

Functional Programming With Lists

Amitabha Sanyal

Department of Computer Science and Engineering
IIT Bombay.

Powai, Mumbai - 400076

`as@cse.iitb.ac.in`

August 2014

Functional Programming with Lists

```
typedef enum {Nil_t,Cons_t}    List* Nil ()  {
                                Tagtype;      List* l = (List*) malloc(sizeof(List));
                                l->tag = Nil_t;
typedef struct L               return(l);}
{
    Tagtype tag;
    struct {
        int head;
        struct L* tail;
    };
} List;

List* Cons (int hd, List* tl)  {
    List* l = (List*) malloc(sizeof(List));
    l->tag = Cons_t;
    l->head = hd;
    l->tail = tl;
    return(l);}

int main ()
{
    int i;    List* l = Nil();
    i = 10;
    while (i > 0) l = Cons (i--, l);}
```

Functional Programming with Lists

Lists in C

```
typedef enum {Nil_t,Cons_t}
            Tagtype;
```

```
typedef struct L
{
    Tagtype tag;
    struct {
        int head;
        struct L* tail;};
} List;
```

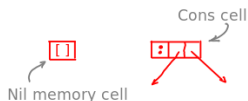
```
List* Cons (int hd, List* tl)
{
    List* l = (List*) malloc(sizeof(l));
    l->tag = Cons_t;
    l->head = hd;
    l->tail = tl;
    return(l);}
```

Lists in Haskell

```
data List a = Nil | Cons a (List a)
```

- Defines a polymorphic type `List a`.
- `List` is a **type constructor**.
- `Nil` and `Cons` are **data constructors**. They also serve the role of tags.
- `|` is disjoint union.

```
data [a] = [] | a : [a]
```

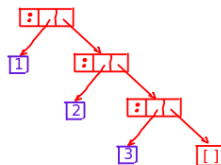


Functional Programming with Lists

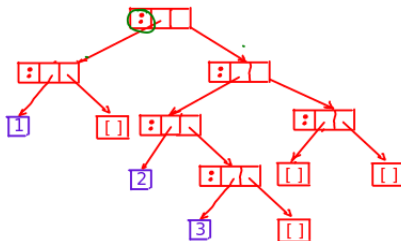
- `[]` is a shorthand for `Nil`.
- `[1,2,3]` is a shorthand for `1:(2:(3:[]))`.
- `1:(2:(3:[]))` is Haskell notation for `Cons 1 (Cons (Cons 3 []))`.

`[1,2,3]`

`['a', 'b', 'c'] :: [Char]`



`[[1], [2,3], []] :: [[Int]]`



Functional Programming with Lists

length

Functional Programming with Lists

```
length [] = 0
length (x:xs) = 1 + length xs

sum
```

Functional Programming with Lists

```
length [] = 0  
length (x:xs) = 1 + length xs
```

```
sum [] = 0  
sum (x:xs) = x + sum xs
```

```
product
```

Functional Programming with Lists

```
length [] = 0
length (x:xs) = 1 + length xs

sum [] = 0
sum (x:xs) = x + sum xs

product [] = 1
product (x:xs) = x * product xs

[] ++ ys
```


Functional Programming with Lists

```
length [] = 0
length (x:xs) = 1 + length xs

sum [] = 0
sum (x:xs) = x + sum xs

product [] = 1
product (x:xs) = x * product xs

[] ++ ys = ys
(x:xs) ++ ys = x:(xs ++ ys)

reverse
```

Functional Programming with Lists

```
length [] = 0
length (x:xs) = 1 + length xs
```

map

```
sum [] = 0
sum (x:xs) = x + sum xs
```

```
product [] = 1
product (x:xs) = x * product xs
```

```
[] ++ ys = ys
(x:xs) ++ ys = x:(xs ++ ys)
```

```
reverse [] = []
reverse (x:xs) = reverse xs ++ x
```

