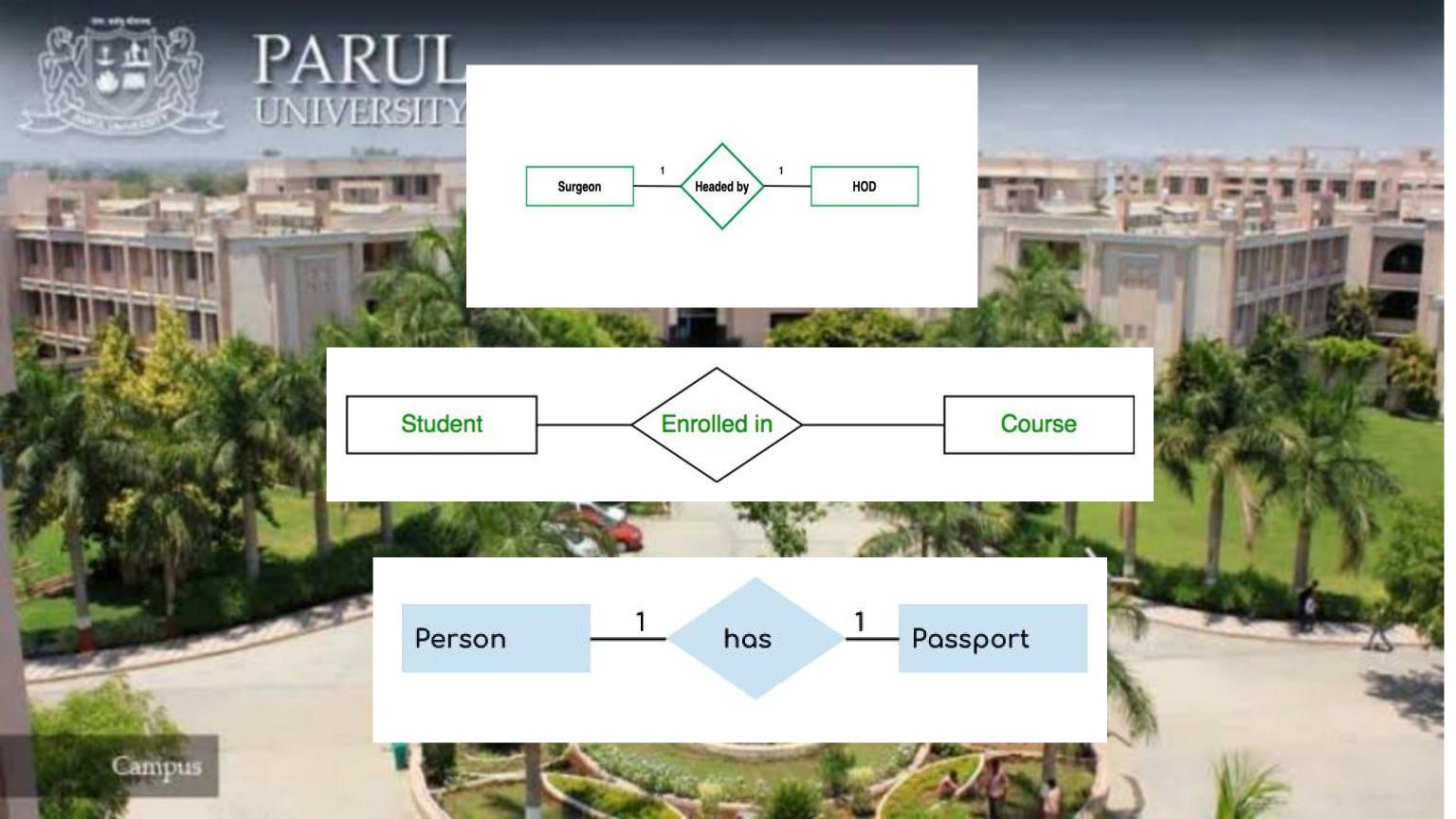
1. One-to-One: When each entity in each entity set can take part only once in the relationship, the cardinality is one-to-one. Let us assume that a male can marry one female and a female can marry one male. So the relationship will be one-to-one.

