

# Code Transformation for Distributed Python ML Code

# 1 Identifying trainig loop

## 1.1 Restrictions

1. Training loop must be defined in only one file.
2. Training loop type is either `distributed gradient tape` or `distributed optimizer`.
3. Each model must have only one type of training loop.
4. Function must not be assigned to the variables.
5. Training loop must not be defined conditionally.

## 1.2 Rules

### 1.2.1 Summary

|           |       |  |                   |
|-----------|-------|--|-------------------|
| $Summary$ | $::=$ | $ModuleSummary\ s\ (SumMap)^*\ tl$     | $(MODULESUMMARY)$ |
|           |       | $FuncSummary\ s\ (SumMap)^*\ tl$       | $(FUNC SUMMARY)$  |
|           |       | $ClassSummary\ s\ arg\ (SumMap)^*\ tl$ | $(CLASS SUMMARY)$ |
| $SumMap$  | $::=$ | $(s\ Summary)^*$                       | $(SUMMAP)$        |
| $arg$     | $::=$ | $???$                                  | $(ARG)$           |
| $tl$      | $::=$ | $GradTape \mid Optimizer \mid NoTl$    | $(TRAINING LOOP)$ |

### 1.2.2 Training Loop

$$\boxed{summary_M : Module \rightarrow ModuleSummary}$$

$$summary_M \llbracket s\ stmt^*\ type\_ignore \rrbracket =$$

$$\text{LET } summary\_list, tl = summary_{\bar{S}} \llbracket stmt^* \rrbracket (\sigma)(\gamma) \text{ IN}$$

$$ModuleSummary\ s\ summary\_list\ tl$$

$$\boxed{summary_{\bar{S}} : Stmt\ list \rightarrow \Sigma \rightarrow SumMap \rightarrow (SumMap \times tl)}$$

$$summary_S \llbracket stmt_1\ stmt_2 \dots stmt_n \rrbracket (\sigma)(\gamma) =$$

$$\text{LET } summary_1, tl_1 = trans_S \llbracket stmt_1 \rrbracket (\sigma) \text{ IN}$$

$$\text{LET } summary_2, tl_2 = trans_S \llbracket stmt_2 \rrbracket (\sigma) \text{ IN}$$

$$\dots$$

$$\text{LET } summary_n, tl_n = trans_S \llbracket stmt_n \rrbracket (\sigma) \text{ IN}$$

$$\text{LET } \gamma = summary_1\ summary_2 \dots summary_n \text{ IN}$$

$$\gamma, ???$$

$$\boxed{summary_S : Stmt \rightarrow \Sigma \rightarrow SumMap \rightarrow (Summary \times \Sigma \times tl)}$$

$$summary_S \llbracket (@expr_1)^*\ \text{def } id\ (args)\ (-> expr_2)^? : (\#type:s)^? stmt^* \rrbracket (\sigma)(\gamma) =$$

$$\text{LET } \gamma', tl = summary_{\bar{S}} \llbracket stmt^* \rrbracket \text{ IN}$$

$$FuncSummary\ id.name\ \gamma'\ tl, ???, ???$$

$$summary_S \llbracket (@expr_1)^*\ \text{async def } id\ (args)\ (-> expr_2)^? : (\#type:s)^? stmt^* \rrbracket (\sigma)(\gamma) =$$

$$\text{LET } \gamma', tl = summary_{\bar{S}} \llbracket stmt^* \rrbracket \text{ IN}$$

$$FuncSummary\ id.name\ \gamma'\ tl, ???, ???$$

$$summary_S \llbracket (@expr_1)^*\ \text{class } id\ (expr_2^*\ keyword^*) : stmt^* \rrbracket (\sigma)(\gamma) =$$

$$\text{LET } arg = summary_{\bar{A}} \llbracket expr_2^*\ keyword^* \rrbracket (\sigma)(\gamma) \text{ IN}$$

$$\text{LET } \gamma', tl = summary_{\bar{S}} \llbracket stmt^* \rrbracket \text{ IN}$$

$$ClassSummary\ arg\ id.name\ \gamma'\ tl, ???, ???$$

$$\boxed{summary_E : Expr \rightarrow \Sigma \rightarrow SumMap \rightarrow tl}$$

## 2 Python Abstract Syntax

|               |  |                |
|---------------|--|----------------|
| <i>module</i> | <code>::= stmt* type_ignore</code>   | (MODULEDEF)    |
| <i>stmt</i>   | <code>::= (@expr)* def id (args) (-&gt; expr)? : (#type:s)? stmt*</code>   | (FUNDEF)       |
|               | <code>(@expr)* async def id (args) (-&gt; expr)? : (#type:s)? stmt*</code> | (ASYNCFUNDEF)  |
|               | <code>(@expr)* class id (expr* keyword*) : stmt*</code>                    | (CLASSDEF)     |
|               | <code>return expr?</code>  | (RETURN)       |
|               | <code>delete expr*</code>  | (DELETE)       |
|               | <code>expr* = expr (#type:s)?</code>                                       | (ASSIGN)       |
|               | <code>expr binop = expr</code>   | (AUGASSIGN)    |
|               | <code>expr : expr (= expr)?</code>   | (ANNASSIGN)    |
|               | <code>(#type:s)? for expr in expr : stmt* (else : stmt*)?</code>           | (FORLOOP)      |
|               | <code>(#type:s)? async for expr in expr : stmt* (else : stmt*)?</code>     | (ASYNCFORLOOP) |
|               | <code>while (expr) : stmt* (else : stmt*)?</code>                          | (WHILELOOP)    |
|               | <code>if (expr) : stmt* (else : stmt*)?</code>                             | (IF)           |
|               | <code>(#type:s)? with with_item* : stmt*</code>                            | (WITH)         |
|               | <code>(#type:s)? async with with_item* : stmt*</code>                      | (ASYNCSWITH)   |
|               | <code>match expr : match_case*</code>                                      | (MATCH)        |
|               | <code>raise expr? (from expr)?</code>                                      | (RAISE)        |
|               | <code>try : stmt* exc_handler* (else : stmt*)? (finally : stmt*)?</code>   | (TRY)          |
|               | <code>assert expr expr?</code>   | (ASSERT)       |
|               | <code>import alias*</code>   | (IMPORT)       |
|               | <code>from i id? import alias*</code>                                      | (IMPORTFROM)   |
|               | <code>global id*</code>  | (GLOBAL)       |
|               | <code>nonlocal id*</code>  | (NONLOCAL)     |
|               | <code>expr</code>  | (EXPRSTMT)     |
|               | <code>pass</code>  | (PASS)         |
|               | <code>break</code>   | (BREAK)        |
|               | <code>continue</code>  | (CONTINUE)     |

