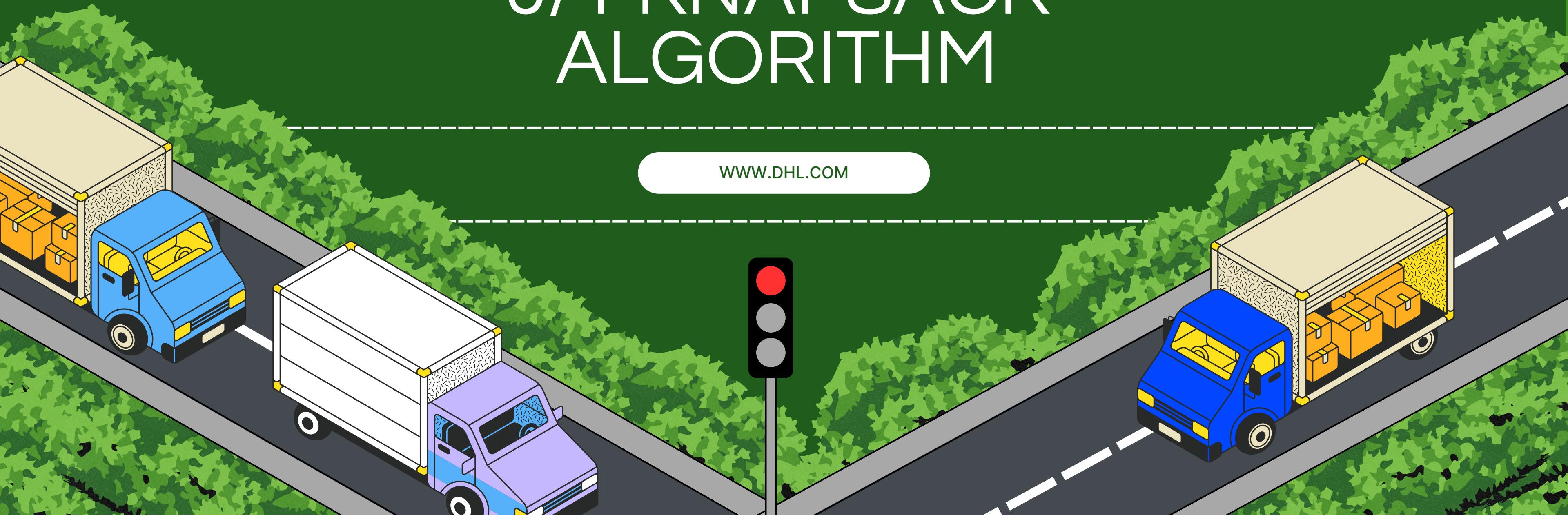


OPTIMIZING CARGO LOADING USING THE 0/1 KNAPSACK ALGORITHM

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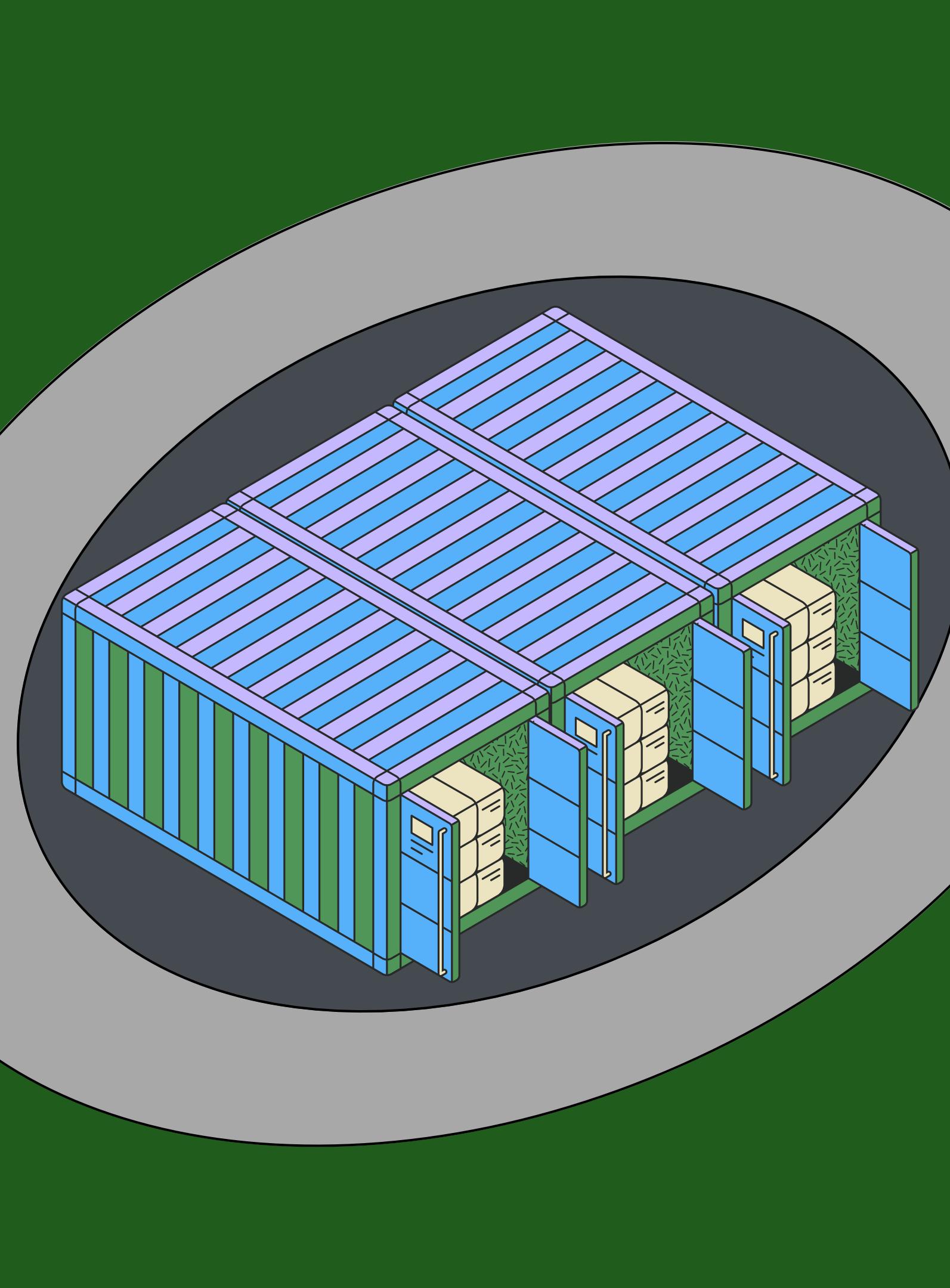
SCENARIO

