

Q

Given a LL. Find the middle node.

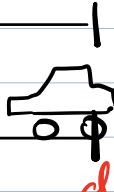
→ Slow & fast pointer.



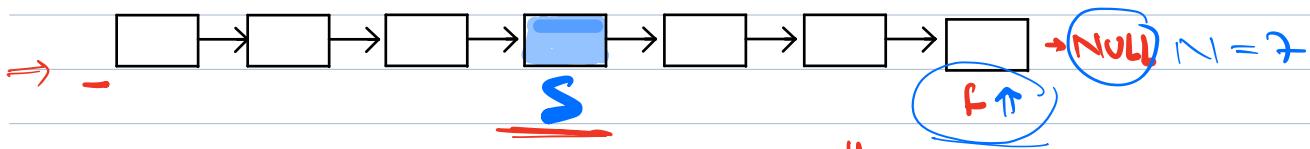
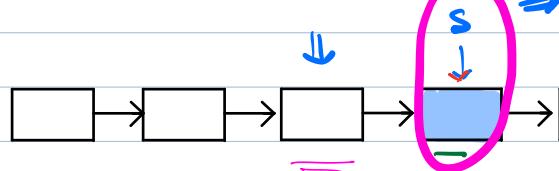
middle



s



$N = 6$



(fast == NULL OR fast Next == NULL) \Rightarrow stop

Node get_Mid (head) &

slow = head;

fast = head

while (fast != NULL && fast! Next != NULL) {

 slow = slow.next;

 fast = fast.next.next;

 next =

val

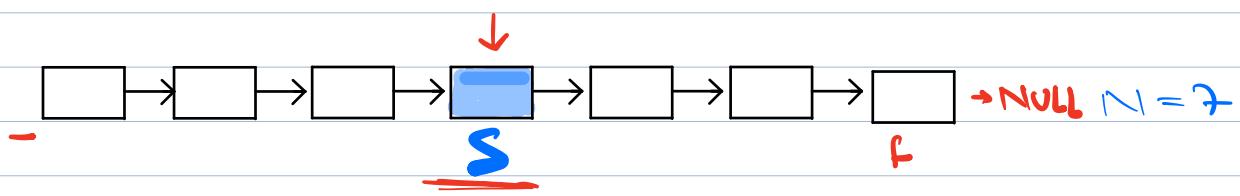
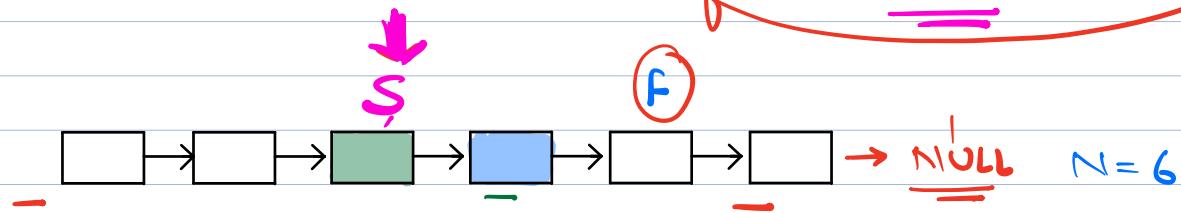
}

return slow;

$!= \underline{\text{NULL}}$

3

fast.next.next = NULL



while (fast.next != NULL && fast.next.next != NULL)



Given 2 sorted linked lists.