Functional Programming with Lists

```
length [] = 0                                map f [] = []
length (x:xs) = 1 + length xs  map f (x : xs) = f x : map f xs

sum [] = 0                                    filter
sum (x:xs) = x + sum xs

product [] = 1
product (x:xs) = x * product xs

[] ++ ys = ys
(x:xs) ++ ys = x:(xs ++ ys)

reverse [] = []
reverse (x:xs) = reverse xs ++ x
```

Functional Programming with Lists

```
length [] = 0
length (x:xs) = 1 + length xs

map f [] = []
map f (x : xs) = f x : map f xs

sum [] = 0
sum (x:xs) = x + sum xs

filter p [] = []
filter p (x:xs)
  | p x = x:(filter p xs)
  | otherwise = (filter p xs)

product [] = 1
product (x:xs) = x * product xs

[] ++ ys = ys
(x:xs)++ys = x:(xs ++ ys)

reverse [] = []
reverse (x:xs) = reverse xs ++ x
```

The foldr function

foldr - the natural abstraction of list processing functions:

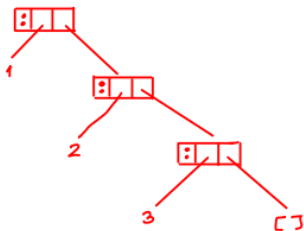
```
foldr f id [] = id
foldr f id (x:xs) = f x (foldr f id xs)
```

```
length l = foldr (\x y -> 1 + y) 0 l
l1 ++ l2 = foldr ...
reverse l = foldr ...
map f l = foldr ...
filter p l = foldr ...
```

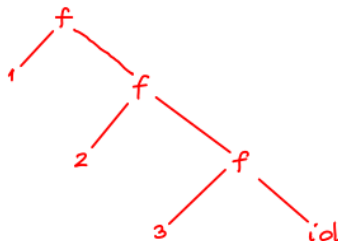
What does foldr do?

```
foldr f id [] = id  
foldr f id (x:xs) = f x (foldr f id xs)
```

`[1,2,3]`



`foldr f id [1,2,3]`



Other List Processing Functions

```
xs !! n | n < 0 = error "negative index"
[] !! n       = error "index too large"
(x:xs) !! 0   = x
(x:xs) !! n   = xs !! (n-1)
```

```
take 0 _ = []
take _ [] = []
take (n + 1) (x:xs) = x : take n xs
```

```
zip [] _ = []
zip _ [] = []
zip (x:xs) (y:ys) = (x,y) : zip xs ys
```

```
takeWhile p [] = []
takeWhile p (x:xs) = if (p x) then x : takeWhile p xs else []
```

```
dropWhile p [] = []
dropWhile p (x:xs) = if (p x) then dropWhile p xs else x:xs
```

List Comprehensions

We wanted a function `dropWhile'` which satisfied

```
dropWhile' p l = (dropWhile p l, l)
```

Now we wanted to express `dropWhile'` as a foldr, i.e.

```
dropWhile' p l = foldr g id l
```

Let us calculate `g` and `id`.

```
id = dropWhile' p [] = (dropWhile p [], [])
```

```
Now dropWhile' p (x:xs)
  = (dropWhile p (x:xs), (x:xs))
  = (if (p x) then dropWhile p xs else x:xs, x:xs)
  = if (p x) then (dropWhile p xs, x:xs) else (x:xs, x:xs)
  = if (p x) then (dropWhile p xs), x:xs) else (x:xs, x:xs)
  = if (p x) then (d, x:l) else (x:xs, x:xs)
    where (d, l) = dropWhile' p xs -- and l = xs
```

Extract the `g`

```
g x (d,l) = if (p x) then (d, x:l) else (x:l, x:l)
```


List Comprehensions

```
[x * x | x <- [1,2,3,4,5]] => [1,4,9,16,25]
[x * x | x <- [1,2,3,4,5], even x] => [4,16]
[x + y | x <- [1,2,3], y <- [6,7]] => [7,8,8,9,9,10]
```

```
qsort [] = []
qsort (x:xs) = qsort lows ++ [x] ++ qsort highs
  where lows = [y | y <- xs, y <= x]
        highs = [y | y <- xs, y > x]
```

```
fib = 0:1:[x + y | (x,y) <- zip fib (tail fib)]
```

```
fib = 0:1:[x + y | x <- fib, y <- (tail fib)]
```

does not work. Why?