|                      |                  |  |                  |
|----------------------|------------------|--|------------------|
| <i>expr</i>          | <code>::=</code> | <code>expr boolop expr</code>  | (BOOLOP)         |
|                      |                  | <code>expr := expr</code>  | (NAMEDEXPR)      |
|                      |                  | <code>expr binop expr</code>   | (BINARYOP)       |
|                      |                  | <code>unop expr</code>   | (UNARYOP)        |
|                      |                  | <code>lambda args : expr</code>  | (LAMBDA)         |
|                      |                  | <code>expr if expr else expr</code>  | (IFEXPR)         |
|                      |                  | <code>{(expr : expr)* (**expr)*}</code>  | (DICTIONARY)     |
|                      |                  | <code>{expr*}</code>   | (SET)            |
|                      |                  | <code>[expr*]</code>   | (LIST)           |
|                      |                  | <code>(expr*)</code>   | (TUPLE)          |
|                      |                  | <code>{expr:expr comprehension*}</code>  | (DICTIONARYCOMP) |
|                      |                  | <code>{expr comprehension*}</code>   | (SETCOMP)        |
|                      |                  | <code>[expr comprehension*]</code>   | (LISTCOMP)       |
|                      |                  | <code>(expr comprehension*)</code>   | (GENERATORCOMP)  |
|                      |                  | <code>await expr</code>  | (AWAIT)          |
|                      |                  | <code>yield expr?</code>   | (YIELD)          |
|                      |                  | <code>yield from expr</code>   | (YIELDFROM)      |
|                      |                  | <code>expr (compop expr)*</code>   | (COMPOP)         |
|                      |                  | <code>expr (expr* keyword*)</code>   | (CALL)           |
|                      |                  | <code>{expr (!i)? (:expr)?}</code>   | (FORMATTEDVALUE) |
|                      |                  | <code>expr*</code>   | (JOINEDSTR)      |
|                      |                  | <code>constant</code>  | (CONSTANT)       |
|                      |                  | <code>expr.id</code>   | (ATTRIBUTE)      |
|                      |                  | <code>expr[expr]</code>  | (SUBSCRIPT)      |
|                      |                  | <code>*expr</code>   | (STARRED)        |
|                      |                  | <code>**expr</code>  | (DOUBLESTARRED)  |
|                      |                  | <code>id</code>  | (NAME)           |
|                      |                  | <code>expr? (:expr)? (:expr)?</code>   | (SLICE)          |
| <i>boolop</i>        | <code>::=</code> | <code>and   or</code>  | (BOOLOPERATOR)   |
| <i>binop</i>         | <code>::=</code> | <code>+   -   *   @   /   **   &lt;&lt;   &gt;&gt;       ^   &amp;   //   %</code> | (BINOPERATOR)    |
| <i>unop</i>          | <code>::=</code> | <code>~   not   +   -</code>   | (UNOPERATOR)     |
| <i>compop</i>        | <code>::=</code> | <code>==   !=   &lt;   &lt;=   &gt;   &gt;=   is   is not   in   not in</code>     | (COMPOP)         |
| <i>comprehension</i> | <code>::=</code> | <code>for expr in expr (if expr)*</code>   | (COMPREHENSION)  |
|                      |                  | <code>async for expr in expr (if expr)*</code>                                     | (ASYNCCOMPREHEN) |
| <i>exc.handler</i>   | <code>::=</code> | <code>except expr? (as id)? : stmt*</code>   | (EXCHANDLER)     |
| <i>args</i>          | <code>::=</code> | <code>(arg (= expr)?)*, (arg (= expr)?)*, arg?, (arg (= expr)?)*, arg?</code>      | (ARGUMENTS)      |
| <i>arg</i>           | <code>::=</code> | <code>id expr? s?</code>   | (ARGUMENT)       |
| <i>keyword</i>       | <code>::=</code> | <code>id? = expr</code>  | (KEYWORD)        |
| <i>alias</i>         | <code>::=</code> | <code>id (.id)* (as id)?</code>  | (ALIAS)          |
| <i>with.item</i>     | <code>::=</code> | <code>expr (as expr)?</code>   | (WITHITEM)       |

