

Week 2, Thursday

Review: root-finding problem: find p st $f(p)=0$

① bisection method: shrinkage the interval including the root

1. find a_1, b_1 s.t. $f(a_1)f(b_1) < 0$, $P_1 = \frac{a_1+b_1}{2}$

2. For $i=1 \dots, N$,

if $f(a_i)f(P_i) < 0$, set $a_{i+1} = a_i, b_{i+1} = P_i$

else if $f(a_i)f(P_i) > 0$, set $a_{i+1} = P_i, b_{i+1} = b_i$

Thm if $f \in C[a, b]$, $f(a)f(b) < 0$, then the sequence $\{P_n\}_{n=1}^{\infty}$ generated by the bisection method converges to p , which is the root of $f(x)$, i.e. $f(p)=0$ and $|P_n - p| \leq \frac{b-a}{2^n}$.

② fixed-point iteration: convert root-finding to fixed-point-finding
 $f(x)=0 \Leftrightarrow x=g(x)$ (or at least $f(p)=0 \Leftrightarrow p=g(p)$)

Thm (existence) $g(x) \in C[a, b]$, $g([a, b]) \subset [a, b]$

(uniqueness) $g'(x) \exists$ on (a, b) , and $|g'(x)| \leq K < 1$, $\forall x \in (a, b)$.

(fixed-point iteration)

The sequence: take $P_0 \in [a, b]$, $P_n = g(P_{n-1})$, $n=1, 2, \dots$.

Thm under the same condition in above theorem, the sequence $\{P_n\}_{n=0}^{\infty}$ converges to the unique fixed point, with speed:

$$\text{i}, |P^n - p| \leq K^n \max\{P_0 - a, b - P_0\}$$

$$\text{ii}, |P^n - p| \leq \frac{K^n}{1-K} |P_1 - P_0|$$

e.g. $f(x) = x^3 + 4x^2 - 10 = 0$, $f(1) = -5 < 0$, $f(2) = 14 > 0$, \exists root $\in [1, 2]$

$$x(3x^2 + 8x) = x(3x^2 + 8x) - (x^3 + 4x^2 - 10)$$

$$\Rightarrow x = x - \frac{x^3 + 4x^2 - 10}{3x^2 + 8x} := g_5(x)$$

$$P_0 = 1.5, P_4 = 1.365230013, \text{ super fast!}$$

compare with bisection method: $P_{27} \approx 1.365230013$

Rk: essentially, the convergence speed depends on $|g'(p)|$.

$$g_5(x) = x - \frac{f(x)}{f'(x)} \Rightarrow g'_5(x) = \frac{f(x)f''(x)}{(f'(x))^2} \Rightarrow g'_5(p) = 0$$

ΔQ : how to set $|g'(p)|$ as small as possible?

A general formula to construct fixed-point function:

$$x = x - \phi(x)f(x) := g(x), \text{ easy to check:}$$

$$p = g(p) \Leftrightarrow f(p) = 0 \text{ if } \phi(p) \neq 0. \text{ To this end:}$$

$$g'(x) = 1 - (\phi'(x)f(x) + \phi(x)f'(x)) \text{ if we want to have:}$$

$$g'(p) = 0 \Leftrightarrow 1 - (\phi'(p)f(p) + \phi(p)f'(p)) = 0 \Rightarrow \phi(p) = \frac{1}{f'(p)}$$

so it is reasonable to take $\phi(x) = \frac{1}{f'(x)}$, i.e.

$$g(x) = x - \frac{f(x)}{f'(x)}, \quad P_n = P_{n-1} - \frac{f(P_{n-1})}{f'(P_{n-1})} \rightarrow (\text{Newton's iteration})$$

Δ another way to derive Newton's iteration: for $P_0 \approx p$

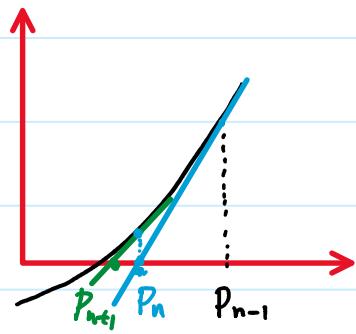
$$f(p) = f(P_0) + (p - P_0)f'(P_0) + \frac{1}{2}(p - P_0)^2 f''(\xi_0)$$

$$\text{when } |P_0 - p| \ll 1, |p - P_0|^2 \ll |p - P_0| \Rightarrow (p - P_0)f'(P_0) \approx -f(P_0)$$

$$P \approx P_0 - \frac{f(P_0)}{f'(P_0)}, \quad P_n = P_{n-1} - \frac{f(P_{n-1})}{f'(P_{n-1})} \rightarrow (\text{Newton's iteration})$$

△ Geometrically: $f'(P_{n-1}) = \frac{f(P_n) - 0}{P_{n-1} - P_n}$

$y - f(P_{n-1}) = f'(P_{n-1})(x - P_{n-1})$: the line tangent $f(x)$ intersect x -axis at P_n at $(P_{n-1}, f(P_{n-1}))$



Input: P_0, ε, N_0

Output: approximation of the root.

