

# Differentially Private Multi-Party Computation

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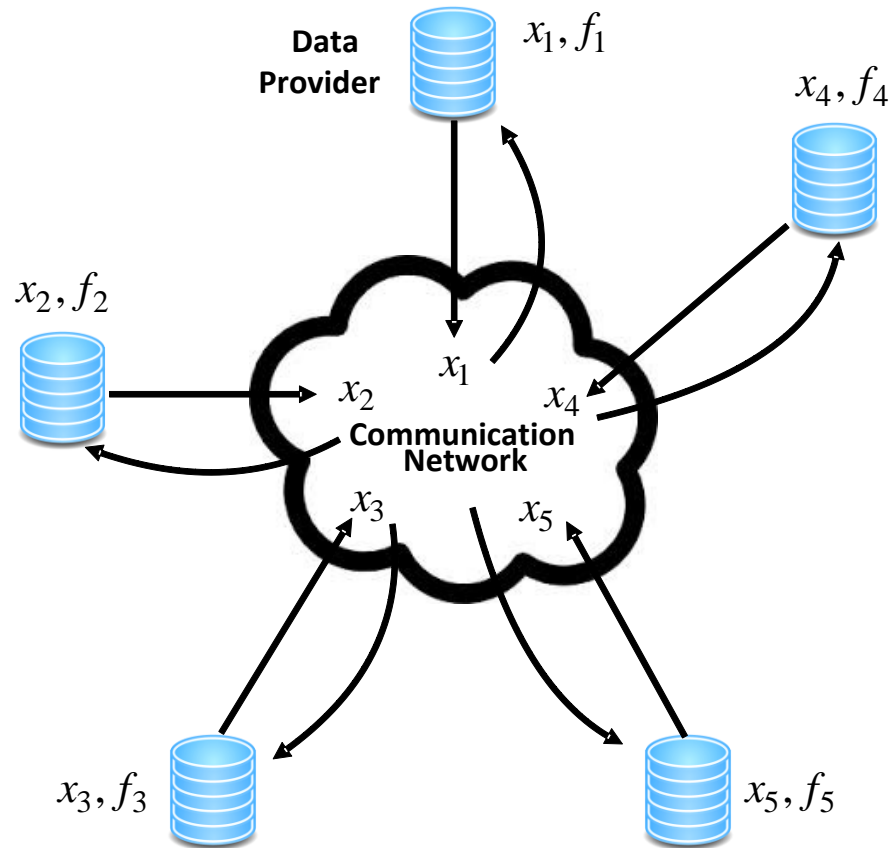


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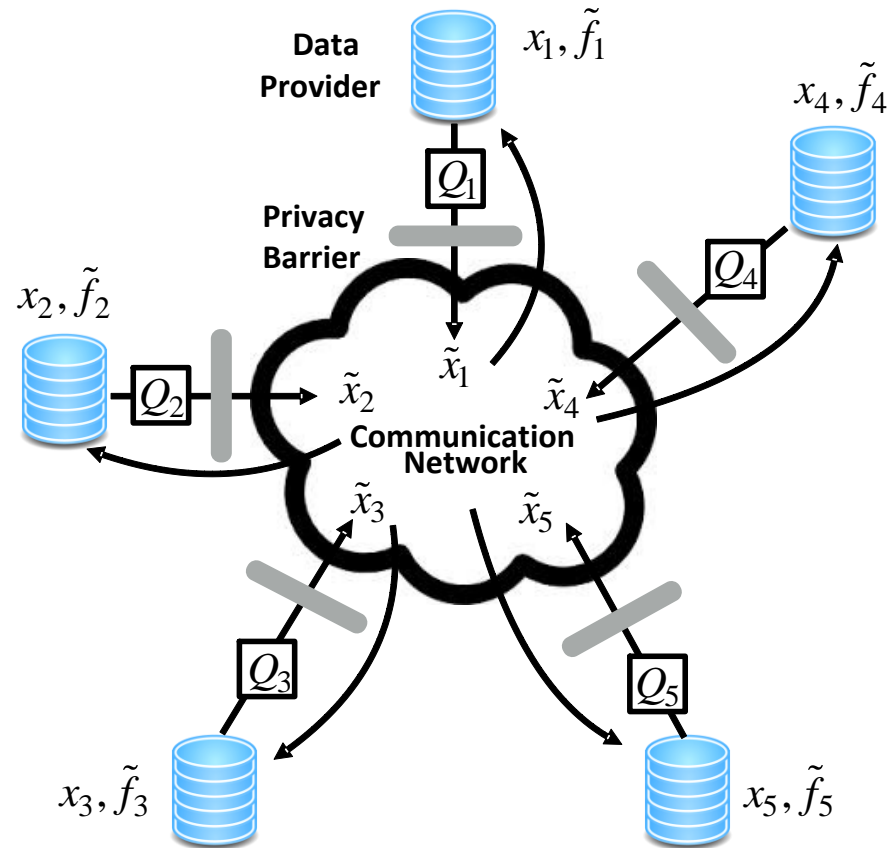


# Multi-party computation



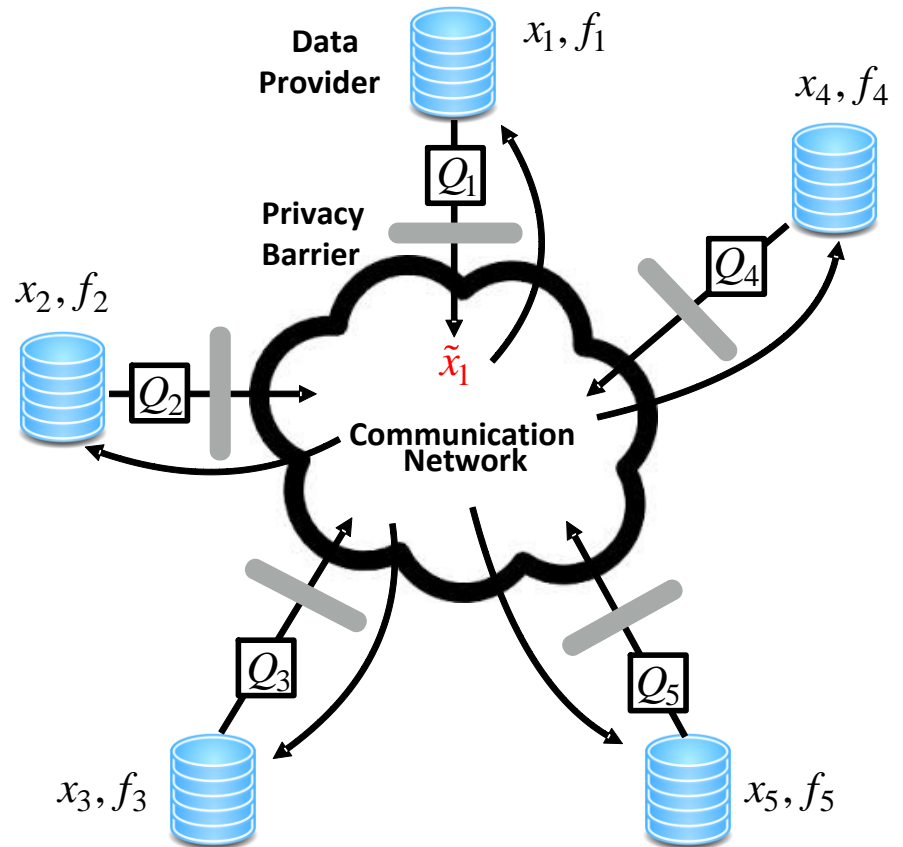
- important setting in distributed systems and cloud computing