the total number of tables that can be u



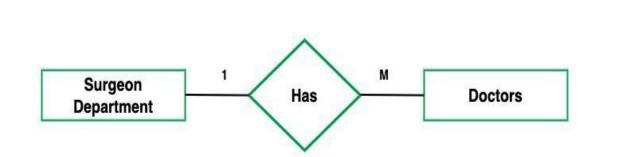
One-to-One Relationship



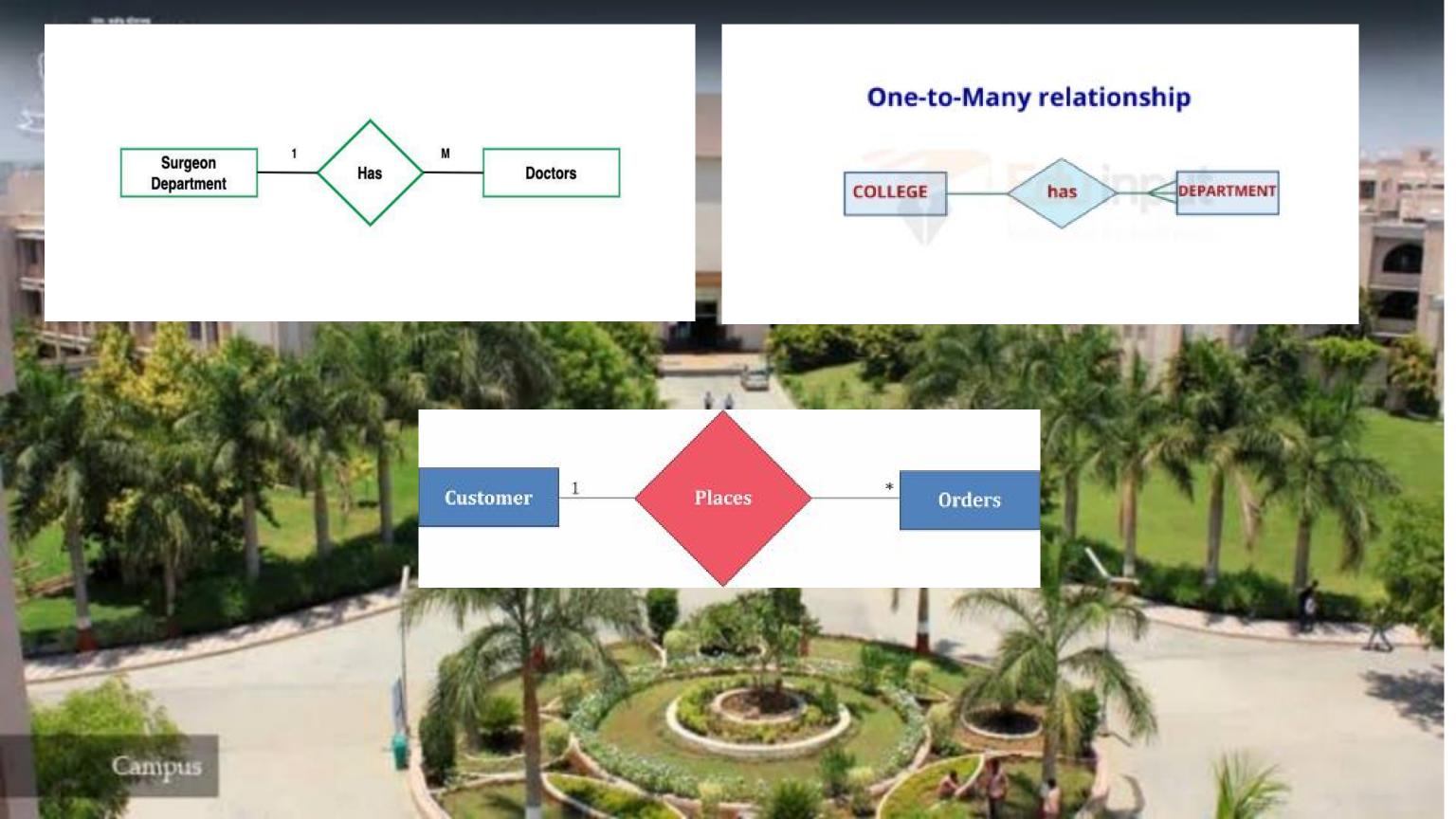




2. One-to-Many: In one-tomany mapping as well where each entity can be related to more than one entity and the total number of tables that can be used in this is 2. Let us assume that one surgeon department can accommodate many doctors. So the Cardinality will be 1 to M. It means one department



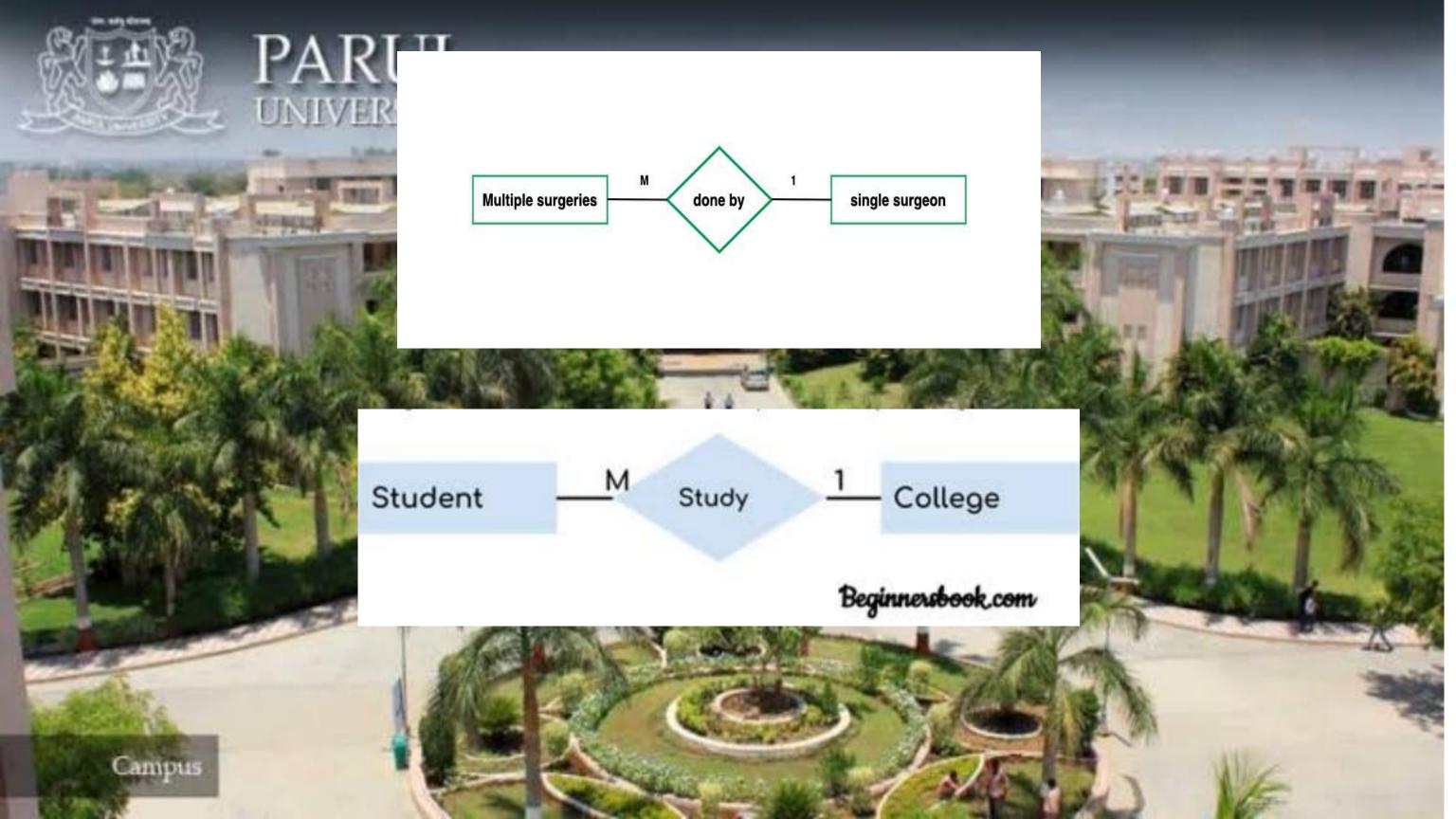






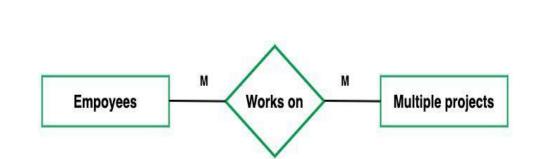
3. Many-to-One: When entities in one entity set can take part only once in the relationship set and entities in other entity sets can take part more than once in the relationship set, cardinality is many to one. Let us assume that a student can take only one course but one course can be taken by many students. So the cardinality will be n to 1. It means that for one course there can be n students but for one student, there will be only one course.