List Comprehension

The eight queens problem

```
queens 0 = [[]]
queens n = [board ++ [pos] | board <- queens (n-1),
                        pos <- [1..8],
                        safeconfig board pos]
where safeconfig board pos = all (safepos (n,pos))
                                (zip [1..n-1] board)
    safepos (n1, pos1) (n, pos) = pos /= pos1 &&
                                abs (n-n1) /= abs (pos-pos1)
```

Reasoning about Functional Programs

Universal property of foldr:

$$\begin{aligned} e [] &= \text{id} & \Leftrightarrow & e = \text{foldr } f \text{ id} \\ e (x:xs) &= f \ x \ (e \ xs) \end{aligned}$$

\leq foldr f id is a solution of the equations on the left

\geq foldr f id is the only solution of the equations on the left

Reasoning about Functional Programs

Universal property of foldr:

$$\begin{aligned} e [] &= \text{id} & \Leftrightarrow & e = \text{foldr } f \text{ id} \\ e (x:xs) &= f \ x \ (e \ xs) \end{aligned}$$

\leq foldr f id is a solution of the equations on the left

\geq foldr f id is the only solution of the equations on the left

One can prove many interesting results using this property

- ① $(+ \ 1).\text{sum} = \text{foldr } (+) \ 1$
- ② if $h \ w = v$ and $h \ (g \ x \ y) = f \ x \ (h \ y)$, then
 $h.(\text{foldr } g \ w) = \text{foldr } f \ v.$
- ③ $\text{map } s.\text{map } t = \text{map } (s.t)$
- ④ $\text{map } s.\text{concat} = \text{concat.map } (\text{map } s)$

Reasoning about Lists

inits: Finds all initial segments of a list.

```
inits [] = [[]]  
inits (x:xs) = [] : (map (x:) inits xs)
```

tails: Finds all tail segments of a list

```
tails [] = [[]]  
tails (x:xs) = (x:head l):l  
    where l = tails xs
```

Can you express **heads** and **tails** using **foldr**?

Reasoning about Functional Programs

`scanl/scanr`: Accumulating `foldl/foldr`. n times the the time taken to apply `f` on a list element.

```
scanl f id l = map (foldl f id ) (inits l)
scanr f id l = map (foldr f id ) (tails l)
```

`foldr1`: Fold without identity element

```
foldr1 f (x:xs) = if null xs then x else f x (foldr1 f xs)
```

Reasoning about Functional Programs

Bookkeeping law: If f is associative with identity a . Then:

$\text{foldr } f \ a \ . \ \text{concat} = \text{foldr } f \ a \ . \ \text{map } (\text{foldr } f \ a)$

Generalized Horner's rule:

$$1 + x_0 + x_0 * x_1 + x_0 * x_1 * x_2 + x_0 * x_1 * x_2 * x_3 \quad - \text{foldr1 } + \ . \ \text{scanl } * \ 1 \\ = 1 + x_0 * (1 + x_1 * (1 + x_2 * (1 + x_3))) \quad - \text{foldr } f \ 1 \text{ where } f \ x \ y = 1 + x * y$$

Under what conditions is:

$\text{foldr1 } \oplus \ . \ \text{scanl } \otimes \ \text{id}$ the same as $\text{foldr } \odot \ \text{id}$

where $x \odot y = \text{id} \oplus (x \otimes y)$?

Answer:

- 1 id should be the identity element of \otimes .
- 2 \otimes should distribute over \oplus .
- 3 \otimes should be associative.

