



Introduction to Mathematical Logic

For CS Students

CS104

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Decision Problems

1 Decidability

Definition

A **decision problem** is a problem which has a yes or no answer on a given input.

Example: The problem “Is an integer x positive?” is a decision problem.



Decidability

1 Decidability

Definition

A decision problem is **decidable** (可判定的) if there exists an algorithm that always halts with a yes/no answer for every input.

Otherwise, we say that the decision problem is **undecidable** (不可判定的).



The Collatz Conjecture

1 Decidability

For even numbers, divide by 2; for odd numbers, multiply by 3 and add 1. With enough repetition, do all positive integers converge to 1?

```
while (n > 1){  
    if ( (n % 2) == 0){  
        n = n/2;  
    } else {  
        n = 3 * n + 1;  
    }  
}
```

The Collatz Conjecture is one of the most famous unsolved problems in math, and is algorithmically undecidable.



The Halting Problem

1 Decidability

The Halting Problem (停机问题) is a decision problem in computability theory.

Given a computer program P and an input I , determine if P halts on input I .

Theorem (Turing, 1936): The Halting Problem is undecidable.

Theorem proved in class.

A nice animation for the proof: <https://www.youtube.com/watch?v=92WHN-pAFCs>



The Partial Correctness Problem

1 Decidability

Given a Hoare Triple $\langle P \rangle C \langle Q \rangle$, does C satisfy the triple under partial correctness?

Theorem: The Partial Correctness Problem is undecidable.

Proof Sketch: Assume towards a contradiction that there is an algorithm B that returns “yes” when $\langle P \rangle C \langle Q \rangle$ is satisfied under partial correctness and “no” otherwise. Then, we can construct an algorithm A that decides the Halting Problem. Contradiction.



Introduction to Mathematical Logic

Thank you for listening!
Any questions?