The total number of tables that can be

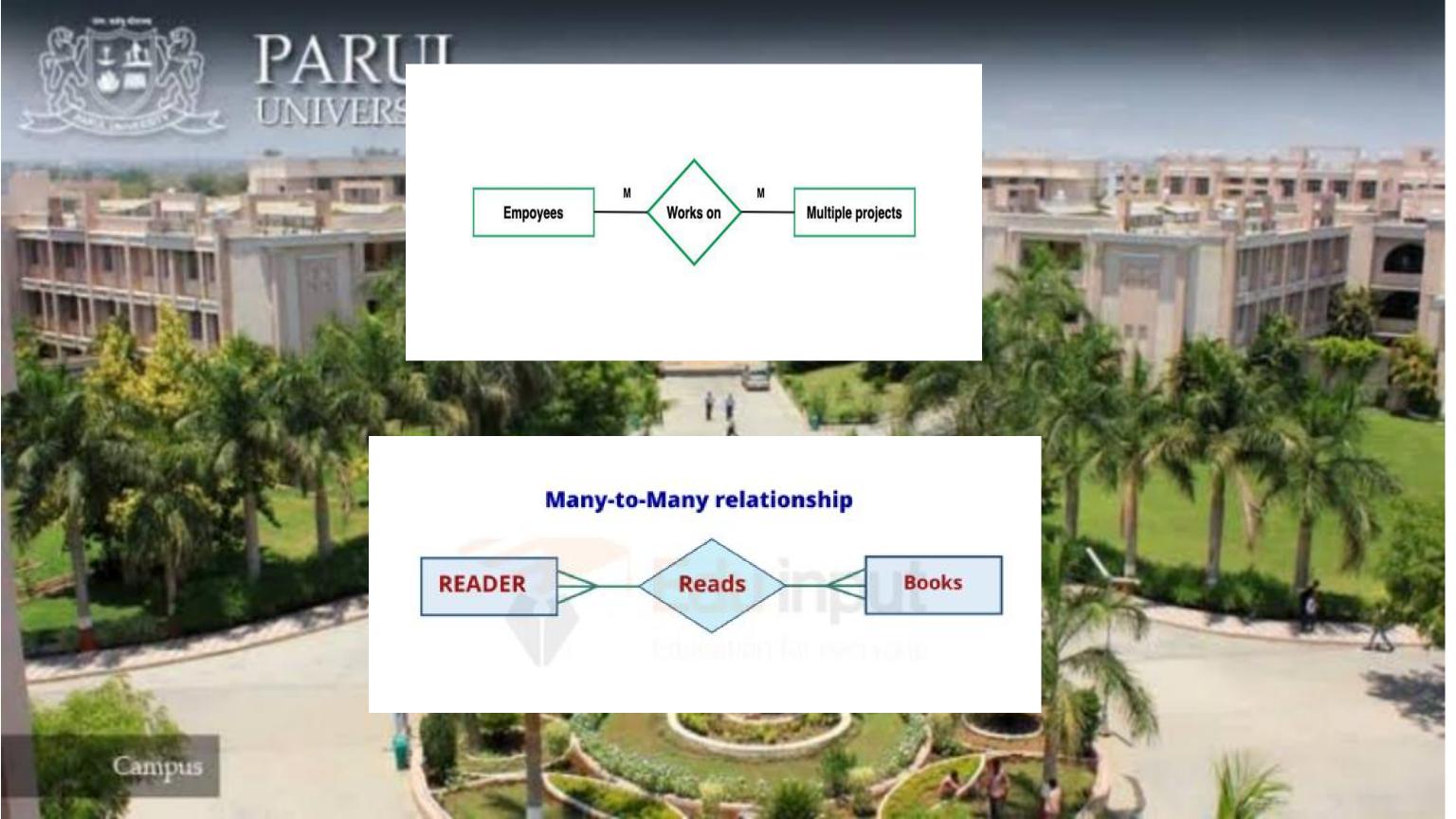


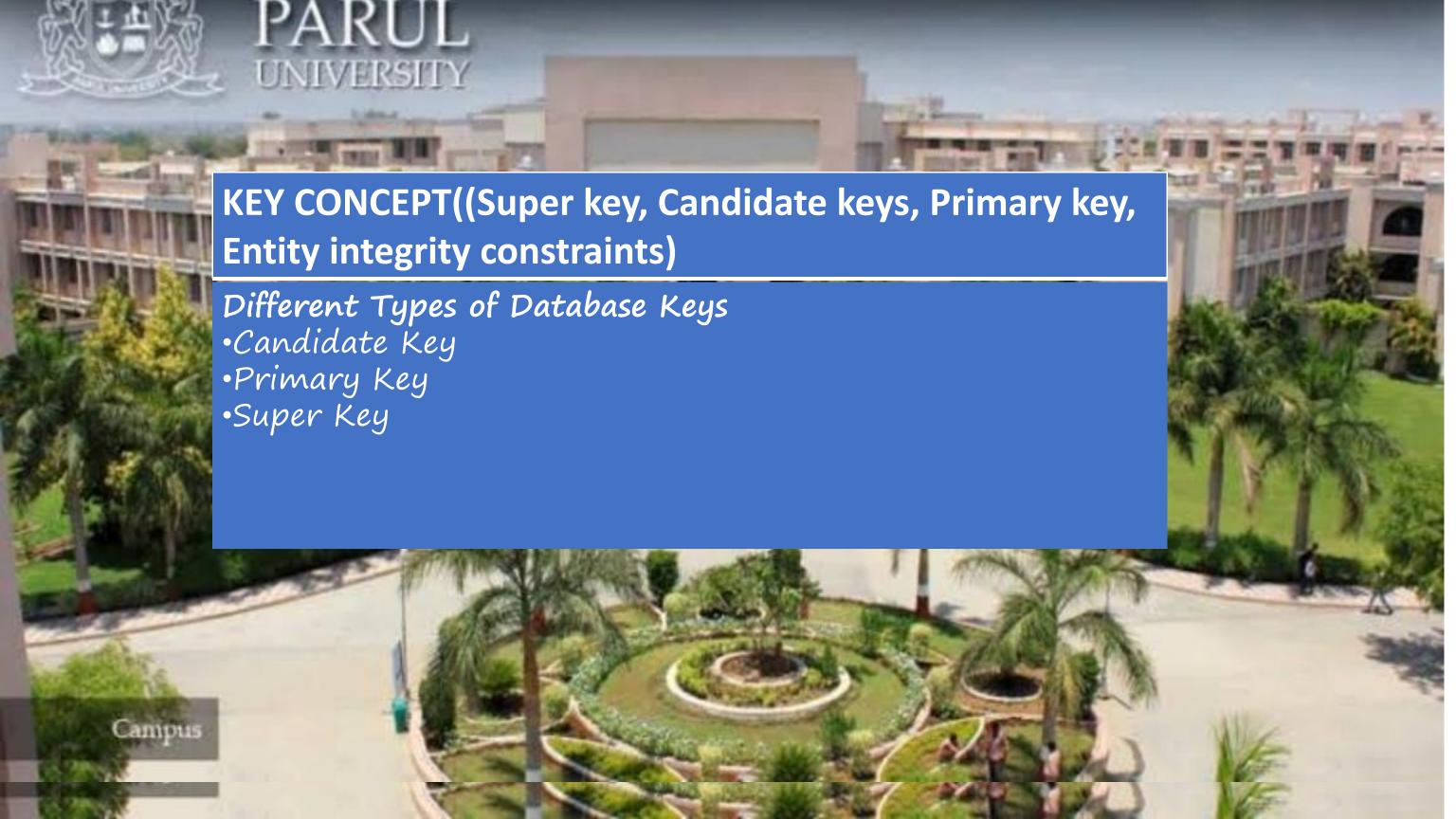
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4. Many-to-Many: When entities in all entity sets can take part more than once in the relationship cardinality is many to many. Let us assume that a student can take more than one course and one course can be taken by many students. So the relationship will be many to many. the total number of tables that can be used in this is 3.