The maximum segment sum problem

Given a sequence of integers, find the maximum of the sum of all (contiguous) segments.

Example:

`mss` `[-1, 2, -3, 5, -2, 1, 3, -2, -2, -3, 6] = 7`

An obviously correct but inefficient definition of `mss` :

```
mss = maximum . map sum . segs
segs = concat . map inits . tails
```

- ❶ `tails` – $O(n)$
- ❷ `map inits` – $O(n^3)$, assuming $O(n)$ sublists, each of length $O(n)$
- ❸ `concat` – $O(n^2)$, assuming $O(n)$ sublists, each of length $O(n)$.
- ❹ `map sum` – $O(n^3)$, assuming $O(n^2)$ sublists, each of length $O(n)$.

Reasoning about Functional Programs

```
mss = maximum . map sum . concat . map inits . tails
      {map f . concat = concat . map (map f)}
= maximum . concat . map (map sum) . map inits . tails
      {map f . map g = map (f . g)}
= maximum . concat . map (map sum . inits) . tails
      {maximum = foldr max -infinity}
= maximum . map maximum . map (map sum . inits) . tails
      {map f . map g = map (f . g)}
= maximum . map (maximum . map sum . inits) . tails
      {map sum . inits = scanl (+) 0}
= maximum . map (maximum . scanl (+) 0) . tails
      {maximum = foldr1 max}
      { 0 is id for +}
      { + is associative}
      { + distributes over max}
= maximum . map (foldr f 0) . tails
      where f x y = 0 'max' (x+y)
= maximum . scanr f 0
```

A Sudoku solver

As an example of

- List processing
- Backtracking in lazy languages

Reference: FUNCTIONAL PEARL - A program to solve Sudoku, RICHARD BIRD

www.cs.tufts.edu/~nr/comp150fp/archive/richard-bird/sudoku.pdf

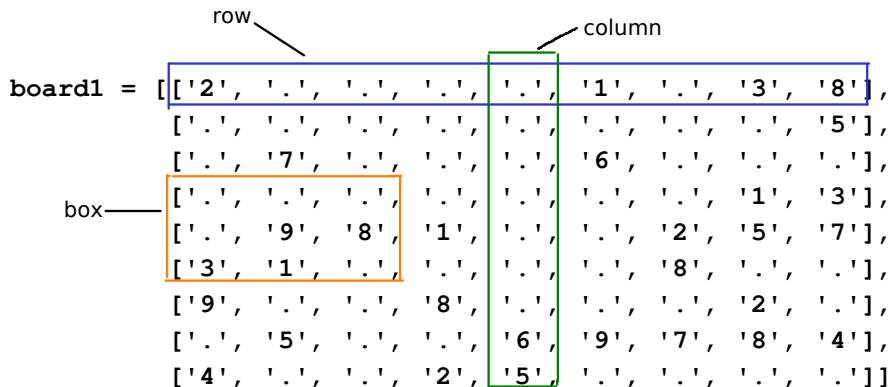
The Board

row

column

```
board1 = [ ['2', '.', '.', '.', '.', '1', '.', '3', '8'],  
            ['.', '.', '.', '.', '.', '.', '.', '.', '5'],  
            ['.', '7', '.', '.', '.', '6', '.', '.', '.'],  
            ['.', '.', '.', '.', '.', '.', '.', '1', '3'],  
            ['.', '9', '8', '1', '.', '.', '2', '5', '7'],  
            ['3', '1', '.', '.', '.', '.', '8', '.', '.'],  
            ['9', '.', '.', '8', '.', '.', '.', '2', '.'],  
            ['.', '5', '.', '.', '6', '9', '7', '8', '4'],  
            ['4', '.', '.', '2', '5', '.', '.', '.', '.'] ]
```

box



```
type Matrix a = [[a]]  
type Board = Matrix Char
```

Characterizing a correct solution

Some constants

```
boxsize = 3:: Int
allvals = "123456789"
blank c = c == '.'
```

