



HACETTEPE UNIVERSITY
COMPUTER ENGINEERING DEPARTMENT

BBM204 SOFTWARE PRACTICUM II - 2024 SPRING

Programming Assignment 1

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1 Problem Definition

Our goal when writing our codes is not only to ensure that the code works correctly, but also to make it take as little time as possible. Because when we process large data sets, it can take a very long time. If it takes too long, its usefulness in daily life applications decreases. As qualified engineers of the future, we must take this aspect of our projects into consideration.

2 Solution Implementation

After implementing some popular sort and search algorithms in our project, we compare the average times we obtained after running these algorithms with data sets of different sizes by creating a graph.

2.1 Merge Sort Algorithm

```
1 public class SortAlgorithms {
2     static int[] mergeSort(int[] datas) {
3         int size = datas.length;
4
5         if (size <= 1) return datas;
6         int[] leftArray = Arrays.copyOfRange(datas, 0, size/2);
7         int[] rightArray = Arrays.copyOfRange(datas, size/2, size);
8
9         leftArray = mergeSort(leftArray);
10        rightArray = mergeSort(rightArray);
11
12        return merge(leftArray, rightArray);
13    }
14
15
16    static int[] merge(int[] left, int[] right) {
17        int[] mergeArray = new int[left.length + right.length];
18
19        int leftIndex = 0, rightIndex = 0, mergeIndex = 0;
20
21        while (leftIndex < left.length || rightIndex < right.length) {
22            if (leftIndex == left.length || (rightIndex < right.length && left
23                [leftIndex] > right[rightIndex])) {
24                mergeArray[mergeIndex++] = right[rightIndex++];
25            } else if (rightIndex == right.length || (leftIndex < left.length
26                && left[leftIndex] <= right[rightIndex])) {
27                mergeArray[mergeIndex++] = left[leftIndex++];
28            }
29
30        }
31
32        return mergeArray;
33    }
34 }
```

2.2 Insertion Sort Algorithm

```
31 public class SortAlgorithms {
32     static int[] insertionSort(int[] datas) {
33         for (int j=1; j < datas.length; j++) {
34             int key = datas[j];
35             int i = j-1;
36
37             while (i >= 0 && datas[i] > key) {
38                 datas[i+1] = datas[i];
39                 i--;
40             }
41             datas[i+1] = key;
42         }
43         return datas;
44     }
45 }
```

2.3 Counting Sort Algorithm

```
46 public class SortAlgorithm {
47     static int[] countingSort(int[] datas) {
48         int maxValue = Utilities.findMax(datas);
49
50         int[] count = new int[maxValue + 1];
51         int[] output = new int[datas.length];
52
53         for (int i : datas) count[i]++;
54         for (int i = 1; i < maxValue+1; i++)
55             count[i] = count[i] + count[i - 1];
56
57         for (int i = datas.length-1; i >= 0; i--) {
58             int j = datas[i];
59             count[j]--;
60             output[count[j]] = datas[i];
61         }
62         return output;
63     }
64 }
```

2.4 Linear Search Algorithm

```
65 public class SearchAlgorithms {  
66     static int linearSearch(int[] datas, int value) {  
67         for (int i = 0; i < datas.length; i++) {  
68             if (datas[i] == value) return i;  
69         }  
70         return -1;  
71     }  
72 }
```

2.5 Binary Search Algorithm

```
73 public class SortAlgorithms {  
74     static int binarySearch(int[] datas, int value) {  
75         int low = 0;  
76         int high = datas.length - 1;  
77  
78         while (high - low > 1) {  
79             int mid = (high + low) / 2;  
80  
81             if (datas[mid] < value) low = mid + 1;  
82             else high = mid;  
83         }  
84  
85         if (datas[low] == value) return low;  
86         else if (datas[high] == value) return high;  
87  
88         return -1;  
89     }  
90 }
```

3 Results, Analysis, Discussion

Table 1: Results of the running time tests performed for varying input sizes (in ms).