1. Set $i=0$, while $i \leq N_0$, do 2-5
2. If $f'(P_0)=0$, STOP, output ("failure")
3. set $p = P_0 - \frac{f(P_0)}{f'(P_0)}$
4. If $|P - P_0| < \varepsilon$ / or $|f(p)| < \varepsilon$. STOP, output (p).
5. Set $i=i+1$, $P_0=p$
6. STOP, output ("failure")

e.g. $f(x) = \cos(x) - x$

(a) fixed point with: $g(x) = \cos x$

$$P_0 = \frac{\pi}{4}; \quad P_n = \cos(P_{n-1}) \quad n=1, 2, \dots$$

(b) Newton's iteration with: $g(x) = x - \frac{\cos x - x}{-\sin x - 1}$

$$P_0 = \frac{\pi}{4}; \quad P_n = P_{n-1} - \frac{P_{n-1} - \cos(P_{n-1})}{1 + \sin(P_{n-1})}$$

(a)

$$n=0 \quad 0.7853$$

$$n=1 \quad 0.7071$$

$$n=2 \quad 0.7602$$

$$0.7247$$

$$0.7487$$

$$0.7325$$

(b)

$$0.7853981635 \leftarrow$$

$$0.7395361337$$

$$0.7390851781$$

$$0.7390851332$$

$$0.7390851332$$

, \vdots \vdots \vdots \vdots \vdots

0.7481

••1210001000

0.7325

much faster!

0.7435

0.7361

Q: How to quantify the speed of convergence?

Def: $\{P_n\} \rightarrow P$ and $P_n \neq P, \forall n, \exists \lambda, \alpha$ st.

$$\lim_{n \rightarrow \infty} \frac{|P_{n+1} - P|}{|P_n - P|^\alpha} = \lambda, \text{ then we say}$$

$\{P_n\}$ converges to P of order α with asymptotic error constant λ .

(i) $\alpha = 1, \lambda < 1$, the sequence is linearly convergent

(ii) $\alpha = 2$, quadratically convergent

e.g. $\{\alpha_n\}, \{\beta_n\} \rightarrow 0$, and

$$\lim_{n \rightarrow \infty} \frac{|\alpha_{n+1}|}{|\alpha_n|} = 0.5, \quad \lim_{n \rightarrow \infty} \frac{|\beta_{n+1}|}{|\beta_n|^2} = 0.5$$

$$|\alpha_n| \approx 0.5 |\alpha_{n-1}| \approx \dots \approx 0.5^n / \alpha_0$$

$$|\beta_n| \approx 0.5 |\beta_{n-1}|^2 \approx 0.5 (0.5 |\beta_{n-2}|^2)^2 \approx \dots \approx 0.5^{2^0 + 2^1 + 2^2 + \dots + 2^{n-1}} / |\beta_0|^{2^n}$$

$$= (0.5)^{2^{n-1}} (\beta_0)^{2^n}$$

$$\text{if } |\alpha_0| = |\beta_0| = 1$$

	α_n	β_n
1	0.5×10^0	0.5×10^0
2	0.25×10^0	0.125×10^{-1}
3	0.125×10^0	0.0625×10^{-2}
4	0.0625×10^{-1}	0.03125×10^{-5}
5	0.03125×10^{-1}	0.015625×10^{-10}
6	0.015625×10^{-1}	$0.0078125 \times 10^{-19}$
7	0.0078125×10^{-2}	$0.00390625 \times 10^{-39}$



Thm let $f \in C^2[a, b]$ $p \in (a, b)$ st. $f(p) = 0$ and $f'(p) \neq 0$

Thm * let $f \in C^2[a, b]$, $p \in (a, b)$ s.t. $f(p) = 0$ and $f'(p) \neq 0$,

then there exist a δ s.t. Newton's method generates a sequence converges to p for any $P_0 \in [p-\delta, p+\delta]$.

RK: basic idea is same as the fixed-point iteration,

i.e. find $(p-\delta, p+\delta)$ s.t. $|g'(x)| \leq k < 1$, $\forall x \in (p-\delta, p+\delta)$

Proof: as $g(x) = x - \frac{f(x)}{f'(x)}$, $g'(x) = \frac{f(x)f''(x)}{(f'(x))^2}$

Since $f \in C^2[a, b]$ and $f'(p) \neq 0$,

$\exists \delta_1 > 0$, s.t. $f'(x) \neq 0$, $\forall x \in [p-\delta_1, p+\delta_1] \subset [a, b]$

on the other hand, $g'(p) = 0$ and $g \in C^1[p-\delta_1, p+\delta_1]$

$\therefore \forall 0 < k < 1$, $\exists 0 < \delta \leq \delta_1$ s.t. $|g'(x)| \leq k$, $\forall x \in [p-\delta, p+\delta]$

Next show $g([p-\delta, p+\delta]) \subset [p-\delta, p+\delta]$: $\forall x \in [p-\delta, p+\delta]$

$$|g(x) - p| = |g(x) - g(p)| = |g'(x)| |x - p| \leq k |x - p| < \delta$$

$\therefore g(x) \in [p-\delta, p+\delta]$, $\forall x \in [p-\delta, p+\delta]$

By the theorem for fixed-point iteration, we have

$p_n \rightarrow p$ as $n \rightarrow \infty$.