DHL LOGISTICS

Handles various **logistics services** from express parcel delivery to large-scale freight movement across land, sea, and air

Reputation for **speed** and **reliability**





PROBLEMS

DHL LOGISTICS

Internal reviews have shown that trucks are often **underloaded** or **inefficiently packed**

Shipments containing high-priority or time-sensitive items may be **delayed** if they are grouped with lower-priority goods in an inefficient manner

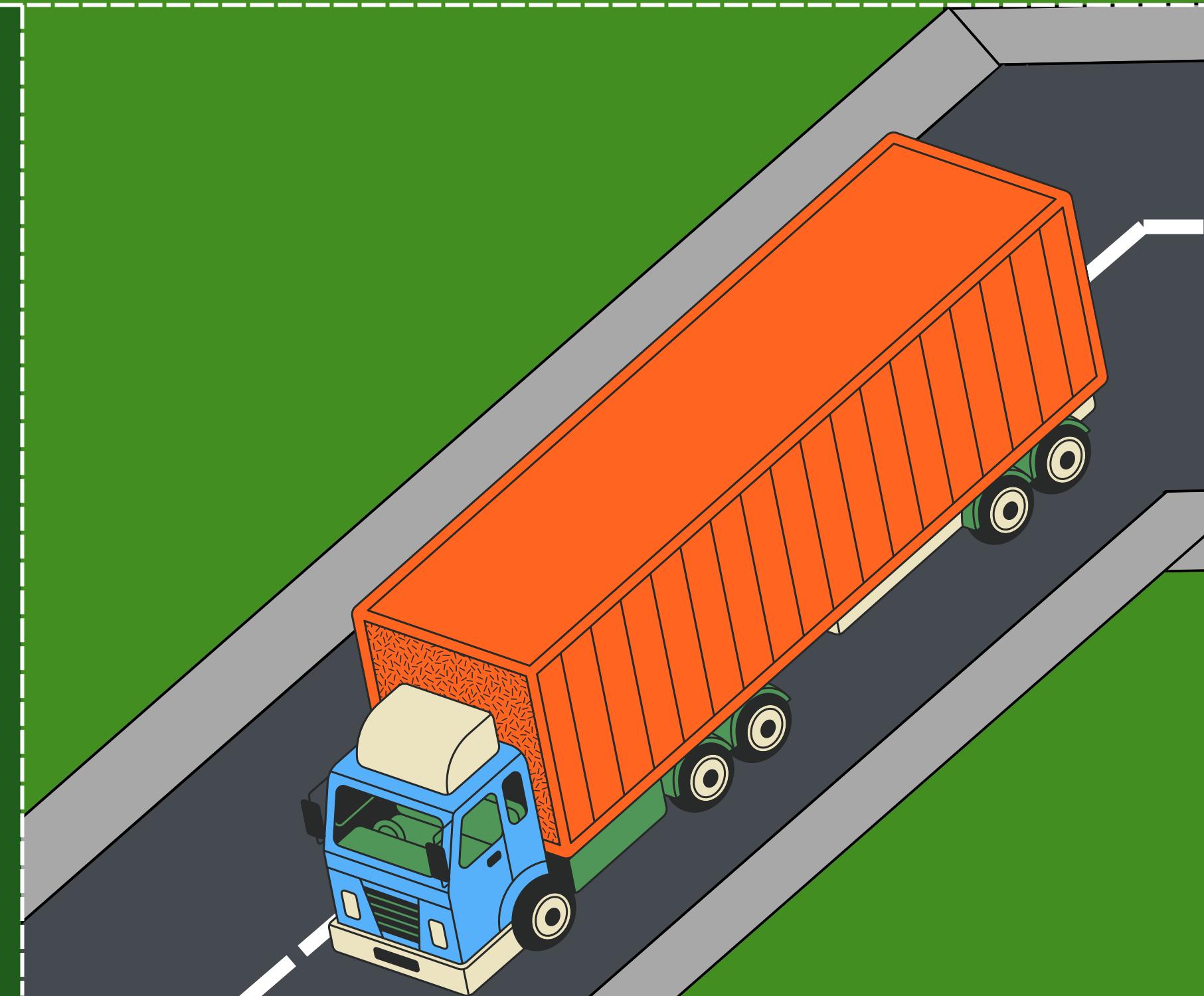
CHALLENGES IN LOGISTICS



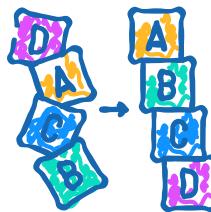
Unhappy
Customer