# Private multi-party computation

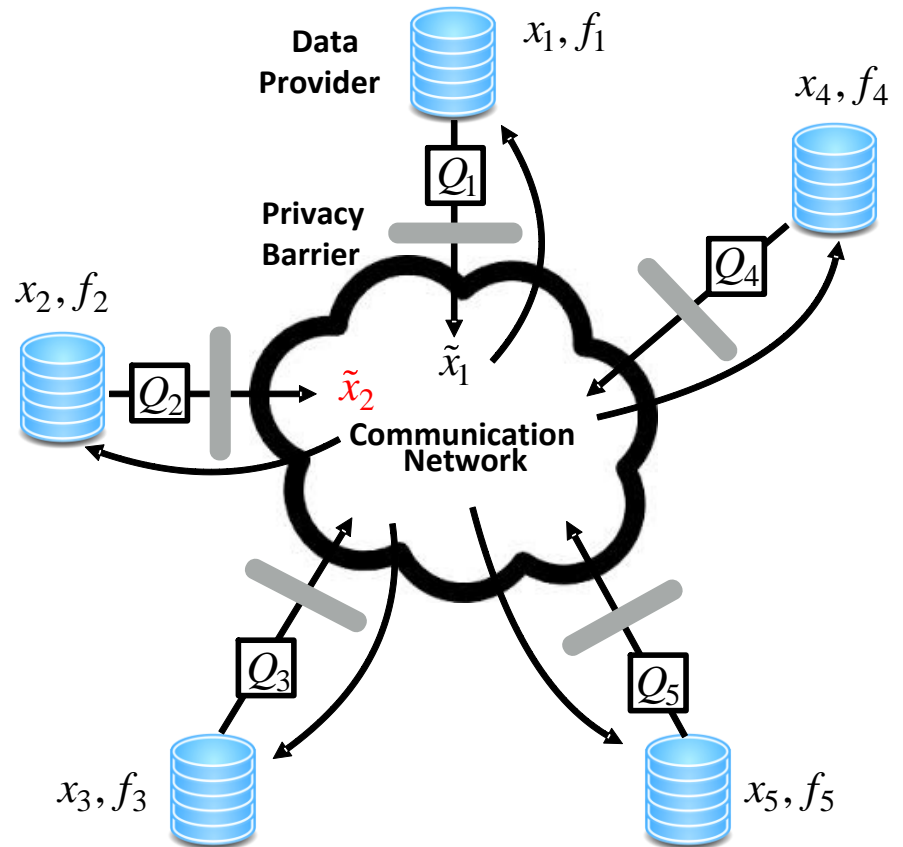
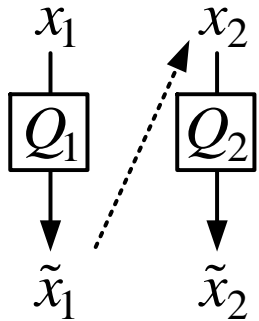


- each party shares a noisy version of its data

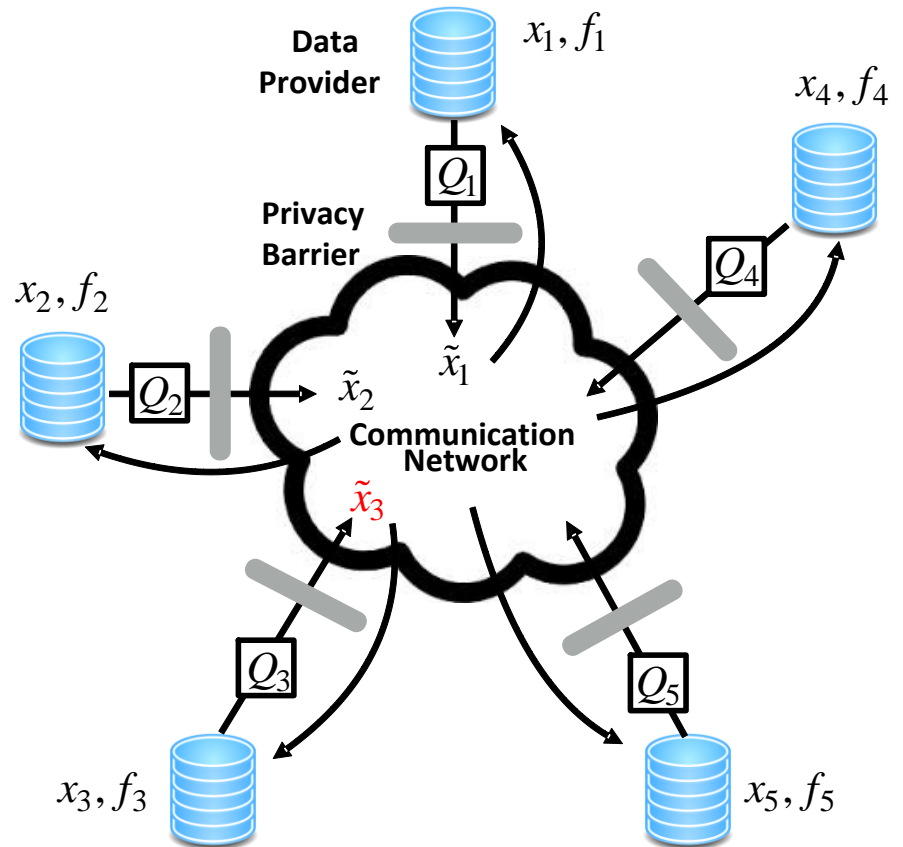
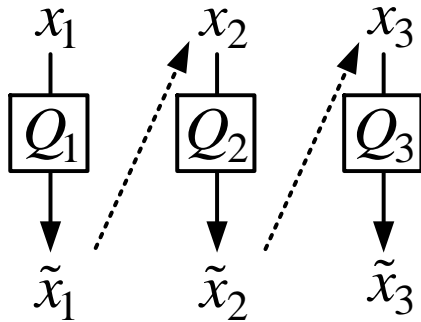
# Interactive mechanisms



