

# Functional Programming With Numbers

The Josephus problem:  $n$  people numbered 1 to  $n$  are made to stand in a circle. Starting from the person numbered 1, every third live person is killed. This is done till only two persons are left. As an example, if  $n$  is 15, then the survivors are the persons who were originally at positions 5 and 14.

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

Write a function (`canSurvive pos n`) that takes as its arguments a position `pos` and the number of people `n`, and returns `True` if the person at position `pos` is one of the last two survivors otherwise it returns `False`.

# Functional Programming With Numbers

- Kill every third person = Repeatedly kill the third person.
- After every killing, renumber the survivors with the person after the dead person counted as 1.

`canSurvive 8 15 ⇒`

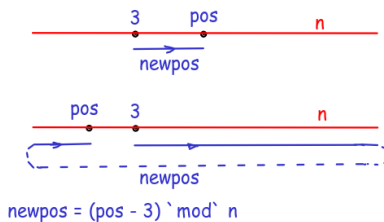
`canSurvive 5 14 ⇒`

`canSurvive 2 13 ⇒`

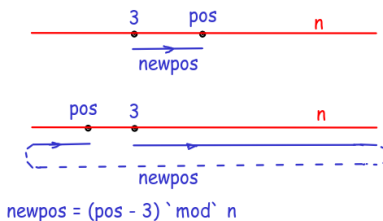
`canSurvive 12 12 ...`

- Stop when `canSurvive` is called with 3 or the number remaining people is less than 2.

# Functional programming with numbers



# Functional programming with numbers



```
canSurvive pos n | n <= 2 = True
                  | pos == 3 = False
                  | otherwise = canSurvive (mod (pos - 3) n) (n - 1)
```

# Functional programming with numbers

The following iterative sequence is defined for the set of positive integers:

$$\begin{aligned} n &\rightarrow n/2, && \text{if } n \text{ is even} \\ n &\rightarrow 3n + 1, && \text{if } n \text{ is odd and } \neq 1 \\ 1, &&& \text{otherwise} \end{aligned}$$

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Write a function `longestchain n` which finds the length of the longest chain starting from any number less than or equal to `n`.



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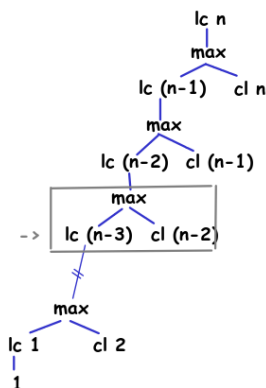
Example: `longestchain 22` is 21.

This corresponds to  $18 \rightarrow 9 \rightarrow 28 \rightarrow 14 \rightarrow 7 \rightarrow 22 \rightarrow 11 \rightarrow 34 \rightarrow 17 \rightarrow 52 \rightarrow 26 \rightarrow 13 \rightarrow 40 \rightarrow 20 \rightarrow 10 \rightarrow 5 \rightarrow 16 \rightarrow 8 \rightarrow 4 \rightarrow 2 \rightarrow 1$ .

# Functional programming with numbers

```
longestChain n | n == 1 = 1
               | otherwise = max (longestChain (n-1)) (chainLength n)
where chainLength n | n == 1 = 1
                   | even n = 1 + chainLength (n `div` 2)
                   | odd  n = 1 + chainLength (3 * n + 1)
```

# Calculating a tail-recursive form



After computing  $lc(n-3)$ , it still remains to compute  $max\ (lc\ (n-3), cl\ (n-2), cl\ (n-1), cl\ n)$  to compute  $(lc\ n)$ .

Suppose we wanted to avoid the computation after the return. A way of doing this is to carry the computation of  $max\ (cl\ (n-2), cl\ (n-1), cl\ n)$  as we go down. At the step shown in the box, we do the computation

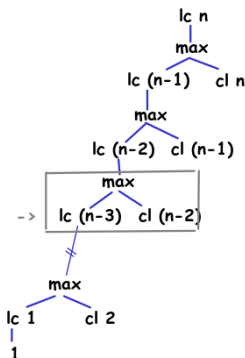
$max\ (lc(n-3), k)$   
where  $k = max\ (cl\ (n-2), cl\ (n-1), cl\ n)$

The boxed computation can be generalized to:

$lc\_tr\ i\ k = max\ (lc\ i, k)$  -- specification of  $lc\_tr$

How does one find a definition of  $lc\_tr$  that is independent of  $lc$ ? How does  $lc\_tr$  get the right  $k$ ?