| Input Size n | | | | | | | | | | |
|--|-------|-------|-------|-------|-------|-------|-------|-------|--------|--------|
| Algorithm | 500 | 1000 | 2000 | 4000 | 8000 | 16000 | 32000 | 64000 | 128000 | 250000 |
| Random Input Data Timing Results in ms | | | | | | | | | | |
| Insertion sort | 0.3 | 0.1 | 0.2 | 0.7 | 2.6 | 11.2 | 41.8 | 162.4 | 672.9 | 2774.4 |
| Merge sort | 0.1 | 0.0 | 0.1 | 0.3 | 0.6 | 1.2 | 2.5 | 5.6 | 11.7 | 25.0 |
| Counting sort | 117.3 | 93.5 | 91.8 | 96.6 | 92.9 | 92.6 | 92.5 | 94.7 | 102.3 | 96.9 |
| Sorted Input Data Timing Results in ms | | | | | | | | | | |
| Insertion sort | 0.0 | 0.0 | 0.1 | 0.4 | 0.1 | 0.3 | 0.0 | 0.0 | 0.2 | 0.5 |
| Merge sort | 0.2 | 0.1 | 0.4 | 0.7 | 0.8 | 2.3 | 4.6 | 6.7 | 10.1 | 26.3 |
| Counting sort | 181.7 | 177.5 | 170.3 | 174.3 | 172.8 | 177.7 | 170.6 | 174.5 | 179.0 | 188.7 |
| Reversely Sorted Input Data Timing Results in ms | | | | | | | | | | |
| Insertion sort | 0.1 | 0.1 | 0.3 | 1.2 | 5.3 | 20.2 | 81.9 | 328.6 | 1338.7 | 4906.8 |
| Merge sort | 0.0 | 0.0 | 0.1 | 0.2 | 0.3 | 0.5 | 1.2 | 2.4 | 4.8 | 11.2 |
| Counting sort | 93.8 | 94.2 | 93.6 | 95.8 | 93.2 | 94.1 | 92.5 | 96.3 | 96.3 | 97.2 |

Table 2: Results of the running time tests of search algorithms of varying sizes (in ns).

| Input Size n | | | | | | | | | | |
|-----------------------------|--------|--------|-------|--------|--------|--------|--------|--------|---------|---------|
| Algorithm | 500 | 1000 | 2000 | 4000 | 8000 | 16000 | 32000 | 64000 | 128000 | 250000 |
| Linear search (random data) | 3434.3 | 2263.4 | 311.9 | 2644.1 | 1148.2 | 2757.8 | 2980.8 | 5779.3 | 8833.5 | 25330.3 |
| Linear search (sorted data) | 790.1 | 963.7 | 531.6 | 521.0 | 910.6 | 2678.9 | 4851.4 | 8498.6 | 16967.3 | 37269.6 |
| Binary search (sorted data) | 359.8 | 205.1 | 281.2 | 108.1 | 84.2 | 87.5 | 107.3 | 86.9 | 89.5 | 101.6 |

| Algorithm | Best Case | Average Case | Worst Case |
|----------------|--------------------|--------------------|---------------|
| Insertion sort | $\Omega(n)$ | $\Theta(n^2)$ | $O(n^2)$ |
| Merge sort | $\Omega(n \log n)$ | $\Theta(n \log n)$ | $O(n \log n)$ |
| Counting Sort | $\Omega(n + k)$ | $\Theta(n + k)$ | $O(n + k)$ |
| Linear Search | $\Omega(1)$ | $\Theta(n)$ | $O(n)$ |
| Binary Search | $\Omega(1)$ | $\Theta(\log n)$ | $O(n \log n)$ |

| Algorithm | Auxiliary Space Complexity |
|----------------|----------------------------|
| Insertion sort | $O(1)$ |
| Merge sort | $O(n)$ |
| Counting sort | $O(k)$ |
| Linear Search | $O(1)$ |
| Binary Search | $O(1)$ |

