CSE 331 Fall 2017

Homework 2: Q3

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1 Algorithm Idea

Begin Algorithm Idea: set n be the number of student s and engineers e, and $m = n^2$. So that m > n will always be satisfied. In this algorithm, I will set each student meet with all the engineer possible. However, it will be ended up pairing them as (s_1, e_n) , (s_2, e_{n-1}) ,, (s_n, e_1) .

2 Algorithm Details

Algorithm: Algorithm1

```
Initially all student and engineer are free for s \ni n, 1...n //loop through all the student for e \ni n, 1...n //loop through all the engineer if e is the last engineer pair (s, e) //pair up each student and engineer as (s_1, e_n), (s_2, e_{n-1}), ....., (s_n, e_1) else student meet with engineer and leave itself endif endfor
```

3 Proof Idea

Begin Proof Idea: In this algorithm, we will be demonstrating an example of n=2, m=4. Whereas m>n. I'll let the first student meet with all the engineer, if he is the last engineer they will match. However, student will meet with another engineer and do nothing. So on so forth, when it iterate all the student the result should be (s_1, e_n) , (s_2, e_{n-1}) ,, (s_n, e_1) , which no one gets stood-up.

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4 Proof Details

Begin Proof Details

CSE Major	Slot 1	Slot 2	Slot 3	Slot 4
S ₁	E ₁	Free	E ₂ (truncate here)	
S ₂	Free	E ₁ (truncate here)		

Based on the schedule we can see that there is n (which is 2) number of students and engineers, and m (n^2 =4) slots. While we iterate through all the students, we check if he meets the last engineer which is E_2 in this case. Before he has met E_2 , he will first meet with E_1 . However, he will just go talk to the E_1 and not do anything until he met E_2 and assigned with E_1 . For E_2 , since our E_1 we will pair E_2 with E_1 which is E_2 and not care about the rest.