Amazon

Merge both of them to create a new sorted linked list

s.c. $\underline{\underline{O(1)}}$

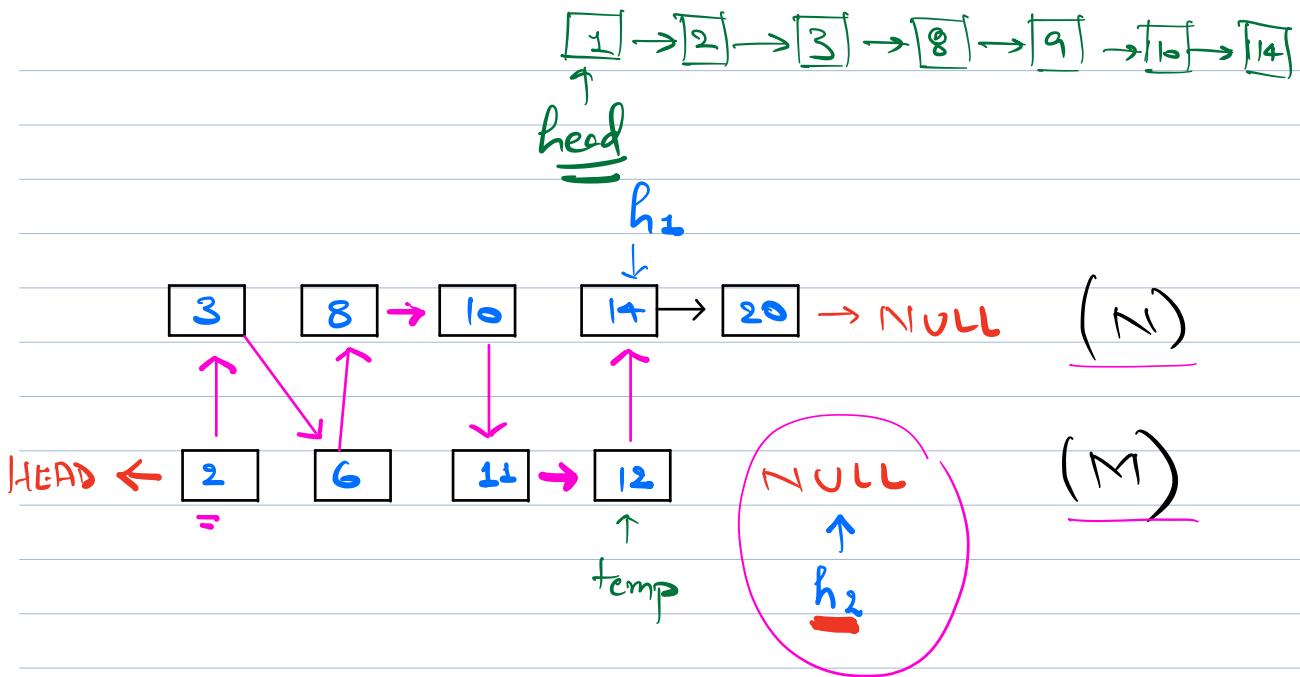
Heads.



Head2



head 1
head 2



Code

```
Node merge ( h1 , h2 ) {
```

```
    if ( h1.val < h2.val ) {
```

Head = h₁;

h₁ = h₁.next;

```
} else {
```

Head = h₂

h₂ = h₂.next;

```
}
```

Initialisation

```
temp = Head;
```

```
while ( h1 != NULL && h2 != NULL ) {
```

```
    if ( h1.val < h2.val ) {
```

temp.next = h₁;

$h_1 = h_1.\text{next};$
 $\text{temp} = \text{temp}.\text{next};$

{
else if

$\text{temp}.\text{next} = h_2;$
 $h_2 = h_2.\text{next};$
 $\text{temp} = \text{temp}.\text{next};$

}

if ($h_1 == \text{NULL}$) {
 $\text{temp}.\text{next} = h_2;$

{

else if

$\text{temp}.\text{next} = h_1;$

}

return Head

{

T.C. = $O(N+M)$

S.C. = $O(1)$

Google
MS
Amazon
LinkedIn

Given a Linked List.

Sort the linked list using
Merge Sort.

Merge Sort () {
mid

mergeSort(low, mid)
mergeSort(mid+1, high)
merge();

}

Code

□ → NULL

Node mergeSort (head) {

Base Case [if (head == NULL || head.next == NULL) {
 return head;
}

Node mid = get 1st Mid (head); O(N)

Node h2 = mid.next;
mid.next = NULL

Node h1 = mergeSort (head);
h2 = mergeSort (h2);

