# ASSIGNMENT 1- SEQUENTIAL SEARCHING AND PARALLEL SEARCHING USING OPENMP REPORT

CSC4005 High Performance Computing

### Abstract

I doing some coding and some tests to solve the problem.

### 1.introduction

OpenMP API have been developed to support parallel application. They provide the parallel application programmer with a model for parallel programming that is portable across architecture from different ventors , and it has been widely accepted as a set of Compiler Directives for multiprocessor programming of shared memory parallel systems . OpenMP supported programming languages include C, C++, and Fortran; and OpenMp-enabled compilers include Sun Compiler, GNU Compiler, and Intel Compiler. OpenMp provides a high-level abstract description of parallel algorithms. The programmer specifies the intent by adding a dedicated pragma to the source code, so that the compiler can automatically parallelize the program and add synchronization mutex where necessary. And communication. When you choose to ignore these pragmas, or if the compiler does not support OpenMp, the program can degenerate into the usual program (usually serial), the code will still work, but you can't use multithreading to speed up program execution.

In this assignment, I focus on the basic use of openMP, and discuss the performance about using openMP to a straightforward pattern matching algorithm with diffident threads. In detail, understand how straightforward pattern matching algorithm works. Generate the test cases, and coding the searching sequential.c into an OpenMP program, also need find the bug and do some optimizations on code to improve the performance.

# 2.1 Part A - Sequential Searching

### About straightforward pattern matching algorithm:

The name of the straightforward pattern matching algorithm is very cool, but the efficiency is not very good, which is a simple, rude way to search.

About the Worst case of straightforward pattern matching algorithm: the time complexity of straightforward pattern matching is O(mn), which also have not a preprocessing stage. So that In the worst case, the straightforward pattern matching algorithm looks for a length of M in a text of length N and requires ~NM character comparisons, until it found the pattern on text at last M.

Therefore, I figured out a way to solve this, for example if the pattern is [aaab], then the text is [aaaaaaaaaaaab], the algorithm will do full pattern length 4 times searches from each letter, and totally do (14-4)\*4 times searches until the pattern found at the end.

# About the code:

### *Time function:*

In some testcases, elapsed wall clock time is "0.000000" because the value returned from the time() function which in the original code couldn't have a high accurate to the 0.000000. After tried a few different timing function and finally a high precision time function - gettimeofday() has been choose to be used on this code because the time it returns accurate to microseconds which meet the requirement perfectly. What's more it's also in the time.h, so that we can introduce it easily without add new head files.

[Length of Pattern and text choose] part:'

Totally 20 patterns and texts which will result in the worst case performance and note the CPU execution time for different length(text)\*length(pattern).

For the selection of pattern and text length. In the first four testcases I design the pattern length is 2 ,4, 5 ,10(the pattern is "ab","aaaab","aaaab","aaaaaaaaab") and the text length is simply the 100/[pattern Length], as there is a rule that the product length(text)\*length(pattern) should be approximately equal to 10^2, 10^4, 10^6, 10^8 and 10^10. So, each 4 test cases is a stage. Pattern Length will multiply by 100 in next stage, length is always following the formula: [text length] = product / [pattern Length]. There is another benefit of the design is that the totally file size is very small, less than 100 MB whatever.

And the Snippet of code as below:

```
long textLength;
long patternLength;// 2 4 5 10
if(testNumber>=0 && testNumber<4){</pre>
       if(testNumber%4==0)/
               patternLength = 2;
        if(testNumber%4==1)//
              patternLength = 4;
        if(testNumber%4==2)/
               patternLength = 5;
        if(testNumber%4==3)//
            patternLength = 10;
        textLength = 100/patternLength;
if(testNumber>=4 && testNumber<8){
        if(testNumber%4==0)/
               patternLength = 20;
        if(testNumber%4==1)//
               patternLength = 40;
        if(testNumber%4==2)//6
               patternLength = 50;
        if(testNumber%4==3)//
               patternLength = 100;
        textLength = 10000/patternLength;
if(testNumber>=8 && testNumber<12){</pre>
        if(testNumber%4==0)//8
               patternLength = 200;
        if(testNumber%4==1)/
               patternLength = 400;
        if(testNumber%4==2)//
               patternLength = 50
        if(testNumber%4==3)//13
            patternLength = 1000;
        textLength = 1000000/patternLength;
```

Folder creation part, generate text and pattern, and write them to each document:

About the text and pattern generation, I used 2 for loop to create them respective:

Put 'a' to index (0~arrayLength-2), then put b to index arrayLength-1

And use open() function and write() function to create text file and write the array of text and pattern to the corresponding file.