# Summation of infinite series



$lc\_tr\ i\ k = \max(lc\ i, k)$  -- specification of  $lc\_tr$

We calculate:

$lc\_tr\ 1\ k = \max(lc\ 1, k)$   
 $= \max(1, k)$  --  $lc\ 1 = 1$

$lc\_tr\ i\ k = \max(lc\ i, k)$   
 $= \max(\max(lc\ (i-1), (cl\ i)), k)$  -- definition of  $lc$   
 $= \max(lc(i-1), \max((cl\ i), k))$  -- assoc. of  $\max$   
 $= lc\_tr\ (i-1)\ (\max((cl\ i), k))$  -- specification of  $lc\_tr$

Also,

$lc\ n = \max(lc\ n, 0) = lc\_tr\ n\ 0$  -- 0 is right identity of  $\max$ .

Initially,  $lc\ n$  passes the right  $k$  to  $lc\_tr$

Once  $lc\_tr$  gets the right  $k$ , it passes the right  $k$  to the next recursive call because of the above calculation.

# Tail recursive form of longestChain

```
longestChain n = longestChain_tr n 0
```

```
longestChain_tr n m  
  | n == 1 =  max 1 m  
  | otherwise = longestChain_tr (n-1) (max (chainLength n) m)
```

# Is the tail-recursive form of `longestChain` better?

# Summation of infinite series

- ①  $\tan(x)$  represented as a continued fraction. Only three terms of the fraction are shown:

$$\frac{x}{1 - \frac{x^2}{3 - \frac{x^2}{5 - \dots}}}$$

- ② The nested expression of square roots

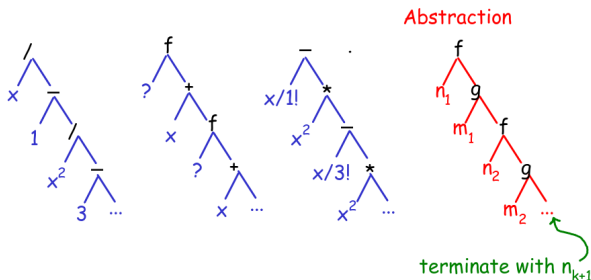
$$\sqrt{x + \sqrt{x + \sqrt{x + \dots}}}$$

- ③ The series expansion of  $\sin$  as a nested expression:

$$\frac{x}{1!} - x^2 * \left( \frac{x}{3!} - x^2 * \left( \frac{x}{5!} - \dots \right) \dots \right)$$

# Higher order functions

Can we see these computations as instances of a common pattern?





# Computing the abstract pattern

```
hof f g n m k i | i <= k = f (n i) (g (m i) (hof f g n m k (i + 1)))  
                | otherwise = n (k+1)
```

```
tan' x k = hof (/) (-) n m k 1  
           where n i = if i == 1 then x else (x ^ 2)  
                 m i = 2 * i - 1
```

```
nested_sqrt x k = hof f (+) n m k 1  
                  where f a b = sqrt b  
                        n i = 0  
                        m i = x
```

```
sin' x k = hof (-) (*) n m k 1  
           where n i = x / factorial (2 * i - 1)  
                 m i = x * x
```

# A better way to write hof

```
hof f g n m k = hof' 1
  where hof' i | i <= k = f (n i) (g (m i)(hof' (i + 1)))
          | otherwise = n (i+1)
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What is the type of hof?

```
hof :: (Ord a, Num a) => (t->t2->t)->
  (t1->t->t2)->
  (a->t1)->
  (a->t)->
  a->
  t
```

# Higer order thinking - Representing a geometric region

Assume that the only use that we shall put a region to is to ask whether a point belongs to it or not.

```
type Point = (Float, Float)
type Region = Point -> Bool
```

`circleMaker` takes a radius and produces a circular region around the origin with the given radius.

```
circleMaker r (x, y) = x ^ 2 + y ^ 2 <= r ^ 2
```

Using lambda notation, we express `circleMaker` as:

```
circleMaker r = \ (x, y) -> x ^ 2 + y ^ 2 <= r ^ 2
```

# Representing a geometric region

Similarly `rectangleMaker` takes a length and a breadth and produces a rectangle around the origin.

```
rectangleMaker :: Float -> Float -> Region
rectangleMaker l b = \ (x,y) -> (abs x) <= l/2 && (abs y) <= b/2
```

Define the regions `notIn`, `intersection`, `union`, `annulus`:

```
notIn :: Region -> Region
notIn r = \p -> not(r p)
```

```
intersection :: Region -> Region -> Region
intersection r1 r2 = \p -> r1 p && r2 p
```

# Representing a geometric region

```
union :: Region -> Region -> Region
union r1 r2 = \p -> r1 p || r2 p

annulus :: Region -> Region -> Region
annulus r1 r2 = intersection r1 (notIn r2)
```

Finally define a function called `translate` which will translate a region to a given distance:

```
type Distance = (Float, Float)

translate :: Region -> Distance -> Region
translate r (x,y) = \(x1, y1) -> r (x1-x, y1-y)
```

# A EMI calculator

Loan amount: 1200000

Interest rate: 10%

Monthly Repayment - 1500

	Balance at beginning of year	Interest for the year	Balance at end of the year
Year 1	1200000	120000	1140000
Year 2	1140000	114000	1074000
Year 3	1074000	107400	1001400
Year 4	1001400	100140	921540
Year 5	921540	92154	833694



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EMI: For what monthly repayment, does the amount become zero at the end of 5 years?

## A EMI (Equated Monthly Installment) calculator.

Loan amount: 1200000

Interest rate: 10%

Monthly Repayment - 26379.58

	Balance at beginning of year	Interest for the year	Balance at end of the year
Year 1	1200000	120000	1003443
Year 2	1003443	100344	787230
Year 3	787230	78723	549396
Year 4	549396	54939	287778
Year 5	287778	28777	0

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A EMI is the monthly payment such that the balance is 0 after the given time period.

The EMI for this example is 26379.58

# A EMI calculator

The function `balance pr r y mp` gives the balance after `y` years for a monthly payment of `mp`, for a principal of `pr` and an interest rate of `r`:

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```
balance pr r 0 mp = pr
balance pr r y mp = balance newpr r (y-1) mp
  where newpr = (pr + pr * r / 100 - 12 * mp)
```

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If `pr`, `r` and `y` are fixed, then `(balance pr r y)` is a function from `mp` to the balance for the given set of values.

```
b = (balance 1200000 10 5)
```

```
b 1500 = 833694
```

```
b 26379 = 0
```

# A EMI calculator

EMI is the value of `mp` for which this function returns 0

```
b mp = 0 -- what value of mp satisfies this equation
```



# A EMI calculator

EMI is the value of `mp` for which this function returns 0

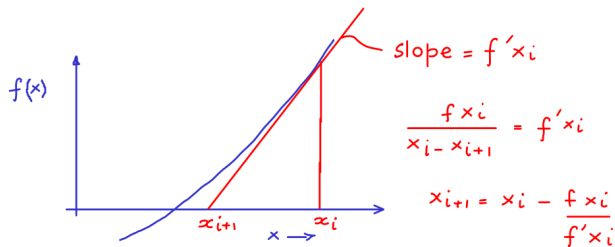
`b mp = 0` -- what value of `mp` satisfies this equation

If we had a function called `zero` that finds for any function the argument value for which the function returns zero:

```
emi_calc pr r y = zero b
  where b = (balance pr r y)
```

# Finding the zero of a function

(zero  $f$ ) is a value  $x$  such that  $f(x) = 0$ .



Approximate using the formula:  $x_{i+1} = x_i - f(x_i)/f'(x_i)$   
 $= \text{improve}(x_i)$

# Finding the zero of a function

First define a function for differentiation:

```
diff f x = (f (x + delta) - f x) / delta
           where delta = 0.0001
```

Next write a function which computes the approximations:

$x_0, f(x_0), f(f(x_0)), f(f(f(x_0))), \dots, f^n(x_0)$

```
until p f x | p x = x
             | otherwise = until p f (f x)
```

Using these, define zero as:

```
zero f = until goodenough improve initial
         where initial = 1.0
               improve xi = xi - (f xi)/diff f xi
               goodenough xi = abs (f xi) < 0.0001
```

# Integration using Simpson's rule

Simpson's rule: The integral of a function  $f$  between  $a$  and  $b$  is:

$$h/3 * \sum_{k=0}^n c_k * y_k$$

$h = (b - a)/n$  and  $n$  is even

$$c_0 = 1$$

$$c_n = 1$$

$c_k = 2$  for  $1 < k < n$  and  $k$  even.

$c_k = 4$  for  $1 < k < n$  and  $k$  odd.

$$y_k = f(a + k * h), 1 < k \leq n$$

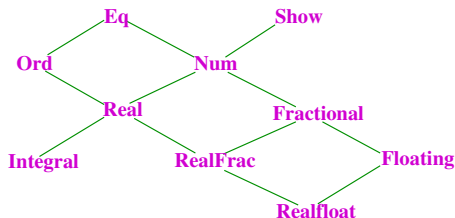
# Integration using Simpson's rule

```
simpson f a b n = (h / 3) * sumseries term n
  where h = (b - a) / n
        c 0 = 1
        c i | i == n = 1
              | even i = 2
              | odd i  = 4
        y k = f (a + k * h)
        term k = (c k) * (y k)
```

```
sumseries term 0 = term 0
sumseries term n = sumseries term (n - 1) + term n
```