|                    |   |                  |
|--------------------|---|------------------|
| <i>match_case</i>  | <code>::= case pattern (if expr)? : stmt*</code>                | (MATCHCASE)      |
| <i>pattern</i>     | <code>::= expr</code>   | (MATCHVALUE)     |
|                    | <i>constant</i>   | (MATCHSINGLETON) |
|                    | [ <i>pattern</i> *]   | (MATCHSEQUENCE)  |
|                    | <i>*(id)?</i>   | (MATCHSTAR)      |
|                    | {( <i>expr</i> : <i>pattern</i> )* <i>id</i> ?}                 | (MATCHMAPPING)   |
|                    | <i>expr</i> ( <i>pattern</i> * ( <i>id</i> = <i>pattern</i> )*) | (MATCHCLASS)     |
|                    | ( <i>pattern</i> as)? <i>id</i>                                 | (MATCHAS)        |
|                    | <i>pattern</i>   <i>pattern</i>                                 | (MATCHOR)        |
|                    | -   | (MATCHWILDCARD)  |
| <i>constant</i>    | <code>::= None</code>   | (NONELITERAL)    |
|                    | <i>i</i>  | (INTLITERAL)     |
|                    | <i>f</i>  | (FLOATLITERAL)   |
|                    | <i>c</i>  | (COMPLEXLITERAL) |
|                    | <i>s</i>  | (STRINGLITERAL)  |
|                    | <i>b</i>  | (BOOLEANLITERAL) |
|                    | ( <i>constant</i> *)  | (TUPLELITERAL)   |
|                    | ...   | (ELLIPSIS)       |
| <i>type_ignore</i> | <code>::= i*</code>   | (TYPEIGNORE)     |
| <i>id</i>          | $\in Id$  |                  |
| <i>s</i>           | $\in Str$   |                  |
| <i>b</i>           | $\in \{\text{True}, \text{False}\}$                             |                  |
| <i>i</i>           | $\in \mathbb{Z}$  |                  |
| <i>f</i>           | $\in \mathbb{R}$  |                  |
| <i>c</i>           | $\in \mathbb{C}$  |                  |

Note: there may be more constant terms such as an immutable container containing only constant elements. Please update the *constant* rule if you know such terms.

## 3 Transformation for TF2 Python Code

### 3.1 Restrictions

1. All import statements must be placed at the top of a module.
2. The `tensorflow` module must be assigned to a variable only using an import statement.
3. Members of the `tensorflow` module must not be aliased.
4. `print` function call expressions must not introduce side-effects.
5. A dataset and an optimizer object must only be created once via an assignment statement with a function call expression and must not be aliased.
6. Variables storing a dataset or a optimizer object must not be reassigned to store another value that is not a dataset or an optimizer object.
7. A dataset and an optimizer object must not be created conditionally.
8. `optimizer.apply_gradients` function call expressions must be expression statements or a direct RHS of an assignment statement.
9. Global variables storing an `optimizer` object and referred in functions must be defined before the functions' definitions and must not be changed after their initializations.
10. A checkpoint object must only be created once via an assignment statement with a function call expression and must not be aliased.

### 3.2 Rules

#### 3.2.1 Types and Auxiliary Functions

|            |                      |                                  |
|------------|----------------------|----------------------------------|
| $\tau ::=$ | <i>Module</i>        | (MODULES)                        |
|            | <i>Stmt</i>          | (STATEMENTS)                     |
|            | <i>Expr</i>          | (EXPRESSIONS)                    |
|            | <i>Comprehension</i> | (COMPREHENSIONS)                 |
|            | <i>ExcHandler</i>    | (EXCEPTION HANDLERS)             |
|            | <i>Alias</i>         | (ALIASES)                        |
|            | <i>WithItem</i>      | (WITHITEMS)                      |
|            | <i>Pattern</i>       | (PATTERNS)                       |
|            | <i>Id</i>            | (IDENTIFIERS)                    |
|            | <i>Str</i>           | (STRINGS)                        |
|            | $\tau$ list          | (LIST OF ELEMENTS TYPED $\tau$ ) |

$\sigma \in \Sigma = Str \xrightarrow{\text{fin}} Id$       ENVIRONMENT STORING MAPPINGS FROM STRINGS TO IDENTIFIERS

|           |   |   |   |
|-----------|---|---|---|
| <b>.1</b> | : | $\tau \times \tau \rightarrow \tau$   | GET THE FIRST ELEMENT OF THE GIVEN PAIR         |
| <b>@</b>  | : | $\tau \text{ list} \rightarrow \tau \text{ list} \rightarrow \tau \text{ list}$ | CONCATENATE TWO LISTS                           |
| <b>::</b> | : | $\tau \rightarrow \tau \text{ list} \rightarrow \tau \text{ list}$              | APPEND AN ELEMENT TO A LIST (RIGHT-ASSOCIATIVE) |

### 3.2.2 Transformation Rules

$$\boxed{trans_M : Module \rightarrow Module}$$

$$trans_M \llbracket stmt^* \text{ type\_ignore} \rrbracket = trans_{\bar{S}} \llbracket stmt^* \rrbracket(\sigma).1 \text{ type\_ignore}$$

$$\boxed{trans_{\bar{S}} : Stmt \text{ list} \rightarrow \Sigma \rightarrow (Stmt \text{ list} \times \Sigma)}$$

$$\begin{aligned} trans_{\bar{S}} \llbracket stmt_1 \ stmt_2 \dots stmt_n \rrbracket(\sigma) = & \text{LET } stmt_1^{*'}, \sigma_1 = trans_S \llbracket stmt_1 \rrbracket(\sigma) \text{ IN} \\ & \text{LET } stmt_2^{*'}, \sigma_2 = trans_S \llbracket stmt_2 \rrbracket(\sigma_1) \text{ IN} \\ & \dots \\ & \text{LET } stmt_n^{*'}, \sigma_n = trans_S \llbracket stmt_n \rrbracket(\sigma_{n-1}) \text{ IN} \\ & (stmt_1^{*'} @ stmt_2^{*'} @ \dots @ stmt_n^{*'}, \sigma_n) \end{aligned}$$