Thm ① $g \in C[a, b]$, $g([a, b]) \subset [a, b]$

② $g'(x)$ is continuous on (a, b) and $|g'(x)| \leq k < 1$, $\forall x \in (a, b)$

③ $g'(p) \neq 0$, p is the fixed point

then $\{p_n\}$ generated by the fixed-point iteration

converges only linearly to the fixed point.

key idea: $\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|} = \lim_{n \rightarrow \infty} \frac{|g(p_n) - g(p)|}{|p_n - p|} = \boxed{\lim_{n \rightarrow \infty} |g'(z_n)|} = |g'(p)| < 1$

g' is continuous

$$\text{key idea: } \lim_{n \rightarrow \infty} \frac{|P_n - p|}{|P_n - p|} = \lim_{n \rightarrow \infty} \frac{|P_n - p|}{|P_n - p|} = \boxed{\lim_{n \rightarrow \infty} |g'(\xi_n)| = |g'(p)| < 1}$$

↓
mean value theorem.

Thm ① $p = g(p)$, $g'(p) = 0$

② g'' is continuous and $|g''(x)| < M$ on an open interval I , $p \in I$, then $\exists \delta > 0$ s.t. $\forall p_0 \in [p-\delta, p+\delta]$, the sequence defined by $P_n = g(P_{n-1}) \quad \forall n \geq 1$, converges to p at least quadratically

key idea: since g'' is continuous on I and $g'(p) = 0$, $\forall 0 < k < 1$, $\exists \delta > 0$, s.t. $[p-\delta, p+\delta] \subset I$ and $|g'(x)| \leq k$, $\forall x \in [p-\delta, p+\delta]$ and $\forall p_0 \in [p-\delta, p+\delta]$, $P_n = g(P_{n-1}) \rightarrow p$ as $n \rightarrow \infty$ (by thm *)

$$g(x) = g(p) + g'(p)(x-p) + \frac{g''(\xi)}{2}(x-p)^2$$

$$= p + \frac{1}{2}g''(\xi)(x-p)^2$$

$$\Rightarrow P_{n+1} = p + \frac{1}{2}g''(\xi_n)(P_n - p)^2$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{|P_{n+1} - p|}{|P_n - p|^2} = \frac{1}{2}|g''(p)| \quad \text{at least quadratically.}$$

$$\leq \frac{1}{2}M$$

Summary:

① bisection simple, always convergent, a good initial data producer for more faster methods, slow

② fixed-point iteration: faster than bisection with suitable choice of g . Usually linearly convergent. not easy to construct g .

③ Newton's method (can be viewed as a special fixed-point iteration with $g'(p) \neq 0$).

drawbacks i) require g to be differentiable, and $g'(p) \neq 0$

2) require g'' to be continuous to get quadratic convergence rate

3) sensitive to initial guess

advantage quadratically convergent, a uniform formula.

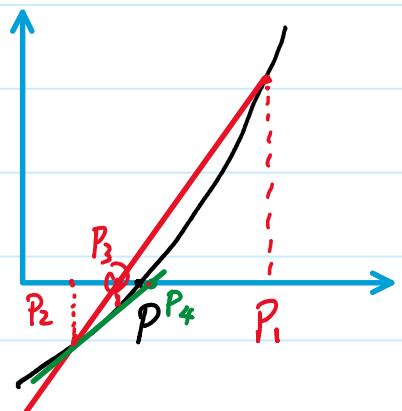
RK: ① what if f is not differentiable or difficult to compute $f'(x)$

(割线法)
Scant method: replace: $f'(P_{n-1}) \approx \frac{f(P_{n-1}) - f(P_{n-2})}{P_{n-1} - P_{n-2}}$, i.e.

$$P_n = P_{n-1} - \frac{P_{n-1} - P_{n-2}}{f(P_{n-1}) - f(P_{n-2})} f(P_{n-1})$$

red line: go through $(P_1, f(P_1)), (P_2, f(P_2))$ i.e. scant line

P_3 is the scant line intersects with x -axis.



② sensitive to initial guess (close enough to the root)
use bisection method to get a good initial.

③ Q: what if $f'(p)=0$? (HW 2.4 10)

HW 2 : 2.3 1, 2, 3 9, 4 9, 31, 32

2.4 8, 9, 10, 12

Using Newton's method, see what will happen:

$$\textcircled{1} f(x) = 1 - x^2, P_0 = 0 \quad \textcircled{2} f(x) = x^3 - 2x + 2, P_0 = 0$$

$$\textcircled{3} f(x) = \sqrt[3]{x}, P_0 = 1 \quad \textcircled{4} f(x) = \sqrt{|x|}, P_0 = 1$$