

SQL-4

Agenda

~~①~~ Views ←

① Indexing ←

② Transactions ←

Naman Bhalu

→ SWE @ Google
Campus
Employee

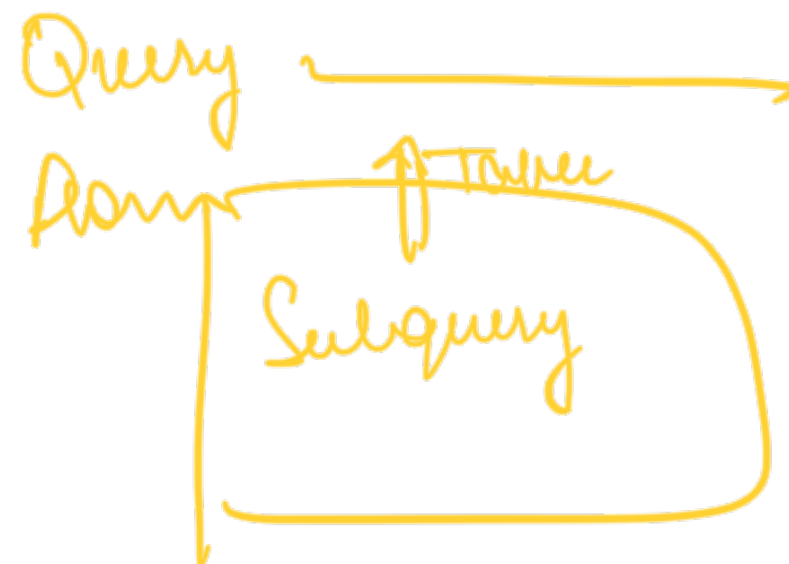
Subquery

→ SELECT

→ FROM

→ WHERE

→ result of any SQL query is like a table



Scalar's DB



Students			
id	name	b-id	pop

Teachers		
b-id	name	t-id

Teaching	
id	name

Qⁿ Give average pop for all students being taught by a teacher

name	pop
Naman	87
Satyasai	95
Tanavay	89
...	...

Module 100

X

<u>b_id</u>	<u>t_name</u>	<u>st_id</u>	<u>t_id</u>	<u>psp</u>
	⇒			

4
Select avg(psp), t_id
from ~~co~~
group by t_id

↓ ↓

Select s.id, t.id, t.name, s.psp, b.id
from students s
join batches b
on s.t_id = b.t_id

on s.b-id = t.id
join teachers t
on b.t-id = t.id

↓ Cross Join

Select

from students teachers

where _____

Q^m

We want to get average psp of every batch.

```
Select avg(bsp), b-id  
from X    
group by b-id
```



View

⇒ View is like a query stored in
my database

⇒ View acts like a table.

Transactions

→ a group of ^{SQL} statements that we want to be run atomically ← Complete / None