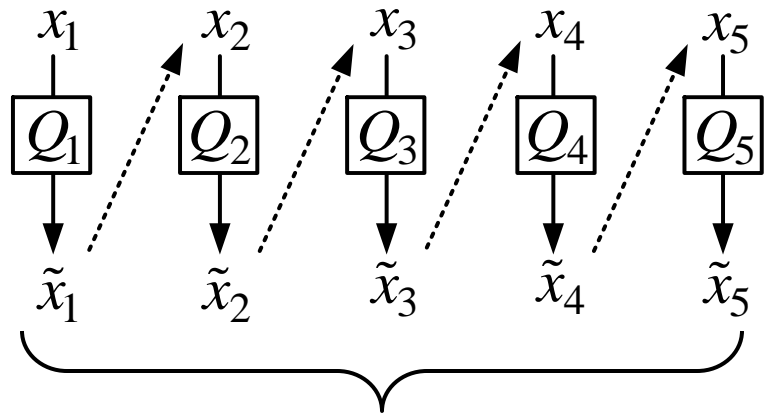
# Interactive mechanisms



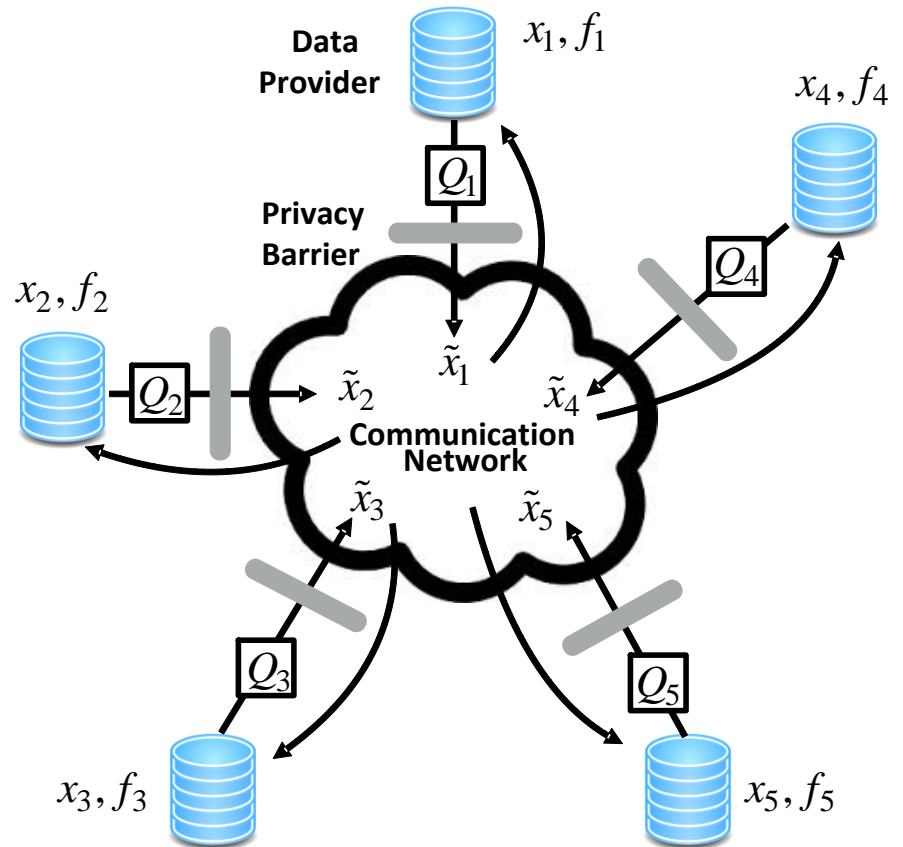
# Interactive mechanisms



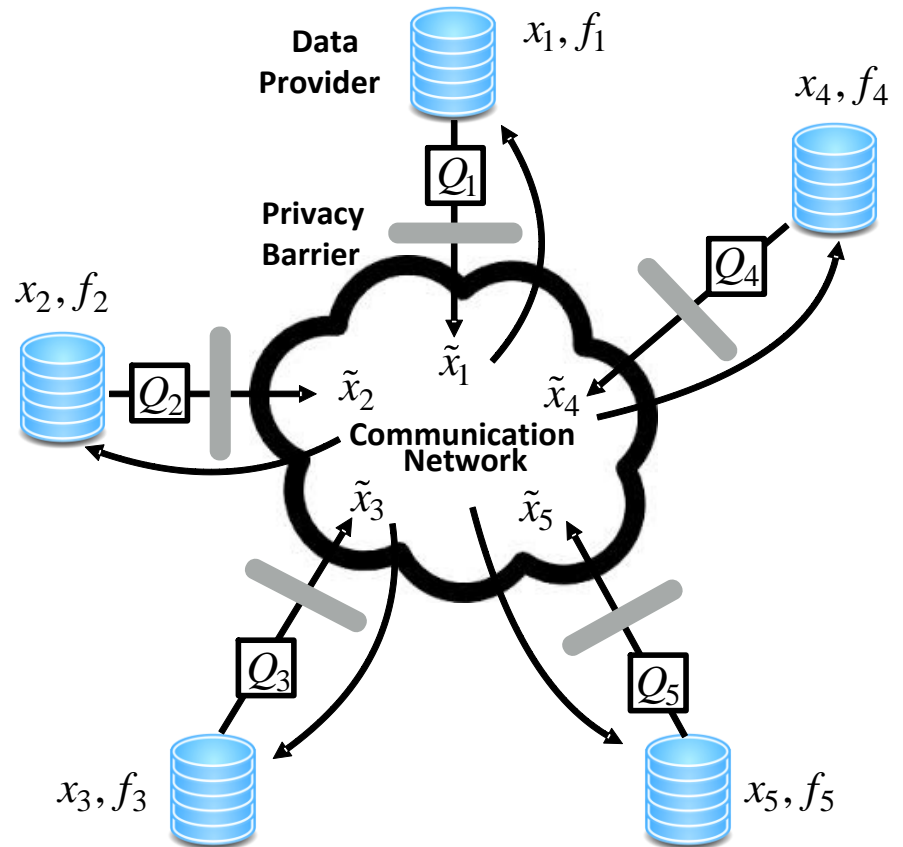
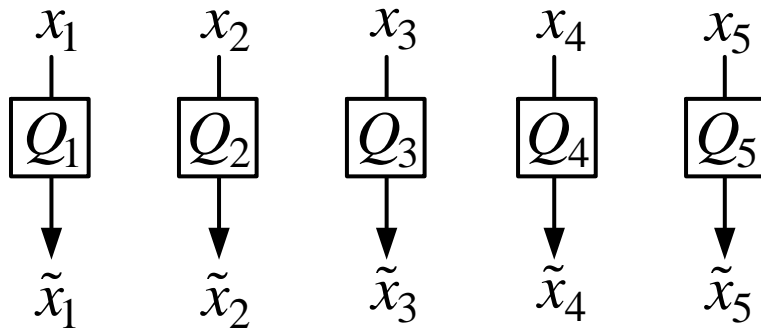
# Interactive mechanisms