Candidate Key

The minimal set of attributes that can uniquely identify a tuple is known as a candidate key. For Example, STUD_NO in STUDENT relation.

- ·It is a minimal super key.
- ·It is a super key with no repeated data is called a candidate key.
- •The minimal set of attributes that can uniquely identify a record.
- ·It must contain unique values.
- ·It can contain NULL values.
- ·Every table must have at least a single candidate key.
- ·A table can have multiple candidate keys but only one primary key.
- •The value of the Candidate Key is unique and may be null for a tuple.
- ·There can be more than one candidate key in a relationship



Example:

STUD_NO is the candidate key for relation STUDENT.

Table STUDENT

STUD_NO	SNAME	ADDRESS	PHONE
1	Shyam	Delhi	123456789
2	Rakesh	Kolkata	223365796
3	Suraj	Delhi	175468965

Example:

{STUD_NO, COURSE_NO} is a composite candidate key for relation STUDENT_COURSE.

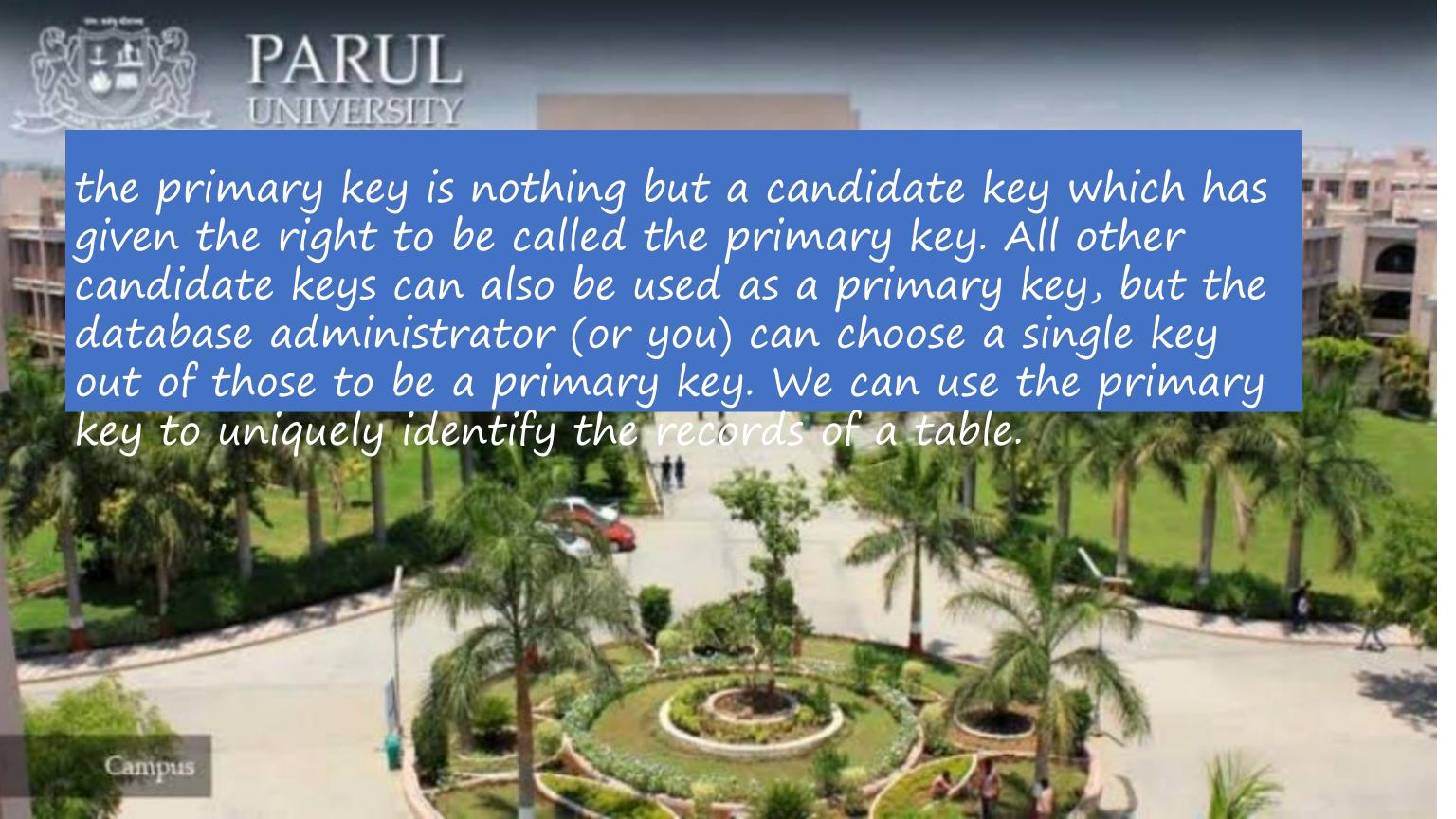


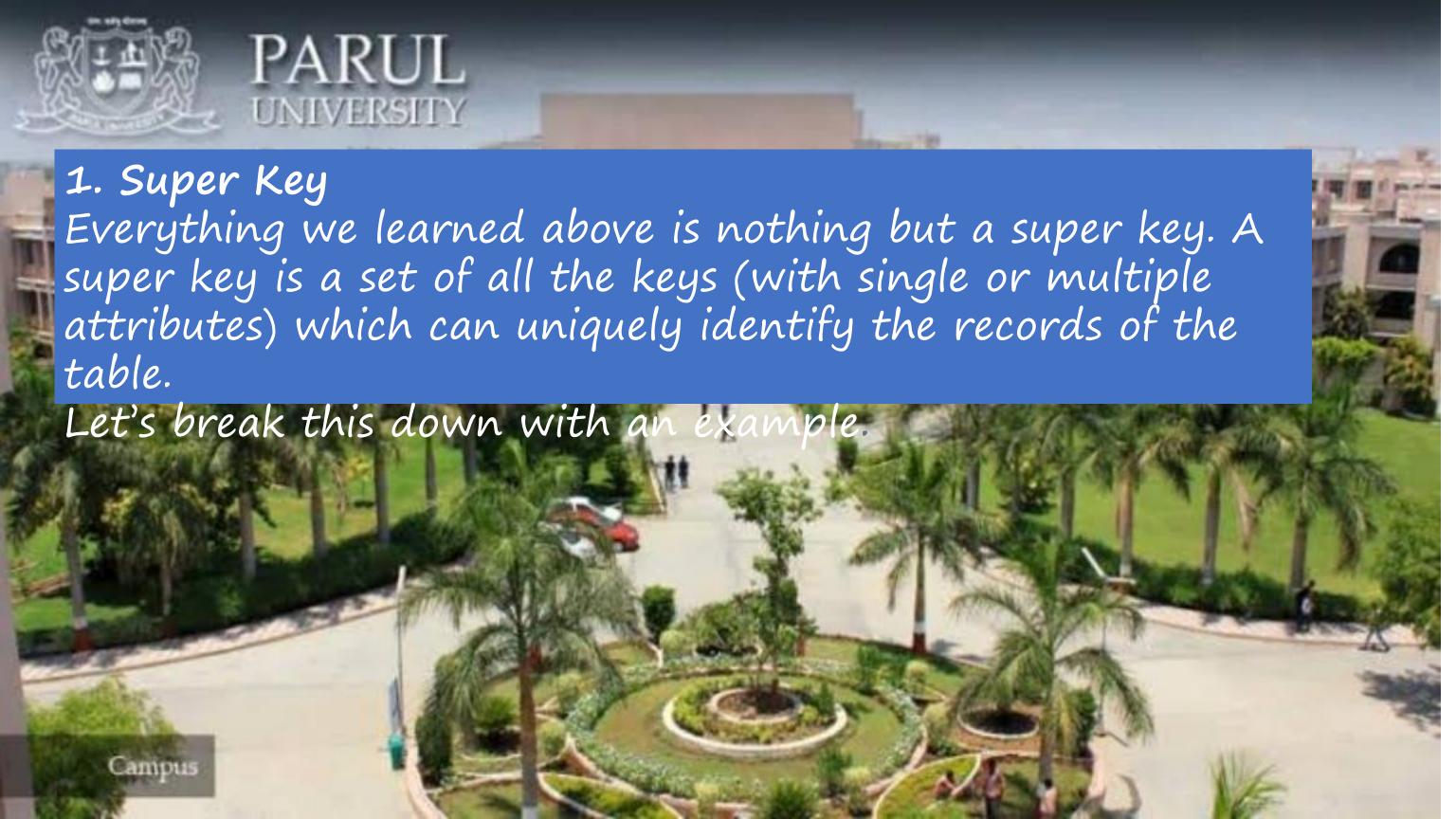
Note: In <u>SQL</u> Server a unique constraint that has a nullable column, allows the value 'null 'in that column only once. That's why the STUD_PHONE attribute is a candidate here, but can not be a 'null'

