

The Dart type system

The Dart language is type safe: it uses a combination of static type checking and [runtime checks](#) to ensure that a variable's value always matches the variable's static type, sometimes referred to as sound typing. Although *types* are mandatory, type *annotations* are optional because of [type inference](#).

For a full introduction to the Dart language, including types, see the [language tour](#).

One benefit of static type checking is the ability to find bugs at compile time using Dart's [static analyzer](#).

You can fix most static analysis errors by adding type annotations to generic classes. The most common generic classes are the collection types `List<T>` and `Map<K, V>`.

For example, in the following code the `printInts()` function prints an integer list, and `main()` creates a list and passes it to `printInts()`.

```
x static analysis: error/warning
void printInts(List<int> a) => print(a);

void main() {
  var list = [];
  list.add(1);
  list.add('2');
  printInts(list);
}
```

The preceding code results in a type error on `list` (highlighted above) at the call of `printInts(list)`:

```
error - The argument type 'List<dynamic>' can't be assigned to the parameter type 'List<int>'. -
argument_type_not_assignable
```

The error highlights an unsound implicit cast from `List<dynamic>` to `List<int>`. The `list` variable has static type `List<dynamic>`. This is because the initializing declaration `var list = []` doesn't provide the analyzer with enough information for it to infer a type argument more specific than `dynamic`. The `printInts()` function expects a parameter of type `List<int>`, causing a mismatch of types.

When adding a type annotation (`<int>`) on creation of the list (highlighted below) the analyzer complains that a string argument can't be assigned to an `int` parameter. Removing the quotes in `list.add('2')` results in code that passes static analysis and runs with no errors or warnings.

```
✓ static analysis: success
void printInts(List<int> a) => print(a);

void main() {
  var list = <int>[];
  list.add(1);
  list.add(2);
  printInts(list);
}
```

[Try it in DartPad](#).

What is soundness?

Soundness is about ensuring your program can't get into certain invalid states. A sound *type system* means you can never get into a state where an expression evaluates to a value that doesn't match the expression's static type. For example, if an expression's static type is `String`, at runtime you are guaranteed to only get a string when you evaluate it.

Dart's type system, like the type systems in Java and C#, is sound. It enforces that soundness using a combination of static checking (compile-time errors) and runtime checks. For example, assigning a `String` to `int` is a compile-time error. Casting an object to a `String` using `as String` fails with a runtime error if the object isn't a `String`.

The benefits of soundness

A sound type system has several benefits:

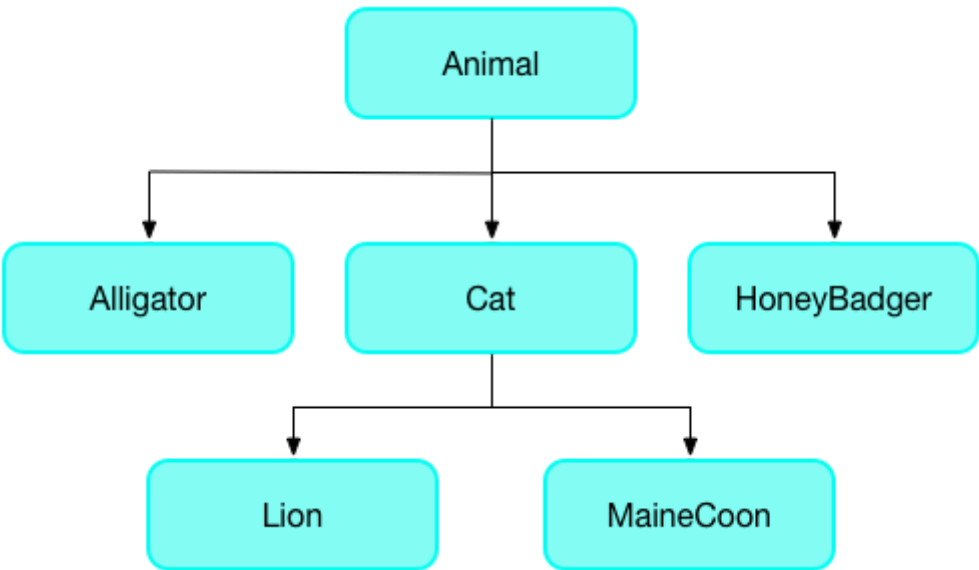
- Revealing type-related bugs at compile time.
A sound type system forces code to be unambiguous about its types, so type-related bugs that might be tricky to find at runtime are revealed at compile time.
- More readable code.
Code is easier to read because you can rely on a value actually having the specified type. In sound Dart, types can't lie.
- More maintainable code.
With a sound type system, when you change one piece of code, the type system can warn you about the other pieces of code that just broke.
- Better ahead of time (AOT) compilation.
While AOT compilation is possible without types, the generated code is much less efficient.

Tips for passing static analysis

Most of the rules for static types are easy to understand. Here are some of the less obvious rules:

- Use sound return types when overriding methods.
- Use sound parameter types when overriding methods.
- Don't use a dynamic list as a typed list.

Let's see these rules in detail, with examples that use the following type hierarchy:



Use sound return types when overriding methods

The return type of a method in a subclass must be the same type or a subtype of the return type of the method in the superclass. Consider the getter method in the `Animal` class:

```
class Animal {  
  void chase(Animal a) { ... }  
  Animal get parent => ...  
}
```

The `parent` getter method returns an `Animal`. In the `HoneyBadger` subclass, you can replace the getter's return type with `HoneyBadger` (or any other subtype of `Animal`), but an unrelated type is not allowed.

```
✓ static analysis: success  
  
class HoneyBadger extends Animal {  
  @override  
  void chase(Animal a) { ... }  
  
  @override  
  HoneyBadger get parent => ...  
}
```

```
x static analysis: error/warning

class HoneyBadger extends Animal {
    @override
    void chase(Animal a) { ... }

    @override
    Root get parent => ...
}
```

Use sound parameter types when overriding methods

The parameter of an overridden method must have either the same type or a supertype of the corresponding parameter in the superclass. Don't "tighten" the parameter type by replacing the type with a subtype of the original parameter.

Note: If you have a valid reason to use a subtype, you can use the [covariant keyword](#).

Consider the `chase(Animal)` method for the `Animal` class:

```
class Animal {
    void chase(Animal a) { ... }
    Animal get parent => ...
}
```

The `chase()` method takes an `Animal`. A `HoneyBadger` chases anything. It's OK to override the `chase()` method to take anything (`Object`).

```
✓ static analysis: success

class HoneyBadger extends Animal {
    @override
    void chase(Object a) { ... }

    @override
    Animal get parent => ...
}
```

The following code tightens the parameter on the `chase()` method from `Animal` to `Mouse`, a subclass of `Animal`.

```
x static analysis: error/warning

class Mouse extends Animal {...}

class Cat extends Animal {
    @override
    void chase(Mouse x) { ... }
}
```

This code is not type safe because it would then be possible to define a cat and send it after an alligator:

```
Animal a = Cat();
a.chase(Alligator()); // Not type safe or feline safe.
```

Don't use a dynamic list as a typed list

A `dynamic` list is good when you want to have a list with different kinds of things in it. However, you can't use a `dynamic` list as a typed list.

This rule also applies to instances of generic types.

The following code creates a `dynamic` list of `Dog`, and assigns it to a list of type `Cat`, which generates an error during static analysis.

```
x static analysis: error/warning

class Cat extends Animal { ... }

class Dog extends Animal { ... }

void main() {
  List<Cat> foo = <dynamic>[Dog()]; // Error
  List<dynamic> bar = <dynamic>[Dog(), Cat()]; // OK
}
```

Runtime checks

Runtime checks in the Dart VM and [dartdevc](#) deal with type safety issues that the analyzer can't catch.

For example, the following code throws an exception at runtime because it is an error to cast a list of dogs to a list of cats:

```
x runtime: error

void main() {
  List<Animal> animals = [Dog()];
  List<Cat> cats = animals as List<Cat>;
}
```

Type inference

The analyzer can infer types for fields, methods, local variables, and most generic type arguments. When the analyzer doesn't have enough information to infer a specific type, it uses the `dynamic` type.

Here's an example of how type inference works with generics. In this example, a variable named `arguments` holds a map that pairs string keys with values of various types.

If you explicitly type the variable, you might write this:

```
Map<String, dynamic> arguments = {'argA': 'hello', 'argB': 42};
```

Alternatively, you can use `var` or `final` and let Dart infer the type:

```
var arguments = {'argA': 'hello', 'argB': 42}; // Map<String, Object>
```

The map literal infers its type from its entries, and then the variable infers its type from the map literal's type. In this map, the keys are both strings, but the values have different types (`String` and `int`, which have the upper bound `Object`). So the map literal has the type `Map<String, Object>`, and so does the `arguments` variable.

Field and method inference

A field or method that has no specified type and that overrides a field or method from the superclass, inherits the type of the superclass method or field.

A field that does not have a declared or inherited type but that is declared with an initial value, gets an inferred type based on the initial value.

Static field inference

Static fields and variables get their types inferred from their initializer. Note that inference fails if it encounters a cycle (that is, inferring a type for the variable depends on knowing the type of that variable).

Local variable inference

Local variable types are inferred from their initializer, if any. Subsequent assignments are not taken into account. This may mean that too precise a type may be inferred. If so, you can add a type annotation.

```
x static analysis: error/warning

var x = 3; // x is inferred as an int.
x = 4.0;
```

```
✓ static analysis: success

num y = 3; // A num can be double or int.
y = 4.0;
```

Type argument inference

Type arguments to constructor calls and [generic method](#) invocations are inferred based on a combination of downward information from the context of occurrence, and upward information from the arguments to the constructor or generic method. If inference is not doing what you want or expect, you can always explicitly specify the type arguments.

```
✓ static analysis: success

// Inferred as if you wrote <int>[].
List<int> listOfInt = [];

// Inferred as if you wrote <double>[3.0].
var listOfDouble = [3.0];

// Inferred as Iterable<int>.
var ints = listOfDouble.map((x) => x.toInt());
```

In the last example, `x` is inferred as `double` using downward information. The return type of the closure is inferred as `int` using upward information. Dart uses this return type as upward information when inferring the `map()` method's type argument: `<int>`.

Substituting types

When you override a method, you are replacing something of one type (in the old method) with something that might have a new type (in the new method). Similarly, when you pass an argument to a function, you are replacing something that has one type (a parameter with a declared type) with something that has another type (the actual argument). When can you replace something that has one type with something that has a subtype or a supertype?

When substituting types, it helps to think in terms of *consumers* and *producers*. A consumer absorbs a type and a producer generates a type.

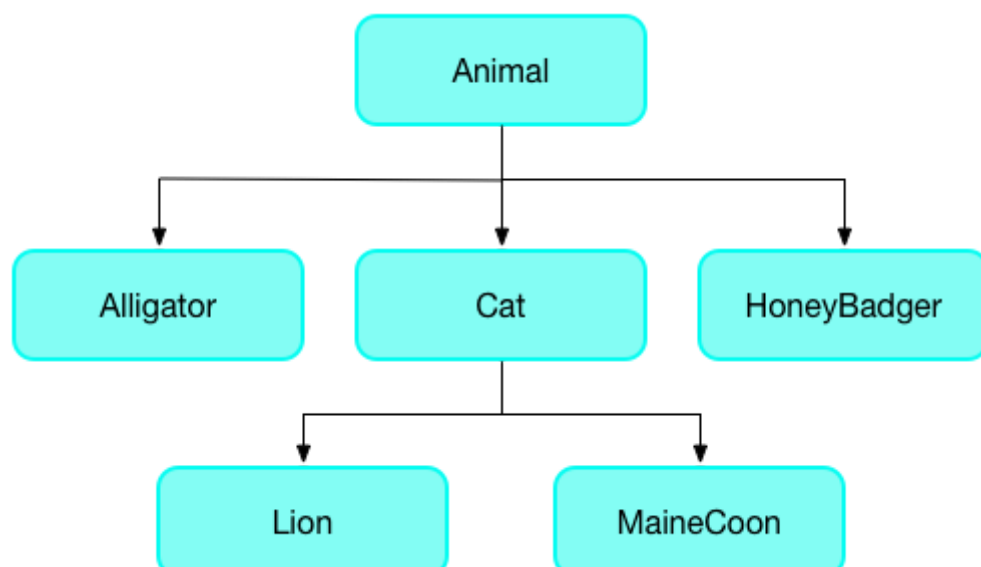
You can replace a consumer's type with a supertype and a producer's type with a subtype.

Let's look at examples of simple type assignment and assignment with generic types.

Simple type assignment

When assigning objects to objects, when can you replace a type with a different type? The answer depends on whether the object is a consumer or a producer.

Consider the following type hierarchy:



Consider the following simple assignment where `Cat c` is a *consumer* and `Cat()` is a *producer*.

```
Cat c = Cat();
```

In a consuming position, it's safe to replace something that consumes a specific type (`Cat`) with something that consumes anything (`Animal`), so replacing `Cat c` with `Animal c` is allowed, because `Animal` is a supertype of `Cat`.

```
✓ static analysis: success
Animal c = Cat();
```

But replacing `Cat c` with `MaineCoon c` breaks type safety, because the superclass may provide a type of `Cat` with different behaviors, such as `Lion`:

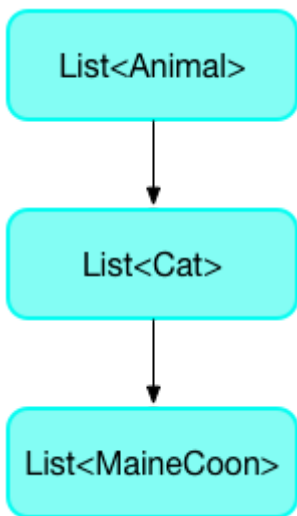
```
✗ static analysis: error/warning
MaineCoon c = Cat();
```

In a producing position, it's safe to replace something that produces a type (`Cat`) with a more specific type (`MaineCoon`). So, the following is allowed:

```
✓ static analysis: success
Cat c = MaineCoon();
```

Generic type assignment

Are the rules the same for generic types? Yes. Consider the hierarchy of lists of animals—a `List` of `Cat` is a subtype of a `List` of `Animal`, and a supertype of a `List` of `MaineCoon`:



In the following example, you can assign a `MaineCoon` list to `myCats` because `List<MaineCoon>` is a subtype of `List<Cat>`:

```
✓ static analysis: success
List<Cat> myCats = <MaineCoon>[];
```

What about going in the other direction? Can you assign an `Animal` list to a `List<Cat>`?

```
✗ static analysis: error/warning
List<Cat> myCats = <Animal>[];
```

This assignment doesn't pass static analysis because it creates an implicit downcast, which is disallowed from non-`dynamic` types such as `Animal`.

🔗 Version note: In packages that use a [language version](#) before 2.12 (when support for [null safety](#) was introduced), code can implicitly downcast from these non-`dynamic` types. You can disallow non-`dynamic` downcasts in a pre-2.12 project by specifying `implicit-casts: false` in the [analysis options file](#).

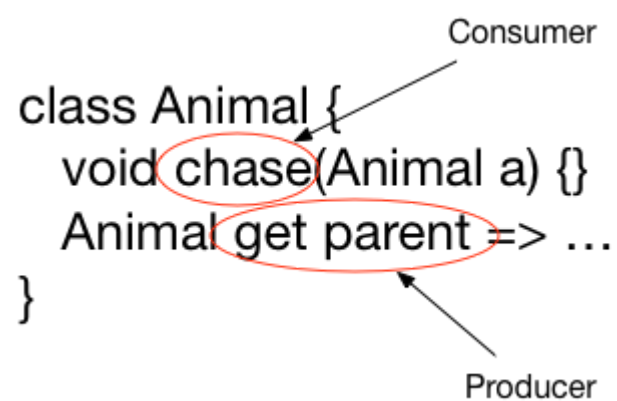
To make this code pass static analysis, use an explicit cast, which might fail at runtime.

```
List<Cat> myCats = <Animal>[] as List<Cat>;
```

Methods

When overriding a method, the producer and consumer rules still apply. For example:

```
class Animal {  
  void chase(Animal a) {}  
  Animal get parent => ...  
}
```



For a consumer (such as the `chase(Animal a)` method), you can replace the parameter type with a supertype. For a producer (such as the `parent` getter method), you can replace the return type with a subtype.

For more information, see [Use sound return types when overriding methods](#) and [Use sound parameter types when overriding methods](#).

Other resources

The following resources have further information on sound Dart:

- [Fixing common type problems](#) - Errors you may encounter when writing sound Dart code, and how to fix them.
- [Fixing type promotion failures](#) - Understand and learn how to fix type promotion errors.
- [Sound null safety](#) - Learn about writing code with sound null safety enabled.
- [Customizing static analysis](#) - How to set up and customize the analyzer and linter using an analysis options file.
- [Dart 2 migration guide](#) - How to update Dart 1.x code to Dart 2.