Increase
Cost



ALGORITHM SUITABILITY REVIEW

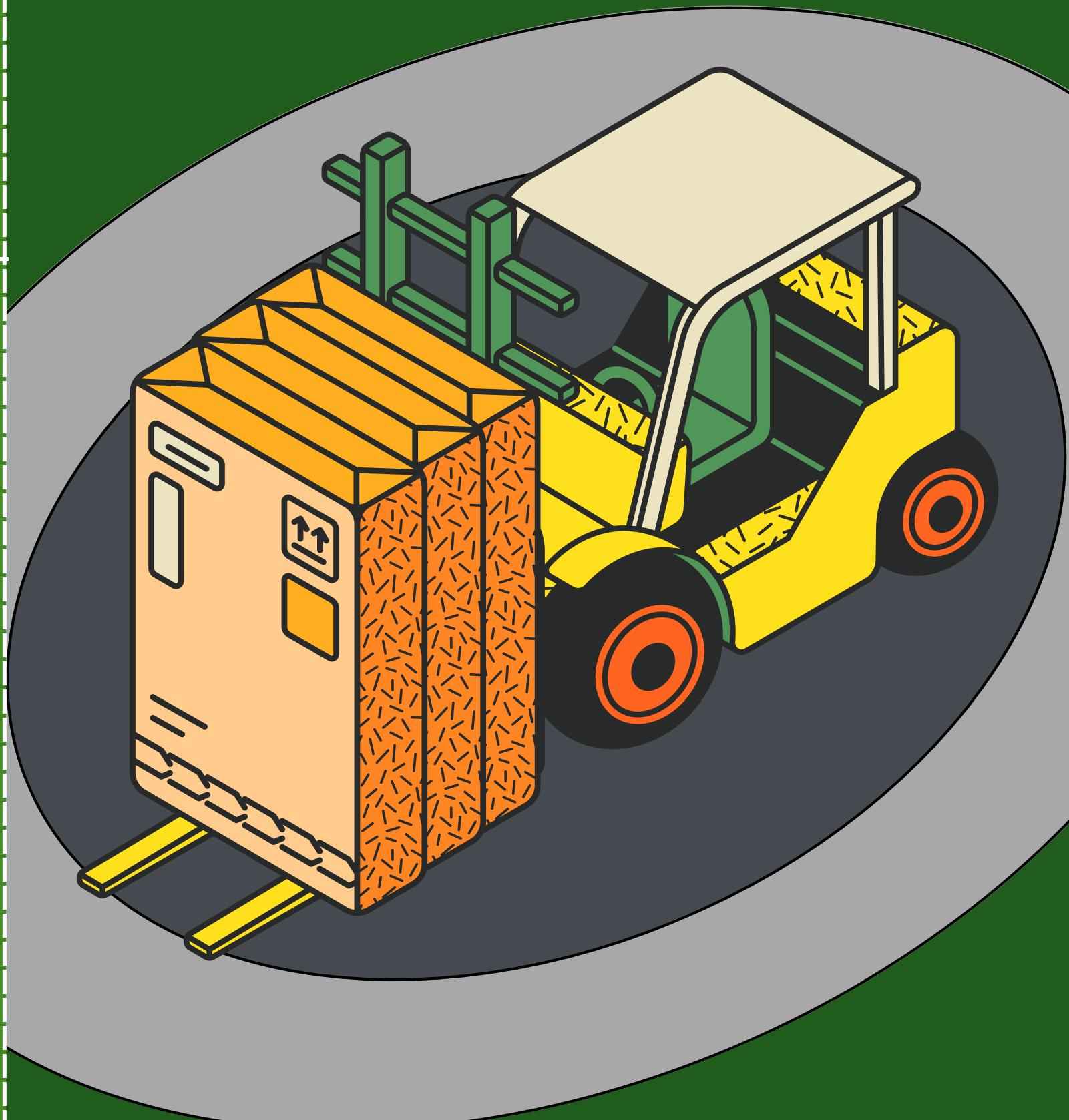


Sorting

- Simple and fast
- Quick ranking by **value-to-weight ratio**



- **Cannot guarantee** optimal selection
- Ignore better combinations of items



ALGORITHM SUITABILITY REVIEW

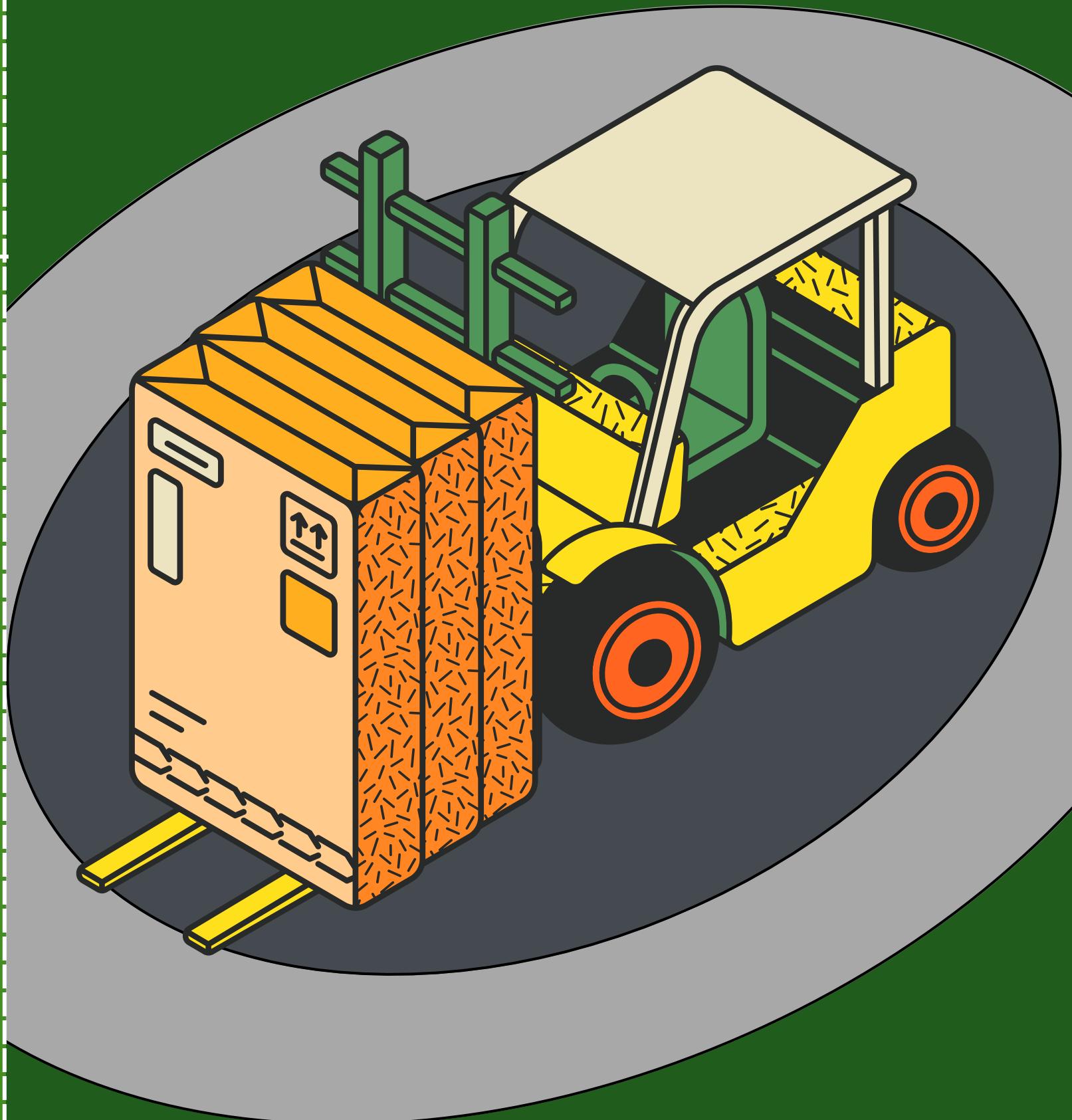


Divide and Conquer

- Allows partial optimization in segmented contexts
- Scales well for distributed cargo management systems



- May not find globally optimal cargo sets unless carefully merged
- Not inherently designed to optimize cumulative cargo value



