Oracle Database Transactions

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- ► A database transaction consists of one of the following:
 - DML statements that constitute one consistent change to the data
 - One DDL statement
 - One data control language (DCL) statement

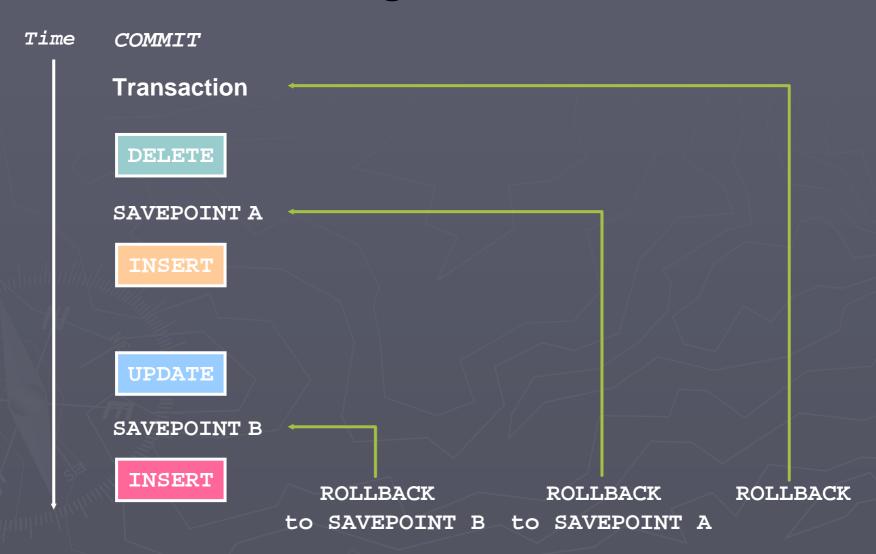
Oracle Database Transactions

- Begin when the first DML SQL statement is executed.
- ► End with one of the following events:
 - A COMMIT or ROLLBACK statement is issued.
 - A DDL or DCL statement executes (automatic commit).
 - The system crashes.

Advantages of COMMIT and ROLLBACK Statements

- ► With COMMIT and ROLLBACK statements, you can:
 - Ensure data consistency
 - Preview data changes before making changes permanent
 - Group logically related operations

Controlling Transactions



Rolling Back Changes to a Marker

- Create a marker in a current transaction by using the SAVEPOINT statement.
- Roll back to that marker by using the ROLLBACK TO SAVEPOINT statement.

```
UPDATE...
SAVEPOINT update_done;
Savepoint created.
INSERT...
ROLLBACK TO update_done;
Rollback complete.
```

Implicit Transaction Processing

- An automatic commit occurs under the following circumstances:
 - DDL statement is issued
 - DCL statement is issued
- An automatic rollback occurs under an abnormal termination of *i*SQL*Plus or a system failure.

State of the Data Before COMMIT or ROLLBACK

- The previous state of the data can be recovered.
- The current user can review the results of the DML operations by using the SELECT statement.
- Other users cannot view the results of the DML statements by the current user.
- The affected rows are *locked*; other users cannot change the data in the affected rows.

State of the Data After COMMIT

- Data changes are made permanent in the database.
- The previous state of the data is permanently lost.
- All users can view the results.
- Locks on the affected rows are released; those rows are available for other users to manipulate.
- All savepoints are erased.

Committing Data

Make the changes:

```
DELETE FROM employees
WHERE employee_id = 99999;
1 row deleted.

INSERT INTO departments
VALUES (290, 'Corporate Tax', NULL, 1700);
1 row created.
```

Commit the changes:

```
COMMIT;
Commit complete.
```

State of the Data After ROLLBACK

- ► Discard all pending changes by using the ROLLBACK statement:
 - Data changes are undone.
 - Previous state of the data is restored.
 - Locks on the affected rows are released.

```
DELETE FROM copy_emp;

22 rows deleted.

ROLLBACK;

Rollback complete.
```

State of the Data After ROLLBACK

```
DELETE FROM test;
25,000 rows deleted.
ROLLBACK;
Rollback complete.
DELETE FROM test WHERE id = 100;
1 row deleted.
SELECT * FROM test WHERE id = 100;
No rows selected.
COMMIT;
Commit complete.
```

Statement-Level Rollback

- If a single DML statement fails during execution, only that statement is rolled back.
- The Oracle server implements an implicit savepoint.
- All other changes are retained.
- The user should terminate transactions explicitly by executing a COMMIT or ROLLBACK statement.

