

Exercise Sheet 3

Handout: Sept 27th — Deadline: October 4th - 4pm

Question 3.1 (0.1 marks)

Consider the following input for MERGESORT:

12	10	4	2	9	6	5	25	8
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Illustrate the operation of the algorithm (follow how it was done in the lecture notes).

Question 3.2 (0.45 marks) Prove using the **substitution method** the runtime of the MERGESORT Algorithm on an input of length n , as follows. Let n be an exact power of 2, $n = 2^k$ to avoid using floors and ceilings. **Use mathematical induction over k** to show that the solution of the recurrence involving positive constants $c, d > 0$

$$T(n) = \begin{cases} d & \text{if } n = 2^0 = 1 \\ 2T(n/2) + cn & \text{if } n = 2^k \text{ and } k \geq 1 \end{cases}$$

is $T(n) = dn + cn \log n$ (we always use \log to denote the logarithm of base 2, so $\log = \log_2$).

Hint: you may want to rewrite the above by replacing n with 2^k . Then the task is to prove that $T(2^k) = d2^k + c2^k \cdot k$ using the recurrence

$$T(2^k) = \begin{cases} d & \text{if } k = 0 \\ 2T(2^{k-1}) + c2^k & \text{if } k \geq 1 \end{cases}$$

Question 3.3 (0.4 marks) Use the Master Theorem to give asymptotic tight bounds for the following recurrences. Justify your answers.

1. $T(n) = 2T(n/4) + 1$
2. $T(n) = 2T(n/4) + \sqrt{n}$
3. $T(n) = 2T(n/4) + \sqrt{n} \log^2 n$
4. $T(n) = 2T(n/4) + n$

Question 3.4 (0.45 marks) Write the pseudo-code of the *recursive* BINARYSEARCH($A, x, \text{low}, \text{high}$) algorithm discussed during the lecture to find whether a number x is present in an increasingly sorted array of length n . Write down its recurrence equation and prove that its runtime is $\Theta(\log n)$ using the Master Theorem.

Question 3.5 (0.6 marks) Solve programming problems "Heybale Feast", "A good problem", "Swiss" and "Bubble Sort II" provided on the Judge system.