Bank

eg of transaction

transfer 200 rs from A to B

1. Get current balance of A

② Temp = A

③ Put balance of A = $A - 200$

④ Get current balance of B

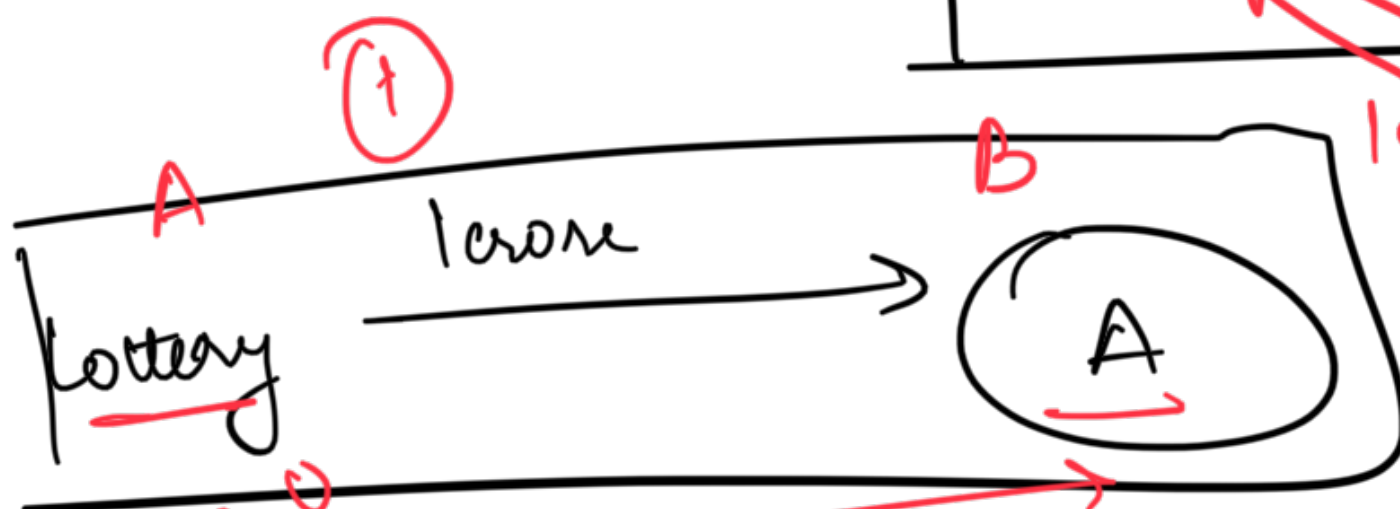
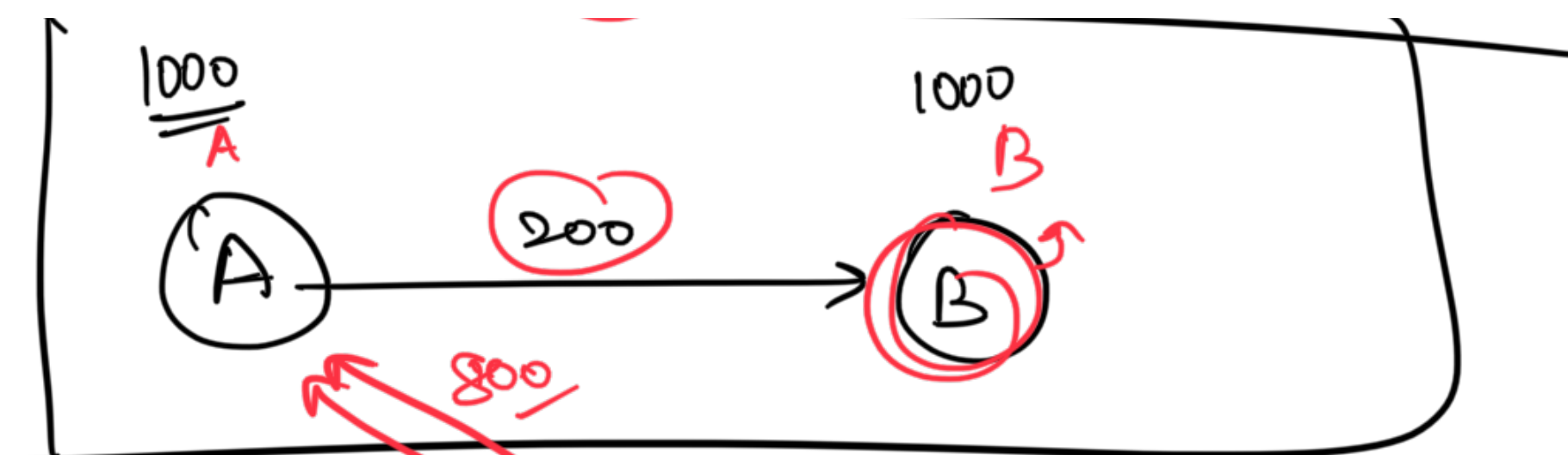
⑤ Temp = B

⑥ Put new balance of B = $B + 200$

Everything
or
nothing

A C I D ^{▷ Durability}
↓ ↓ ↓
Atomcity Consistency Isolation

②



transfer (from to, to amount)