Read Consistency

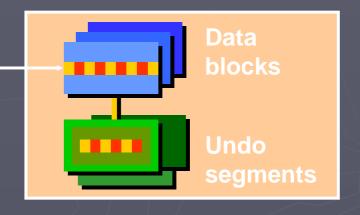
- ► Read consistency guarantees a consistent view of the data at all times.
- Changes made by one user do not conflict with changes made by another user.
- Read consistency ensures that on the same data:
 - Readers do not wait for writers
 - Writers do not wait for readers

Implementation of Read Consistency

User A

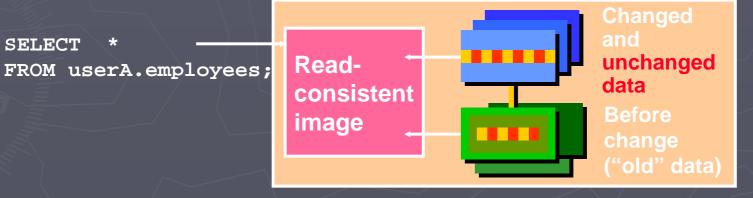


```
UPDATE employees
SET salary = 7000 —
WHERE last name = 'Grant';
```





User B

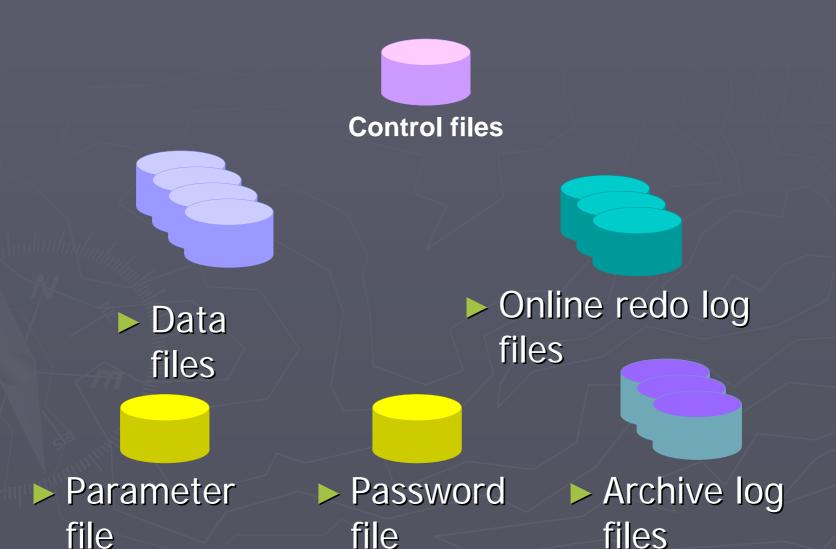


Oracle Architectural Components

Oracle DB Architecture: Overview

- ► The Oracle database consists of two main components:
 - The database or the physical structures
 - The instance or the memory structures

Database Physical Architecture



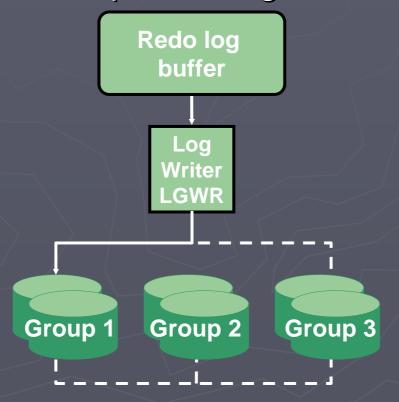
Control Files

- Contains physical database structure information
- Multiplexed to protect against loss
- Read at mount stage