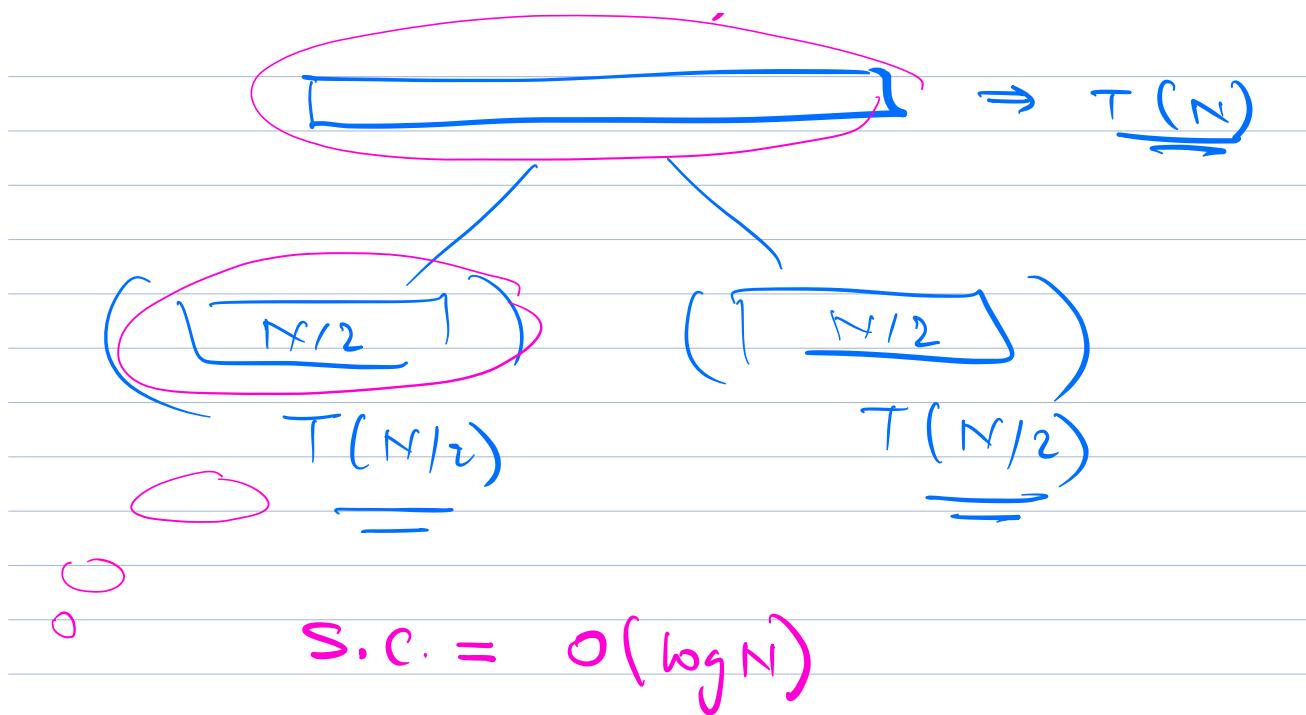
head = merge (h1, h2); O(N)

return head;

}

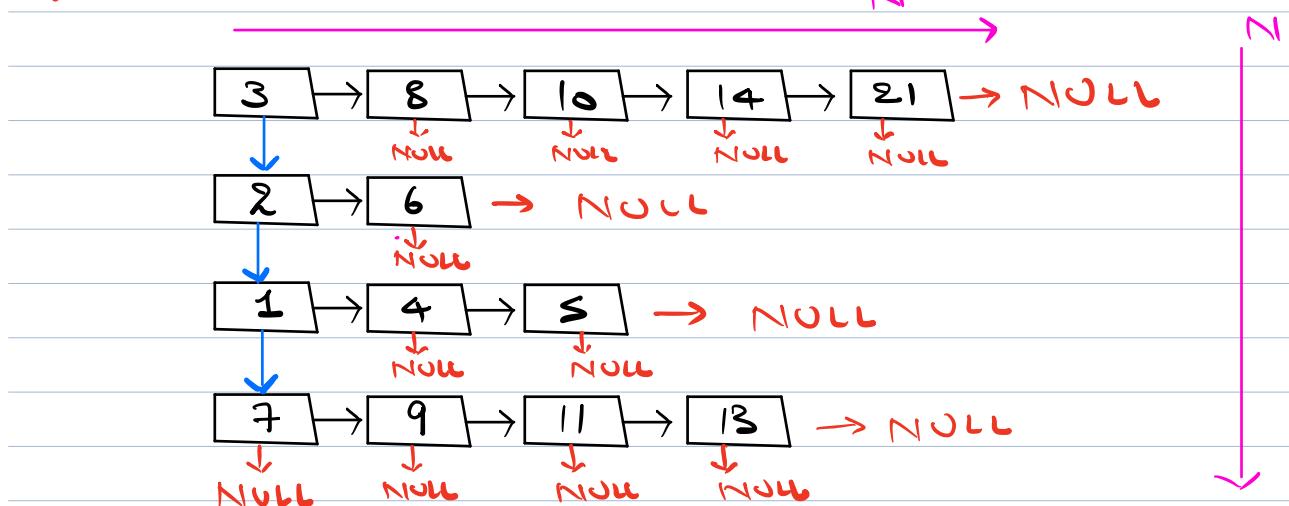
T.C. = O(N log N)
↓

T(N) = 2T(N/2) + O(N)



Break till 10.58

~~Google~~ Given a $2D$ list. flatten it to a single list (sorted) $\xrightarrow{N \times N}$ (sorted horizontally). S.C. $O(1)$