$$\boxed{trans_S : Stmt \rightarrow \Sigma \rightarrow (Stmt \text{ list} \times \Sigma)}$$

$$\begin{aligned} trans_S \llbracket (@expr_1)^* \text{ def } id \ (args) \ (-> \ expr_2)^? : (\#type:s)? \ stmt^* \rrbracket(\sigma) = \\ ([(@expr_1)^* \text{ def } id \ (args) \ (-> \ expr_2)^? : (\#type:s)? \ trans_{\bar{S}} \llbracket stmt^* \rrbracket(\sigma).1], \sigma) \end{aligned}$$

$$\begin{aligned} trans_S \llbracket (@expr_1)^* \text{ async def } id \ (args) \ (-> \ expr_2)^? : (\#type:s)? \ stmt^* \rrbracket(\sigma) = \\ ([(@expr_1)^* \text{ async def } id \ (args) \ (-> \ expr_2)^? : (\#type:s)? \ trans_{\bar{S}} \llbracket stmt^* \rrbracket(\sigma).1], \sigma) \end{aligned}$$

$$\begin{aligned} trans_S \llbracket (@expr_1)^* \text{ class } id \ (expr_2^* \ keyword^*) : stmt^* \rrbracket(\sigma) = \\ ([(@expr_1)^* \text{ class } id \ (expr_2^* \ keyword^*) : trans_{\bar{S}} \llbracket stmt^* \rrbracket(\sigma).1], \sigma) = \end{aligned}$$

$$trans_S \llbracket \text{return } expr^? \rrbracket(\sigma) = ([\text{return } (trans_E \llbracket expr \rrbracket(\sigma))^?], \sigma)$$

$$trans_S \llbracket \text{delete } expr^* \rrbracket(\sigma) = ([\text{delete } expr^*], \sigma)$$

A strict form of assignment statements

```

transS[[ idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)? ]](σ) =
  IF σ("tensor_flow") = idt AND expr1 = idt.data.Dataset.expr3 THEN
    ([idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)?],
     σ["dataset" ↦ idr])
  ELIF σ("tensor_flow") = idt AND expr1 = idt.train.Checkpoint THEN
    ([idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)?],
     σ["checkpoint" ↦ idr])
  ELIF σ("tensor_flow") = idt AND expr1 = idt.optimizers.Adam THEN
    IF idi = learning_rate WHEN 1 ≤ i ≤ k THEN
      ([idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... idi = expr2i * hvd.size()
        ... (idk = )? expr2k) (#type:s)?], σ["optimizer" ↦ idr])
    ELSE
      ([idr = expr1 (expr11 * hvd.size() ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)?],
       σ["optimizer" ↦ idr])
  ELIF σ("optimizer") = idt AND expr1 = idt.apply_gradients THEN
    IF idi = grads_and_vars WHEN 1 ≤ i ≤ k THEN
      LET idz = NewID() IN
        ([idz = expr2i,
          idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... idi = idz ... (idk = )? expr2k) (#type:s)?,
          global hvd.broadcast_done,
          if not hvd.broadcast_done: [ hvd.broadcast_variables([x[1] for x in idz], root_rank=0),
            hvd.broadcast_variables(idt.variables(), root_rank=0),
            hvd.broadcast_done = True ], σ)
        ELSE
          LET idz = NewID() IN
            ([idz = expr11,
              idr = expr1 (idz expr12 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)?,
              global hvd.broadcast_done,
              if not hvd.broadcast_done: [ hvd.broadcast_variables([x[1] for x in idz], root_rank=0),
                hvd.broadcast_variables(idt.variables(), root_rank=0),
                hvd.broadcast_done = True ], σ)
      ELIF σ("checkpoint") = idt AND expr1 = idt.save THEN
        ([if hvd.rank() == 0: [idr = expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) (#type:s)?],
          σ)
        ELIF σ("keras") = idk AND expr1 = idk.models.Sequential THEN
          ([idr = expr1 (expr11 ... expr1n (id2 = )? expr21 ... (idk = )? expr2k)], σ["model" ↦ idk])
        ELSE ([idr = transE[[ expr1 (expr11 ... expr1n (id2 = )? expr21 ... (idk = )? expr2k) ]](σ) (#type:s)?],
          σ)

```

**A general form of assignment statements except for the strict form**

$trans_S[[ expr_1^* = expr_2 \text{ (\#type:s)? } ]](\sigma) = ([expr_1^* = trans_E[[ expr_2 ]](\sigma) \text{ (\#type:s)?}], \sigma)$

$trans_S[[ expr_1 \text{ binop} = expr_2 ]](\sigma) = ([expr_1 \text{ binop} = trans_E[[ expr_2 ]](\sigma)], \sigma)$



$trans_S \llbracket expr_1 : expr_2 (= expr_3)? \rrbracket(\sigma) = ([expr_1 : expr_2 (= trans_E \llbracket expr_3 \rrbracket(\sigma))?], \sigma)$   
**IF**  $expr_1 = id_1$  **AND**  $\sigma(\text{"tensor\_flow"}) = id_2$   
**AND**  $expr_3 = id_2.data.Dataset.expr_4 (expr_5^* keyword^*)$  **THEN**  
 $([expr_1 : expr_2 = expr_3 (\#type:s)?], \sigma[\text{"dataset"} \mapsto id_1])$   
**ELSE**  $([expr_1 : expr_2 (= trans_E \llbracket expr_3 \rrbracket(\sigma))?], \sigma)$

