## Algorithm 1: Online Association Algorithm

```
Input: User i with time demand \{T_{i1}, T_{i2}, ..., T_{iN}\} and priority w_i
Output: user association \{x_{i1}, x_{i2}, ... x_{iN}\}

Initialisation: x_{ij} = 0, \quad \forall 1 \leq j \leq N;

-if WaitingWork().empty = False and WaitingWork().QueueHead.priority \geq w_i then
      WaitingWork.add(T_{i1}, T_{i2}, ..., T_{iN}; w_i)
       j_{min} = argmin_j([\sum_{k=1}^M x_{kj}.T_{kj}] + T_{ij}) /* choose the less loaded AP if [(\sum_{k=1}^M x_{kj_{min}}.T_{kj_{min}}) + T_{ij_{min}}] \le 1 then
               /* Time demand exceeds the remaining available time
                ChooseAP=1, chooseDetail=priorityAllowDesassociate(AP1,T_{i1}, w_i);
                for j=1 to N do
                       k=priorityAllowDesassociate(AP_i, T_{ij}, w_i);
                       if k.response==True then

/* The current candidate AP has less priority clients to be removed
                               chooseAP=j, chooseDetail=k;
                               else
                                       \mathbf{if}\ \mathit{k.directDisassociate} \! = \! \mathit{True}\ \mathbf{then}
                                               \mathbf{if}\ chooseDetail.directDisassociate {==} False\ \mathbf{then}
                                                      /* The First available AP allowing direct disassociation
                                                      chooseAP=j, chooseDetail=k;
                                               else
                                                      /* choose the better AP through the less priority disassociated clients if last\text{-}elem.prio(k) > last\text{-}elem.prio(chooseDetail) then
                                                        _ chooseAP=j, chooseDetail=k;
                                       else
                                               /* choose the better AP through the less priority disassociated clients if (chooseDetail.directDisassociate == False) and
                                                  (last-elem.prio(k) > last-elem.prio(chooseDetail)) then
                                                 chooseAP=j, chooseDetail=k;
                       else
                               R_1 = k.period * c[i][j];
                               R_2 = chooseDetail.period * c[i][chooseAP];
                               if (chooseDetail.response = False \ and \ R_1 > R_2) then

/* The total demand of the current client is not available but currently, this is the
                                       chooseAP=j, chooseDetail=k;
               \mathbf{if}\ \mathit{chooseDetail.response} \! = \! \mathit{False}\ \mathbf{then}
                       */
* The current client could not be associated immediately, because higher priority clients could not be disassociated, and should be put on the waiting queue. The client could choose to start working with the available current period */
WaitingWork(T_{i1}, T_{i2}, ... T_{iN}, w_i) /* The waiting clients are put in a priority based queue, for the same priority clients a FIFO is used */
                else
                       \mathbf{if}\ chooseDetail.DirectDisassociate == False\ \mathbf{then}
                               */
* less priority clients exist, but should be disassociated when they finish the current critical task, the client is put in the queue and associated as possible, or could choose to start working with the chosen AP with the current available period

*/
                               WaitingWork(T_{i1}, T_{i2}, ... T_{iN}, w_i)
                               Disassociate(chooseAP,ChooseDetail.index) /* Disassociate clients from the current index until the tail of the Queue of clients j=chooseAP, x_{ij}=1;
```