And the Snippet of code as below:

```
char textName[10
                             test%d/text.txt",testNumber);
sprintf(textName,
int handle_text;
char *text = (char*)malloc(textLength * sizeof(char));
handle text = open(textName, O CREAT | O RDWR, 0
if(handle_text ==-1) {
    printf("errno=%d\n",errno);
    printf("couldn't create text
                                           for test%d\n", testNumber);
         printf('
for(i=0;i<textLength-1;i++)</pre>
       text[i] = ch;
text[textLength-1] = 'b';
write(handle_text,text,textLength);
free(text);
char patternName[1000];
                             uts/test%d/pattern.txt",testNumber);
sprintf(patternName,
int handle_pattern;
char pattern[patternLength];
handle pattern = open(patternName, O CREAT | O RDWR, 0777);
if(handle_pattern ==-1) {
    printf("errno=%d\n",errno);
    printf("couldn't create pattern for test%d\n",testNumber);
pattern[patternLength-1] = 'b';
for(i=patternLength-2;i>=0;i--)
        pattern[i] = 'a';
write(handle pattern, pattern, patternLength);
if(handle_text!=-1 && handle_pattern!=-1)
                                                    test%d successfully\n", testNumber);
         printf("
```

```
Results(output):
Read test number 0
Text length = 50
Pattern length = 2
Pattern found at position 48
# comparisons = 98
Test 0 elapsed wall clock time = 0.000005
Test 0 elapsed CPU time = 0.000000
Read test number 1
Text length = 25
Pattern length = 4
Pattern found at position 21
# comparisons = 88
Test 1 elapsed wall clock time = 0.000005
Test 1 elapsed CPU time = 0.000000
Read test number 2
Text length = 20
Pattern length = 5
Pattern found at position 15
# comparisons = 80
Test 2 elapsed wall clock time = 0.000004
Test 2 elapsed CPU time = 0.000000
Read test number 3
Text length = 10
Pattern length = 10
Pattern found at position 0
# comparisons = 10
```