A Board is correct, if each row, each column and each box is free of duplicates.

```
correct :: Board -> Bool
```

```
correct b = all nodups (rows b) &&
             all nodups (cols b) &&
             all nodups (boxes b)
```

```
nodups [] = True
```

```
nodups (x:xs) = notElem x xs && nodups xs
```

Characterizing a correct solution

```
rows = id
```

cols makes rows of columns:

```
cols [] = replicate 9 []  
cols (r:rs) = zipWith (:) r (cols rs)
```

boxes makes rows of boxes:

```
board1 = [['2', '.', '.', '.', '.', '.', '1', '.', '3', '8'],  
          ['.', '.', '.', '.', '.', '.', '2', '.', '5'],  
          ['.', '7', '.', '.', '.', '.', '6', '.', '4'],  
          ['.', '9', '8', '1', '.', '2', '5', '7'],  
          ['3', '1', '.', '.', '.', '8', '2', '4'],  
          ['9', '.', '.', '8', '.', '1', '2', '.'],  
          ['.', '5', '.', '6', '9', '7', '8', '4'],  
          ['4', '.', '.', '2', '5', '1', '3', '.']]
```

Characterizing a correct solution

```
boxes = map unchop . unchop . map cols . chop . map chop
```

```
chop :: [a] -> [[a]]
```

```
chop = chopBy boxsize
```

```
  where chopBy bsize [] = []
```

```
        chopBy bsize l = (take bsize l) :
```

```
                        (chopBy bsize (drop bsize l))
```

```
unchop = concat
```

Notice that rows, cols or boxes done twice gives identity.

```
rows . rows = id
```

```
cols . cols = id
```

```
boxes . boxes = id
```

Choices

A Choice is a list of characters, that represent the choices for a cell

- Initially, the choices for a blank cell are all possible characters, and the choices for a non-blank cell is the only character in the cell.

```
type Choices = [Char]

initialChoices :: Board -> Matrix Choices
initialChoices = map (map fillin)
  where fillin initialChar = if blank initialChar then allvals
                             else [initialChar]
```

From a Matrix of Choices, we want to generate all possible boards.

How does one do that?

Easier problem: From a list of choices, how does one generate all possible list?

Choices

```
cp :: [[a]] -> [[a]]  -- cp for cartesian product
cp [] = [[]]
cp (x:xs) = [h:t | h <- x, t <- (cp xs)]
```

How can one use cp, to calculate the matrix cartesian product, mcp?

```
mcp :: Matrix [a] -> [Matrix a]
```

Surprisingly, mcp is easy to define using cp?

```
mcp = cp (map cp)
```

map cp converts a Matrix of choices into

- [list of possible first rows
- list of possible second rows
- ...
- list of possible ninth rows]

cp converts it into possible matrix of Boards

Choices

A `Sudoku` solver takes a board and returns all possible completions of the board. Returns `[]` if there are none.

```
sudokusolver1 :: Board -> [Board]
```

```
--sudokusolver - first attempt
```

```
sudokusolver1 = filter correct . mcp . initialChoices
```

Pruning

We would like to do pruning of the following form:

24	2	34	12
34	234	134	13
124	23	13	4
14	123	123	3

4	2	34	1
34	34	134	1
12	3	13	4
14	1	12	3

This is one time pruning.

Given a row, column or a box, we collect all the fixed choices and remove these from the non-fixed choices.

```
fixed :: [Choices] -> Choices -- fixed identifies fixed choices
fixed = concat . filter single
  where single [_] = True
        single _   = False
```

Pruning

pruneList takes a list of choices and prunes it into a list of choices:

```
pruneList :: [Choices] -> [Choices]
```

```
pruneList css = map (remove (fixed css)) css
```

```
  where remove fs cs = if single cs then cs else delete fs cs
```

```
      delete fs cs = filter (\c -> not (c `elem` fs)) cs
```