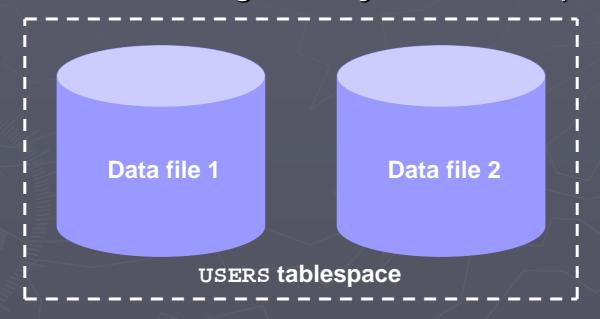
Redo Log Files

- Record changes to the database
- Multiplexed to protect against loss



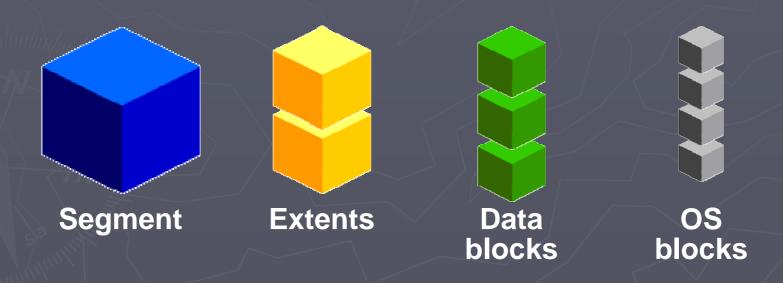
Tablespaces and Data Files

- Tablespaces consist of one or more data files.
- Data files belong to only one tablespace.

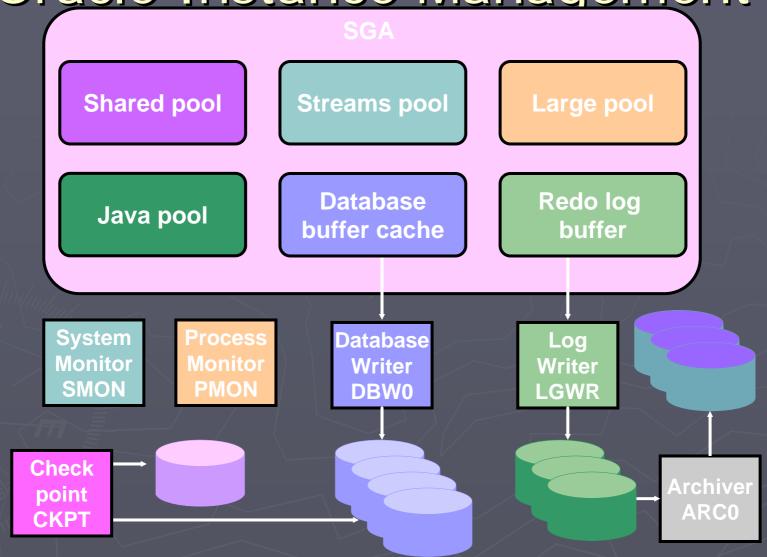


Segments, Extents, and Blocks

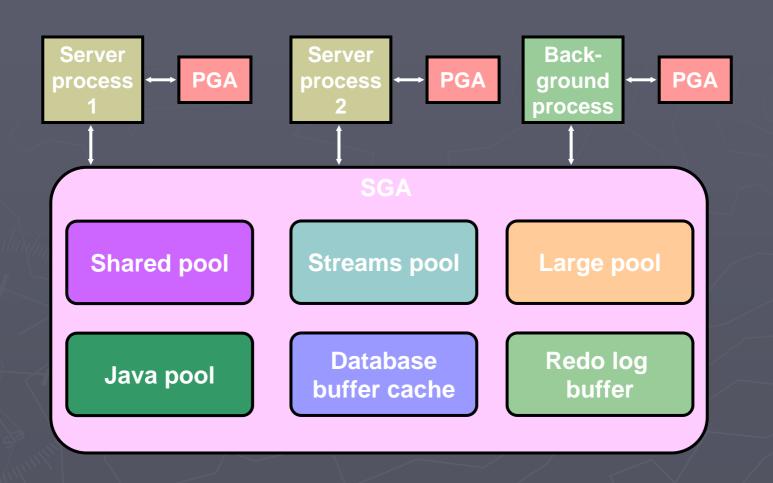
- Segments exist within a tablespace.
- Segments consist of a collection of extents.
- Extents are a collection of data blocks.
- Data blocks are mapped to OS blocks.



Oracle Instance Management



Oracle Memory Structures



Oracle Processes

Server process

Server process

Server process

Server process

System Global Area SGA

System monitor SMON

Process monitor PMON

Database writer DBW0

Check point CKPT

Log writer LGWR

Archive

Background processes

Other Key Physical Structures



Processing a SQL Statement

- Connect to an instance using:
 - ► The user process
 - ► The server process
- The Oracle server components that are used depend on the type of SQL statement:
 - Queries return rows
 - ► DML statements log changes
 - ► Commit ensures transaction recovery
- Some Oracle server components do not participate in SQL statement processing.