ALGORITHM SUITABILITY REVIEW

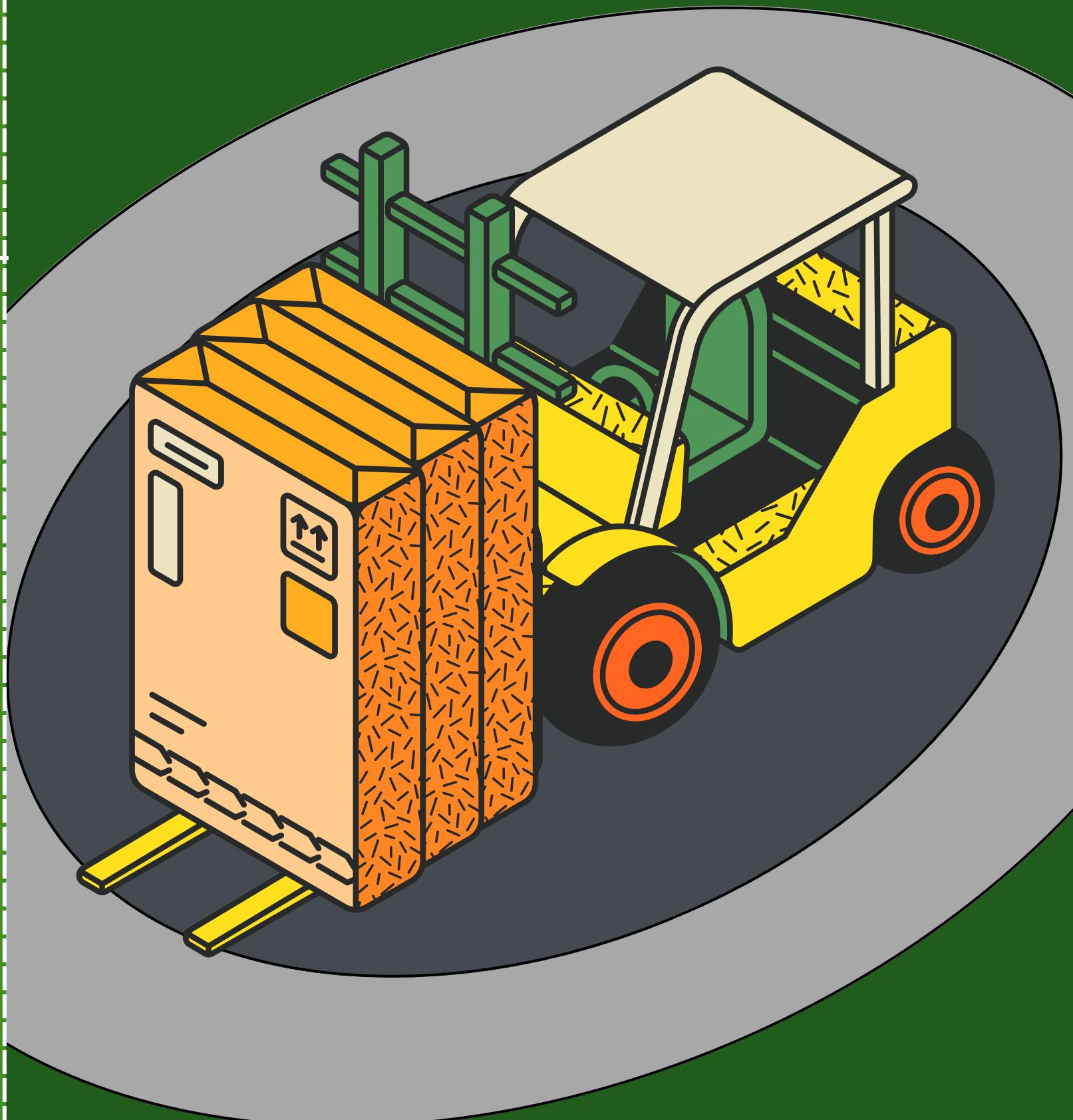


Dynamic Programming (DP)

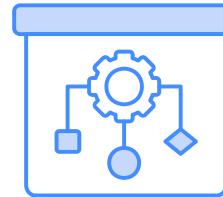
- Provides an optimal solution using the **0/1 Knapsack Problem**
- Evaluates all viable combinations of cargo within the weight limit



- Computationally intensive for large item sets
- Requires careful modeling of item weights and values