Test 3 elapsed wall clock time = 0.000004

Test 3 elapsed CPU time = 0.000000

Read test number 4

Text length = 500

Pattern length = 20

Pattern found at position 480

# comparisons = 9620

Test 4 elapsed wall clock time = 0.000033

Test 4 elapsed CPU time = 0.000000

Read test number 5

Text length = 250

Pattern length = 40

Pattern found at position 210

# comparisons = 8440

Test 5 elapsed wall clock time = 0.000029

Test 5 elapsed CPU time = 0.000000

Read test number 6

Text length = 200

Pattern length = 50

Pattern found at position 150

# comparisons = 7550

Test 6 elapsed wall clock time = 0.000026

Test 6 elapsed CPU time = 0.000000

Read test number 7

Text length = 100

Pattern length = 100

Pattern found at position 0

# comparisons = 100

Test 7 elapsed wall clock time = 0.000005

Test 7 elapsed CPU time = 0.000000

Read test number 8

Text length = 5000

Pattern length = 200

Pattern found at position 4800

# comparisons = 960200

Test 8 elapsed wall clock time = 0.002755

Test 8 elapsed CPU time = 0.000000

Read test number 9

Text length = 2500

Pattern length = 400

Pattern found at position 2100

# comparisons = 840400

Test 9 elapsed wall clock time = 0.002289

Test 9 elapsed CPU time = 0.000000

Read test number 10

Text length = 2000

Pattern length = 500

Pattern found at position 1500

# comparisons = 750500

Test 10 elapsed wall clock time = 0.002062

Test 10 elapsed CPU time = 0.000000

Read test number 11

Text length = 1000

Pattern length = 1000

Pattern found at position 0

# comparisons = 1000

Test 11 elapsed wall clock time = 0.000008

Test 11 elapsed CPU time = 0.000000

Read test number 12

Text length = 50000

Pattern length = 2000

Pattern found at position 48000

# comparisons = 96002000

Test 12 elapsed wall clock time = 0.261211

Test 12 elapsed CPU time = 0.260000

Read test number 13

Text length = 25000

Pattern length = 4000

Pattern found at position 21000

# comparisons = 84004000

Test 13 elapsed wall clock time = 0.228414

Test 13 elapsed CPU time = 0.230000

Read test number 14

Text length = 20000

Pattern length = 5000

Pattern found at position 15000

# comparisons = 75005000

Test 14 elapsed wall clock time = 0.210776

Test 14 elapsed CPU time = 0.200000

Read test number 15

Text length = 10000

Pattern length = 10000

Pattern found at position 0

# comparisons = 10000

Test 15 elapsed wall clock time = 0.000032

Test 15 elapsed CPU time = 0.000000

Read test number 16

Text length = 500000

Pattern length = 20000

Pattern found at position 480000

# comparisons = 9600020000

Test 16 elapsed wall clock time = 25.954374

Test 16 elapsed CPU time = 25.900000

Read test number 17

Text length = 250000

Pattern length = 40000

Pattern found at position 210000

# comparisons = 8400040000

Test 17 elapsed wall clock time = 22.770283

Test 17 elapsed CPU time = 22.709999

Read test number 18

Text length = 200000

Pattern length = 50000

Pattern found at position 150000

# comparisons = 7500050000

Test 18 elapsed wall clock time = 20.294216

Test 18 elapsed CPU time = 20.190001

Read test number 19

Text length = 100000			
Pattern length = 100000			
Pattern found at position 0			
# comparisons = 100000			
Test 19 elapsed wall clock time = 0.000276			
Test 19 elapsed CPU time = 0.000000			
Points selection:			
Results(graphs):			

# Part B - Parallel Searching using OpenMP

### the data below will be used as a reference point(sequential) for test0,1,2:

[40120405@login1(kelvin) assignment-1b]\$ \$time ./searching\_sequential

Read test number 0

Text length = 2000000

Pattern length = 2000

Pattern found at position 1998000

# comparisons = 3996002000

Test 0 elapsed wall clock time = 11

Test 0 elapsed CPU time = 10.680000

Read test number 1

Text length = 10000000

Pattern length = 1000

Pattern found at position 1001

# comparisons = 1002000

Test 1 elapsed wall clock time = 0

Test 1 elapsed CPU time = 0.000000

Read test number 2

Text length = 10000000

Pattern length = 1000

Pattern found at position 5001001

# comparisons = 5001002000

Test 2 elapsed wall clock time = 14

Test 2 elapsed CPU time = 13.380000

# Test0:

# 3.result:

When using openMP instead of sequential:

3a. Table of elapsed times and numbers of comparisons for 3 diffident scheduling strategies:

### Static:

threads	elapsed wall clock time	# comparisons
1	15	3996002000
2	11	2108807405
4	11	1084215980
8	11	855766177
16	12	603062187
32	11	483029111
64	11	410961034

# Dynamic:

threads	elapsed wall clock time	# comparisons
1	14	3996002000
2	21	675219377
4	14	252413741
8	15	300126896
16	16	332466081
32	13	252801616
64	14	178058564

# Guided:

threads	elapsed wall clock time	# comparisons
1	15	3996002000
2	12	2130151583
4	11	1034482481
8	12	759857499
16	13	487172607
32	13	330788813
64	12	218670030

### 3b.

### Static:

parallel speedup(1) = 15/15=1

parallel speedup(2) = 15/11=1.364

parallel speedup(4) = 15/11=1.364

parallel speedup(8) = 15/11=1.364

parallel speedup(16) = 15/12=1.25

parallel speedup(32) = 15/11=1.364

parallel speedup(64) = 15/11=1.364

### dynamic:

parallel speedup(1) = 14/14=1

parallel speedup(2) = 14/21=0.667

parallel speedup(4) = 14/14=1

parallel speedup(8)=14/15=0.933

parallel speedup(16)=14/16=0.875

parallel speedup(32)=14/13=1.077

parallel speedup(64)=14/14=1

### guided:

parallel speedup(1) = 15/15=1

parallel speedup(2) = 15/12=1.25

parallel speedup(4) = 15/11=1.364

parallel speedup(8)=15/12=1.25

parallel speedup(16)=15/13=1.154

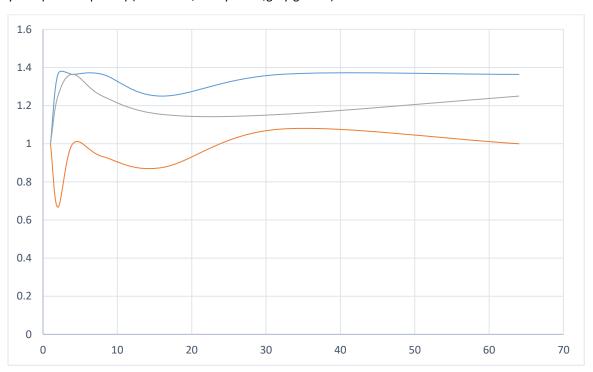
parallel speedup(32)=15/13=1.154

parallel speedup(64)=15/12=1.25

graph of the parallel speedup vs the number of threads:

x: the number of threads

y: the parallel speedup(blue:static,red:dynamic,gray:guided)