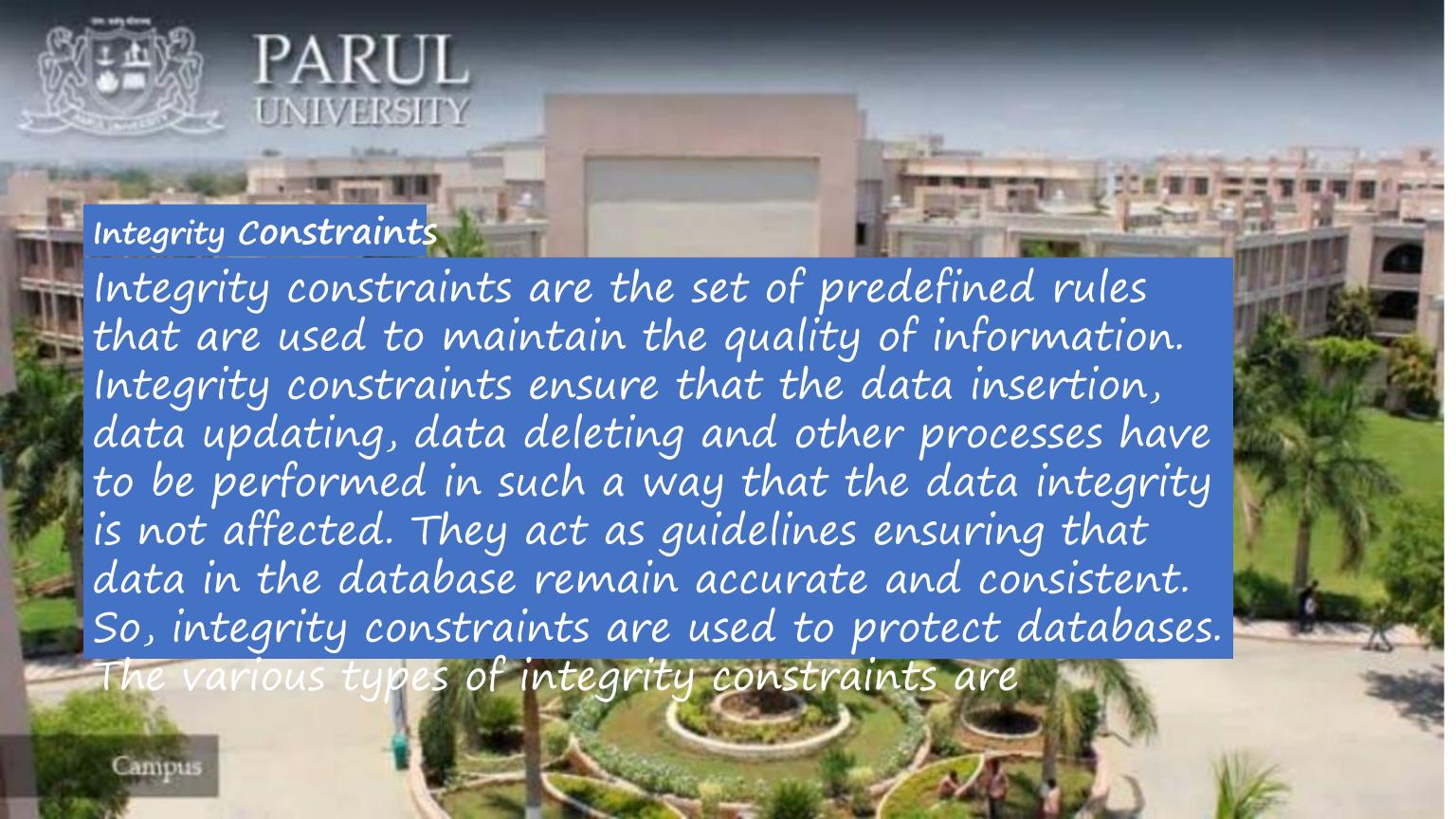
Primary Key

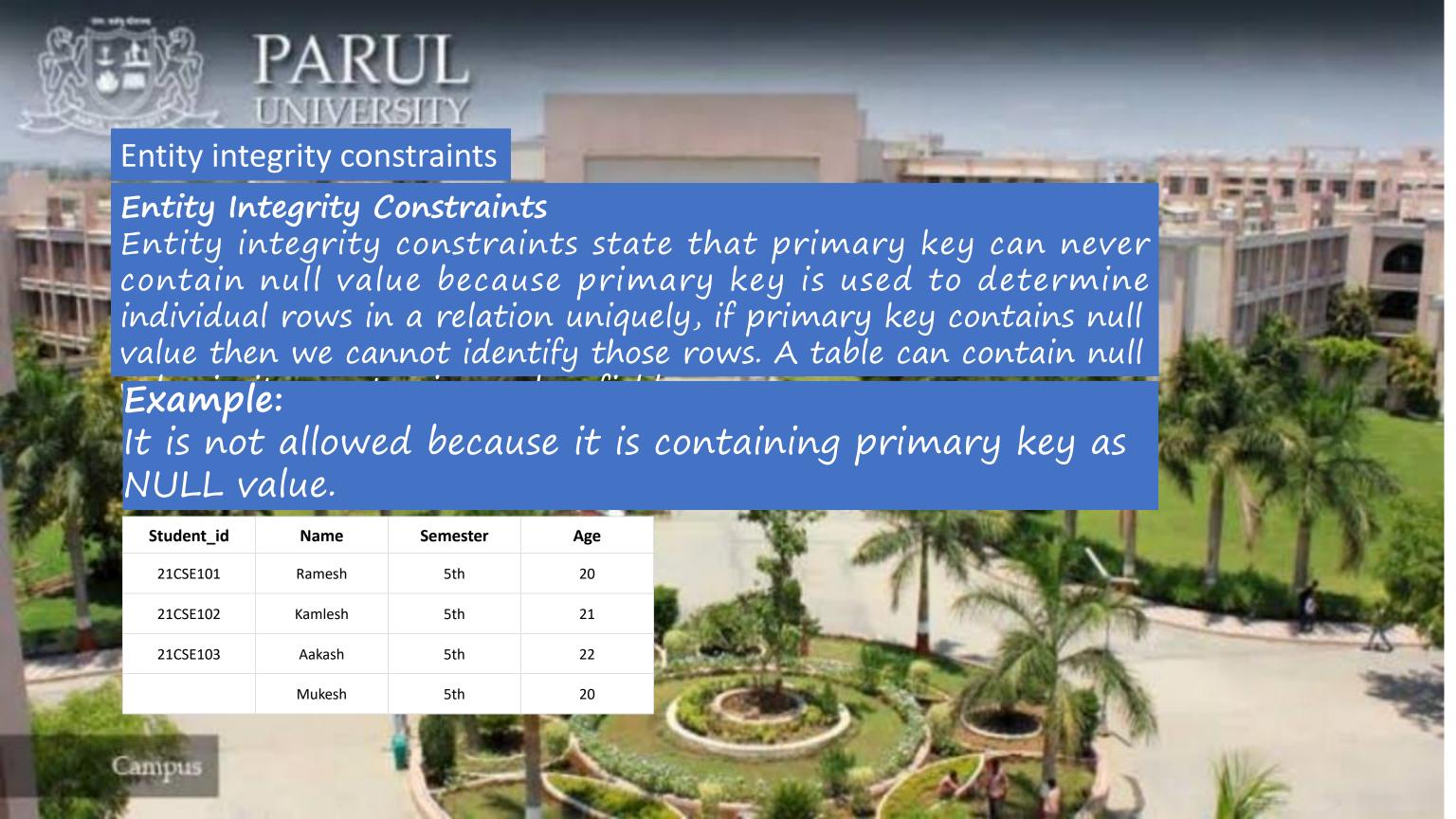
There can be more than one candidate key in relation out of which one can be chosen as the primary key. For Example, STUD_NO, as well as STUD_PHONE, are candidate keys for relation STUDENT but STUD_NO can be chosen as the <u>primary key</u> (only one out of many candidate keys).

- •It is a unique key.
- ·It can identify only one tuple (a record) at a time.
- ·It has no duplicate values, it has unique values.
- •It cannot be NULL.
- •Primary keys are not necessarily to be a single column; more than one column can also be a primary key for a table.