ALGORITHM SUITABILITY REVIEW

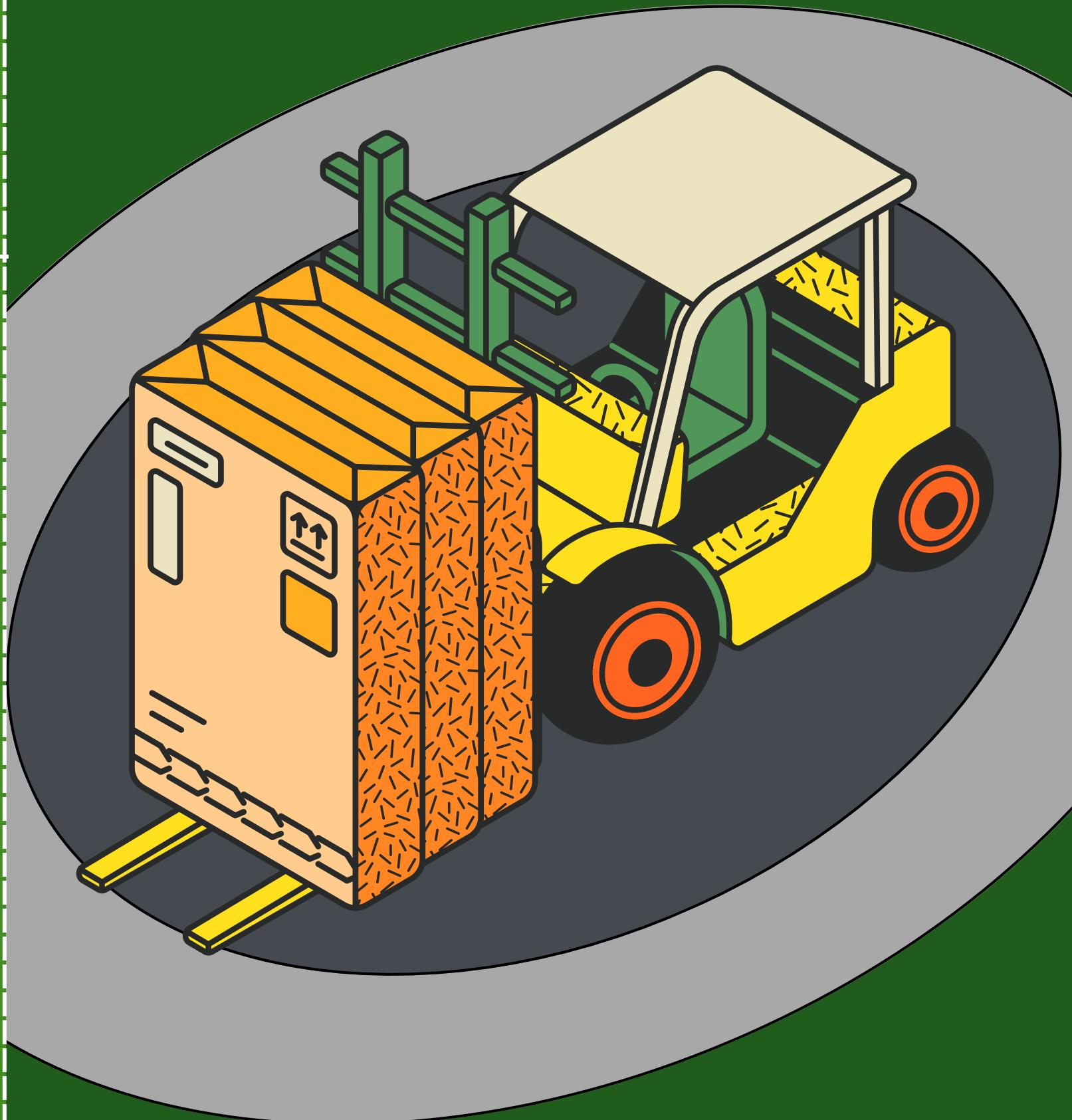


Greedy Algorithms

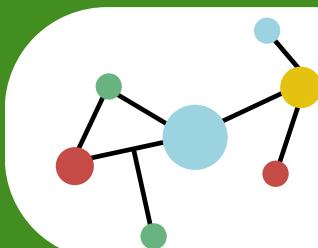
- Very fast and easy to implement
- Useful for **quick approximations** or when time/resources are limited



- Does **not guarantee** optimal cargo value
- Cannot revise earlier decisions, leading to suboptimal total loads



ALGORITHM SUITABILITY REVIEW

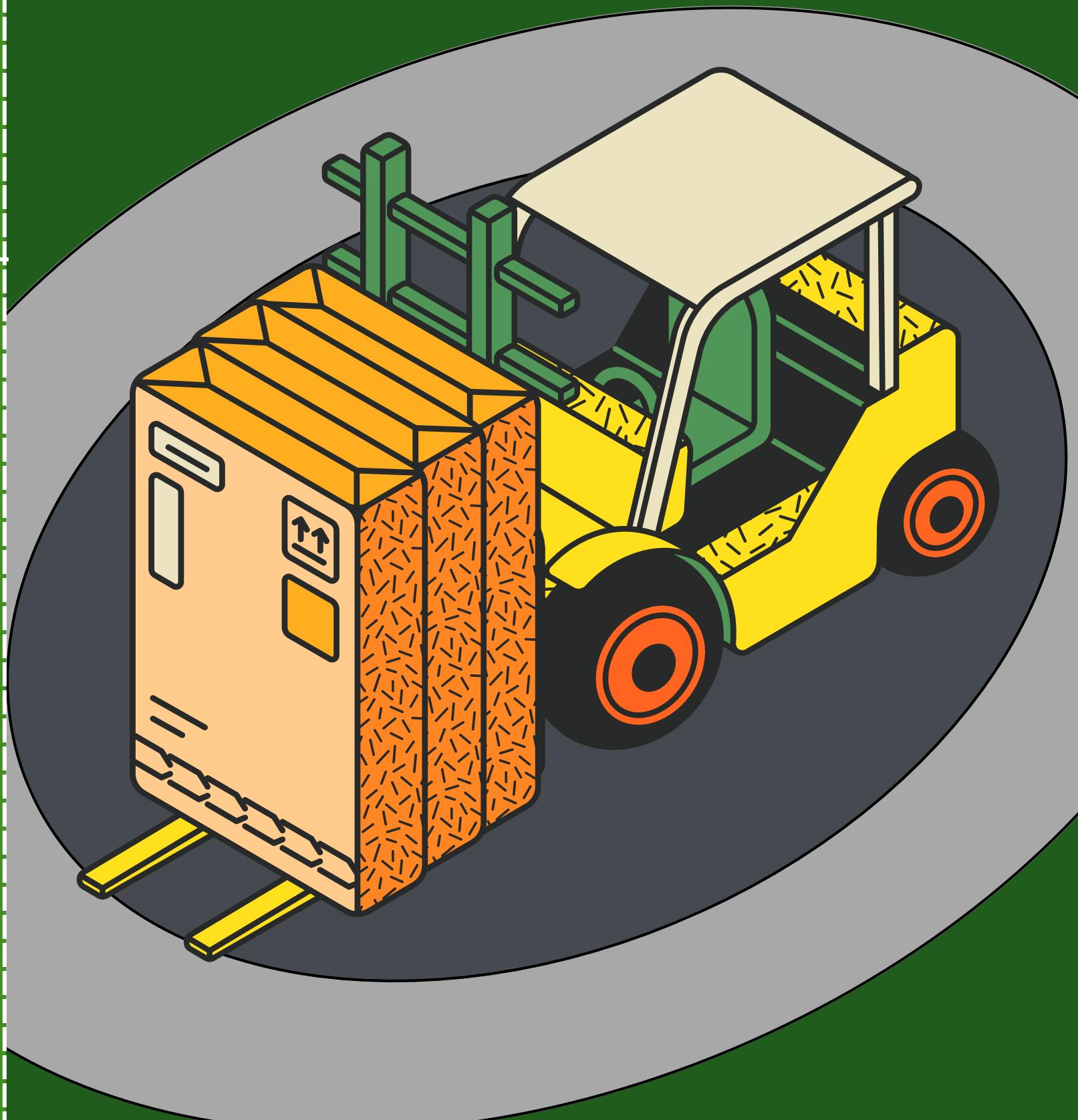


Graph Algorithms

- Can model complex dependencies or constraints
- Useful for visualizing cargo groupings or routing with weight limits