# 3c.

# Static:

parallel efficiency (1) = 1/1=1

parallel efficiency (2) = 1.364/2=0.682

parallel efficiency (4) = 1.364/4=0.341

parallel efficiency (8) = 1.364/8=0.175

parallel efficiency (16) = 1.25/16=0.078

parallel efficiency (32) = 1.364/32=0.043

parallel efficiency (64) = 1.364/64=0.021

# dynamic:

```
parallel efficiency (1) = 1/1=1
```

parallel efficiency (2) = 0.667/2 = 0.334

parallel efficiency (4) = 1/4 = 0.25

parallel efficiency (8) = 0.933/8=0.117

parallel efficiency (16) = 0.875/16=0.055

parallel efficiency (32) = 1.077/32=0.034

parallel efficiency (64) = 1/64=0.016

# guided:

parallel efficiency (1) = 1/1=1

parallel efficiency (2) = 1.25/2=0.625

parallel efficiency (4) = 1.364/4 = 0.341

parallel efficiency (8) = 1.25/8=0.156

parallel efficiency (16) = 1.154/16=0.072

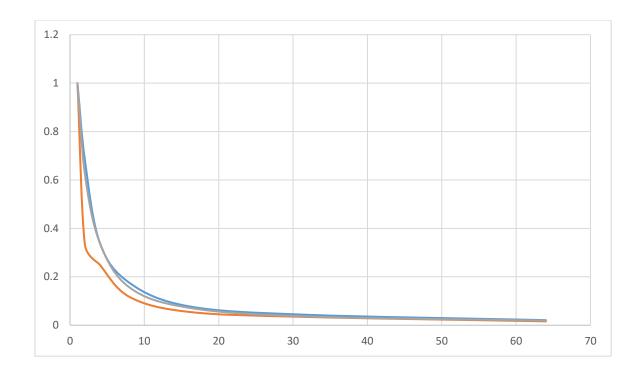
parallel efficiency (32) = 1.154/32=0.036

parallel efficiency (64) = 1.25/64=0.019

graph of the parallel efficiency vs the number of threads:

x: number of threads:

y: parallel efficiency (blue: static,red:dynamic,gray:guided)



Conclusion: According to the 2 graphs of PS vs P and PE vs P, there is no too much diffidence regarding to 3 scheduling strategies, and the performance of static schedule is slightly better than other 2's. However, while the threads increase, the parallel efficiency is decline, the relationship of them is nearly y = 1/x.

### Test1:

# 2.

Result of using searching OMP 0.c on the test case 1 using two threads:

using 2 threads:

Text length = 10000000

Pattern length = 1000

Pattern found at position 1001

# comparisons = 5154093467

using 2 threads: on test0 elapsed wall clock time = 29

using 2 threads: elapsed CPU time = 54.869999

After observe the result, the wall clock time is nearly 30s, which is much longer than running on test case 0. The number of comparisons is also much longer. That's because in test case 1, the pattern is in the front part of text. If it's in the sequential program, the pattern would be found very quickly then the program quit. But on openMP there is no exit function to break the loop, it must run through all the letters in the text even if the pattern has been found, that's where the "bug" is.

4.

Combined with the reason of the performance problem, I simply add a conditional statement:

Which means when pattern has been found, the position of the pattern will assign to the variable [position], after condition checking, the program passed at loop start. then the program will not do any search until all the threads finish. The performance is very good:

```
sing 2 threads:
ext length = 10000000
attern length = 1000
attern found at position 1001
comparisons = 1002000
sing 2 threads: on test0 elapsed wall clock time = 0
sing 2 threads: elapsed CPU time = 0.020000
sing 4 threads:
ext length = 10000000
attern length = 1000
attern found at position 1001
comparisons = 1002000
using 4 threads: on test0 elapsed wall clock time = 0
sing 4 threads: elapsed CPU time = 0.020000
sing 8 threads:
ext length = 10000000
attern length = 1000
attern found at position 1001
comparisons = 1002000
sing 8 threads: on test0 elapsed wall clock time = 0
sing 8 threads: elapsed CPU time = 0.020000
sing 16 threads:
```