

# Performance

## CS 540- High Performance Computing

*Spring 2017*

# Performance Analysis

- Measure the gain in application performance, leveraging multiple processors/threads
- Performance gain is not perfect and not always guaranteed
- Limitations of parallel computing
- Very simple techniques for calculating performance gain (i.e. speedup, Amdahl's law)

# Speedup

- The performance gain achieved by running your algorithm/application in parallel (i.e. multiple cores)
- The length of time it takes a program to run on a single processor divided by the time it takes to run on multiple processors.
- speedup ranges between 0 – p, p is the number of processors

# Speedup

- Very simple method for calculating speedup

$$S_p = \frac{\text{time for a single processor run}}{\text{time using } p \text{ processors}}$$

# Linear Speedup

- If one processor executes a specific task in time  $t$  and the number processors,  $p$ , can do the task in time  $t / p$
- For example: if 4 processors improve runtime by a factor of 4.



# Scalability

- The ability to achieve performance proportional to the number of processors used is called, *scalability*.
- Ideally, performance should improve as your parallel code leverages multiple processor
- Computing speedup is a good indicator of how efficient a program scales as the number of processor cores increases.

# Efficiency

- The average speedup per processor core .
- Efficiency is a fraction that ranges between 0 and 1

$$E_p = \frac{S_p}{p}$$

# Amdahl's Law

- An alternative formula for calculating speedup
- Named after Gene Amdahl, well respected computer scientists
- The law states speedup will be limited by its percentage (fraction) of sequential code, regardless of the number of processor cores used in a parallel application.



# Amdahl's Law cont.

- Defines speedup with  $p$  processors as follows:

$$S_p = \frac{1}{f + (1 - f) / p}$$

# Amdahl's Law cont.

- The term  $f$  stands for the fraction of instructions executed sequentially with one processor
- The term  $(1 - f)$  stands for the fraction of instructions executed in perfect parallelism with  $p$  processors.
- The sequential fraction of code,  $f$ , ranges between 0 and 1.
- If  $f$  is 0, there is no sequential code.
- If  $f$  is 1, there is no parallel code, then speedup is 1.
- Amdahl's speedup ranges between 1 and  $p$ , where  $p$  is the number of processors used in a parallel processing application.

# Poor Performance

- When the calculated speedup is less than one, it's a good indicator of poor performance.
- the parallel code runs slower than the sequential code.
- When there isn't enough processing, or work, to be completed by each processor core.
- The computing overhead of creating and managing the concurrent threads offsets the performance gain of parallel computing.

# Speedup Limitations

- **Too much I/O**, when the code is I/O bound. too much I/O compared to the amount of computation.
- **Wrong algorithm**, when the numerical algorithm is not appropriate for a parallel computer.
- **Too much memory contention**, when there is too much memory contention.
- **Wrong problem size**, when the data set is too small to take advantage of a parallel computer architecture.

# Speedup Limitations

- **Too much sequential code**, when there's too much sequential code. This is proved by Amdahl's Law.
- **Too much parallel overhead**, when there is too much parallel overhead compared to the amount of computation. Examples are the additional CPU cycles accumulated in creating parallel regions, creating threads, synchronizing threads, blocking threads.
- **Load imbalance**, when the processors have different workloads. The processors that finish early will be idle while they wait for other processors cores to complete or catch up.

# References

1. Blaise Barney, Lawrence Livermore National Laboratory,  
[https://computing.llnl.gov/tutorials/parallel\\_comp/#Overview](https://computing.llnl.gov/tutorials/parallel_comp/#Overview)
2. <https://www.citutor.org/index.php>, Parallel Computing Explained