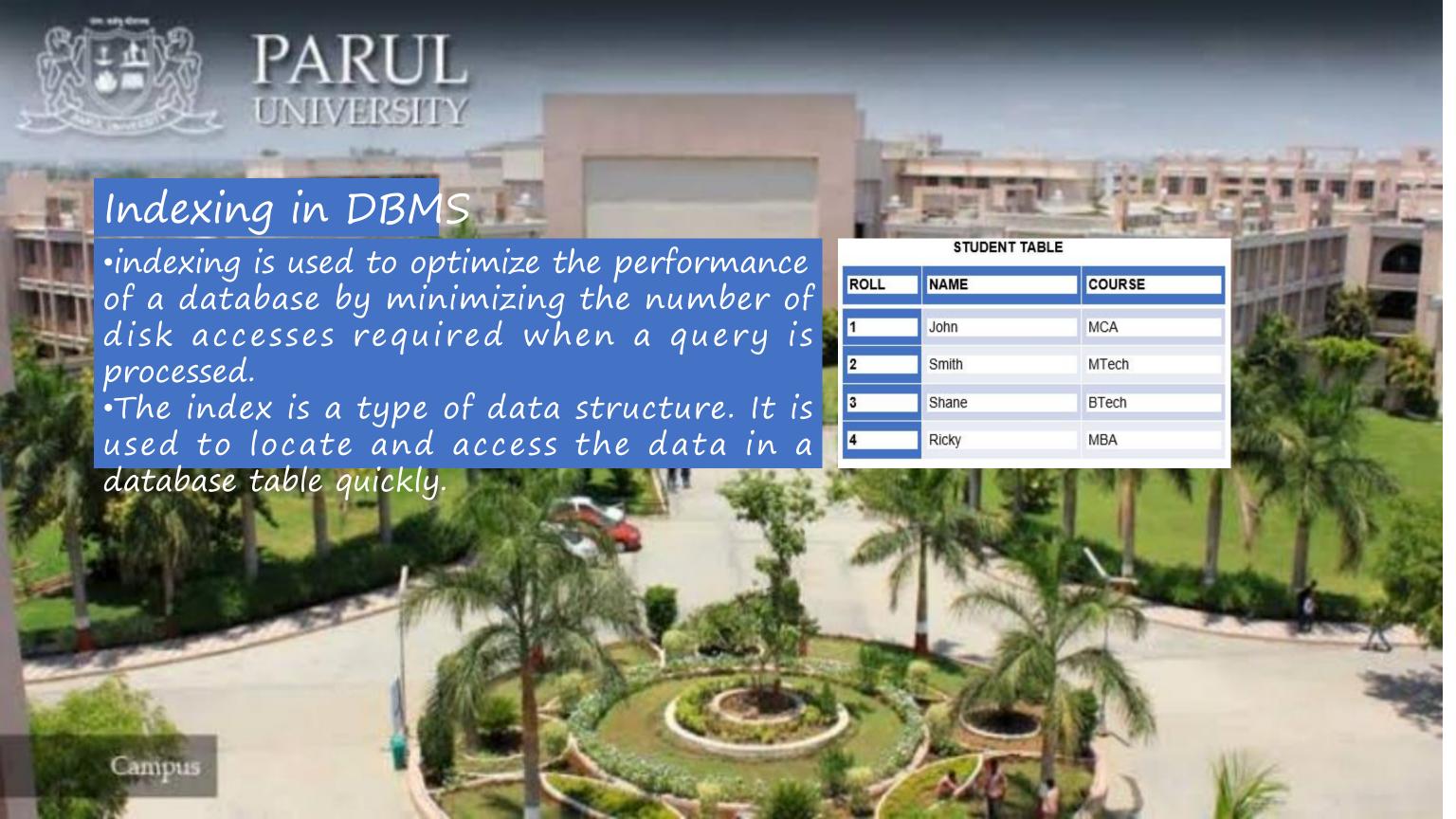






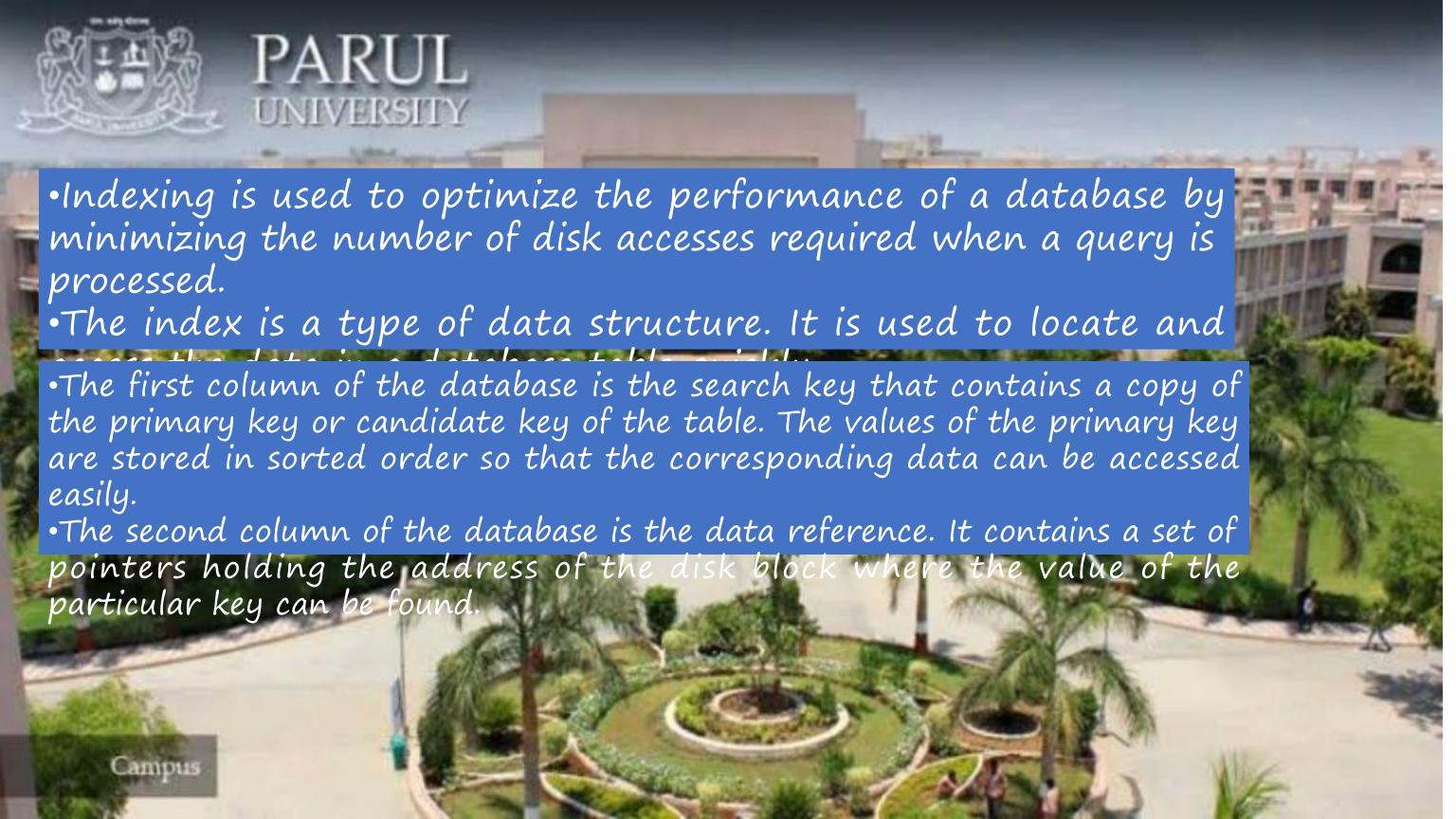












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Ordered indices

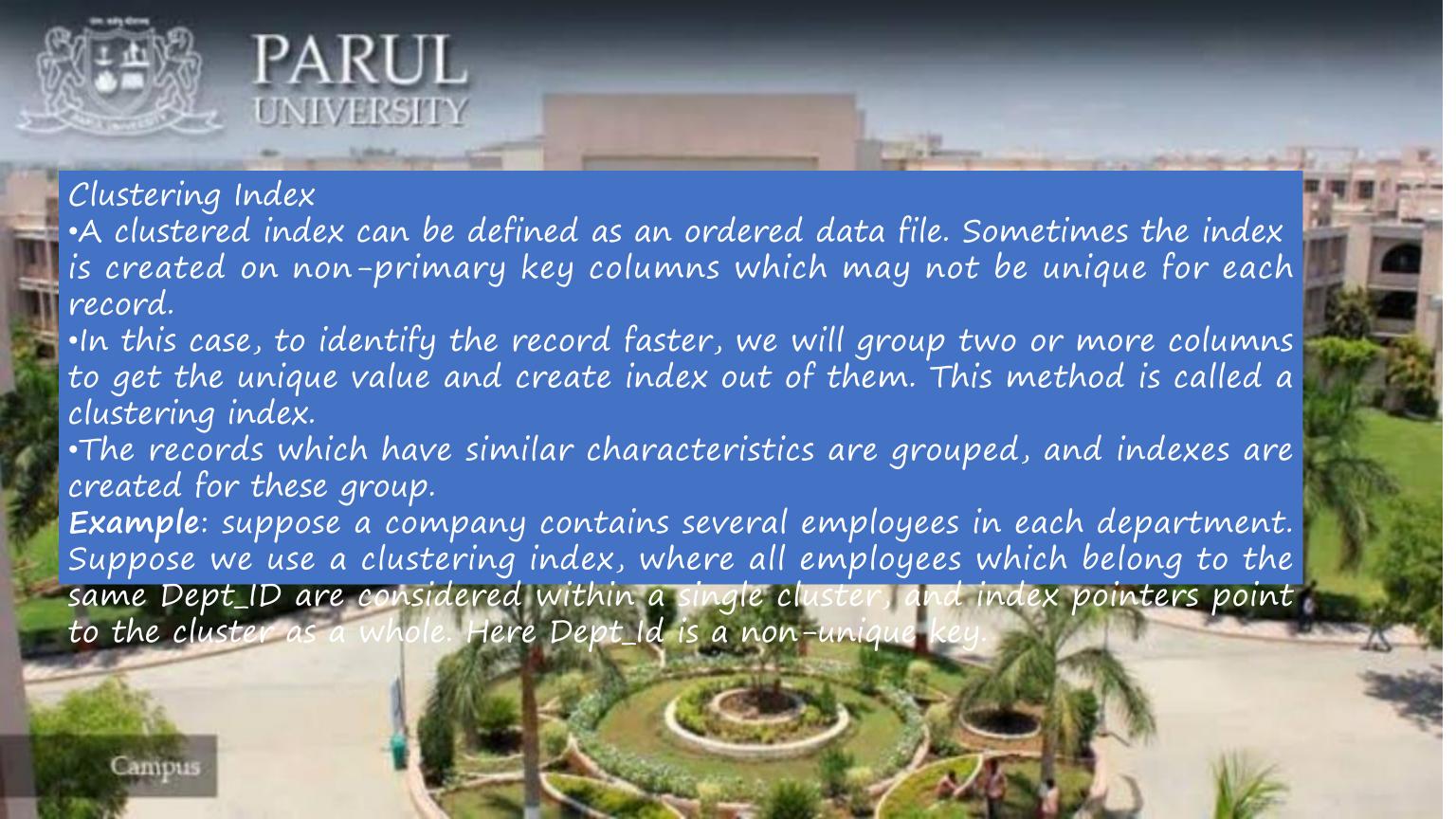
The indices are usually sorted to make searching faster. The indices which are sorted are known as ordered indices.

Example: Suppose we have an employee table with thousands of record and each of which is 10 bytes long. If their IDs start with 1, 2, 3....and so on and we have to search student with ID-543.

•In the case of a database with no index, we have to search the disk block from starting till it reaches 543. The DBMS will read the record after reading 543*10=5430 bytes.









In the sparse indexing, as the size of the table grows, the size of mapping also grows. These mappings are usually kept in the primary memory so that address fetch should be faster. Then the secondary memory searches the actual data based on the address got from mapping. If the mapping size grows then fetching the address itself becomes slower. In this case, the sparse index will not be efficient. To overcome this problem, secondary indexing is introduced.

In secondary indexing, to reduce the size of mapping, another level of indexing is introduced. In this method, the huge range for the columns is selected initially so that the mapping size of the first level becomes small. Then each range is further divided into smaller ranges. The mapping of the first level is stored in the primary memory, so that address fetch is faster. The mapping of the second level and actual data are stored in the

secondary memory (hard disk).