$\tau = \text{communication transcript}$

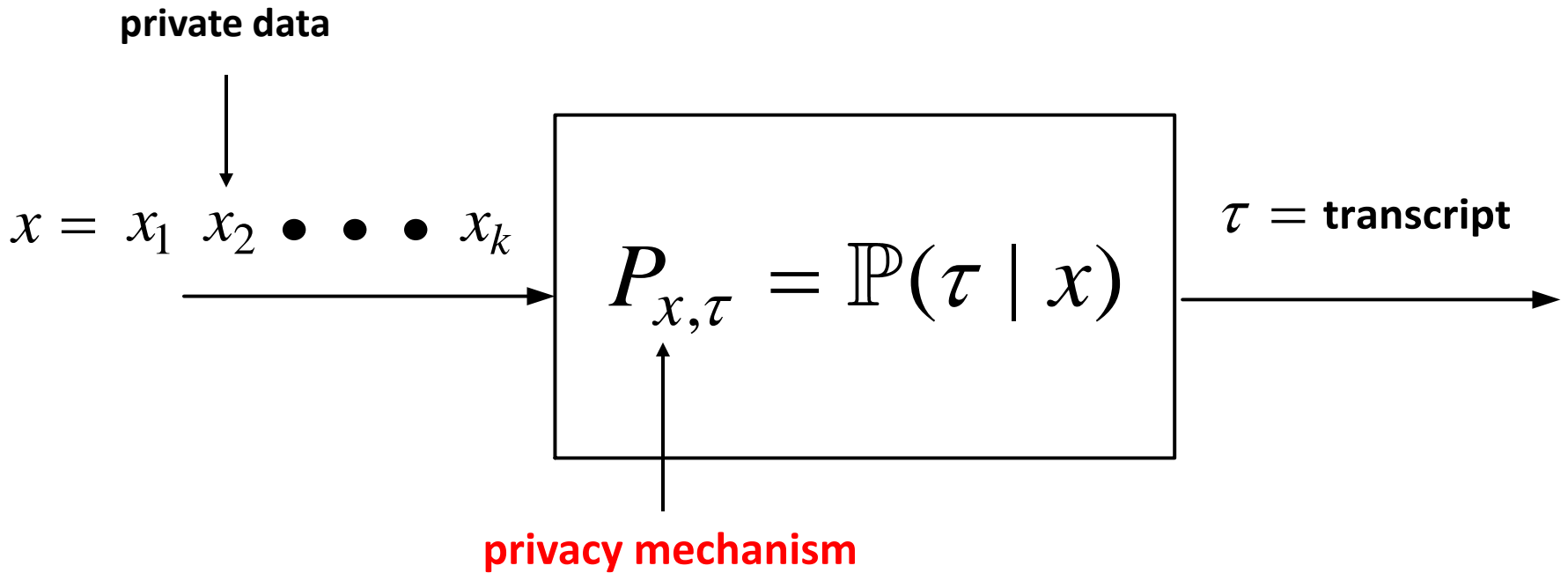


# Non-interactive mechanisms

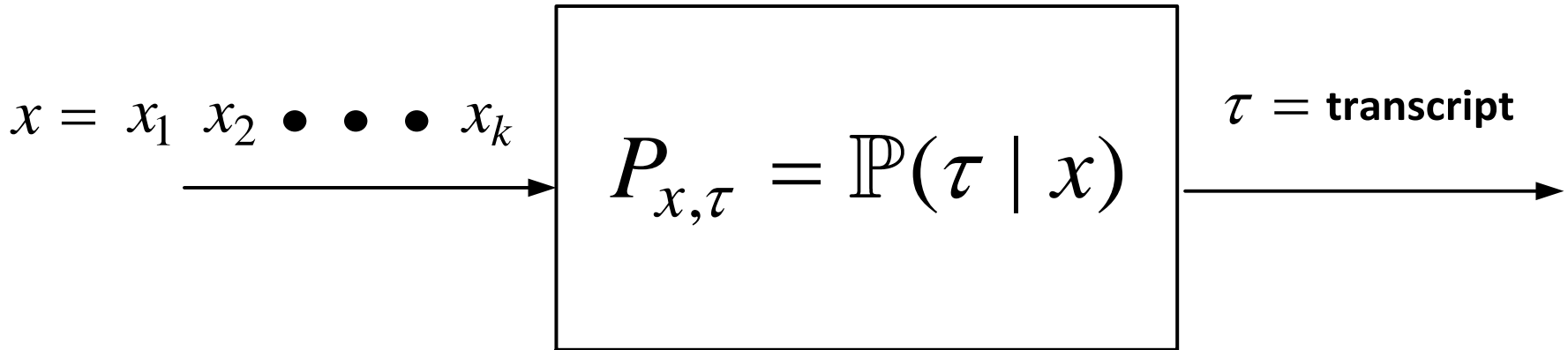




# General representation



# Multi-party differential privacy

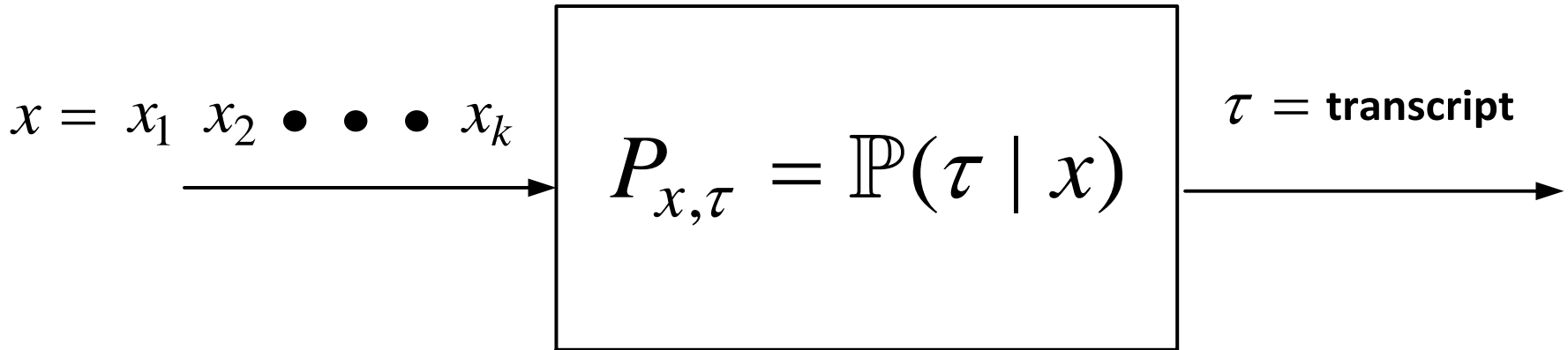


$$e^{-\varepsilon_i} \leq \frac{\mathbb{P}(\tau \mid x_i = 0, x_{-i})}{\mathbb{P}(\tau \mid x_i = 1, x_{-i})} \leq e^{\varepsilon_i}$$

$$x_{-i} = (x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_k)$$

- bounded likelihood even when all parties but one collude

# Multi-party differential privacy



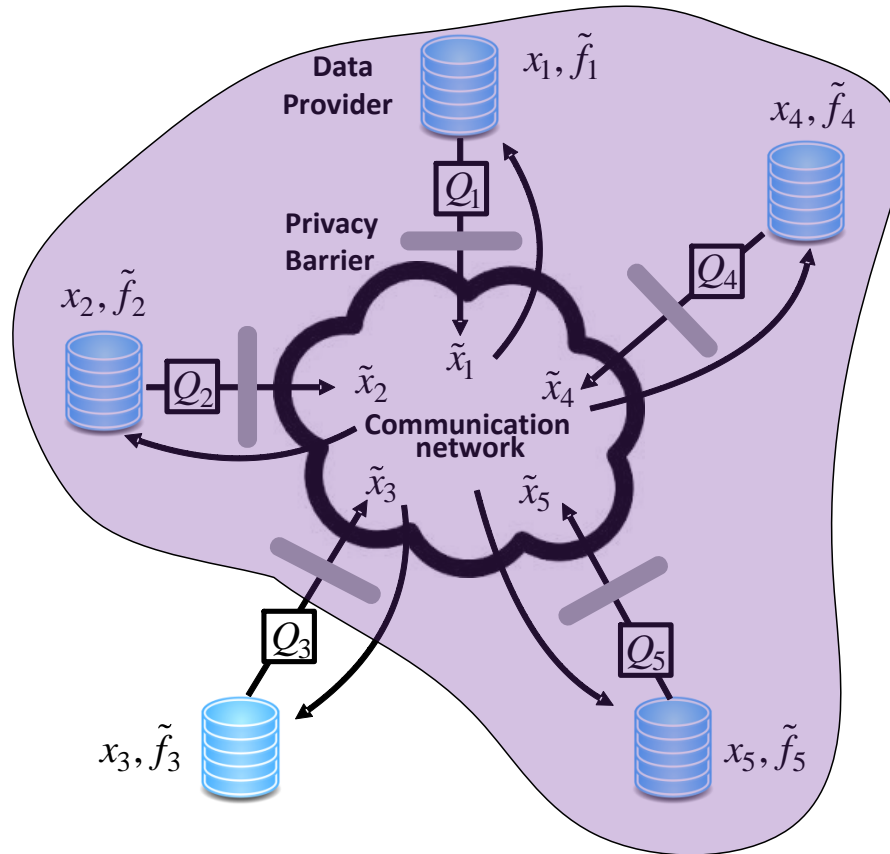
$$e^{-\varepsilon_i} \leq \frac{\mathbb{P}(\tau \mid x_i = 0, x_{-i})}{\mathbb{P}(\tau \mid x_i = 1, x_{-i})} \leq e^{\varepsilon_i}$$