- **Overly complex** for simple weight-based value maximization



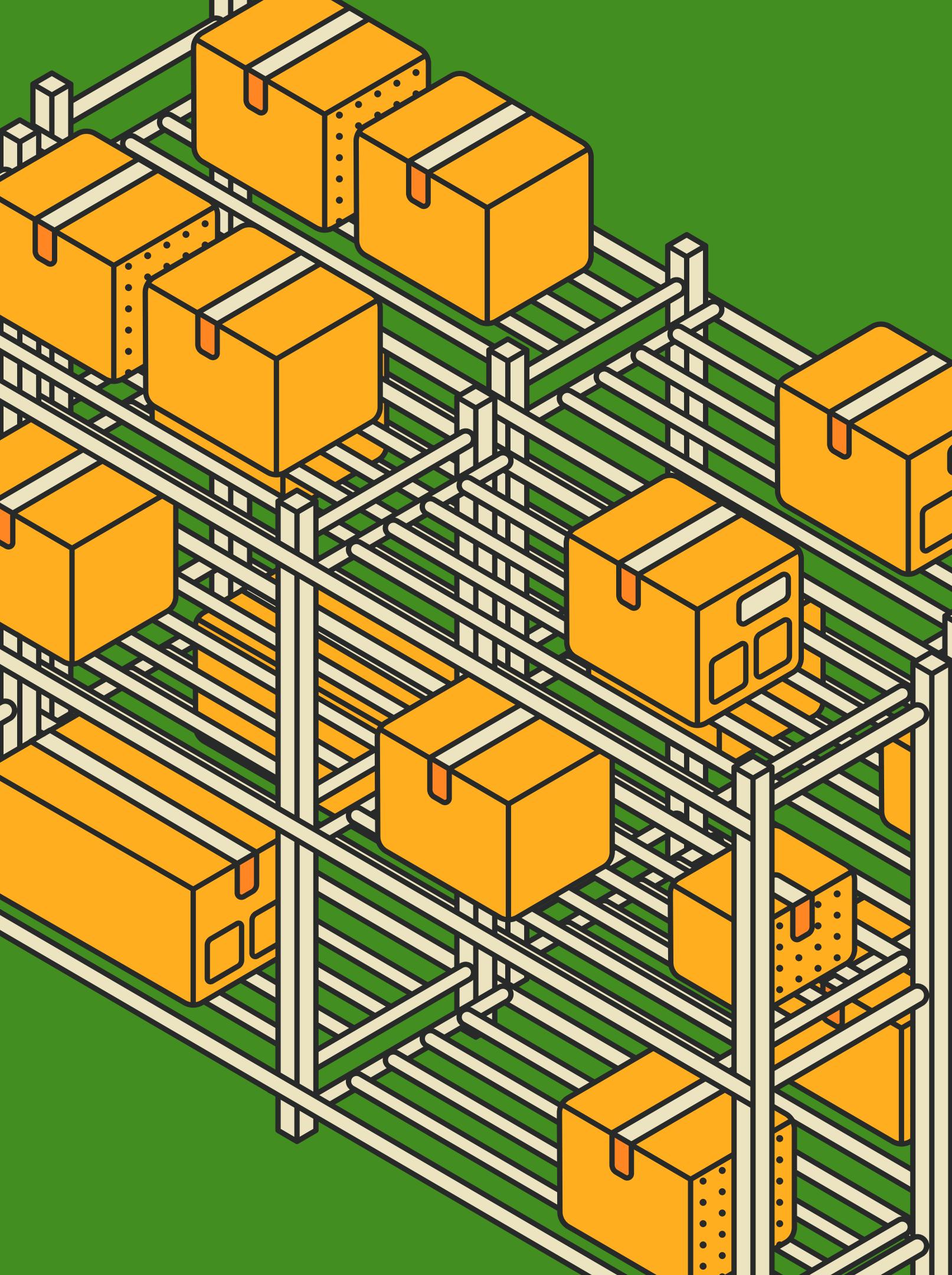
THE OPTIMIZED CHOICE:

DYNAMIC PROGRAMMING (0/1 KNAPSACK)

Maximization of total cargo value
within truck's weight capacity

Best combination of feasible items

Systematic & reproducible results



ALGORITHM DESIGN

OBJECTIVE

Maximize the total value of shipments loaded onto a truck without exceeding the weight limit.

MAIN COMPONENTS

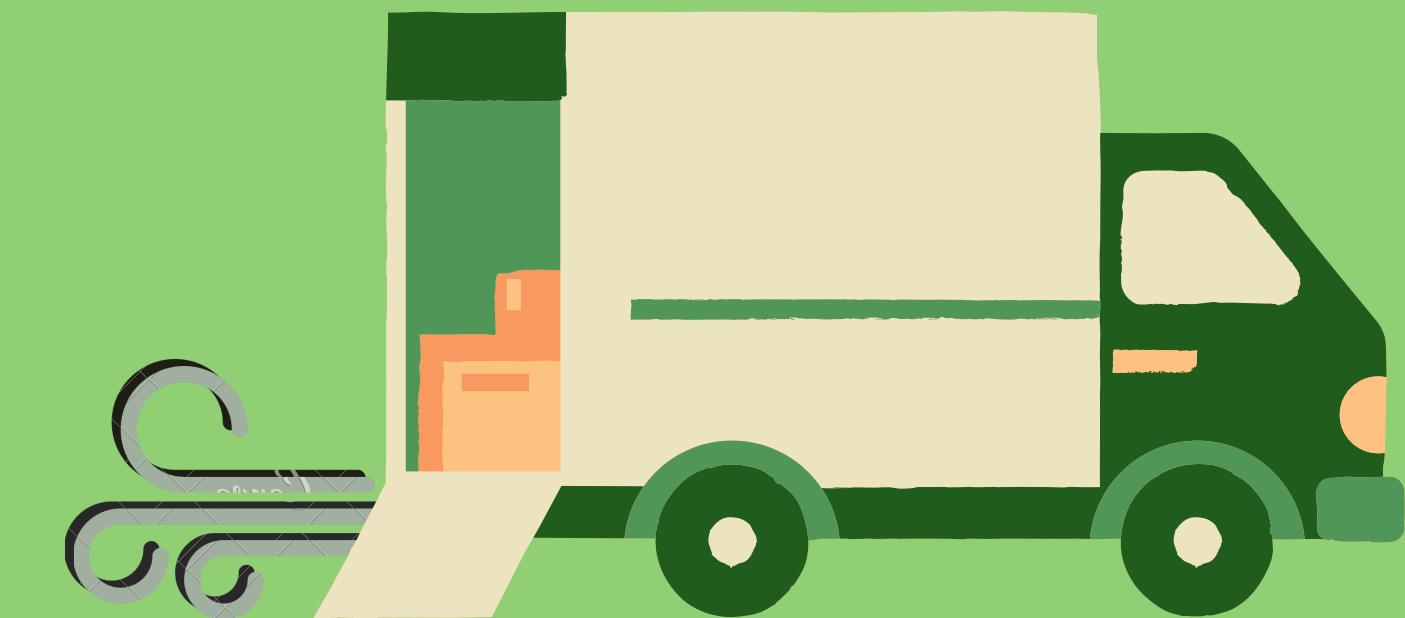
① Shipment Prioritization

- Rank shipments by value score
- Value reflects:
 - Urgency
 - Profitability
 - Delivery deadlines