Node of

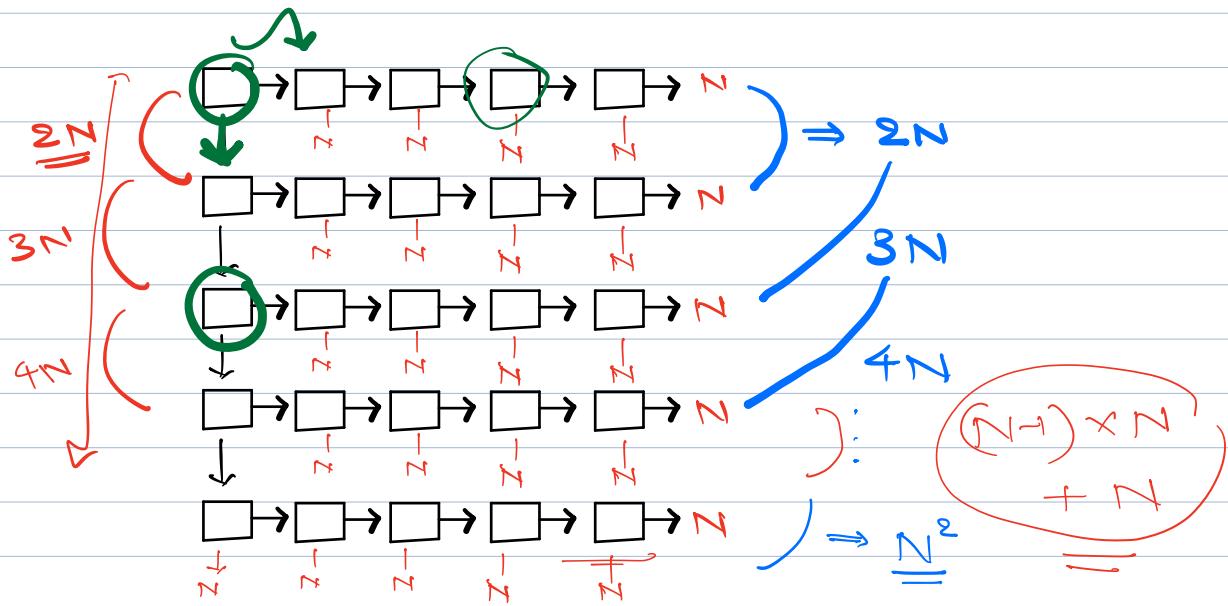
$1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7 \rightarrow 8 \rightarrow 9 \rightarrow 10$
 $11 \rightarrow 13 \rightarrow 14 \rightarrow 21$

```

int val;
Node next;
Node down;

```

k



⇒ Merge 2 lists at a time.

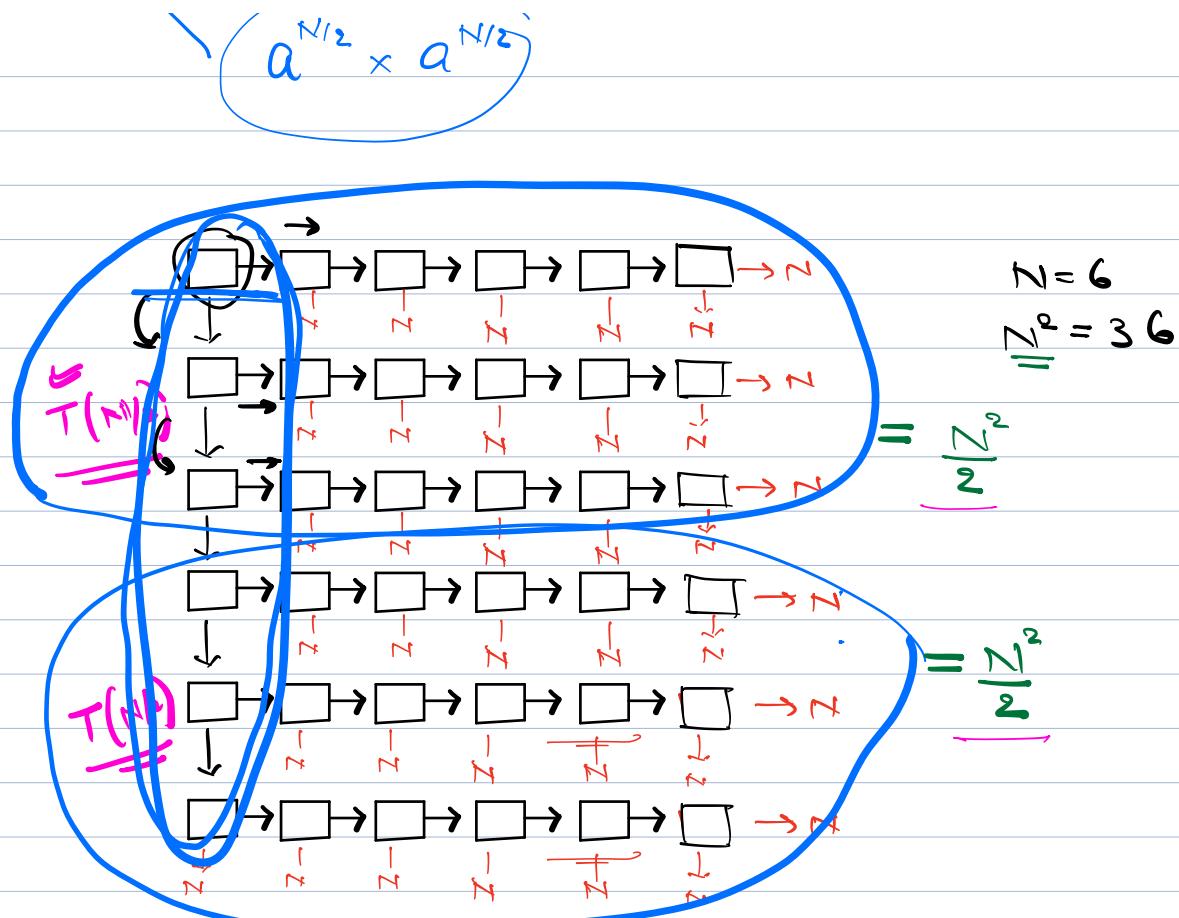
$$\# \text{ iterations} = 2N + 3N + 4N \dots \dots \quad (N \times N)$$