$\varepsilon_i$  controls the level of privacy

large  $\varepsilon_i$ , low privacy

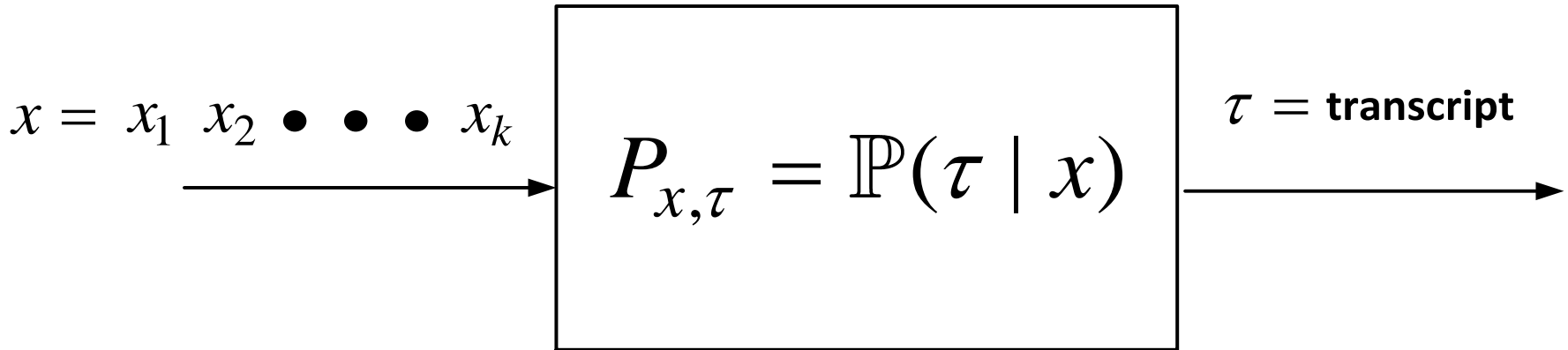
small  $\varepsilon_i$ , high privacy

# Can't say much even if...



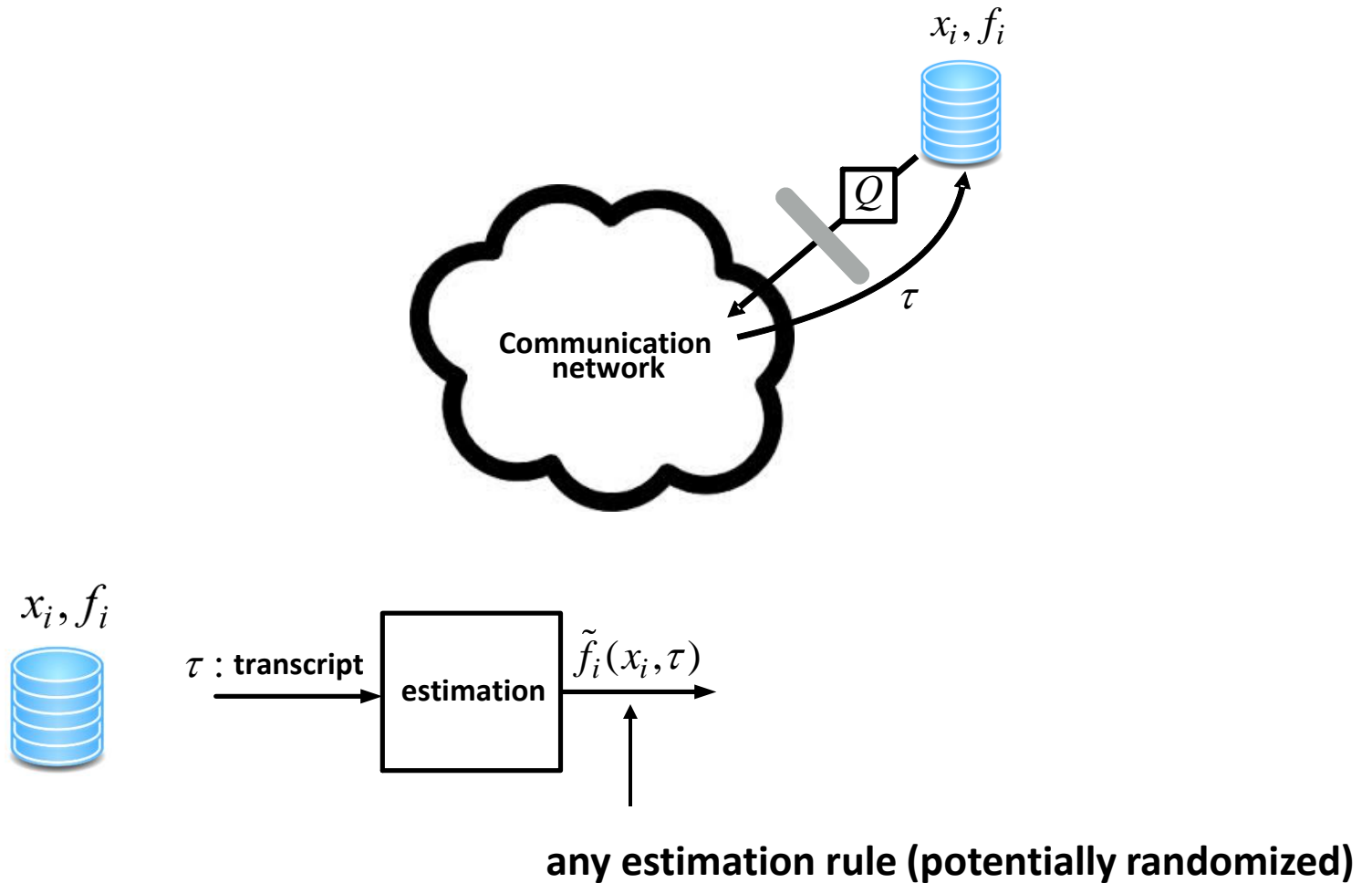
- all parties but one collude to figure out a party's data