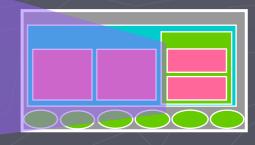
Processing a Query

- Parse:
 - ► Search for identical statement
 - ► Check syntax, object names, and privileges
 - ► Lock objects used during parse
 - Create and store execution plan
- Execute: Identify rows selected
- Fetch: Return rows to user process

The Shared Pool

- The library cache contains the SQL statement text, parsed code, and execution plan.
- The data dictionary cache contains table, column, and other object definitions and privileges.
- The shared pool is sized by SHARED_POOL_SIZE.

Shared pool Library cache Data dictionary cache



Program Global Area (PGA)

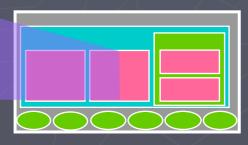
- Not shared
- Writable only by the server process
- Contains:
 - ► Sort area
 - **▶** Session information
 - **►** Cursor state
 - ► Stack space



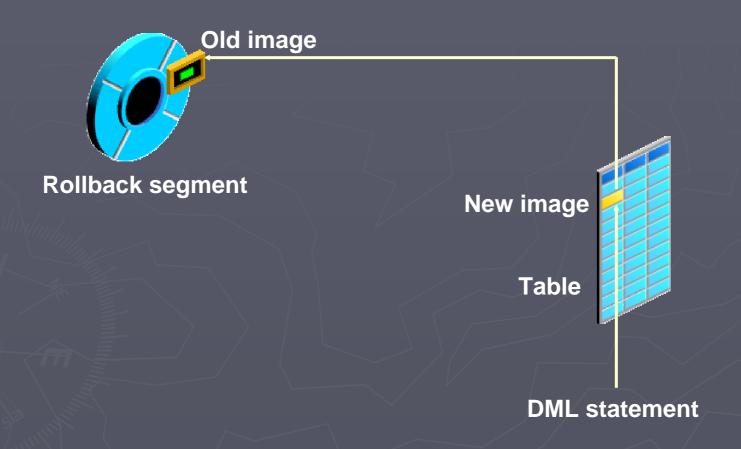
Redo Log Buffer

- Has its size defined by LOG_BUFFER
- Records changes made through the instance
- Is used sequentially
- Is a circular buffer

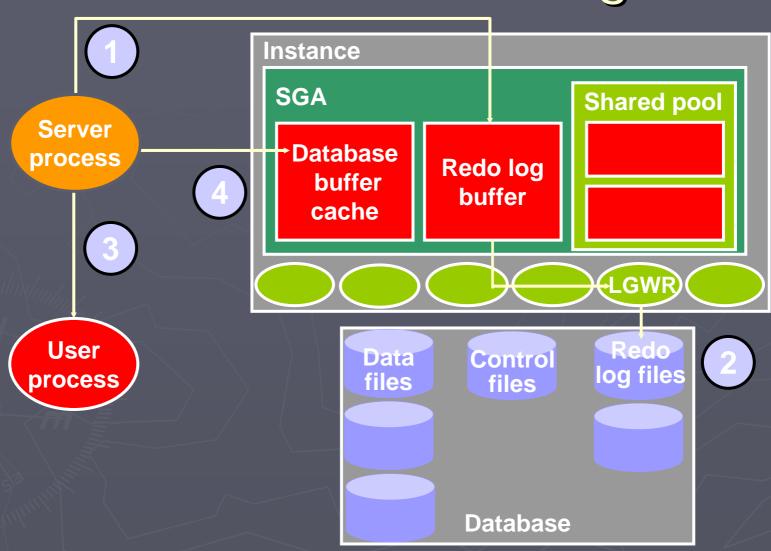
Database buffer cache



Rollback Segment



COMMIT Processing



CAP Theorem

- Three properties of a system
 - Consistency (all copies have same value)
 - Availability (system can run even if parts have failed)
 - Partitions (network can break into two or more parts)
- Brewer's CAP "Theorem": You can have at most two of these three properties for any system