Now pruneMatrix prunes each row, each column and each box using pruneList

24	2	34	12
34	234	134	13
124	23	13	4
14	123	123	3

24	2	34	234
----	---	----	-----

4	2	34	12
34	34	134	13
124	23	13	4
14	123	123	3

Pruning

The rows pruning can be done by

```
rows . map pruneList . rows
```

Similarly for pruning by columns and boxes. We therefore abstract:

```
pruneBy f = f . map pruneList . f
pruneMatrix = pruneBy boxes . pruneBy cols . pruneBy rows

sudokusolver2 :: Board -> [Board]
sudokusolver2 = filter correct.mcp.pruneMatrix.initialChoices
```

↑
plug in your own
pruning strategy here

Expand → Prune → Expand → Prune

24	2	4	12
34	234	134	13
124	23	13	4
14	123	123	3

Expand



24	2	34	12
34	234	134	13
124	23	13	4
14	123	123	3

Expand



24	2	3	12
34	234	134	13
124	23	13	4
14	123	123	3

Prune



	2	4	1
34	34	13	1
12	3	1	4
14	1	12	3

Blocked

4	2	3	1
34	34	14	1
12	3	1	4
14	1	12	3

Blocked

Expand → Prune → Expand → Prune

Take a Choice Matrix that has a cell with at least two (say x) choices, and generate x Choice Matrices with fixed choice for this cell.

```
expand :: Matrix Choices -> [Matrix Choices]
```

Sometimes a Choice Matrix can be blocked. Conditions are

1. No choices for a cell,
2. Same fixed choices in more than one cells in row, col or box.

We shall discard blocked matrices during expansion and pruning.

Expand → Prune → Expand → Prune

```
blocked :: Matrix Choices -> Bool
blocked cm = void cm || not (safe cm)

void :: Matrix Choices -> Bool
void cm = any (any null) cm

safe :: Matrix Choices -> Bool
safe cm = all (nodups . fixed) (rows cm) &&
          all (nodups . fixed) (cols cm) &&
          all (nodups . fixed) (boxes cm)
```

Expand → Prune → Expand → Prune

To expand, we select the first cell that has the minimum number of choices amongst all cell which have more than one choices.

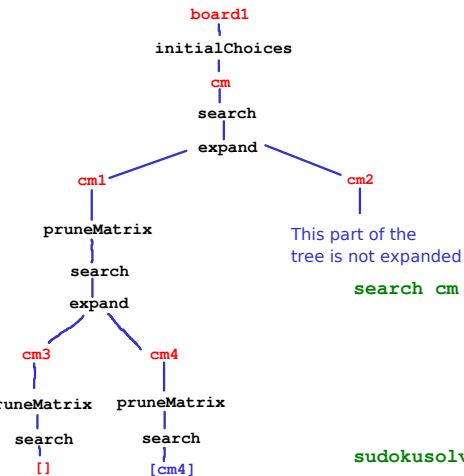
```
minchoice = minimum . filter ( > 1 ) . concat . map (map length)
```

A choice list is a candidate for expansion if its length is the same as the minimum. We pick the first candidate for-- expansion. This goes as follows:

```
expand cm = [rows1 ++ [row1 ++ [c] : row2] ++ rows2 | c <- cs]
  where (rows1, row:rows2) = break (any isCandidate) cm
        (row1, cs:row2) = break isCandidate row
        isCandidate cs = (length cs == n)
        n = minchoice cm
```


The Final Solution

```
search cm | blocked cm = []  
  | all (all single) cm = [cm]  
  | otherwise = (concat .  
                  map (search . pruneMatrix) .  
                  expand) cm  
  
sudokusolver3 = map (map head) . head . search . initialChoices
```



```

search cm | blocked cm = []
          | all (all single) cm = [cm]
          | otherwise = (concat .
                        map (search . pruneMatrix) .
                        expand) cm
  
```

```

sudokusolver3 = map (map head) .
                  head . search .
                  initialChoices
  
```