# Approximate differential privacy

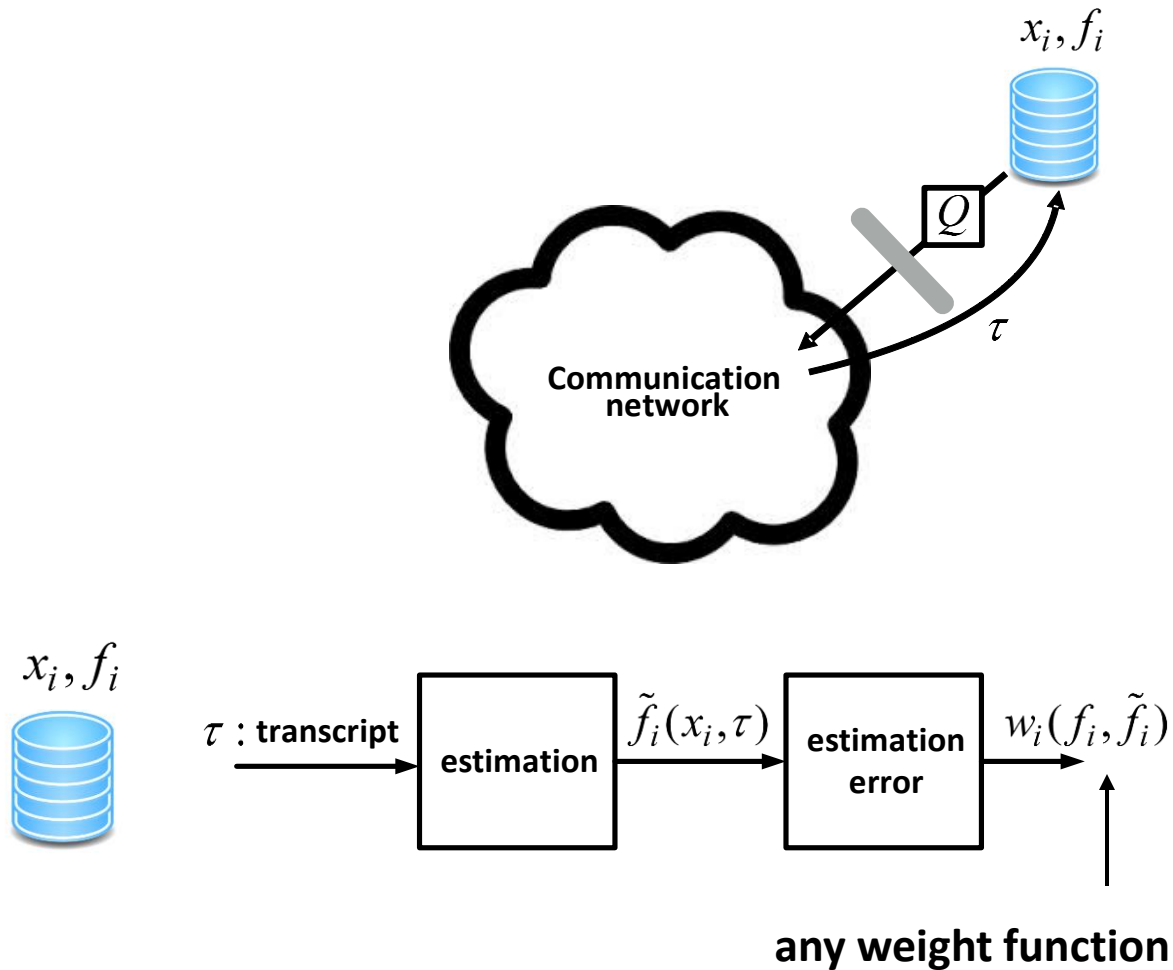


  $\delta_i$   
provides some slack

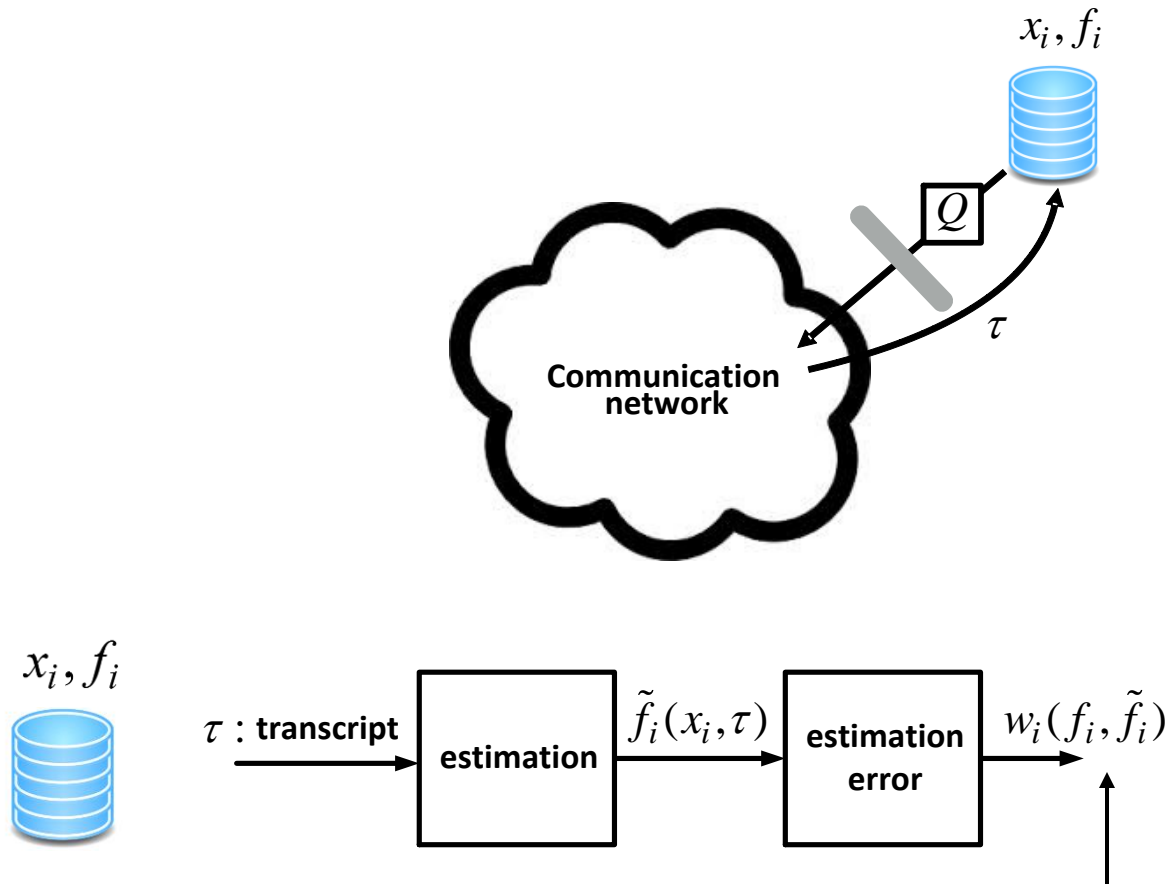
# Function estimation



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# Function estimation



any weight function

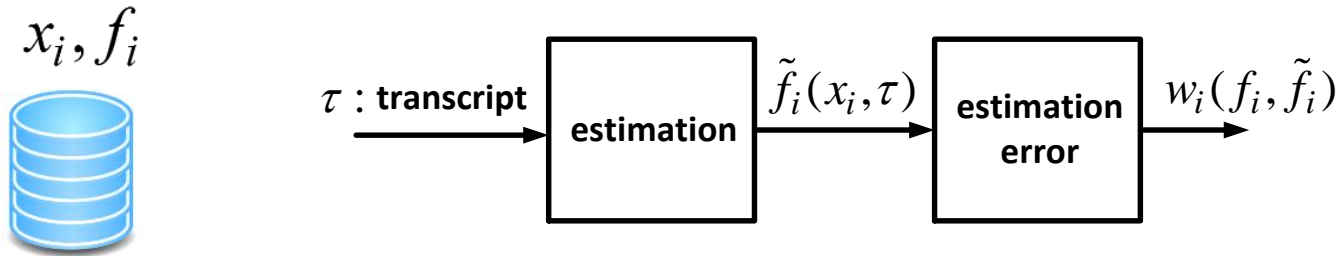
examples:

$$w(f, \tilde{f}) = 1_{(f=\tilde{f})}$$

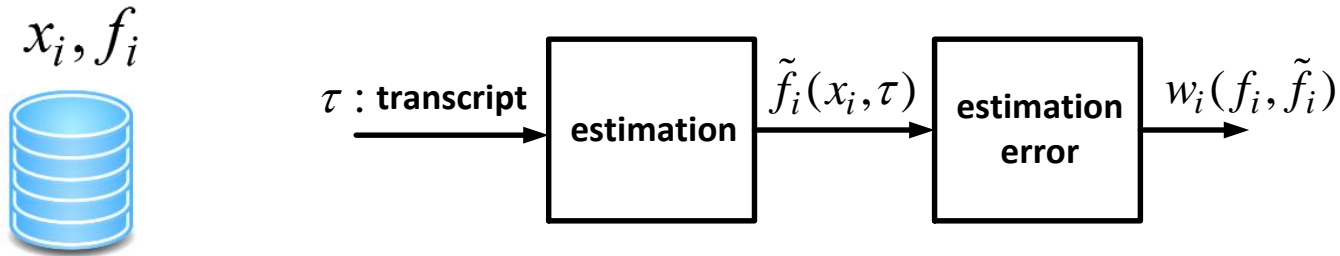
$$w(f, \tilde{f}) = |f - \tilde{f}|$$



# Utility: average accuracy



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$$\text{ACC}_{\text{ave}} \equiv \frac{1}{2^k} \sum_{x \in \{0,1\}^k} \mathbb{E}_{\hat{f}_i, P_{x,\tau}} [w_i(f_i(x), \tilde{f}_i(\tau, x_i))]$$

average over all possible inputs

# Privacy-utility tradeoff

$$\underset{P, \tilde{f}_i}{\text{maximize}} \quad \text{ACC}_{\text{ave}}(P, w_i, f_i, \tilde{f}_i),$$

subject to  $P$  and  $\tilde{f}_i$  are row-stochastic matrices

$$P_{(x_i, x_{-i}), \tau} \leq e^{\varepsilon_i} P_{(x'_i, x_{-i}), \tau} + \delta_i \quad \forall i, x_i, x'_i, x_{-i}, \tau$$

# Privacy-utility tradeoff

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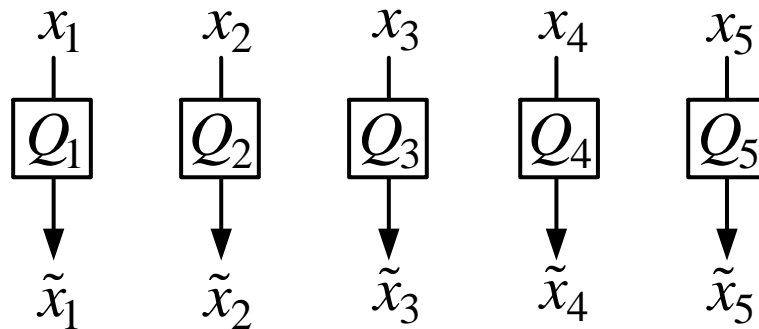
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$$P_{(x_i, x_{-i}), \tau} \leq e^{\varepsilon_i} P_{(x'_i, x_{-i}), \tau} + \delta_i \quad \forall i, x_i, x'_i, x_{-i}, \tau$$

- **heterogeneous privacy levels** across users
- each party possesses **a single bit**
- the **functions can vary** from one party to the other
- the **weight functions can vary** from one party to the other
- **interactive & non-interactive** mechanisms