$$N \left(\underline{2+3+4+5} \dots \dots \underline{N} \right)$$

$$N \times N^2$$

$$= O(N^3)$$

$$a^N \xrightarrow{\text{ } \atop \circlearrowleft} @ \circlearrowleft @^{N^2} \Rightarrow O(N)$$



$$T(N) = 2 \underbrace{T(N)}_{\text{H.W.}} + \underbrace{O(N^2)}_{\text{H.W.}}$$

Node merge 2D List (head) {

```
if (head == NULL || head.down == NULL) {
    return head;
}
```

// Ass: merge 2D list \Rightarrow merges all lists whose heads are connected + a single list.

θ
Node mid = get1stMid(head) // Replace
next with down.

$h_2 = \text{mid}.down$

$\text{mid}.down = \text{NULL};$

$\text{head} = \text{merge2Dlist}(\text{head});$

$\underline{h_2} = \text{merge2Dlist}(h_2);$

$\text{head} = \underline{\text{merge}}(\text{head}, h_2);$

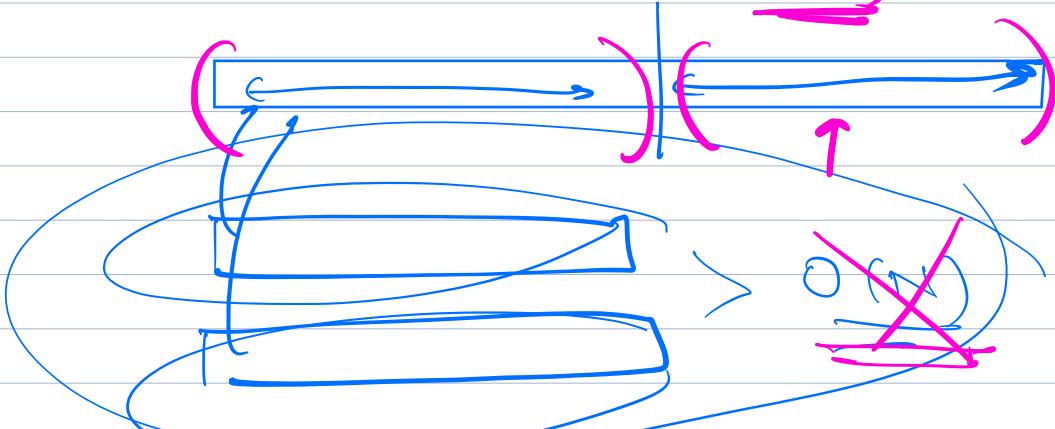
return head;

↓

Doubt

log n

O(1)



~~O(n)~~ + O(log n)

while (_____.next) = 1 - O(1)) $\Rightarrow N + \frac{N}{2}$

—

—

$$\cup = \underline{\left(\frac{3N}{n} \right)}$$

Iterations \Rightarrow

$$\underline{\left(\frac{N}{n} \right)}$$