# Integration using Simpson's rule

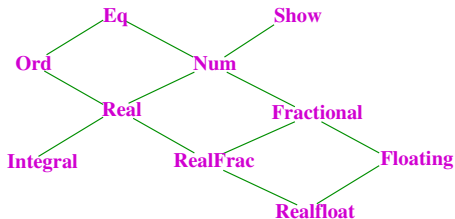
The numeric class hierarchy in Haskell:



- 1 *Num*: Any numeric type. Coercable
- 2 *Fractional*: Support non-integral division. Coercable from `Rational`.
- 3 *Floating*: Supports trigonometric and hyperbolic functions.
- 4 *Real*: Should be expressible as a ratio.
- 5 *RealFloat*: Supports operations specific to floats.

# Integration using Simpson's rule

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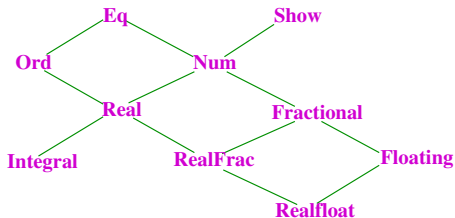


```
class (Eq a, Show a) => Num a where
    (+), (-), (*)      :: a -> a -> a
    negate             :: a -> a
    abs, signum        :: a -> a
    fromInteger        :: Integer -> a

    -- Minimal complete definition: All, except negate or (-)
    x - y              = x + negate y
    negate x           = 0 - x
```

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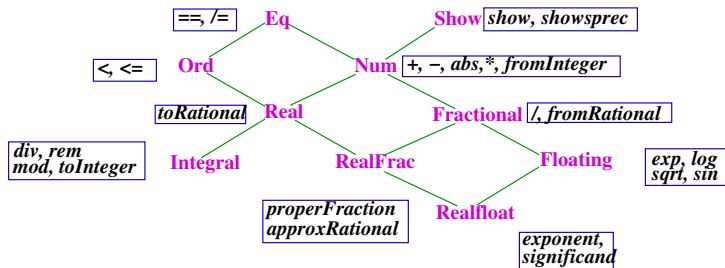


```
instance Num Integer where
  a + b = ...
  a - b = ...
  a * b = ...
  abs a = ...
  signum a = ...
  fromInteger a = a
```



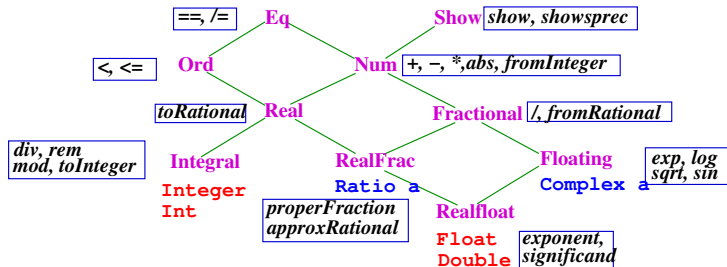
# Integration using Simpson's rule

Some of the operators/functions supported by the classes:



# Integration using Simpson's rule

The types which belong to the classes:



# Integration using Simpson's rule

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  where h = (b - a) / fromInteger n
        c 0 = 1
        c i | i == n = 1
              | i `mod` 2 == 0 = 2
              | i `mod` 2 == 1 = 4
        y k = f (a + k * h)
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```
sumseries term 0 = term 0
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```

# Overloading resolution using the *default* rule

```
*Main> :t 6          -- What is the type of 6?  
6 :: Num a => a
```

```
*Main> let x = 6      -- Serious business, what is the type of the result  
*Main> :t x          -- of this program?  
x :: Integer
```

In the second case `x` is assumed to be the result of a program. Similarly:

```
*Main> :t sin 3.1415  
sin 3.1415 :: Floating a => a  
  
*Main> let val = sin 3.1415  
*Main> :t val  
val :: Double
```

# Overloading resolution using the *default* rule

- If the type of the result of a program is a overloaded value (ambiguous), it is resolved using the *default rule*.
- *defaults* are limited to numeric classes.
- The default *default* is given by the declaration:  
`default{Integer, Double}`.
- If the type of the result is overloaded, it defaults to the first type in the default list that satisfies all the class contexts in the overloaded type.

# Overloading resolution using the *default* rule

Example:

- The overloaded type of the result is:  
`(Show a), (Fractional a) => a -> Int.`
- Assume that the default list is `{Integer, Double}`.
- `Integer` does not satisfy the `(Fractional a)` context.
- `Double` satisfies both `(Show a)` and `(Fractional a)` contexts.
- The type of the result resolves to `Double a -> Int.`
- The default can be altered by a `default` declaration.