CS61C Spring 2025

Precheck: TLP and Concurrency Discussion 11

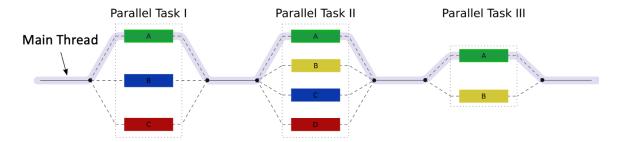
1 Pre-Check: T/F?

- 1.1 Multithreaded code will always outperform single-threaded code.
- 1.2 Race conditions occur when multiple threads attempt to modify the same resource at the same time.
- 1.3 Every thread has its own set of registers, program counter, heap and global variables
- 1.4 Multithreaded programs can still execute on single-core processors

2 Thread-Level Parallelism

Multithreading improves performance by utilizing a form of parallelism called **thread-level par-allelism**. The key idea behind thread-level parallelism is splitting from a single line of execution to multiple lines executing concurrently.

Multithreaded programs will start with a main thread that will spawn multiple threads when parallelism is required (See Lecture 31).



In this class, we use the OpenMP library to create and manage threads. Consider the sample hello world program, which prints "hello world from thread #" from each thread:

```
int main() {
    #pragma omp parallel
    {
        int thread_id = omp_get_thread_num();
        printf("hello world from thread %d\n", thread_id);
    }
}
```

This program will create a team of parallel threads. Each thread prints out a hello world message, along with its own thread number.

Let's break down the pragma omp parallel line:

- pragma tells the compiler that the rest of the line is a directive.
- omp declares that the directive is for OpenMP.
- parallel says that the following block statement the part inside the curly braces ({/}) should be executed in parallel by different threads.

Note that each thread has its own registers (including stack pointer) and program counter (PC). However, memory in the heap or global variables are shared amongst all threads.

3 OpenMP

OpenMP provides an easy interface for using multithreading within C programs. Some examples of OpenMP directives:

• The **parallel** directive indicates that each thread should run a copy of the code within the block. If a for loop is put within the block, **every** thread will run every iteration of the for loop.

```
#pragma omp parallel
{
    ...
}
```

NOTE: The opening curly brace needs to be on a newline or else there will be a compile-time error!

• The parallel for directive will split up iterations of a for loop over various threads. Every thread will run different iterations of the for loop. The exact order of execution across all threads, as well as the number of iterations each thread performs, are both non-deterministic, as the OpenMP library load balances threads for performance. The following two code snippets are equivalent.

```
#pragma omp parallel for
for (int i = 0; i < n; i++) {
    ...
}

#pragma omp parallel
{
    #pragma omp for
    for (int i =0; i < n; i++) { ... }
}</pre>
```

The critical directive only allows a single thread to access the following line / block of code at
once. It is useful to prevent race conditions when modifying resources shared between threads.

```
// Assume arr has length n
int fast_sum(int *arr, int n) {
   int result = 0;
   int result2 = 0;
   #pragma omp parallel for
   for (int i = 0; i < n; i++) {
        #pragma omp critical
        result += arr[i]; // limited to one thread access at once
   }
   return result;
}</pre>
```

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- The parallel for reduction (operation: var) directive creates and optimizes the critical section for a loop, given a variable that should be in the critical section and the operation being performed on that variable. An example is given below.

```
// Assume arr has length n
int fast_sum(int *arr, int n) {
   int result = 0;
   #pragma omp parallel for reduction(+: result)
   for (int i = 0; i < n; i++) {
      result += arr[i];
   }
   return result;
}</pre>
```

Additionally, there are two OpenMP functions that may be useful to you:

- int omp_get_thread_num() will return the number of the thread executing the code
- int omp_get_num_threads() will return the number of total hardware threads executing the code

4 Race Conditions

Two threads attempting to access the same memory can result in a race condition. Consider the following:

```
void example(int n, int *a) {
  int result = 0;
  #pragma omp parallel for
  for (int i = 0; i < n; i += 1) {
    result += a[i];
  }
}</pre>
```

For a single thread to compute the line result += a[i]:

- (a) Read the value a[i]
- (b) Compute the value result + a[i]
- (c) Store the new value back to result

Different threads often have their instructions interleaved. Without defining a critical section, these threads may be executing steps (a), (b), and (c) at any time and the value of **result** may be indeterminate.

To fix this, we can introduce a critical section to limit one thread to executing **result += a[i]** at any given time:

```
void example(int n, int *a) {
  int result = 0;
  #pragma omp parallel for
  for (int i = 0; i < n; i += 1) {
    #pragma omp critical
    result += a[i];
  }
}</pre>
```

It can also be fixed with the **reduction** directive (see above)!