$trans_S \llbracket (\#type:s)? \text{ for } expr_1 \text{ in } expr_2 : stmt_1^* (\text{else} : stmt_2^*)? \rrbracket(\sigma) =$   
 $([(\#type:s)? \text{ for } expr_1 \text{ in } trans_E \llbracket expr_2 \rrbracket(\sigma) :$   
 $trans_{\overline{S}} \llbracket stmt_1^* \rrbracket(\sigma).1 (\text{else} : trans_{\overline{S}} \llbracket stmt_2^* \rrbracket(\sigma).1)?], \sigma)$

$trans_S \llbracket (\#type:s)? \text{ async for } expr_1 \text{ in } expr_2 : stmt_1^* (\text{else} : stmt_2^*)? \rrbracket(\sigma) =$   
 $([(\#type:s)? \text{ async for } expr_1 \text{ in } trans_E \llbracket expr_2 \rrbracket(\sigma) :$   
 $trans_{\overline{S}} \llbracket stmt_1^* \rrbracket(\sigma).1 (\text{else} : trans_{\overline{S}} \llbracket stmt_2^* \rrbracket(\sigma).1)?], \sigma)$

$trans_S \llbracket \text{while } (expr) : stmt_1^* (\text{else} : stmt_2^*)? \rrbracket(\sigma) =$   
 $([\text{while } (trans_E \llbracket expr \rrbracket(\sigma)) : trans_{\overline{S}} \llbracket stmt_1^* \rrbracket(\sigma).1 (\text{else} : trans_{\overline{S}} \llbracket stmt_2^* \rrbracket(\sigma).1)?], \sigma)$

$trans_S \llbracket \text{if } (expr) : stmt^* (\text{else} : stmt^*)? \rrbracket(\sigma) =$   
 $([\text{if } (trans_E \llbracket expr \rrbracket(\sigma)) : trans_{\overline{S}} \llbracket stmt^* \rrbracket(\sigma).1 (\text{else} : trans_{\overline{S}} \llbracket stmt^* \rrbracket(\sigma).1)?], \sigma)$

$trans_S \llbracket (\#type:s)? \text{ with } with\_item^* : stmt^* \rrbracket(\sigma) =$   
**LET**  $with\_item^{*'}, \sigma_1 = trans_{\overline{W}} \llbracket with\_item^* \rrbracket(\sigma)$  **IN**  
**LET**  $stmt^{*'}, \sigma_2 = trans_{\overline{S}} \llbracket stmt^* \rrbracket(\sigma_1)$  **IN**  
**IF**  $\sigma_1 \setminus \sigma = [\text{"gradient\_tape"} \mapsto id]$  **THEN**  
 $([(\#type:s)? \text{ with } with\_item^{*'} : stmt^{*'},$   
 $id = \text{hvd.DistributedGradientTape}(id)], \sigma_2)$   
**ELSE**  $([(\#type:s)? \text{ with } with\_item^{*'} : stmt^{*'}], \sigma_2)$

$trans_S \llbracket (\#type:s)? \text{ async with } with\_item^* : stmt^* \rrbracket(\sigma) =$   
**LET**  $with\_item^{*'}, \sigma_1 = trans_{\overline{W}} \llbracket with\_item^* \rrbracket(\sigma)$  **IN**  
**LET**  $stmt^{*'}, \sigma_2 = trans_{\overline{S}} \llbracket stmt^* \rrbracket(\sigma_1)$  **IN**  
**IF**  $\sigma_1 \setminus \sigma = [\text{"gradient\_tape"} \mapsto id]$  **THEN**  
 $([(\#type:s)? \text{ async with } with\_item^{*'} : stmt^{*'},$   
 $id = \text{hvd.DistributedGradientTape}(id)], \sigma_2)$   
**ELSE**  $([(\#type:s)? \text{ async with } with\_item^{*'} : stmt^{*'}], \sigma_2)$

$trans_S \llbracket \text{match } expr : match\_case^* \rrbracket(\sigma) =$   
 $([\text{match } trans_E \llbracket expr \rrbracket(\sigma) : (trans_C \llbracket match\_case \rrbracket(\sigma))^*], \sigma)$

$trans_S \llbracket \text{raise } expr_1? (\text{from } expr_2)? \rrbracket(\sigma) = ([\text{raise } expr_1? (\text{from } expr_2)?], \sigma)$

$trans_S \llbracket \text{try} : stmt_1^* exc\_handler^* (\text{else} : stmt_2^*)? (\text{finally} : stmt_3^*)? \rrbracket(\sigma) =$   
 $([\text{try} : trans_{\overline{S}} \llbracket stmt_1^* \rrbracket(\sigma).1 (trans_H \llbracket exc\_handler \rrbracket(\sigma))^*$   
 $(\text{else} : trans_{\overline{S}} \llbracket stmt_2^* \rrbracket(\sigma).1)? (\text{finally} : trans_{\overline{S}} \llbracket stmt_3^* \rrbracket(\sigma).1)?], \sigma)$