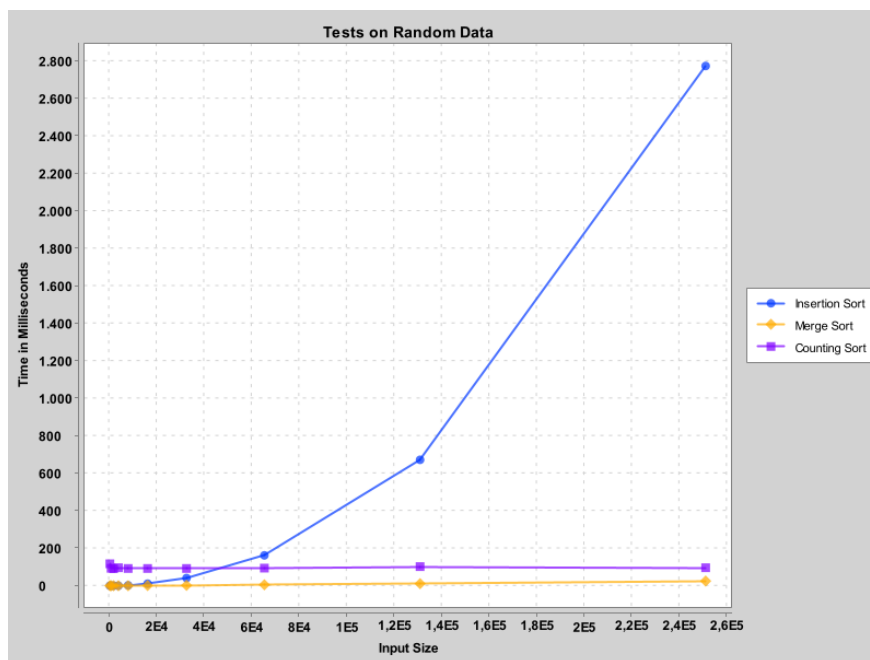


Figure 1: Tests on Random Data

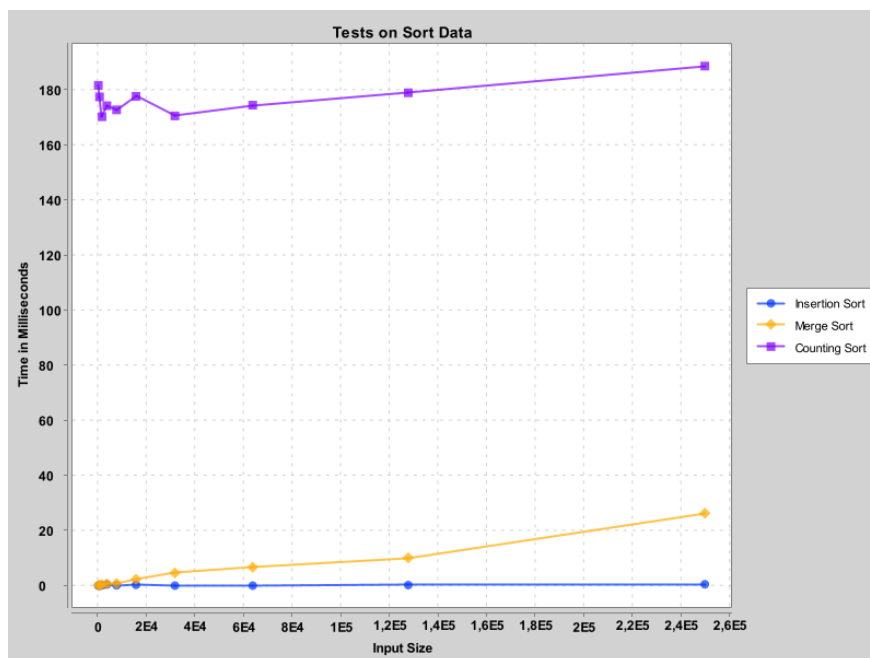


Figure 2: Tests on Sorted Data

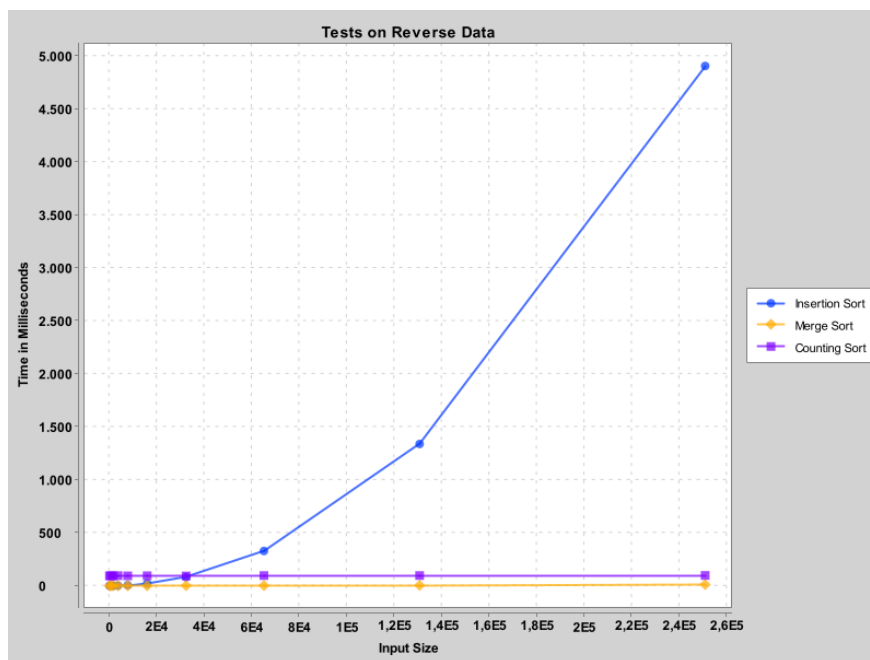


Figure 3: Tests on Reverse Data

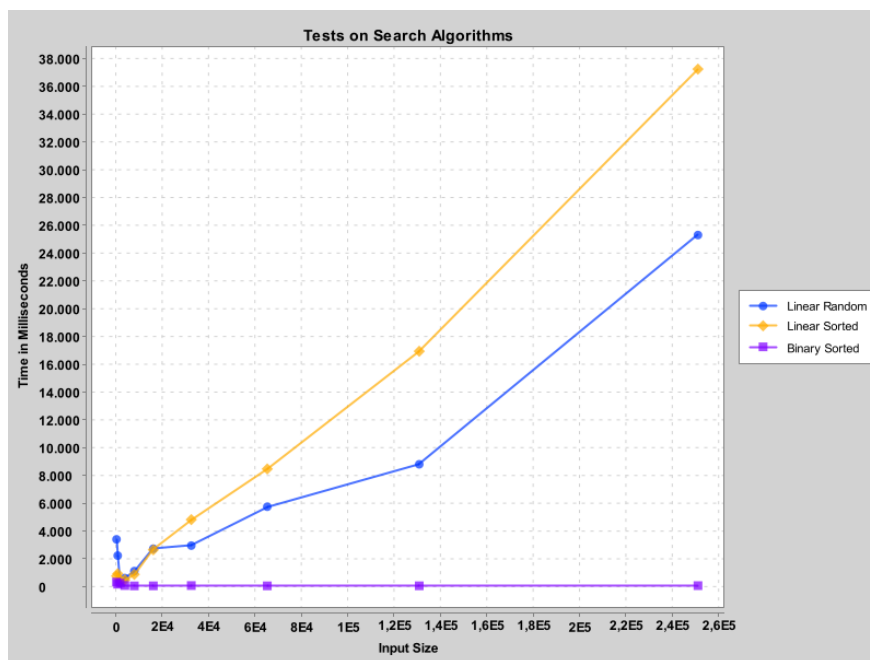


Figure 4: Times of Search Algorithms

4 Notes

- When working with random data or reverse data, working with the insertion sort algorithm takes a lot of time for large data sizes. Therefore, it would be more useful to use merge sort or counting sort algorithms to improve performance when working with this data.
- Algorithms will take much less time when working with already sorted data. However, in our chart, the counting sort algorithm takes longer than others. I guess this may be due to the data range we are working with being too high.
- While analyzing the search algorithms, I tried to choose the value to be searched in a balanced way so that the results would be more consistent. For this reason, I chose 10 values to search and these values are evenly distributed in the data. Additionally, as can be seen from our graph, the binary search algorithm is a much more efficient algorithm compared to linear.