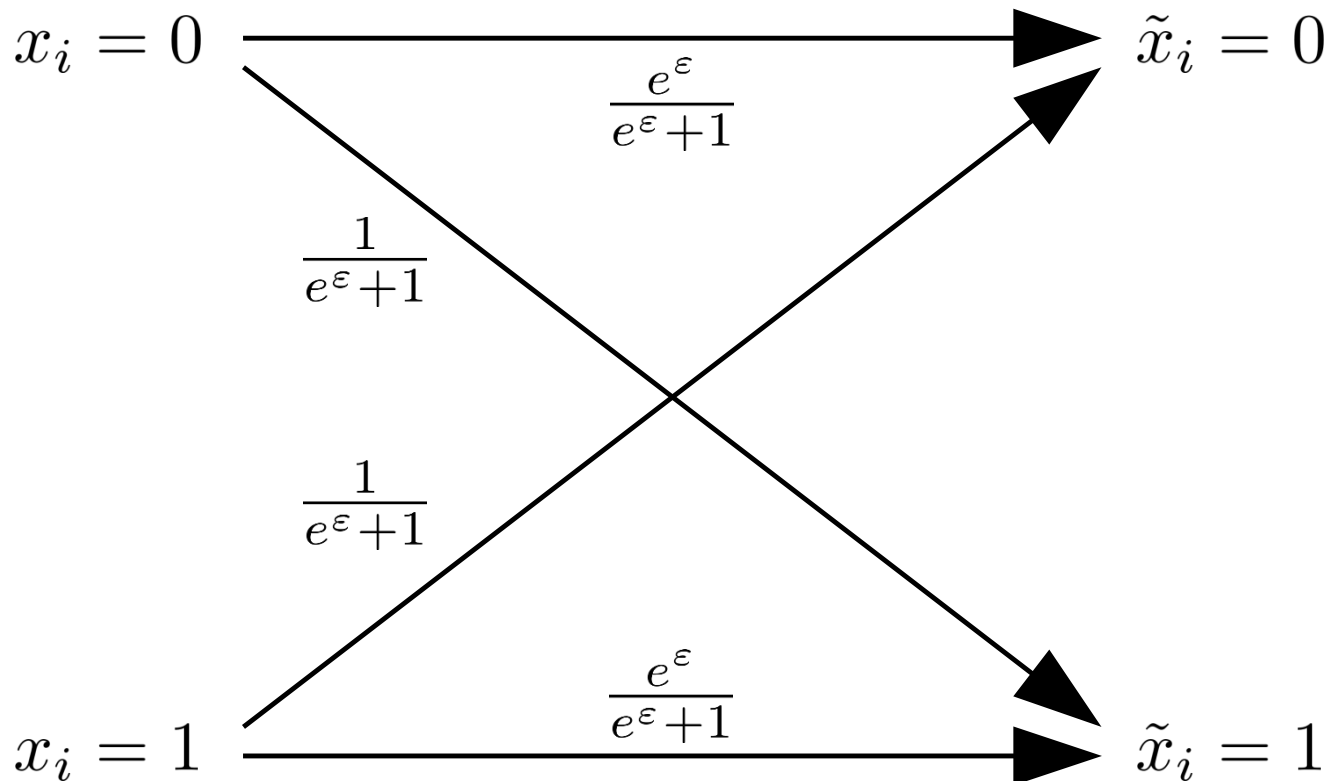
# Main result: differential privacy

**Non-interactive mechanisms are optimal**



# Main result: differential privacy

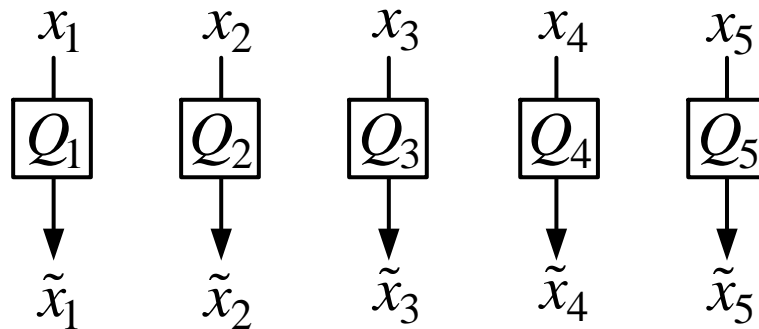
The randomized response is optimal!



# Approximate differential privacy?

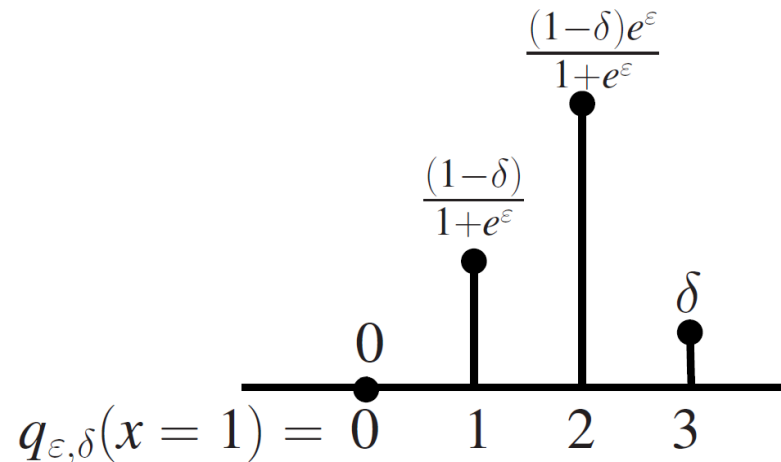
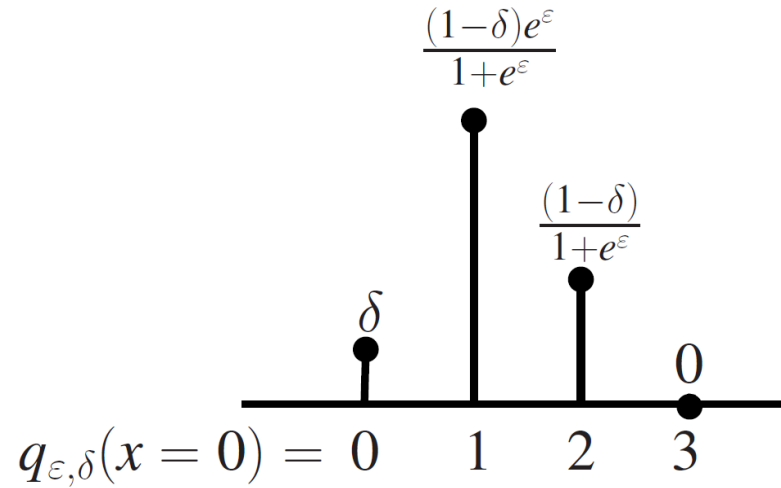
# Approximate differential privacy?

**Non-interactive mechanisms are optimal**

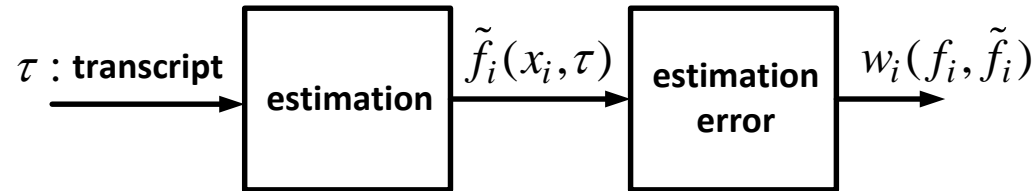




# Approximate differential privacy



# Optimal estimation rule

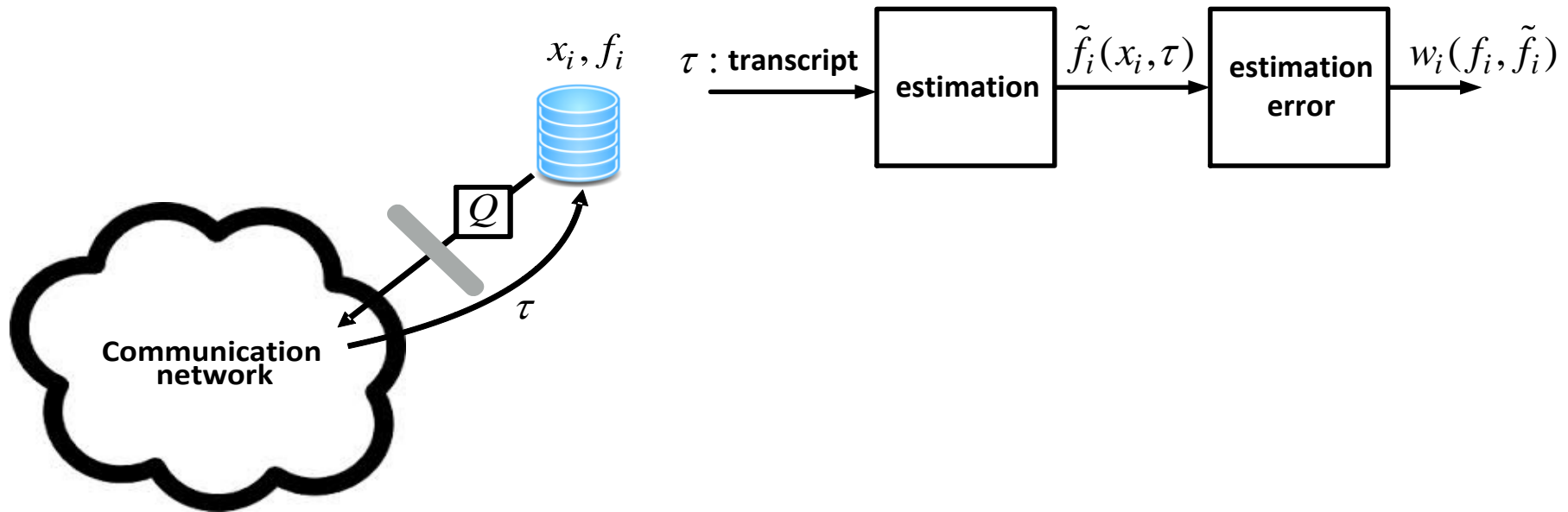


$$\text{ACC}_{\text{ave}} \equiv \underbrace{\frac{1}{2^k} \sum_{x \in \{0,1\}^k} \mathbb{E}_{\hat{f}_i, P_{x,\tau}} [w_i(f_i(x), \tilde{f}_i(\tau, x_i))]}_{\text{average over all possible inputs}}$$

average over all possible inputs

$$\tilde{f}_{i,\text{opt}}(\tau, x_i) = \arg \max_y \sum_{x_{-i} \in \{0,1\}^{k-1}} P_{x,\tau} w_i(f_i(x), y)$$