$trans_S \llbracket \text{assert } expr_1 \text{ } expr_2? \rrbracket(\sigma) = ([\text{assert } trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ } expr_2?], \sigma)$

```

transS[[ import alias* ]](σ) =
  LET σ1 = transA[[ alias* ]](σ) IN
  IF σ1 \ σ = [“tensorflow” ↦ id] THEN
    ([import alias*,
     import horovod.tensorflow as hvd,
     hvd.broadcast_done = False,
     hvd.init(),
     gpus = id.config.experimental.list_physical_devices(‘GPU’),
     for gpu in gpus: id.config.experimental.set_memory_growth(gpu, True),
     if gpus: id.config.experimental.set_visible_devices(gpus[hvd.local_rank()], ‘GPU’)], σ1)
  ELSE ([import alias*], σ1)

transS[[ from i id? import alias* ]](σ) =
  LET σ1 = transA[[ alias* ]](σ) IN
  IF id = tensorflow AND σ1 \ σ = [“keras” ↦ id] THEN
    ([from i id? import alias*,
     import horovod.tensorflow.keras as hvd,
     hvd.init(),
     gpus = id.config.experimental.list_physical_devices(‘GPU’),
     for gpu in gpus: id.config.experimental.set_memory_growth(gpu, True),
     if gpus: id.config.experimental.set_visible_devices(gpus[hvd.local_rank()], ‘GPU’)], σ1)
  ELSE ([from i id? import alias*], σ)

transS[[ global id* ]](σ) = ([global id*], σ)

transS[[ nonlocal id* ]](σ) = ([nonlocal id*], σ)

```

**A strict form of expr statements**

```

transS[[ expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) ]](σ) =
  IF σ("optimizer") = idt AND expr1 = idt.apply_gradients THEN
    IF idi = grads_and_vars WHEN 1 ≤ i ≤ k THEN
      LET idz = NewID() IN
        ([idz = expr2i,
          expr1 (expr11 ... expr1n (id1 = )? expr21 ... idi = idz ... (idk = )? expr2k),
          global hvd.broadcast_done,
          if not hvd.broadcast_done: [ hvd.broadcast_variables([x[1] for x in idz], root_rank=0),
            hvd.broadcast_variables(idt.variables(), root_rank=0),
            hvd.broadcast_done = True ]], σ)
    ELSE
      LET idz = NewID() IN
        ([idz = expr11,
          expr1 (idz expr12 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k),
          global hvd.broadcast_done,
          if not hvd.broadcast_done: [ hvd.broadcast_variables([x[1] for x in idz], root_rank=0),
            hvd.broadcast_variables(idt.variables(), root_rank=0),
            hvd.broadcast_done = True ]], σ)
    ELIF expr1 = print THEN
      ([if hvd.rank() == 0:,
        [expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k)]], σ)
    ELIF σ("checkpoint") = idt AND expr1 = idt.save THEN
      ([if hvd.rank() == 0: [expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k)]], σ)
    ELIF σ("model") = idt AND expr1 = idt.fit THEN
      IF idi = callbacks WHEN 1 ≤ i ≤ k THEN
        ([expr1 (expr11 ... expr1n (id1 = )? expr21 ...
          idi = expr2i + [hvd.callbacks.BroadcastGlobalVariablesCallback(0)]
          ... (idk = )? expr2k)], σ)
      ELSE
        ([expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k
          callbacks = [hvd.callbacks.BroadcastGlobalVariablesCallback(0)]], σ)
    ELSE
      [transE[[ expr1 (expr11 ... expr1n (id1 = )? expr21 ... (idk = )? expr2k) ]](σ)], σ)

```

**A general form of expr statements except for the strict form**

```
transS[[ expr ]](σ) = ([transE[[ expr ]](σ)], σ)
```

```
transS[[ pass ]](σ) = ([pass], σ)
```

```
transS[[ break ]](σ) = ([break], σ)
```

```
transS[[ continue ]](σ) = ([continue], σ)
```

$$\boxed{trans_E : Expr \rightarrow \Sigma \rightarrow Expr}$$

$$trans_E \llbracket expr_1 \text{ boolop } expr_2 \rrbracket(\sigma) = trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ boolop } trans_E \llbracket expr_2 \rrbracket(\sigma)$$

$$trans_E \llbracket expr_1 := expr_2 \rrbracket(\sigma) = expr_1 := trans_E \llbracket expr_2 \rrbracket(\sigma)$$

$$trans_E \llbracket expr_1 \text{ binop } expr_2 \rrbracket(\sigma) = trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ binop } trans_E \llbracket expr_2 \rrbracket(\sigma)$$

$$trans_E \llbracket unop \text{ expr } \rrbracket(\sigma) = unop \text{ } trans_E \llbracket expr \rrbracket(\sigma)$$

$$trans_E \llbracket \text{lambda } args : expr \rrbracket(\sigma) = \text{lambda } args : trans_E \llbracket expr \rrbracket(\sigma)$$

$$trans_E \llbracket expr_1 \text{ if } expr_2 \text{ else } expr_3 \rrbracket(\sigma) = \\ trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ if } trans_E \llbracket expr_2 \rrbracket(\sigma) \text{ else } trans_E \llbracket expr_3 \rrbracket(\sigma)$$

$$trans_E \llbracket expr_1 \text{ if } expr_2 \text{ else } expr_3 \rrbracket(\sigma) = \\ trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ if } trans_E \llbracket expr_2 \rrbracket(\sigma) \text{ else } trans_E \llbracket expr_3 \rrbracket(\sigma)$$

$$trans_E \llbracket \{(expr_1 : expr_2)^* (**expr_3)^*\} \rrbracket(\sigma) = \{(expr_1 : trans_E \llbracket expr_2 \rrbracket(\sigma))^* (**expr_3)^*\}$$

$$trans_E \llbracket \{expr^*\} \rrbracket(\sigma) = \{(trans_E \llbracket expr \rrbracket(\sigma))^*\}$$

$$trans_E \llbracket [expr^*] \rrbracket(\sigma) = [(trans_E \llbracket expr \rrbracket(\sigma))^*]$$

$$trans_E \llbracket (expr^*) \rrbracket(\sigma) = ((trans_E \llbracket expr \rrbracket(\sigma))^*)$$

$$trans_E \llbracket \{expr_1 : expr_2 \text{ comprehension}^*\} \rrbracket(\sigma) = \{expr_1 : trans_E \llbracket expr_2 \rrbracket(\sigma) (trans_O \llbracket comprehension \rrbracket(\sigma))^*\}$$

$$trans_E \llbracket \{expr \text{ comprehension}^*\} \rrbracket(\sigma) = \{trans_E \llbracket expr \rrbracket(\sigma) (trans_O \llbracket comprehension \rrbracket(\sigma))^*\}$$

$$trans_E \llbracket [expr \text{ comprehension}^*] \rrbracket(\sigma) = [(trans_E \llbracket expr \rrbracket(\sigma) (trans_O \llbracket comprehension \rrbracket(\sigma))^*)]$$

$$trans_E \llbracket (expr \text{ comprehension}^*) \rrbracket(\sigma) = (trans_E \llbracket expr \rrbracket(\sigma) (trans_O \llbracket comprehension \rrbracket(\sigma))^*)$$

$$trans_E \llbracket \text{await } expr \rrbracket(\sigma) = \text{await } trans_E \llbracket expr \rrbracket(\sigma)$$

$$trans_E \llbracket \text{yield } expr? \rrbracket(\sigma) = \text{yield } (trans_E \llbracket expr \rrbracket(\sigma))?$$

$$trans_E \llbracket \text{yield from } expr \rrbracket(\sigma) = \text{yield from } trans_E \llbracket expr \rrbracket(\sigma)$$

$$trans_E \llbracket expr_1 (compop expr_2)^* \rrbracket(\sigma) = trans_E \llbracket expr_1 \rrbracket(\sigma) (compop trans_E \llbracket expr_2 \rrbracket(\sigma))^*$$

$$\begin{aligned} trans_E \llbracket expr_1 (expr_{11} \dots expr_{1n} (id_1 = )? expr_{21} \dots (id_k = )? expr_{2k}) \rrbracket(\sigma) = \\ \text{IF } \sigma(\text{"dataset"}) = id_t \text{ AND } expr_1 = id_t.\text{take} \text{ THEN} \\ \quad \text{IF } id_i = \text{count} \text{ WHEN } 1 \leq i \leq k \text{ THEN} \\ \quad \quad expr_1 (expr_{11} \dots expr_{1n} (id_1 = )? expr_{21} \dots id_i = expr_{2i} // \text{hvd.size}() \dots (id_k = )? expr_{2k}) \\ \text{ELSE} \\ \quad expr_1 (expr_{11} // \text{hvd.size}() \dots expr_{1n} (id_1 = )? expr_{21} \dots (id_k = )? expr_{2k}) \\ \text{ELSE} \\ \quad trans_E \llbracket expr_1 \rrbracket(\sigma) (trans_E \llbracket expr_{11} \rrbracket(\sigma) \dots trans_E \llbracket expr_{1n} \rrbracket(\sigma) \\ \quad (id_1 = )? trans_E \llbracket expr_{21} \rrbracket(\sigma) \dots (id_k = )? trans_E \llbracket expr_{2k} \rrbracket(\sigma)) \end{aligned}$$

$$trans_E \llbracket \{expr_1 \quad (!i)? \quad (:expr_2)?\} \rrbracket(\sigma) = \{expr_1 \quad (!i)? \quad (:expr_2)?\}$$

$$trans_E \llbracket expr^* \rrbracket(\sigma) = expr^*$$

$$trans_E \llbracket constant \rrbracket(\sigma) = constant$$

$$trans_E \llbracket expr.id \rrbracket(\sigma) = (trans_E \llbracket expr \rrbracket(\sigma)).id$$

$$trans_E \llbracket expr_1[expr_2] \rrbracket(\sigma) = trans_E \llbracket expr_1 \rrbracket(\sigma)[trans_E \llbracket expr_2 \rrbracket(\sigma)]$$

$$trans_E \llbracket *expr \rrbracket(\sigma) = *expr$$

$$trans_E \llbracket **expr \rrbracket(\sigma) = **expr$$

$$trans_E \llbracket id \rrbracket(\sigma) = id$$

$$trans_E \llbracket expr_1? (:expr_2)? (:expr_3)? \rrbracket(\sigma) = trans_E \llbracket expr_1 \rrbracket(\sigma)? (:trans_E \llbracket expr_2 \rrbracket(\sigma))? (:trans_E \llbracket expr_3 \rrbracket(\sigma))?$$

$$\boxed{trans_O : Comprehension \rightarrow \Sigma \rightarrow Comprehension}$$

$$trans_O \llbracket \text{for } expr_1 \text{ in } expr_2 \text{ (if } expr_3)^* \rrbracket(\sigma) = \text{for } expr_1 \text{ in } trans_E \llbracket expr_2 \rrbracket(\sigma) \text{ (if } trans_E \llbracket expr_3 \rrbracket(\sigma))^*$$