② Optimal Selection via 0/1 Knapsack Problem

- Each shipment = item with:
 - Weight
 - Value
- **Goal: Maximize total value under max weight limit**



ALGORITHM APPROACH

Dynamic Programming is used to:

- Explore all possible combinations
- Track the best subset of shipments within capacity
- Ensure no shipment is partially selected (0/1 constraint)

CARGO VALUE MODELLING

PRIORITY SCORE FORMULA

$$\text{priority_score} = \alpha \cdot \text{revenue} + \beta \cdot \text{urgency} + \gamma \cdot \text{service_level}$$

REASON PRIORITIZE

- Single shipment: Direct delivery
- Multiple shipments: Weight may exceed capacity → Need to select the best ones

RESULT

Priority score gives us a systematic way to balance profit, urgency, and customer tier."

COMPONENTS

Attribute	Description
Revenue	Income from shipment
Urgency	1/days remaining
Service Level	10 for Premium, 0 for Standard

DEFAULT WEIGHT

Parameter	Value
α	0.5
β	0.3
γ	0.7

EXAMPLE

SCENARIO OVERVIEW

DHL Shipment Optimization

- Goal: Maximize priority score without exceeding truck's 30kg weight capacity
- Shipment Orders: 5 incoming orders with different:
 - Weight (kg)
 - Revenue (RM)
 - Time Remaining (Days)
 - Service Level (Premium/Standard)

STEP 1: PREPARE SHIPMENT DATA & CALCULATE PRIORITY SCORE

Shipment ID	Weight (kg)	Revenue (RM)	Days Left	Service Level	Priority score
A	10	200	1	Premium (10)	107.3 ≈ 107
B	8	180	3	Standard (0)	90.1 ≈ 90
C	5	120	2	Premium (10)	67.15 ≈ 67
D	12	220	4	Standard (0)	110.075 ≈ 110
E	7	150	1	Standard (0)	75.3 ≈ 75

EXAMPLE

STEP 2: APPLY 0/1 KNAPSACK ALGORITHM

		Capacity																												
		1	0	...	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
V	W	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
107	10	0	0	0	0	0	0	0	0	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107		
90	8	0	0	0	0	0	90	90	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	
67	5	0	0	0	67	67	67	90	90	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	
110	12	0	0	0	67	67	67	90	90	107	107	110	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	
75	7	0	0	0	67	67	75	90	90	107	107	142	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	107	

STEP 3: EXTRACTING THE OPTIMAL SHIPMENT SET

- Max Priority Score: 339
- Selected Shipments (by backtracking):
 - $339 - 75(\text{index } 5\text{'s priority score}) = 264$
 - $264 - 67 (\text{index } 3\text{'s priority score}) = 197$
 - $197 - 90 (\text{index } 2\text{'s priority score}) = 107$
 - $107 - 107 (\text{index } 1\text{'s priority score}) = 0$

EXAMPLE

STEP 4: FINAL SHIPMENT SELECTION

Shipment ID	Weight	Priority Score
A	10	107
B	8	90
C	5	67
E	7	75
<i>Total</i>	30	339

ALGORITHM IMPLEMENTATION

DATASET

A = Shipment ID
 B = Revenue
 C = Time Remaining
 D = Service Level
 E = Weight

	A	B	C	D	E
1	1	202	2.85	10	18
2	2	370	1.13	0	65
3	3	206	4.86	0	14
4	4	171	3.93	10	8
5	5	288	4.71	0	16
6	6	120	4.5	10	18
7	7	202	3.09	0	11
8	8	221	4.63	10	56
9	9	314	0.67	0	27
10	10	187	1.18	0	35
11	11	199	0.46	10	89
12	12	251	1.8	0	5
13	13	230	2.1	10	12
14	14	249	1.54	0	5
15	15	357	4.19	10	15
16	16	393	1.94	0	65
17	17	291	1.58	0	76
18	18	376	2.83	0	78
19	19	260	0.92	10	16
20	20	121	4.06	10	10
21	21	352	0.6	0	11
22	22	335	4.94	10	9
23	23	148	3.92	0	35
24	24	158	1.19	10	5
25	25	269	0.28	0	37

	A	B	C	D	E
26	26	287	4.12	10	18
27	27	370	3.61	0	6
28	28	289	3.71	0	49
29	29	274	3.91	0	16
30	30	150	0.6	0	54
31	31	154	1.95	10	80
32	32	343	0.8	10	61
33	33	230	4.35	10	8
34	34	234	3.21	10	41
35	35	120	1.82	0	15
36	36	266	0.55	10	25
37	37	373	1.73	0	95
38	38	188	1.79	0	5
39	39	113	3.72	10	19
40	40	341	3.28	10	10
41	41	364	4.46	10	82
42	42	152	2.49	10	39
43	43	191	0.82	0	38
44	44	363	3.64	0	76
45	45	134	3.86	0	10
46	46	305	2.92	0	10
47	47	180	3.91	10	15
48	48	149	2.6	10	49
49	49	101	2.73	10	5
50	50	153	2.28	10	13

RESULT

Final Shipment Selection Based on Optimization:		
ShipmentID	Weight (kg)	Priority Score
001	18	108
003	14	103
004	8	93
005	16	144
006	18	67
007	11	101
009	27	157
012	5	126
013	12	122
014	5	125
015	15	186
016	65	197
019	16	137
020	10	68
021	11	177
022	9	175
024	5	86
025	37	136
026	18	151
027	6	185
029	16	137
033	8	122
034	41	124
035	15	60
036	25	141
038	5	94
040	10	178
045	10	67
046	10	153
047	15	97
049	5	58
050	13	84
Total	499	3959

32 selected shipment

CORRECTNESS OF ALGORITHM

PROBLEM OVERVIEW

- Objective: Maximize total priority score of selected shipments without exceeding weight limit $W = 500$.
- Approach: Bottom-up Dynamic Programming.
- DP Table Definition: $dp[i][w]$: Maximum priority score using first i shipments with capacity w .