← read current value of from

← read current value of to

⇒ update to = to + amount

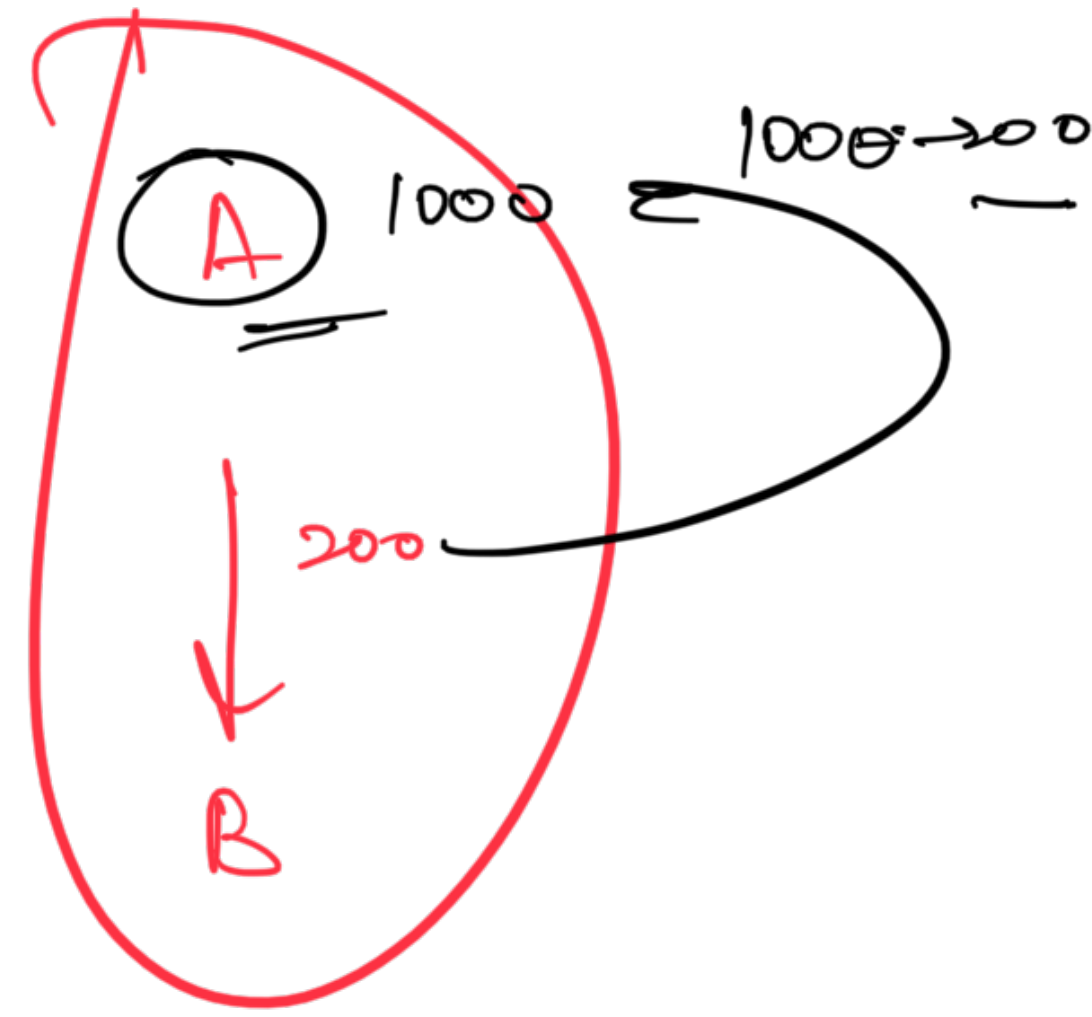
update from = from - amount

1 cr
~~1000~~

1000

1000
1000

Ans
↓
800



→ Cyclic request for locks

T1

|

T2

|

→ gets lock over 3

gets lock over 7

request lock over 7

→ T1 is waiting for
t2 to complete

request lock over 3

→ T2 is waiting for T1
to complete

→ deadlock

→ In MySQL whenever detects deadlock it
fails the earlier transaction



A → B fail

B → C success

T1 has lock ~~X~~

T2 requests lock which leads to deadlock



CN → TTS → Naman
→ MANG → Town