

# Quantum Computing

## Lab 2. Algorithms

Frank C Langbein

```
[1]: from qutip import *  
import numpy as np
```

## Oracles

Create a quantum circuit to implement the following oracles as sign and bit oracle:

1.  $f_1$  given by

$x_1$	$x_2$	$f_1(x_1, x_2)$
0	0	0
0	1	0
1	0	1
1	1	1

- 2.

$$f_2(x_1, x_2, x_3, x_4) = (x_1 \text{ XOR } x_2) \text{ AND } (x_3 \text{ AND } x_4)$$

- 3.

$$f_3(x_1, x_2, x_3) = \text{NOT}(x_1 \text{ AND } x_2 \text{ AND } x_3) \text{ OR } (x_1 \text{ AND } x_2 \text{ AND } x_3)$$

## Solution

1. Bit Oracles
  - $f_1$  Bit Oracle
  - $f_2$  Bit Oracle
  - $f_3$  Bit Oracle
2. Sign Oracles
  - $f_1$  Sign Oracle
  - $f_2$  Sign Oracle
  - $f_3$  Bit Oracle

## Grover's Search

Create a quantum circuit to implement Grover's search to find  $|1011\rangle$ .

## Solution

Grover circuit for  $|1011\rangle$

## Reconstructing an Oracle

Assume you are given a sign oracle  $O_f$  for a function of the type

$$f(x_1, \dots, x_n) = x_1 a_1 \oplus x_2 a_2 \oplus \dots \oplus x_n a_n$$

where  $x_l$  and  $a_l$  are  $n$ -bit strings and  $\oplus$  is the addition modulo 2. I.e.  $f$  calculates the modulo-two sum of the  $x_l$  for which  $a_l = 1$ .

Create a quantum algorithm that can find the bit-string  $a_l$  from evaluating  $O_f$  for four bits. The solution circuits below contain an (unknown) oracle which you should try to reconstruct.

Show that a quantum algorithm requires less evaluations of  $f$  than a classical algorithm to solve this problem and explain how your algorithm works.

## Solution

### Circuit for Oracle 1

Note, oracle 1 ( $O1$ ) is across all four qubits, but does not act on the first qubit, so in quirk only spans 3 qubits.

### Circuit for Oracle 2

[ ]: