

# Exercise Sheet 3

Handout: Sept 27th — Deadline: October 4th - 4pm

**Question 3.1** (0.1 marks)

Consider the following input for MERGESORT:

12	10	4	2	9	6	5	25	8
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Illustrate the operation of the algorithm (follow how it was done in the lecture notes).

**Question 3.2** (0.45 marks) Prove using the **substitution method** the runtime of the MERGESORT Algorithm on an input of length  $n$ , as follows. Let  $n$  be an exact power of 2,  $n = 2^k$  to avoid using floors and ceilings. Use **mathematical induction over  $k$**  to show that the solution of the recurrence involving positive constants  $c, d > 0$

$$T(n) = \begin{cases} d & \text{if } n = 2^0 = 1 \\ 2T(n/2) + cn & \text{if } n = 2^k \text{ and } k \geq 1 \end{cases}$$

is  $T(n) = dn + cn \log n$  (we always use  $\log$  to denote the logarithm of base 2, so  $\log = \log_2$ ).

**Hint:** you may want to rewrite the above by replacing  $n$  with  $2^k$ . Then the task is to prove that  $T(2^k) = d2^k + c2^k \cdot k$  using the recurrence

$$T(2^k) = \begin{cases} d & \text{if } k = 0 \\ 2T(2^{k-1}) + c2^k & \text{if } k \geq 1 \end{cases}$$

**Question 3.3** (0.4 marks) Use the Master Theorem to give asymptotic tight bounds for the following recurrences. Justify your answers.

1.  $T(n) = 2T(n/4) + 1$
2.  $T(n) = 2T(n/4) + \sqrt{n}$
3.  $T(n) = 2T(n/4) + \sqrt{n} \log^2 n$
4.  $T(n) = 2T(n/4) + n$

**Question 3.4** (0.45 marks) Write the pseudo-code of the *recursive* BINARYSEARCH( $A, x, \text{low}, \text{high}$ ) algorithm discussed during the lecture to find whether a number  $x$  is present in an increasingly sorted array of length  $n$ . Write down its recurrence equation and prove that its runtime is  $\Theta(\log n)$  using the Master Theorem.

**Question 3.5** (0.6 marks) Solve programming problems "Heybale Feast", "A good problem", "Swiss" and "Bubble Sort II" provided on the Judge system.