$$\boxed{trans_H : ExcHandler \rightarrow \Sigma \rightarrow ExcHandler}$$

$$trans_H \llbracket \text{except } expr? \text{ (as } id)? : stmt^* \rrbracket(\sigma) = \text{except } expr? \text{ (as } id)? : trans_{\bar{S}} \llbracket stmt^* \rrbracket(\sigma).1$$

$$\boxed{trans_{\bar{A}} : Alias \text{ list} \rightarrow \Sigma \rightarrow \Sigma}$$

$$trans_{\bar{A}} \llbracket alias_1 \dots alias_{n-1} alias_n \rrbracket(\sigma) = trans_A \llbracket alias_n \rrbracket(trans_A \llbracket alias_{n-1} \rrbracket(\dots trans_A \llbracket alias_1 \rrbracket(\sigma)))$$

$$\boxed{trans_A : Alias \rightarrow \Sigma \rightarrow \Sigma}$$

$trans_A \llbracket id \rrbracket(\sigma) =$   
**LET**  $id = \text{tensorflow}$  **THEN**  $\sigma["\text{tensorflow}" \mapsto id]$   
**LET**  $id = \text{keras}$  **THEN**  $\sigma["\text{keras}" \mapsto id]$   
**ELSE**  $\sigma$

$trans_A \llbracket id_1 \text{ as } id_2 \rrbracket(\sigma) =$   
**LET**  $id_1 = \text{tensorflow}$  **THEN**  $\sigma["\text{tensorflow}" \mapsto id_2]$   
**LET**  $id_1 = \text{keras}$  **THEN**  $\sigma["\text{tensorflow}" \mapsto id_2]$   
**ELSE**  $\sigma$

$trans_A \llbracket id_1 . id_2 (.id_3)^* (\text{as } id_2)? \rrbracket(\sigma) = \sigma$

|   |
|---|
| $trans_{\overline{W}} : WithItem \text{ list} \rightarrow \Sigma \rightarrow (WithItem \text{ list} \times \Sigma)$ |
|---|

$trans_{\overline{W}} \llbracket with\_item_1 with\_item_2 \dots with\_item_n \rrbracket(\sigma) =$   
**LET**  $with\_item_1', \sigma_1 = trans_W \llbracket with\_item_1 \rrbracket(\sigma)$  **IN**  
**LET**  $with\_item_2', \sigma_2 = trans_W \llbracket with\_item_2 \rrbracket(\sigma_1)$  **IN**  
 $\dots$   
**LET**  $with\_item_n', \sigma_n = trans_W \llbracket with\_item_n \rrbracket(\sigma_{n-1})$  **IN**  
 $(with\_item_1' :: with\_item_2' :: \dots :: [with\_item_n'], \sigma_n)$

|  |
|--|
| $trans_W : WithItem \rightarrow \Sigma \rightarrow (WithItem \times \Sigma)$ |
|--|

$trans_W \llbracket expr \rrbracket(\sigma) = (trans_E \llbracket expr \rrbracket(\sigma), \sigma)$

$trans_W \llbracket expr_1 \text{ as } expr_2 \rrbracket(\text{smodenv}) =$   
**IF**  $\sigma["\text{tensorflow}"] = id_1$  **AND**  $expr_1 = id_1.\text{GradientTape}()$  **AND**  $expr_2 = id_2$  **THEN**  
 $(expr_1 \text{ as } expr_2, \sigma["\text{gradient\_tape}" \mapsto id_2])$   
**ELSE**  $(trans_E \llbracket expr_1 \rrbracket(\sigma) \text{ as } expr_2, \sigma)$

|  |
|--|
| $trans_C : MatchCase \rightarrow \Sigma \rightarrow MatchCase$ |
|--|

$trans_C \llbracket \text{case pattern (if expr)? : stmt}^* \rrbracket(\sigma) =$   
 $\text{case } trans_P \llbracket pattern \rrbracket(\sigma) \text{ (if } trans_E \llbracket expr \rrbracket(\sigma))? : trans_{\overline{S}} \llbracket stmt^* \rrbracket(\sigma).1$

|  |
|--|
| $trans_P : Pattern \rightarrow \Sigma \rightarrow Pattern$ |
|--|

$trans_P \llbracket expr \rrbracket(\sigma) = trans_E \llbracket expr \rrbracket(\sigma)$

$trans_P \llbracket constant \rrbracket(\sigma) = constant$

$trans_P \llbracket [pattern]^* \rrbracket(\sigma) = [trans_P \llbracket pattern \rrbracket(\sigma)]^*$

$trans_P \llbracket *(id)? \rrbracket(\sigma) = *(id)?$

$trans_P \llbracket \{(expr : pattern)^* id?\} \rrbracket(\sigma) = \{(expr : trans_P \llbracket pattern \rrbracket(\sigma))^* id?\}$

$$trans_P \llbracket expr \ (pattern_1^* \ (id = pattern_2)^*) \rrbracket(\sigma) = expr \ (trans_P \llbracket pattern_1 \rrbracket(\sigma)^* \ (id = trans_P \llbracket pattern_2 \rrbracket(\sigma))^*)$$

$$trans_P \llbracket (pattern \ \mathbf{as})? \ id \rrbracket(\sigma) = (trans_P \llbracket pattern \rrbracket(\sigma) \ \mathbf{as})? \ id$$

$$trans_P \llbracket pattern_1 \mid pattern_2 \rrbracket(\sigma) = trans_P \llbracket pattern_1 \rrbracket(\sigma) \mid trans_P \llbracket pattern_2 \rrbracket(\sigma)$$

$$trans_P \llbracket \_ \rrbracket(\sigma) = \_$$