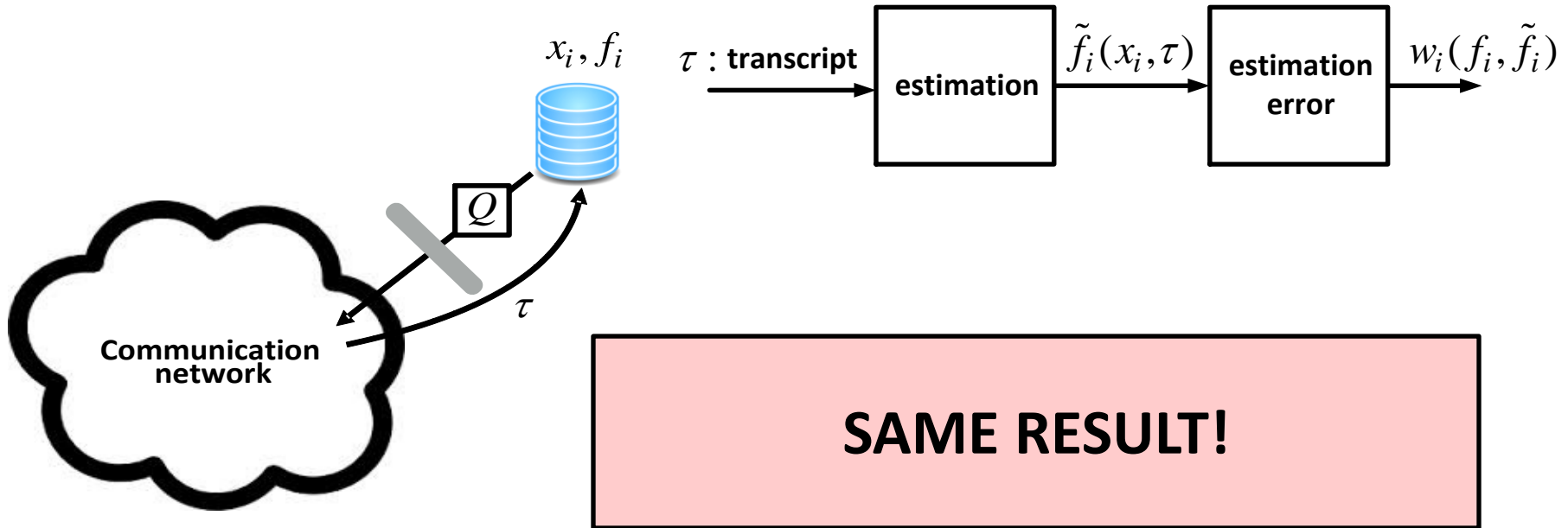
# Worst case accuracy?



$$\text{ACC}_{\text{wc}} \equiv \min_{x \in \{0,1\}^k} \mathbb{E}_{\hat{f}_i, P_{x,\tau}} [w_i(f_i(x), \hat{f}_i(\tau, x_i))]$$

worst case over all possible inputs

# Worst case accuracy

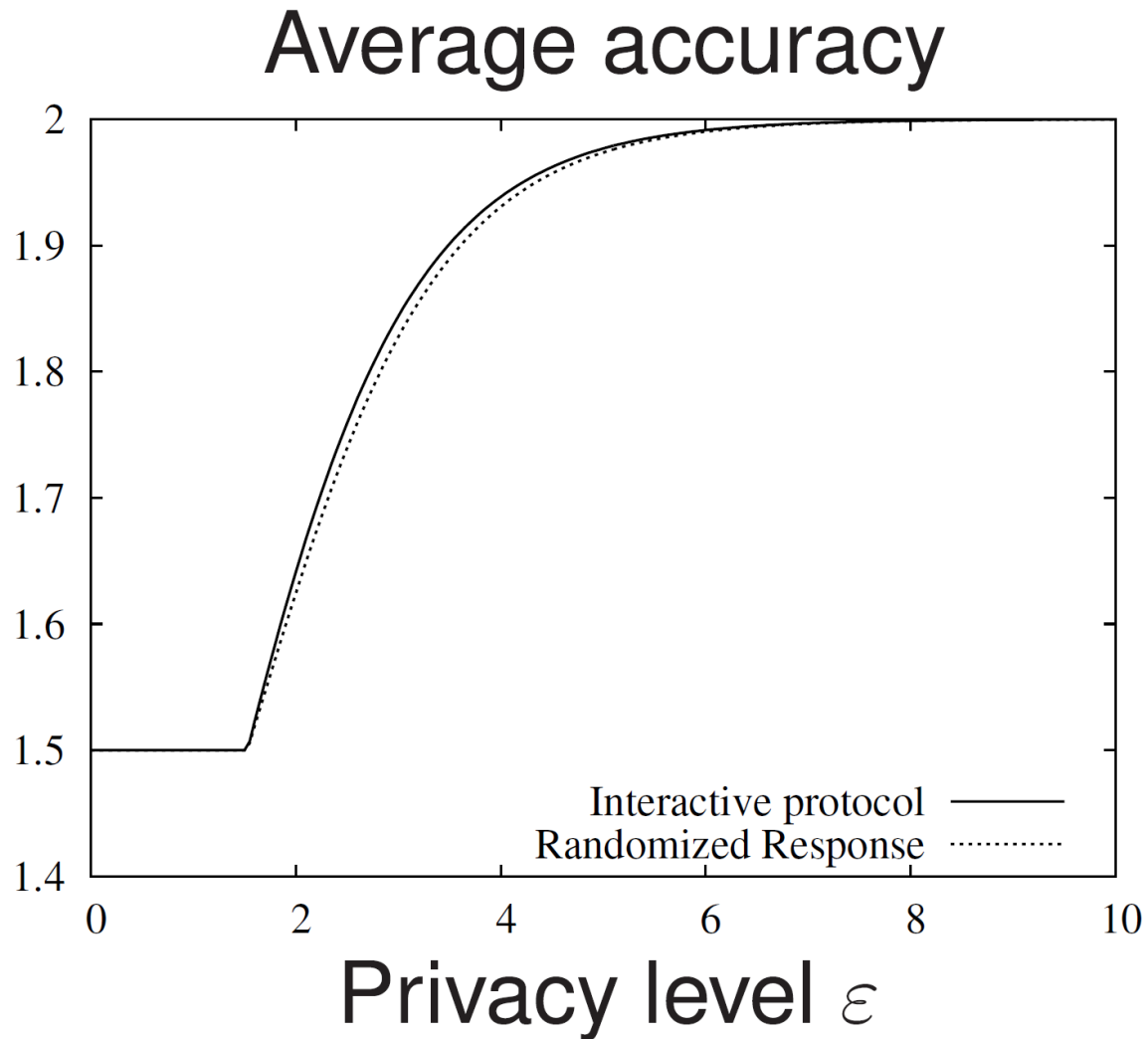


$$\text{ACC}_{\text{wc}} \equiv \min_{x \in \{0,1\}^k} \mathbb{E}_{\hat{f}_i, P_{x,\tau}} [w_i(f_i(x), \tilde{f}_i(\tau, x_i))]$$

worst case over all possible inputs

# Non-binary data?

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# Acknowledgments



**Pramod Viswanath**



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