LOOP INVARIANT DEFINITION

- Invariant: At the start of each iteration, $dp[i][w]$ stores the maximum total priority score using the first i shipments and capacity w .

INITIALIZATION

- All $dp[0][w] = 0$: No items \rightarrow score = 0.
- All $dp[i][0] = 0$: Zero capacity \rightarrow no items can be added.

CORRECTNESS OF ALGORITHM

MAINTENANCE (NESTED LOOPS)

```
for (int i = 1; i <= n; i++) {  
    for (int w = 0; w <= W; w++) {  
        if (weight[i-1] <= w)  
            dp[i][w] = max(dp[i-1][w], priority[i-1] + dp[i-1][w - weight[i-1]]);  
        else  
            dp[i][w] = dp[i-1][w];  
    }  
}
```

- Include Case: Item fits → consider added score.
- Exclude Case: Skip item → carry forward previous best.
- Invariant Maintained: $dp[i][w]$ always stores best score for current state.

TERMINATION & CORRECTNESS

- $dp[n][W]$ holds the final optimal score.
- Backtracking using $keep[][]$ reconstructs the selected shipment IDs.
- Conclusion:
 - Initialization, Maintenance, and Termination follow the loop invariant method.
 - The algorithm is correct and optimal.

TIME COMPLEXITY ANALYSIS

PRIORITY SCORE CALCULATION

```
for (int i = 0; i < n; i++) {  
    // Compute priority score  
}
```

The time complexity for this part is $O(n)$.

DYNAMIC PROGRAMMING TABLE CONSTRUCTION

```
for (int i = 1; i <= n; i++) {  
    for (int w = 0; w <= W; w++) {  
        // Fill DP table using recurrence  
    }  
}
```

The time complexity for filling the table is $O(n \times W)$

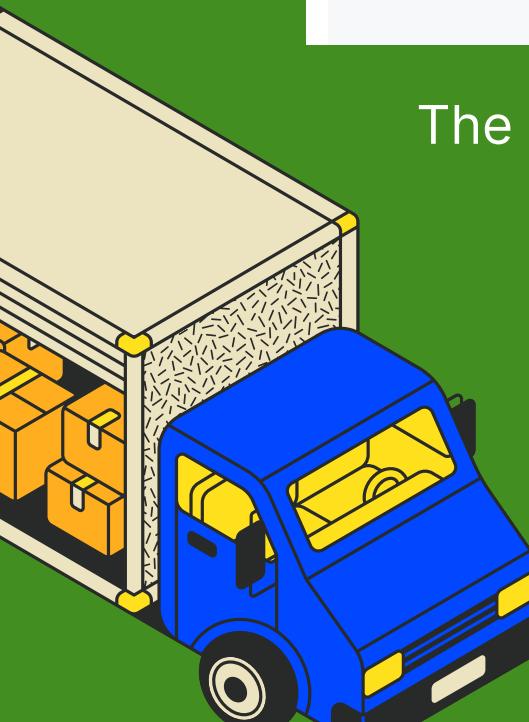


TIME COMPLEXITY ANALYSIS

BACKTRACKING TO FIND SELECTED SHIPMENTS

```
for (int i = n; i > 0; i--) {  
    if (dp[i][w] != dp[i - 1][w]) {  
        // Track selected shipments  
    }  
}
```

The time complexity for this part is $O(n)$.



OUTPUT OF SELECTED SHIPMENTS

```
for (Shipment s : selected) {  
    // Print details  
}
```

The time complexity for filling the table is $O(n \times W)$

BEST CASE

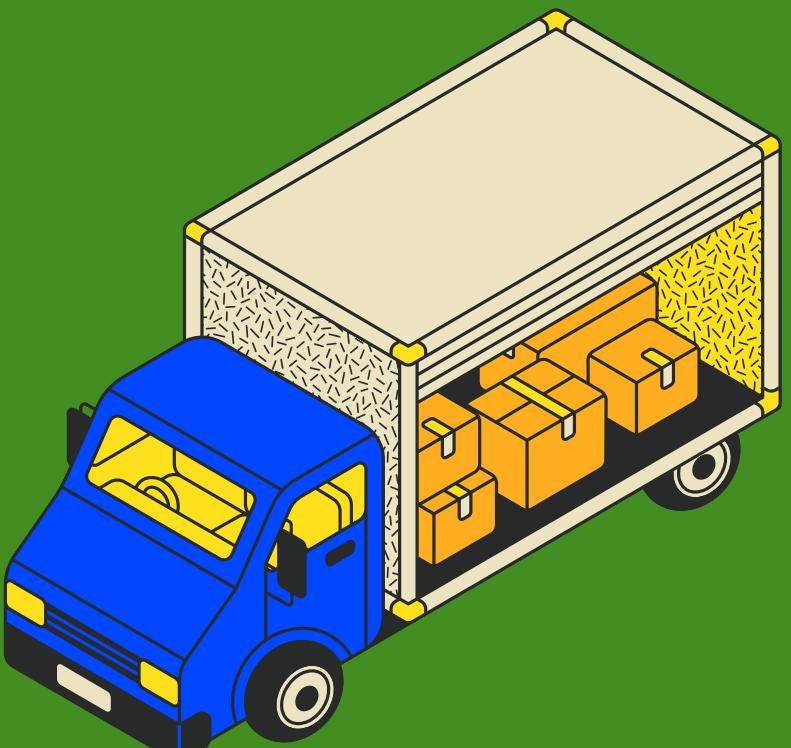
- Time Complexity: $O(nW)$
- The DP table of size $(n+1) \times (W+1)$ is always fully computed.
- Even if all shipments are too heavy to be included, every cell in the table is processed.
- The algorithm does not use early termination or shortcuts.

WORST CASE

- Time Complexity: $O(nW)$
- Every cell in the DP table is processed.
- For each (i, w) , a constant-time operation is performed.
- No input configuration can avoid full-table computation.

AVERAGE CASE

- Time Complexity: $O(nW)$
- Some shipments are selected, others are not.
- Still, all $n \times W$ combinations of items and weights are considered.
- Variations in results do not affect the number of operations.



DHL LOGISTICS

THANKYOU

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